

COMENIUS UNIVERSITY IN BRATISLAVA
FACULTY OF MATHEMATICS, PHYSICS AND INFORMATICS

CYCLIC EDGE-CONNECTIVITY OF CUBIC GRAPHS
MASTER'S THESIS

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COMENIUS UNIVERSITY IN BRATISLAVA
FACULTY OF MATHEMATICS, PHYSICS AND INFORMATICS

CYCLIC EDGE-CONNECTIVITY OF CUBIC GRAPHS
MASTER'S THESIS

Study Programme: Computer Science
Field of Study: Computer Science
Department: Department of Computer Science
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Cyklická hranová súvislosť kubických grafov

Anotácia: Cyklická hranová súvislosť sa ukázala ako dôležitý invariant grafov, obzvlášť tých kubických, nakoľko je previazaná s mnohými vlastnosťami grafov a možno ju použiť v rôznych dôkazoch. Častým prístupom je rozklad väčšieho kubického grafu pozdĺž jeho najmenšieho cyklu separujúceho rezu na dva menšie grafy. Navyše bolo ukázané, že najmenšie protipríklady k mnohým široko otvoreným hypotézam majú vysokú cyklickú hranovú súvislosť popri mnohých iných vlastnostiach, ako absencia hranového 3-farbenia. Práve preto mnohí autori publikovali konštrukcie rôznych kubických grafov bez hranového 3-farbenia s vysokou cyklickou hranovou súvislosťou. Napriek častému používaniu nie sú k dispozícii žiadne ustálené nástroje pre dokazovanie požadovanej cyklickej hranovej súvislosti grafov skonštruovaných týmto prístupom.

Cieľ: Cieľom tejto práce je vyvinúť teoretické nástroje na dokazovanie, že rôzne konštrukcie kubických grafov z menších stavebných blokov majú požadovanú cyklickú hranovú súvislosť. To by malo vychádzať z analýzy publikovaných konštrukcií nekonečných tried kubických grafov a identifikácie kľúčových vlastností, ktoré treba klásť na stavebné bloky. Vyvinutá teória bude demonštrovaná poskytnutím dôkazov požadovanej cyklickej hranovej súvislosti pre príslušné triedy.

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Annotation: Cyclic edge-connectivity proved to be an important invariant of graphs, especially cubic ones, since it is tied to many graph properties and can be used in various proofs. A common approach involves decomposing a larger cubic graph along its smallest cycle-separating edge-cut into two smaller ones. Moreover, it has been shown that the smallest counter-examples to many widely-open conjectures have high cyclic edge-connectivity, besides many other properties like the absence of a 3-edge-colouring. For this reason, several authors provided constructions of various not 3-edge-colourable cubic graphs with high cyclic edge-connectivity. The most common approach is to construct a larger graph by interconnecting several building blocks obtained from smaller graphs. Despite its common usage, there are no established tools for proving the desired cyclic edge-connectivity of graphs obtained through this approach.

Aim: This work aims to develop theoretical tools for proving that various constructions of cubic graphs from smaller building blocks have the desired cyclic edge-connectivity. This should emerge from an analysis of published constructions of infinite families of cubic graphs and identifying key properties that should be imposed on building blocks. The developed theory should be demonstrated by providing proofs of the claimed cyclic edge-connectivity for the respective classes.

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Abstrakt

Cyklická hranová súvislosť je zaujímavá vlastnosť kubických grafov, ktorá sa ukázala ako významná v teórii grafov. Výskum ukázal, že vysoká cyklická hranová súvislosť koreluje s určitými štrukturálnymi vlastnosťami kubických grafov a najmenešie možné protipríklady k viacerým otvoreným problémom musia mať vysokú cyklickú hranovú súvislosť. Teoretických nástrojov na určovanie cyklickej hranovej súvislosti však nie je veľa. V tejto práci skúmame kubické grafy zostrojené z menších stavebných blokov a ukazujeme ako môžeme určiť cyklickú hranovú súvislosť výsledného grafu kladením určitých podmienok na tieto bloky. Naše teoretické výsledky ilustrujeme na praktických príkladoch.

Kľúčové slová: kubické grafy, cyklická hranová súvislosť, multipóly, inflácie

Abstract

The cyclic edge-connectivity is an interesting property of cubic graphs, that has proven to be significant in graph theory. Research has shown, that high cyclic edge-connectivity correlates with certain structural properties of cubic graphs, and notably, smallest potential counterexamples to several open problems need to have high cyclic edge-connectivity. However, theoretical tools for determining cyclic edge-connectivity remain limited. In this work, we study cubic graphs built from smaller building blocks and demonstrate how we can determine the cyclic edge-connectivity of the resulting graph, by imposing certain conditions for these blocks. We also illustrate our theoretical results through practical examples.

Keywords: cubic graphs, cyclic edge-connectivity, multipoles, inflations

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Introduction

In the study of cubic, that is 3-regular graphs, the traditional edge-connectivity faces a natural limitation: it cannot exceed 3, the minimum degree of such graphs. This constraint significantly restricts the exploration of properties of cubic graphs with different values of edge-connectivity.

Cyclic edge-connectivity, on the other hand, offers a wider range of values by imposing additional condition on edge-cuts. Specifically, the cyclic edge-connectivity of a graph G is defined as the size of a smallest cycle-separating edge-cut – an edge-cut whose removal yields two components, each containing a cycle. Unlike traditional edge-connectivity, this property has no fixed upper bound in cubic graphs, enabling the study of graphs with high cyclic edge-connectivity.

Research has shown, that high cyclic edge-connectivity correlates with various structural properties of cubic graphs [1, 4, 21, 30], and several conjectures about high cyclic edge-connectivity have been proposed [7, 14, 31]. Moreover, it has been proven that smallest possible counterexamples to several open problems must have high cyclic edge-connectivity [17, 22].

Our work focuses on cubic graphs constructed from smaller building blocks. To formalize these building blocks, we use the concept of *multipoles*, which can be explained as an extended type of graphs, where each edge has two ends that do not necessarily need to be incident with vertices. These special edges can be connected together to form classical edges between vertices, allowing the construction of larger multipoles and ultimately complete graphs. These connections are formally called *junctions*.

Such constructions of graphs with specific cyclic edge-connectivity remain relatively unexplored. Rajník [29] and Andersen et al. [1] have explored how multipoles with specific properties can be extended to form graphs with desired cyclic edge-connectivity. However, these works focused primarily on the existence of such graphs, rather than determining the cyclic edge-connectivity of graphs built from more than two building blocks.

Our aim is to develop theoretical tools for proving that such constructions of cubic graphs achieve desired cyclic edge-connectivity, when the used multipoles satisfy specific conditions. We demonstrate these results through concrete examples, and provide formal proofs of cyclic edge-connectivity for certain graphs, where such proofs were

previously missing.

In Chapter 1, we introduce the necessary definitions and concepts fundamental to this thesis. Chapter 2 further extends the theory of cyclic edge-connectivity and introduces concepts essential to our results. In Chapter 3, we present our findings on junctions of multipoles, establishing helpful properties of multipoles, and provide theorems that prove these junctions achieve desired cyclic edge-connectivity, along with their practical applications. Chapter 4 introduces the concept of *cyclic part inflations*, where vertices in an original graph are replaced by multipoles satisfying specific conditions. We present several theorems for this concept and demonstrate their applications. Finally, in Chapter 5, we provide potential problems for future research.

Chapter 1

Preliminaries

In this chapter, we establish the fundamental definitions essential for understanding our thesis. We begin with basic graph theory concepts and specify the types of graphs considered in our work, as well as present the definitions of (vertex)-connectivity and edge-connectivity. We then introduce the theoretical framework for our building blocks through the concept of multipoles, along with the formal definitions of connecting these blocks. Following this, we introduce the concept of cyclic edge-connectivity – a key concept further developed in Chapter 2. At the end, we discuss transitive graphs, which prove valuable in proving the existence of multipoles with specific properties required for our theoretical results.

1.1 Basic notions

Definitions not provided in our work can be found in the book “Graph Theory” by Diestel [3]. All graphs considered in this work are undirected and finite. We permit multiple edges, as well as loops.

The *girth* of a graph G , denoted by $g(G)$, is the minimum length of a cycle in G . If G does not contain a cycle, we set the girth to ∞ .

The *distance* between two vertices x and y in a graph G , denoted by $d_G(x, y)$, is defined as the length of a shortest path between x and y in G . If no such path exists, we set $d_G(x, y) = \infty$. If the specific graph G is evident from the context, we shall only use $d(x, y)$.

Although our work focuses on cyclic connectivity, we shall start with a complete definition of connectivity and edge-cuts to establish the necessary background for our main results.

1.2 Edge-cuts and connectivity

As the name suggests, an *edge-cut* decomposes a graph into multiple connected components. It is defined as a set of edges of a graph G , such that if we remove every edge in the set, the resulting graph will be disconnected. Edge-cuts can be created by partitioning vertices into two components and selecting the edges connecting them.

Definition 1. Let G be a graph and $\{V_1, V_2\}$ a partition of $V(G)$. By $E(V_1, V_2)$ we denote the set of edges which have one end in V_1 and the other in V_2 .

Note that $E(V_1, V_2)$ is indeed an edge-cut, when we remove all of these edges, there will not be any edge between a vertex in V_1 and V_2 , splitting the graph into two components. The sets V_1 and V_2 are called the *sides* of the cut. This can help visualize how vertices will be partitioned after the removal of the edges.

In our work, however, we do not remove these edges, but instead introduce the concept of edge severing, which is formally defined in Section 1.4.

A graph which results from a graph G by removing edges from a set S is denoted by $G - S$. Note that this does not require S to be an edge-cut, the result can still be connected.

A graph G is *k-edge-connected* for $k \in \mathbb{N}$, if $|G| > k$ and $G - S$ is connected for every subset S of $E(G)$, such that $|S| < k$. In other words, there is no edge-cut of size less than k .

Definition 2. Let G be a graph. The *edge-connectivity* of G , denoted by $\lambda(G)$, is the largest integer k for which G is k -edge-connected.

Informally, this is the size of a smallest edge-cut in the graph. It holds that for each non-trivial graph G , we have $\lambda(G) \leq \delta(G)$ [3]. The proof comes from the fact that all edges incident with a vertex are indeed edge-cuts, separating the vertex from the rest of the graph, meaning there is an edge-cut of size $\delta(G)$ in each graph. Also it holds that if G is disconnected then $\lambda(G) = 0$.

1.3 Vertex-connectivity

In contrast to edge-connectivity, as the name suggests, vertex-connectivity focuses on the removal of vertices rather than edges. A *vertex-cut* is therefore defined as a subset of vertices, whose removal results in a disconnected graph. A graph G is *k-(vertex)-connected* for $k \in \mathbb{N}$, if $|G| > k$ and $G - X$ is connected for every subset X of $V(G)$ with $|X| < k$. In other words, there is no vertex-cut in G of size less than k .

Similarly to edge-connectivity, the *(vertex)-connectivity* of G , denoted by $\kappa(G)$, is the largest integer k for which G is k -connected.

Note that an argument similar to the one used for edge-connectivity applies to the upper bound of $\kappa(G)$ being $\delta(G)$. We can select a vertex with minimal degree, and its neighbors form a vertex-cut (excluding loops to avoid removing the selected vertex). In complete graphs, this selection does not form a vertex-cut, since only the selected vertex would remain, but the relationship still holds: the connectivity of K_n is $n - 1$ by definition, and $\delta(K_n) = n - 1$. For any nontrivial graph G , these parameters are related by the inequality $\kappa(G) \leq \lambda(G) \leq \delta(G)$ [3].

1.4 Multipoles

In our work, we allow an extended type of graph, which allows ends of its edges to be incident with no vertex, resulting in a graph with semiedges. Such structures are called *multipoles*. These structures can serve as the mentioned building blocks to form larger graphs, as explained in this section. This term was first used by Fiol in 1991 [6], however we will be following a definition by Nedela and Škoviera [26].

Definition 3. A *multipole* is a pair $M = (V, E)$ of distinct finite sets of vertices V and edges E , where every edge e in E has two edge ends, which may or need not be incident with a vertex.

According to the incidence of edge ends with vertices, we define three types of edges:

1. a *link* is an edge whose ends are incident with two distinct vertices,
2. a *loop* is an edge whose ends are incident with the same vertex,
3. a *semiedge* is an edge whose one end is incident with a vertex, and the other is not.

Note that we have not mentioned the case when both edge ends are not incident with a vertex. We do not allow such edges in our work, however in literature these are referred to as *isolated edges*.

A multipole is *connected* if there exists a path between any pair of edge ends. A multipole M is *empty* if $|V(M)| = 0$, otherwise it is *nonempty*.

The semiedges in a multipole M are endowed with a linear order and we denote the tuple of its semiedges by $S(M)$. A multipole M with n semiedges and $S(M) = (a_1, \dots, a_n)$ can also be denoted by $M(a_1, \dots, a_n)$. Similarly to graphs, the *order* of a multipole M , denoted by $|M|$, is the number of its vertices. Also the *degree* of a vertex v of a multipole, denoted by $\deg(v)$, is the number of edge ends incident with v . The minimum degree among all vertices in a multipole M is denoted by $\delta(M)$. A multipole with k semiedges is usually called a *k-pole*. It can be seen, that a graph can be defined as a 0-pole.

Now we define how multipoles can be used as building blocks to form larger structures. Let e and f be semiedges, such that $e \neq f$. The result of the *junction* of e and f is a new edge between the vertices incident with e and f , and a deletion of e and

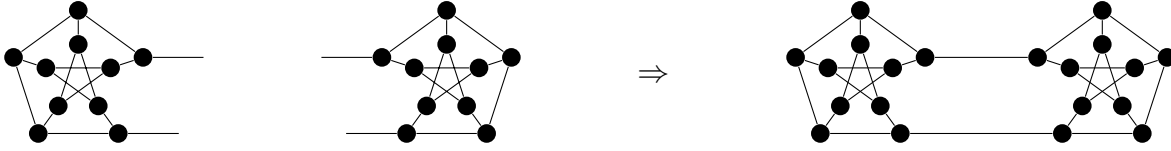


Figure 1.1: Junction of two multipoles

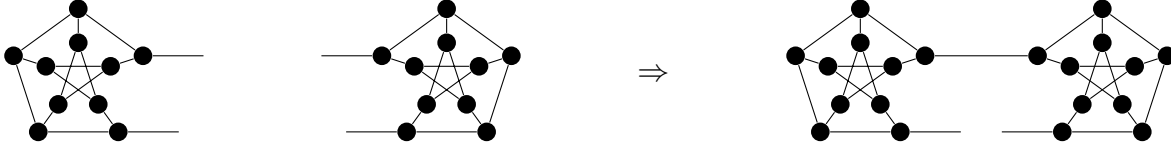


Figure 1.2: Partial junction of two multipoles

f . The junction of two k -poles $M(a_1, \dots, a_k)$ and $N(b_1, \dots, b_k)$ using a bijection $f : \{1, \dots, k\} \rightarrow \{1, \dots, k\}$, denoted by $M *_{f} N$, consists of k individual junctions of semiedges a_i and $b_{f(i)}$ for i from 1 to k . If we use an arbitrary bijection we simply denote it by $M * N$. An example of a junction can be seen in Figure 1.1.

Consider two multipoles $M(a_1, \dots, a_n)$ and $N(b_1, \dots, b_m)$. Their *partial junction of size k* is a junction of some semiedges $(a_{i_1}, \dots, a_{i_k})$ and $(b_{j_1}, \dots, b_{j_k})$, where $k \leq n$ and $k \leq m$. Note that the vertices of the resulting multipole will be $V(M) \cup V(N)$, for clarity we can require the multipoles to have distinct vertices. In contrast to a normal junction, which results in a graph, the partial junction can still result in a multipole. However, we allow in partial junction the junction of all semiedges, thus a junction is a special case of partial junction where the result is a graph. An example of a partial junction can be seen in Figure 1.2, note that the result is still a multipole with some semiedges and not a graph as in the example before.

We say that a multipole M *contains* a multipole N (or N is in M), if there exists a multipole P and a partial junction $*$, such that $M = N * P$. Note that each multipole contains itself, as this can be achieved using an empty partial junction with an empty multipole.

Same as in graphs, a multipole is *cyclic* if it contains a cycle. The *distance* between two semiedges a and b is equal to:

- 0 if $a = b$,
- $d(x, y)$ where x and y are the vertices with which are incident the semiedges a and b respectively.

By *cutting* (or *severing*) a link $e = v_1v_2$, we mean the removal of e , along with

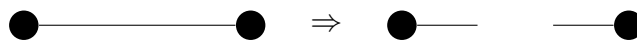
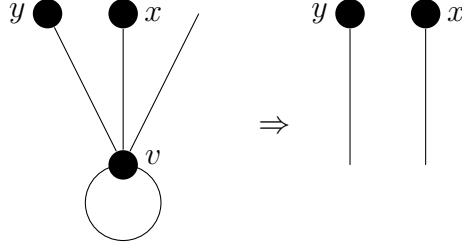


Figure 1.3: Severing an edge

Figure 1.4: Removing a vertex v

adding a semiedge to v_1 and v_2 , as depicted in Figure 1.3.

When we *remove* a vertex v , we delete v along with all its incident semiedges and loops. Additionally, for each link connecting v to another vertex w , we replace that link with a new semiedge incident with w , as depicted in Figure 1.4.

In multipoles, as in graphs, an *edge-cut* is a set of links that, when severed, disconnects the multipole. We also use edge-cuts defined as the set of edges $E(V_1, V_2)$, which have one end in V_1 and the other in V_2 , where $\{V_1, V_2\}$ is a partition of $V(M)$. Note that all of the edges in these edge-cuts are links, since they connect two distinct vertices.

Two edges are *adjacent* if they share a vertex. An edge-cut is *independent* if no two edges in it are adjacent.

We now define edge subdivisions in multipoles, following Gross et al. [11].

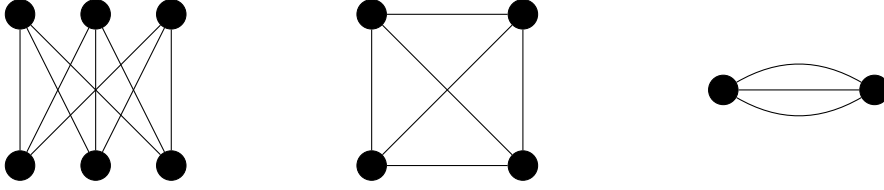
Definition 4. Let e be an edge in a multipole M . *Subdividing* the edge e means that a new vertex w is added to $V(G)$, and the edge e is replaced according to its type:

- a loop on vertex v is replaced by two edges between v and w ,
- a link between vertices x and y is replaced by edges xw and wy ,
- a semiedge incident to vertex v is replaced by an edge vw and two semiedges incident to w .

Conversely, we can define *smoothing* in a multipole.

Definition 5. Let M be a multipole and let v be a vertex of degree 2 in M . By smoothing v in M we obtain a new multipole M' , such that $V(M') = V(M) - \{v\}$ and the edge set is updated according to edges incident with v :

- if v is incident with two semiedges, we remove both from the edge set,
- if v is incident with one semiedge and the other edge is a link vx , we remove both edges and add a new semiedge from x ,
- if v is adjacent to two different vertices x and y , we remove both edges and add a link xy ,

Figure 1.5: Graphs $K_{3,3}$, K_4 and Θ_2

- if v is connected by two edges to a vertex x , we remove both edges and add a loop at x .

Finally, we define when a multipole is induced by a vertex subset of another multipole.

Definition 6. Let M be a multipole and W a subset of its vertices. Then $M[W]$ is an *induced multipole* of M by W , if $V(M[W]) = W$ and there exists a multipole N , such that $M = M[W] * N$.

In other words, if W is empty then $M[W]$ is an empty multipole, if $W = V(M)$ then $M[W] = M$, and otherwise $M[W]$ is the multipole containing the vertices from W , obtained from the edge-cut $E(W, V(M) - W)$ in M .

1.5 Cyclic edge-connectivity

In cubic graphs, the upper bound of the edge-connectivity is 3, which can be a huge limitation in exploring graphs with different connectivity values. Because of this, a new measure regarding edge-cuts was introduced, requiring both of the components resulting after severing edges from an edge-cut to contain cycles. The notation is similar as before, just with a keyword *cyclic* before it. To maintain a reasonable level of formalism, we shall define these properties more precise.

An edge-cut $E(V_1, V_2)$ in a graph G is called *cycle-separating*, if both induced multipoles $G[V_1]$ and $G[V_2]$ are cyclic. In contrast to edge-cuts, there are non-trivial graphs which do not contain a cycle-separating edge-cut [18, 19, 25], for example trees. However, in cubic graphs, there are only three of such graphs: $K_{3,3}$, K_4 , Θ_2 [25] depicted in Figure 1.5.

Similarly to edge-connectivity, we say that a graph G is *cyclically k -edge-connected*, if there is no cycle-separating edge-cut with fewer than k edges.

The *cycle rank* of G , denoted by $\beta(G)$, equals $|E| - |V| + 1$. If we remove at least $\beta(G)$ edges from G , the resulting graph is either disconnected or acyclic. Therefore, if G contains a cycle-separating edge-cut, then it contains one with no more than $\beta(G) - 1$ edges [25].

Definition 7 (Nedela and Škoviera [25]). Let G be a graph. The *cyclic edge-connectivity* of G , denoted by $\zeta(G)$, is the largest integer k , such that $k \leq \beta(G)$ and G is cyclically k -edge-connected.

This means, that if a graph G does not contain a cycle-separating edge-cut, its cyclic edge-connectivity is equal to $\beta(G)$. For the three cubic graphs without cycle-separating edge-cuts mentioned above, we have $\zeta(K_{3,3}) = 4$, $\zeta(K_4) = 3$ and $\zeta(\Theta_2) = 2$.

More properties of cyclic edge-connectivity are regarded in Chapter 2.

1.6 Transitive graphs

In proving the results of our work, we have used *transitive graphs*, for which there are some proved theorems and can be helpful with their properties for our results. To properly define them, we need to dive a bit into an algebraic graph theory and prove what it means for a graph to be transitive. All of these definitions were taken from a book by Godsil and Royle [8].

An *automorphism* on a graph G is an isomorphism from G to itself, in other words a permutation σ of the vertices of G , such that $vw \in E(G)$ if and only if $\sigma(v)\sigma(w) \in E(G)$. In other words, it is a permutation that maps edges to edges and nonedges to nonedges. A trivial case of an automorphism is an identity permutation.

A graph G is called *vertex-transitive*, if for any two vertices x and y in $V(G)$ there is an automorphism of G mapping x to y . Similarly, it is called *edge-transitive*, if for any two edges e and f there is an automorphism of G that maps e to f .

An important property of transitive graphs is that their cyclic edge-connectivity coincides with their girth. This result proves particularly useful in our work, specifically in establishing the existence of multipoles with desired properties.

Theorem 1 (Nedela and Škoviera [25, Theorem 17]). *Let G be a cubic vertex-transitive or edge-transitive graph of girth g . Then $\zeta(G) = g$.*

For now we shall mention the following theorem about the existence of k -regular vertex-transitive graph with arbitrary girth and k , which we will use to construct cubic vertex-transitive graphs.

Theorem 2 (Nedela and Škoviera [27]). *For every pair of integers $k \geq 3$ and $g \geq 3$ there exists a k -regular vertex-transitive graph of girth g .*

Thus there is a cubic vertex-transitive graph of arbitrary girth, implying arbitrarily large cyclic edge-connectivity.

Chapter 2

Cyclic edge-connectivity

From this point forward, unless explicitly stated otherwise, all discussed graphs and multipoles will be cyclic. While fundamental concepts of cyclic edge-connectivity were introduced in Section 1.5, this chapter extends these foundations and presents important theoretical developments for our thesis. We first examine cyclic vertex-connectivity and show how it relates to cyclic edge-connectivity in cubic graphs. We then discuss important theorems, conjectures and other concepts essential for our work. Following this, we introduce a fundamental concept of cyclic parts and conclude this chapter with a concept of inflations, which is an interesting method for constructing larger graphs.

As mentioned, we can also examine cyclic connectivity in graphs from an alternative perspective by using vertex-cuts. A vertex set X of a graph G is called a *cycle-separating* vertex-cut, if G can be expressed as a union of two edge-disjoint cyclic subgraphs G_1 and G_2 , such that $V(G_1) \cap V(G_2) = X$. We say that a graph is *cyclically k -vertex-connected*, if it contains no cycle-separating vertex-cut with fewer than k vertices. Analogous to cyclic edge-connectivity, the *cyclic vertex-connectivity* of a graph G is defined as the largest integer k , such that $k \leq \beta(G)$ and G is cyclically k -vertex-connected.

While this property might appear to offer a distinct perspective for analyzing cubic graphs, the following theorem demonstrates that cyclic vertex-connectivity and cyclic edge-connectivity are actually equivalent properties in cubic graphs.

Theorem 3 (Nedela and Škoviera [25]). *Let G be a connected cubic graph not isomorphic to $K_{3,3}$, K_4 or Θ_2 . Then the following conditions are equivalent:*

- (i) *G is cyclically k -edge-connected,*
- (ii) *G is cyclically k -vertex-connected,*
- (iii) *each independent edge-cut has at least k edges.*

Consequently, any theorem concerning cyclic edge-connectivity in cubic graphs automatically yields corresponding results about cyclic vertex-connectivity and vice-versa.

Note the equivalence with the statement about independent edge-cuts. It holds, that in cubic graphs each independent edge-cut is cycle-separating [25].

Cyclic connectivity is interesting to study, due to numerous important results concerning structural properties of graphs with high cyclic connectivity. For instance:

Theorem 4 (McCuaig [21]). *Let F be a set of k edges in a cyclically $(k + 1)$ -connected cubic graph G . Suppose that for every pair of vertices u and v which are incident with distinct edges in F , $k - 2 \leq d(u, v)$. Then the edges in F lie on a cycle.*

Proposition 1 (Thomassen [30]). *If G has minimum degree at least 3 and is cyclically k -vertex-connected for $k > 3$, and A, B are disjoint vertex sets of G each of cardinality at least $2k - 3$, then G contains a collection of k disjoint paths from A to B .*

Theorem 5 (Thomassen [30]). *If A is a set of k independent edges in a cyclically 2^{k+1} -vertex-connected graph G of minimum degree at least 3, then G has a cycle through A .*

Furthermore, several conjectures involve graphs with high cyclic edge-connectivity:

Conjecture 1 (Thomassen [31]). *Every cyclically 8-edge-connected cubic graph is Hamiltonian.*

Conjecture 2 (Fleischner and Jackson [7]). *Every cyclically 4-edge-connected cubic graph G has a circuit C , such that $G - V(C)$ is acyclic.*

Conjecture 3 (Fleischner and Jackson [7]). *Every cyclically 4-edge connected cubic graph G has a cycle C , such that $G - V(C)$ is an independent set of vertices.*

Conjecture 4 (Jaeger [14]). *Every cyclically 7-edge-connected cubic graph is 3-edge-colourable.*

Moreover, it has been proven that smallest possible counterexamples to several open problems must have high cyclic edge-connectivity [17, 22]. These results suggest that studying cyclic connectivity is not only interesting in the area of exploring structural properties of graphs, but also gives us valuable insights for exploring open questions in graph theory.

We examine a result by Nedela and Škoviera regarding the upper bound of cyclic edge-connectivity in cubic graphs.

Proposition 2 (Nedela and Škoviera [25, Proposition 2]). *For every connected cubic graph G it holds that $\zeta(G) \leq g(G)$.*

Notably, this presents an interesting contrast with edge-connectivity: while regular edge-connectivity in cubic graphs is bounded from above by 3, this upper bound of cyclic edge-connectivity can achieve arbitrarily large values, even within the restricted class of cubic graphs.

2.1 Cyclic parts

Cyclic parts can be described as cyclic multipoles that result from severing a smallest cycle-separating edge-cut in a graph. The formal definition follows:

Definition 8 (Nedela and Škoviera [25]). Let G be a graph and let X be a minimum cycle-separating edge-cut in G . Then the multipoles which arise after cutting X in G are called *cyclic parts* of G . More specifically, if $|X| = k$, then we can refer to them as cyclic k -parts.

This concept serves as a valuable tool for studying connectivity properties, as they help us understand both the structure of edge-cuts and the resulting multipoles that emerge after their severing. Furthermore, in Chapter 4, we develop a method for constructing graphs with desired cyclic edge-connectivity through junctions of cyclic parts. This makes it particularly important to determine whether a given cyclic multipole is as a cyclic part.

According to the definition, determining whether a multipole M is a cyclic part is equivalent to determining whether there exists a graph G with a minimum cycle-separating cut X , such that one of the components of $G - X$ is isomorphic to M .

Researchers have explored methods for completing cyclic k -parts into cyclically k -connected graphs. Andersen et al. [1] demonstrated, that for $k = 4$ this can be achieved through a junction with a 4-pole consisting of two adjacent vertices, each with two semiedges. Máčajová and Rajník [23] extended this research to $k = 5$.

Theorem 6 (Máčajová and Rajník [23, Theorem 9]). *Let H be a cyclic 5-part. If H is not a cycle of length 5, then H can be extended to a cyclically 5-connected cubic graph by performing a junction with a specific 5-pole, which consists of three vertices arranged in a path of length two, with additional semiedges attached to form a cubic multipole.*

A cyclic part of G minimal under inclusion is called an *atom*. Two atoms of a single graph may differ in the number of vertices [25], those with the minimum number of vertices are called *proper* atoms.

Let C_k denote the k -pole consisting of the k -cycle where each vertex is additionally incident with one semiedge. A cyclic part, which is a single cycle C_k is called *trivial*, otherwise it is *nontrivial*. Several estimations regarding a lower bound of the order of these cyclic parts were proven.

Lemma 1 (Rajník [29, Lemma 5.1]). *Let M be an s -pole with girth at least k different from a k -cycle for some $s \leq k$. Then $|M| \geq 2k - 4$.*

Theorem 7 (Nedela and Škoviera [25, Theorem 7]). *Let P be a nontrivial cyclic k -part. Then either $|P| \geq 2k - 3$, or $k = 6$ and $|P| = 2k - 4 = 8$. Moreover, if P is an atom, then $|P| \geq 2k$, or $k = 3$ and $|P| = 2k - 1 = 5$.*

Cyclic parts can be used as potential building blocks for building graphs with desired cyclic connectivity, as proved in the following theorem.

Theorem 8 (Rajník [29, Theorem 5.2]). *Let M_1 and M_2 be cyclic k -parts, such that both M_1 and M_2 are different from C_k or $k \leq 5$. Let G be a junction of M_1 and M_2 . Then G is cyclically k -connected unless the girth of G is smaller than k .*

An interesting area of cyclic parts studied by Rajník [29] are so-called adjuncts. A k -pole A_k is called a k -adjunct, if for each cyclic k -part M the graph $M *_f A_k$ is cyclically k -connected for a bijection f . If $M *_f A_k$ is cyclically k -connected for every bijection f , then we call A_k a *universal k -adjunct*.

We use this concept for proving an important theorem regarding cyclic parts, however, we do not study it in depth, just find a construction for creating universal k -adjunct for an arbitrary k .

Rajník [29] also identified a problem in creating an algorithm to determine whether a cyclic multipole is a cyclic part. Currently, for arbitrary k , this process requires finding a universal k -adjunct, performing a junction with the k -pole, and applying existing algorithms, such as that of Dvořák et al. [5], to verify whether the result has cyclic edge-connectivity equal to k .

Problem 1 (Rajník [29, Problem 2.6]). Find an algorithm that for a given k -pole M decides if M is a cyclic k -part or not.

In Section 3.2 we present an alternative approach to solve this problem by examining the cyclic multipole directly. Our method involves constructing an infinite family of universal adjuncts with specific properties from transitive graphs, using which we can solve this problem by exploring whether the cyclic multipole contains another cyclic multipole with a small number of semiedges or not. While this approach does not provide an explicit algorithm, it offers a different perspective on the problem

2.2 Junctions of multipoles

Previous research in this area has primarily focused on specific cases. For instance, the junction of cyclic k -parts with adjuncts, where $k \in \{4, 5\}$, has been studied by several researchers [1, 29], as previously mentioned.

In the study of snarks, that is cubic graphs which are not 3-edge-colourable, Kochol introduced the concept of *superposition*. This technique involves replacing edges with multipoles called superedges and vertices with supervertices. He denoted such superpositions of a graph G by $S(G)$ and used this method to construct larger cyclically 5-connected snarks from smaller cyclically 5-connected graphs:

Theorem 9 (Kochol [16, Theorem 1]). *Let G be a cyclically 5-edge-connected cubic graph of order n and with girth c equal to 5. Then there exists a cyclically 5-edge-connected snark $S(G)$ of order $22n + 14$ and with girth c .*

Several other specific constructions have been explored in snarks. Goedgebeur, Máčajová, and Škoviera [9, 10] investigated constructions of snarks with cyclic connectivity equal to 4 up to order 44, while Rajník [28] generated all snarks with girth 6 and cyclic connectivity equal to 4 and 5 up to order 40.

Using dot product of graphs, which is similar to junction of multipoles, Chladný and Škoviera [2] proved that each irreducible snark, that is snark where the removal of every nontrivial edge-cut yields 3-edge-colourable components, can be decomposed into cyclically 5-connected irreducible snarks:

Theorem 10 (Chladný and Škoviera [2, Theorem C]). *Every irreducible snark G different from the dumbbell graph can be decomposed into a collection $\{H_1, \dots, H_n\}$ of cyclically 5-connected irreducible snarks, such that G can be reconstructed from them by repeated dot products.*

Our approach differs from these previous works by developing a more general theoretical framework. Rather than focusing on specific values of k in cyclically k -connected graphs, we aim to create a versatile framework applicable to arbitrary values of k .

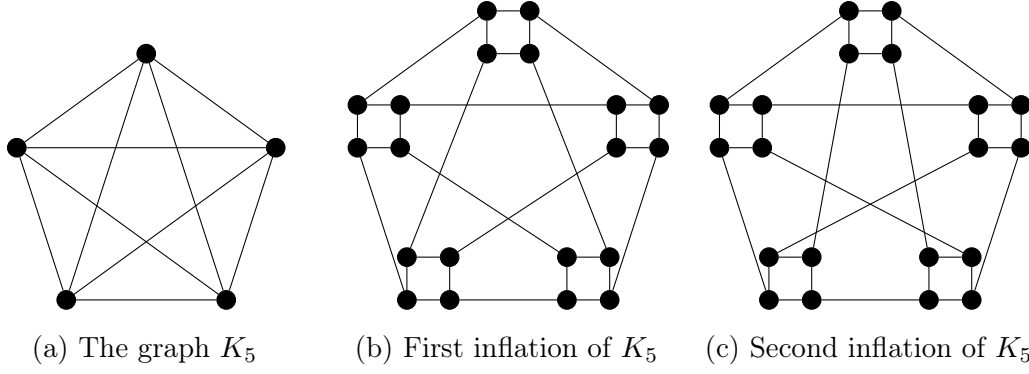
2.3 Inflation

The concept of inflation informally represents a graph transformation where each vertex is replaced by a cycle whose length corresponds to the degree of the vertex. Note that in this concept we consider arbitrary graphs, meaning each vertex can have any degree, however the resulting inflation is always a cubic graph, regardless of the original vertex degrees.

To formalize this concept, we approach the definition from a different perspective: describing how to obtain the original graph from its inflation rather than the transformation process itself. This requires us to first establish the definition of vertex contraction.

Let G be a graph and V a subset of its vertices. The *contraction* of V in G yields a new graph G' where:

- The vertex set is defined as $V(G') = (V(G) - V) \cup \{v\}$, where v represents a new vertex.
- For the edge set, we define:

Figure 2.1: K_5 and two its distinct inflations

- $E = \{xy \mid xy \in E(G), x \in V, y \notin V\}$ (edges connecting V to other vertices),
- $E' = \{xy \mid xy \in E(G), x \in V, y \in V\}$ (edges within V),
- $E'' = \{vy \mid xy \in E, x \in V\}$ (new edges incident to v).

The edge set of G' is then defined as $(E(G) - (E' \cup E)) \cup E''$.

This can be understood as taking the set of vertices and merging them into one while keeping all of the edges, and removing loops which arose after merging.

A k -factor of a graph G is a k -regular subgraph G' of G , such that $V(G') = V(G)$, thus spanning all vertices of G . Specifically, a 2-factor is a subgraph spanning all vertices of the original graph that is a union of disjoint cycles. A *chord* of a cycle is an edge connecting two vertices of a cycle that is not itself an edge of the cycle. A *chordless* cycle of a graph is a cycle containing no chords.

We can now formally define inflations:

Definition 9 (Hoffmann-Ostenhof et al. [12, Definition 10]). Let G be a graph and let H be a cubic graph. Then H is an *inflation* of G if H contains a 2-factor F consisting of chordless cycles, such that the graph obtained from H by contracting each cycle of F to a vertex is isomorphic to G . By $I(G)$ we denote the set of all inflations of G .

It is important to note that a single graph can have multiple distinct inflations, as the connections between cycles can vary. This multiplicity is illustrated in Figure 2.1, which shows two different cyclic part inflations of the same graph.

The following theorem is particularly relevant to our work, as it establishes conditions for cyclic edge-connectivity in inflations:

Theorem 11 (Hoffmann-Ostenhof et al. [12, Theorem 11]). *Let $k \geq 3$ and let G be an arbitrary k -connected graph with girth at least k . Then every $H \in I(G)$ is cyclically k -edge-connected.*

Our research extends this understanding of cyclic edge-connectivity in inflations by developing additional theorems and introducing a novel approach, where vertices are replaced by arbitrary cyclic parts rather than just cycles. These developments are explored in detail in Chapter 4.

Chapter 3

Junctions of multipoles

This chapter investigates junctions of multipoles and identifies conditions under which they yield graphs with desired cyclic edge-connectivity.

First, in Section 3.1 we introduce a new property for multipoles that determines which ones are potentially suitable for constructing graphs with desired cyclic connectivity. In Section 3.2, we develop methods to determine whether a cyclic multipole is a cyclic part without needing to find the original graph from which this cyclic multipole could be obtained as a cyclic part. Moreover, we establish conditions for pairs of multipoles under which their junction has desired cyclic edge-connectivity. Section 3.3 examines graphs that emerge from junctions of multiple multipoles. We demonstrate that when building graphs iteratively from smaller multipoles, the final structure has desired cyclic edge-connectivity if each intermediate step satisfies the conditions of our theoretical framework. To demonstrate practical applications, we conclude with specific examples. Furthermore, we analyze graphs from publications [15, 24] where cyclic edge-connectivity was claimed without formal proof, and provide formal proofs validating these claims using our theoretical framework.

3.1 Potential cyclic connectivity

Now, we establish the necessary conditions for multipoles that can be used to construct graphs with desired cyclic edge-connectivity. Two key restrictions emerge: first, the multipole cannot contain any independent edge-cut yielding a cyclic multipole with a small number of semiedges, as this could imply the existence of a small independent edge-cut in the resulting graph as well, contradicting the desired cyclic edge-connectivity. Second, if the multipole itself is cyclic and not a graph, meaning some partial junction must still be performed to obtain a graph, it cannot have a small number of semiedges. The reason is that the resulting graph would contain this cyclic multipole with small number of semiedges, again possibly contradicting the desired cyclic edge-connectivity.

For this reason we introduce a new property of multipoles, which proves to be fundamental in our subsequent theorems.

Definition 10. Let M be a multipole. We say that M is *potentially cyclically k -connected*, if it has at least k semiedges when cyclic and not a graph, and there exists no independent edge-cut whose severing yields a cyclic multipole with fewer than k semiedges.

The term *potentially* is used because these multipoles can serve as building blocks for constructing graphs with desired cyclic edge-connectivity, provided they meet certain additional requirements. This is visible in the subsequent theorems concerning graph construction from such multipoles.

Since a graph can be viewed as a special case of a multipole, we can examine the relationship between their property of being cyclically k -connected and potentially cyclically k -connected. As the following lemma demonstrates, these properties are, in fact, equivalent.

Lemma 2. *A graph G is cyclically k -connected for an integer k if and only if it is potentially cyclically k -connected.*

Proof. We prove both implications indirectly. First, assume that G is not potentially cyclically k -connected. This means there exists an independent edge-cut yielding a cyclic component with fewer than k semiedges. Since the edge-cut is independent, it is cycle-separating and of size less than k , meaning G is not cyclically k -connected.

Conversely, assume that G is not cyclically k -connected. This means there exists an independent edge-cut of size $l < k$, separating G into two cyclic l -poles. Therefore, G is not potentially cyclically k -connected. \square

The next lemma demonstrates that in a cyclic multipole, the cut separating another cyclic multipole with minimal number of semiedges must be independent. The idea of the proof comes from the observation that if such a cut was not independent, it could be modified to create a new cut yielding a cyclic multipole with fewer semiedges, thus contradicting our assumption about the minimality of the number of semiedges.

Lemma 3. *Let M be a cyclic k -pole containing a cyclic multipole with fewer than k semiedges. Among all such multipoles, let X be one with the minimum number l of semiedges. Then the edge-cut $E(V(X), V(M) - V(X))$ yielding X is independent.*

Proof. Suppose to the contrary that the edge-cut is not independent, meaning there exists a vertex v incident with two edges e_1 and e_2 that are in the cut. If e_3 is as well in the cut, we remove all three from the cut. Otherwise, modify this edge-cut by removing e_1 and e_2 and adding the third edge e_3 incident with v (unless e_3 is a semiedge, then just remove e_1 and e_2).

These modifications ensure the multipole stays cyclic. Since any vertex incident with two or more edges in the cut cannot be part of a cycle after their severing, modifying the cut, resulting either in adding the vertex into X or removing it from X , preserves the cyclic property.

Also, we do not increase the number of semiedges in either multipole, since we remove two edges from the cut and either add one edge to the cut, or add one semiedge to one of the components.

Therefore, this modification yields a cyclic multipole with fewer than l semiedges, contradicting the minimality of l . \square

Using this lemma, we now prove an equivalence lemma for multipoles with at least one semiedge, where we show that the property of being potentially cyclically k -connected is equivalent to them not containing a cyclic multipole with a small number of semiedges.

Lemma 4. *A multipole M with at least one semiedge is potentially cyclically k -connected if and only if there exists no cyclic multipole in M with fewer than k semiedges.*

Proof. We prove both implications indirectly. First, assume that M contains a cyclic l -pole X , where $l < k$ and l is minimal. By Lemma 3, the edge-cut yielding X is independent and therefore M is not potentially cyclically k -connected.

For the converse implication, assume that M is not potentially cyclically k -connected. This can mean two things:

- M is cyclic and has fewer than k semiedges. Then M contains a cyclic multipole with fewer than k semiedges – itself, therefore is not potentially cyclically k -connected.
- There is an independent edge-cut in M , which yields a cyclic multipole X with fewer than k semiedges. Then M contains X , which is cyclic and has fewer than k semiedges, therefore is not potentially cyclically k -connected.

In both cases we have obtained the fact that M is not potentially cyclically k -connected, completing the proof. \square

This lemma is widely used in our proofs, where we come to contradictions with the fact that a multipole is potentially cyclically k -connected, by coming to a case that it contains a cyclic l -pole, such that $l < k$. Given its frequent application, we will use this result implicitly without explicit reference.

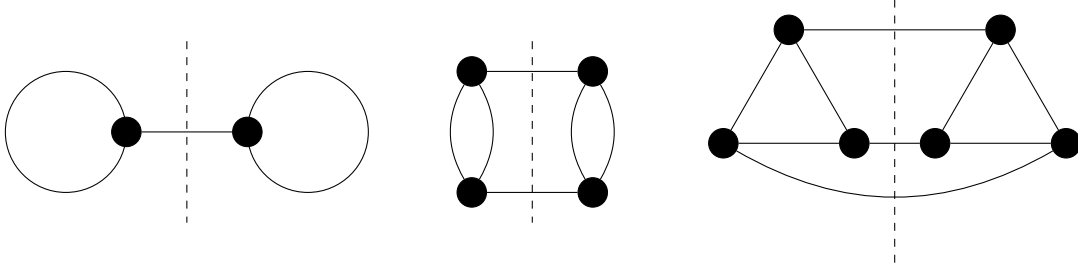


Figure 3.1: Junctions (marked by the dashed lines) $C_1 * C_1$, $C_2 * C_2$ and $C_3 * C_3$

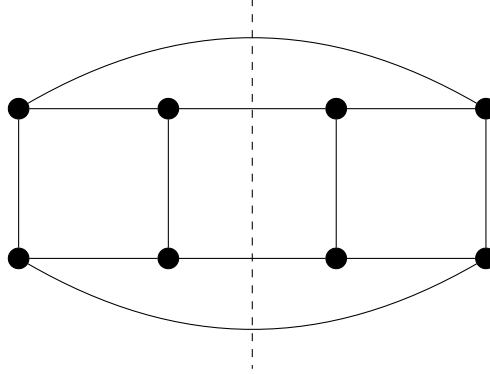


Figure 3.2: Junction (marked by the dashed line) $C_4 * C_4$

3.2 Cyclic parts

In this section, we further explore the concept of cyclic parts. Our main goal, as discussed earlier, is to offer a different perspective on determining whether a cyclic multipole is a cyclic part, specifically by examining only the properties of the given multipole. First, we prove that all k -cycles are cyclic k -parts.

Lemma 5. *The k -pole C_k is a cyclic k -part for every integer $k \geq 1$.*

Proof. First, we tackle the small cases of C_1, C_2, C_3 and C_4 . As shown in Figures 3.1 and 3.2, for each $i \in \{1, 2, 3, 4\}$, the junction $C_i * C_i$ yields a graph whose cyclic edge-connectivity equals i , as can be verified.

Now assume $k > 4$. For a given k , we can construct a cubic vertex-transitive graph with girth k according to Theorem 2. Let us denote this graph by G_k . Then $\zeta(G_k) = k$ by Theorem 1. Since the girth of G_k is k , it contains a cycle of length k and thus $G_k = C_k * P$ for some other k -pole P .

Let P' be a modification of P , such that all semiedges are removed, thus the result is a graph. We prove that $\delta(P') \geq 2$, meaning P' is cyclic, and this implies that P is cyclic as well, since by adding more edges to a cyclic multipole it remains cyclic.

Suppose $\delta(P') < 2$, this would mean that there is a vertex in P incident with more than one semiedge. Let us denote two of these semiedges by s_1 and s_2 (note that there can be three semiedges, it is however sufficient to explore two of them). In the process

of the junction $C_k * P$, separate junctions are performed between the semiedges s_1 and s_2 , and some semiedges in C_k , say t_1 and t_2 . Let X be the shortest path between t_1 and t_2 in C_k . We see that $|X| \leq \frac{k}{2}$, since C_k is a k -cycle. Let e_1 and e_2 be the links in G_k resulting from the junction of C_k and P , such that e_i resulted from a junction of semiedges s_i and t_i , for both $i \in \{1, 2\}$. Then there is a cycle $e_1 X e_2$ in G_k of length at most $\frac{k}{2} + 2$. The girth of G_k is k , therefore $k \leq \frac{k}{2} + 2$. From this we obtain $k \leq 4$, contradicting our assumption that $k > 4$.

Thus P contains a cycle, so the edge-cut $E(V(C_k), V(P))$ in G_k is cycle-separating. Since both P and C_k are k -poles, $|E(V(C_k), V(P))| = k$. Furthermore, since $\zeta(G_k) = k$, this implies that C_k is a cyclic k -part. \square

Now we present a helpful lemma, which formally establishes that a cyclic k -part is potentially cyclically k -connected. The intuition behind this result is straightforward: since a cyclic k -part originates from a graph with cyclic edge-connectivity equal to k , the cyclic part cannot contain a cyclic multipole with fewer than k semiedges, as this would contradict the cyclic edge-connectivity of the graph.

Lemma 6. *Let P be a cyclic k -part. Then P is potentially cyclically k -connected.*

Proof. Suppose to the contrary that P is not potentially cyclically k -connected, therefore it contains a cyclic l -pole Q where $l < k$. Since P is a cyclic k -part, there exists a graph G with $\zeta(G) = k$ formed by the junction of P and a cyclic k -part P' . Also, since Q is in P , there exists an l -pole Q' , such that $G = Q * Q'$. However, since Q is in P , it holds that P' is in Q' , which means that Q' is cyclic, and the edge-cut in G yielding Q and Q' is cycle-separating and of size less than k , which contradicts the fact that $\zeta(G) = k$. \square

Following is a lemma that determines the upper bound on the number of vertices in an acyclic multipole.

Lemma 7. *Let M be an acyclic k -pole for an integer k . Then $|M| \leq k - 2$.*

Proof. For each semiedge of M , perform a junction with a multipole consisting of a single vertex with one semiedge, therefore creating leafs. Denote the resulting forest by T . The number of vertices in T is $k + |M|$, since k is the number of leafs and $|M|$ the number of internal vertices. The total edge count e is thus at most $k + |M| - 1$. Counting edges by vertex degrees, where each internal vertex contributes three edges and each leaf one edge, gives us $k + 3|M|$. Since this method counts each edge twice, we have:

$$k + 3|M| = 2e \leq 2k + 2|M| - 2.$$

This inequality yields $|M| \leq k - 2$. \square

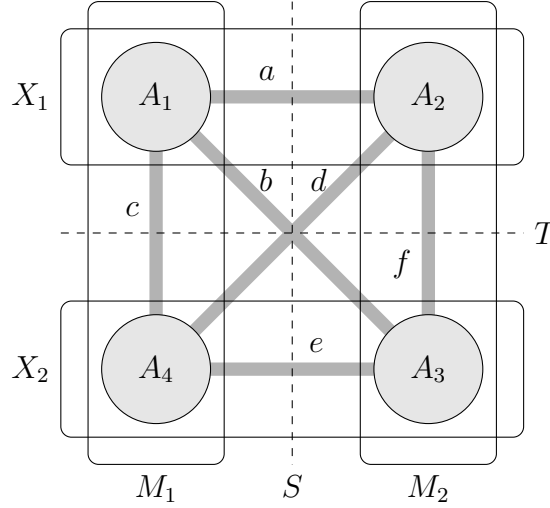


Figure 3.3: [29] Intersecting cuts S and T in the graph G

The following theorem, one of our main results, establishes necessary conditions for achieving desired cyclic edge-connectivity in graphs formed by a junction of two multipoles. While this theorem is a modification of a theorem proved by Rajník [29, Theorem 5.2] and it follows a similar proof structure, it required a modification to address a broader class of multipoles. The original proof used cyclic k -parts, whereas our version uses cyclic potentially cyclically k -connected k -poles. Specifically, where the original proof reached a contradiction by finding a cyclic multipole with few semiedges in a cyclic part, our proof reaches the same contradiction by finding such a multipole in a potentially cyclically k -connected multipole. For completeness, we present the entire proof.

Theorem 12. *Let M_1 and M_2 be cyclic potentially cyclically k -connected k -poles. Let both M_1 and M_2 be different from the cycle of length k or $k \leq 5$. Let G be the junction of M_1 and M_2 . Then $\zeta(G) = k$ unless the girth of G is smaller than k .*

Proof. Let us denote the k -cut obtained after the junction $M_1 * M_2$ by S . Suppose to the contrary that G has girth at least k and contains a cycle-separating edge-cut T of minimal size, such that $|T| = l < k$ and it separates G into two cyclic multipoles X_1 and X_2 . Note that in this case $\zeta(G) = l$. Let A_1, A_2, A_3 and A_4 denote the induced multipoles of G by $V(X_1) \cap V(M_1)$, $V(X_1) \cap V(M_2)$, $V(X_2) \cap V(M_2)$ and $V(X_2) \cap V(M_1)$ respectively. Let a, b, c, d, e and f denote the number of edges connecting vertices from A_1 and A_2 , A_1 and A_3 , A_1 and A_4 , A_2 and A_4 , A_3 and A_4 , A_2 and A_3 respectively as illustrated in Figure 3.3.

Next, let $s_i = |S(A_i)|$ for each $i \in \{1, 2, 3, 4\}$. We see that each of such s_i is determined by the values of a, b, c, d, e and f , for instance $s_1 = a + b + c$. We prove by contradiction that each of A_1, A_2, A_3 and A_4 is non-empty. Suppose that some A_i for $i \in \{1, 2, 3, 4\}$ is empty. Then M_1 or M_2 would contain the whole cyclic l -pole X_1 or X_2 ,

contradicting the assumption that M_1 and M_2 are potentially cyclically k -connected. From this point onward, the proof follows the one from [29, Theorem 5.2]. However, for completeness, we present the entire proof.

From the sizes of S and T we know that

$$\begin{aligned} a + b + d + e &= k, \\ l = c + b + d + f &< k. \end{aligned}$$

We now show that for each $i \in \{1, 2, 3, 4\}$ at least one of the inequalities $s_i \geq k$ or $s_{i+1} \geq k$ is true, taking the indices modulo 4. Suppose to the contrary, that there exists such $i \in \{1, 2, 3, 4\}$, for which $s_j \leq k - 1$ for both $j \in \{i, i + 1\}$. Let us look at the properties of one of such A_j . We know that it contains fewer than k semiedges, say m . It must be acyclic, otherwise A_j would be a cyclic multipole with fewer than k semiedges in M_1 or M_2 . Therefore, for both $j \in \{i, i + 1\}$, Lemma 7 implies $|A_j| \leq s_j - 2$, and since $s_j \leq k - 1$, we obtain $|A_j| \leq s_j - 2 \leq k - 3$. By denoting $G[A_i \cup A_{i+1}]$ by M we obtain $|M| \leq 2k - 6$ (A_i and A_{i+1} are vertex disjoint). Note that M is by its definition one of the multipoles X_1 , X_2 , M_1 , or M_2 . This means that M is cyclic and $|M| \geq k$ due to the girth of G , which is at least k . From $k \leq |M| \leq 2k - 6$ we obtain $k \geq 6$, so by the assumption of the theorem, M_1 and M_2 are different from k -cycles. Also, both X_1 and X_2 are different from k -cycles, since they have fewer than k semiedges. Therefore, M is not a k -cycle and it is a cyclic multipole with at most k semiedges and girth at least k , so by Lemma 1, $|M| \geq 2k - 4$, contradicting the fact that $|M| \leq 2k - 6$.

Thus for each $i \in \{1, 2, 3, 4\}$ we have $s_i \geq k$ or $s_{i+1} \geq k$, from which it follows that

$$s_1 = a + b + c \geq a + b + d + e = k \text{ and } s_3 = b + e + f \geq a + b + d + e = k$$

or

$$s_2 = a + d + f \geq a + b + d + e = k \text{ and } s_4 = c + d + e \geq a + b + d + e = k$$

is true. After summing, the first inequalities imply that $c + f \geq a + 2d + e \geq a + e$ and the second imply that $c + f \geq a + 2b + e \geq a + e$. In both cases we have obtained $c + f \geq a + e$. However, we have

$$l = b + c + d + f < a + b + d + e = k,$$

so $c + f < a + e$, leading to a contradiction.

We see that there is no cycle-separating edge-cut in G with fewer than k edges. Also, there is the cycle-separating k -edge-cut S , therefore it must be smallest and $\zeta(G) = k$. \square

Note that the main property around which this theorem operates is the girth of a graph. A small girth implies small cyclic edge-connectivity, as mentioned earlier. We demonstrate, that when performing a partial junction of multipoles, if the distances between pairs of semiedges in these multipoles are sufficiently large, the resulting multipole will also have large girth. This proves valuable in constructing infinite families of universal adjuncts. In such constructions, we aim to maintain large distances between the pairs of semiedges, ensuring that after junction with any multipole, the resulting graph will have large girth, regardless of the distances in that multipole.

Lemma 8. *Let M_1 and M_2 be multipoles with girth at least k . For both $i \in \{1, 2\}$, let d_i be the minimal distance between a pair of distinct semiedges in M_i . Let M be a result of a partial junction $M_1 * M_2$. If $d_1 + d_2 \geq k - 2$, then $g(M) \geq k$.*

Proof. Let C be a cycle in M . We prove that its length must be at least k . If C lies entirely within either M_1 or M_2 , its length must be at least k , since both multipoles have girth at least k .

If C spans both M_1 and M_2 , it must contain edges from both multipoles, as well as at least two edges connecting them. Note that these edges between the pair of multipoles are created by the partial junction.

Let e_1, \dots, e_j denote the edges in M that were created by the partial junction between M_1 and M_2 . Then C must contain:

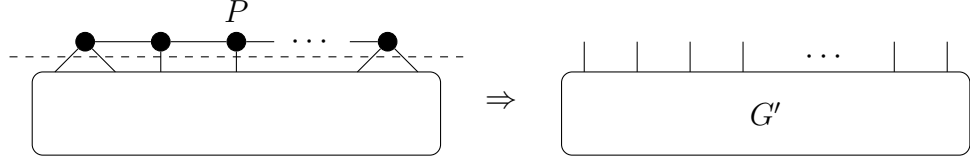
- (i) two distinct edges e_{i_1} and e_{i_2} connecting M_1 and M_2 ,
- (ii) a path P in M_1 between a pair of distinct edges connecting M_1 and M_2 ,
- (iii) a path Q in M_2 between a pair of distinct edges connecting M_1 and M_2 .

Note that the sets of edges in the three cases are pairwise disjoint. Also, each edge between M_1 and M_2 was created by a junction of semiedges from M_1 and M_2 . Now, by the definition of d_1 and d_2 , we obtain:

- (i) path P has length at least d_1 ,
- (ii) path Q has length at least d_2 ,
- (iii) the edges e_1 and e_2 contribute length 2.

Therefore, the total length of this path is at least $2 + d_1 + d_2$. Since $d_1 + d_2 \geq k - 2$, we have $2 + d_1 + d_2 \geq 2 + (k - 2) = k$, implying C has length at least k . \square

We now address the existence of infinite families of universal k -adjuncts. Rather than providing specific constructions, our approach just demonstrates their existence for any value of k . These universal k -adjuncts must satisfy several properties. They

Figure 3.4: Obtaining the multipole G' by removing P

must possess sufficiently large girth and should not create short cycles when joined with other multipoles. Moreover, they must be potentially cyclically k -connected, meaning they cannot contain cyclic multipoles with a small number of semiedges.

While an explicit construction of such multipoles would have practical value, our theoretical framework requires only the proof of their existence. The following lemma provides this existence proof using properties of transitive graphs.

Lemma 9. *For each $k \geq 1$, there exists a cyclic k -pole, denoted by W_k , with the following properties:*

- (i) *it is potentially cyclically k -connected,*
- (ii) *the distance between any pair of semiedges is at least $k - 2$.*

Proof. For k equal to 1 and 2, we can verify that C_k fulfils the requirements, we shall therefore assume $k \geq 3$.

For a given k , we construct a cubic vertex-transitive graph G with girth $2k$ according to Theorem 2. Then $\zeta(G) = 2k$ by Theorem 1. Next, let P be a path of length $k - 3$ in G . Since $k \geq 3$ and $|V(G)|$ is at least $2k$ because of its girth, there will always be such a path. Sever the edges in the edge-cut $E(V(P), V(G) - V(P))$ and denote the multipole not containing the path P by G' . This process is depicted in Figure 3.4.

We prove that the distance between any pair of semiedges in G' is at least $k - 2$. Suppose to the contrary that there are two distinct semiedges s_1 and s_2 , such that $d(s_1, s_2) \leq k - 3$. Let us denote the shortest path between them by Q . These semiedges were created by severing links between vertices in G' and vertices in P , not necessarily distinct. Let us denote these edges by e_1 and e_2 , and the vertices from P by v_1 and v_2 . Since v_1 and v_2 are on P , it holds that $d(v_1, v_2) \leq k - 3$, let us denote the shortest path by P' . Then $e_1 Q e_2 P'$ is a closed path of length at most $2 + (k - 3) + (k - 3) = 2k - 4$, contradicting the fact that G has girth $2k$.

Now we shall prove that G' is a cyclic k -pole, meaning it contains a cycle and has k semiedges. Let H' be a modification of G' , such that all semiedges are removed, therefore the result is a graph. We prove that $\delta(H') \geq 2$, meaning H' is cyclic, and this implies that G' is cyclic as well, since by adding edges to a cyclic multipole it remains cyclic.

Suppose there is a vertex v in H' with degree less than 2. In G' , this vertex was incident with two semiedges, therefore at least two distinct edges vw_1 and vw_2 were severed during the creation of G' , such that w_1 and w_2 are vertices on P . Let us denote these edges by f_1 and f_2 respectively. Since w_1 and w_2 are on a path P of length $k - 3$, the distance between them is at most $k - 3$. Let us denote the shortest path between them by P'' . Then $f_1P''f_2$ is a cycle of length at most $k - 1$ in G , contradicting the fact that the girth of G is at least $2k$.

We prove G' has exactly k semiedges by demonstrating that the size of the edge-cut $E(V(P), V(G'))$ is k . Note that no two non-adjacent vertices on the path P are connected by an edge, as this would create a cycle of length at most $k - 2$, contradicting our assumption.

On path of length $k - 3$ there are $k - 2$ vertices. If $k = 3$, we just cut three edges incident with one vertex (a path of length 0) and the result has three semiedges. If $k > 3$, there are two end vertices of the path and $k - 4$ internal. Each internal vertex contributes by one edge to the edge-cut, the end vertices by two. This means that the size of the edge-cut is $2 \cdot 2 + (k - 4) = k$, proving that G' will have exactly k semiedges.

Finally, we prove that G' is potentially cyclically k -connected by showing that it contains no cyclic multipole with fewer than k semiedges. Suppose to the contrary that it contains a cyclic l -pole M , such that $l < k$. We see that M is in G as well, thus $G = M * N$ for some other l -pole N . Suppose that N is acyclic. By Lemma 7, we know that $|N| \leq l - 2$. However, N contains all vertices of path P , since no vertex from the path P is in G' , thus not in M . This means, that $|N| \geq k$, which yields $k \leq l - 2$, contradicting our assumption that $l < k$.

Therefore, N must be cyclic, meaning the edge-cut in G separating M and N is cycle-separating and of size $l < k$. This contradicts $\zeta(G) = 2k$, proving that G' cannot contain any cyclic multipole with fewer than k semiedges. \square

The following theorem, which represents our main result regarding cyclic parts, builds on our previous results. We demonstrate that the problem of determining whether or not a cyclic multipole is a cyclic part does not have to rely on finding an original graph from which the cyclic part could be obtained, but rather only examine the properties of the multipole itself, specifically, by checking whether it contains a cyclic multipole with a small number of semiedges.

Theorem 13. *A cyclic k -pole M is a cyclic k -part if and only if it is potentially cyclically k -connected.*

Proof. If M is a cyclic k -part, then by Lemma 6 it is potentially cyclically k -connected.

For the converse implication, we can assume that M is different from a k -cycle,

since each k -cycle C_k is a cyclic part, as proven in Lemma 5. Let W_k be a cyclic k -pole satisfying Lemma 9 and let G be a junction $M * W_k$.

Since M is potentially cyclically k -connected, its girth must be at least k . Same holds for W_k . Also, the minimal distance between a pair of semiedges in W_k is $k - 2$, thus by Lemma 8 it holds that $g(G) \geq k$, regardless of the distances in M .

For $k \geq 4$, the multipole W_k cannot be a k -cycle. If it were, it would contain a pair of semiedges with distance 1 between them, contradicting the requirement that the distance between any pair of semiedges must be at least $k - 2$.

Thus we see that G is a cyclic cubic graph with girth at least k , resulting from the junction $M * W_k$. Since both M and W_k are potentially cyclically k -connected, both are different from a cycle of length k or $k < 4$, and their junction G has girth at least k , we can apply Theorem 12 and conclude that $\zeta(G) = k$. Because of this fact, we have proved that M is in fact a cyclic k -part. \square

This result shows that determining whether a cyclic multipole is a cyclic part can be reduced to checking if it contains a cyclic multipole with a small number of semiedges. While this does not completely solve the aforementioned Problem 1, it offers a new perspective: instead of searching for adjuncts, performing junctions, and determining the cyclic edge-connectivity of the resulting graph, we can focus on examining the structure of the multipole itself. However, we should note that this does not necessarily lead to a more efficient algorithm, as we still face the challenge of finding a cyclic multipole with the minimum number of semiedges. This challenge is addressed in Chapter 5 regarding problems for further research, where we propose both theoretical and algorithmic approaches for future investigation.

3.3 Junctions of multiple multipoles

In this section, we explore how to construct graphs with desired cyclic connectivity through junctions of multiple multipoles, establishing the necessary properties these multipoles must possess. As our work frequently involves discussing smallest independent edge-cuts in multipoles, we introduce the following notation:

Definition 11. For a multipole M , let $i(M)$ denote the size of its smallest independent edge-cut, or ∞ if M contains no such cut.

The following lemma establishes a fundamental property regarding independent edge-cuts in partial junctions. This result proves important for analyzing the structural properties of junctions of multipoles and provides a lower bound for the size of their smallest independent edge-cuts.

Lemma 10. *Let M_1 and M_2 be multipoles and let M be the result of a partial junction $M_1 * M_2$ of size j . Then $i(M) \geq \min\{i(M_1), i(M_2), j\}$.*

Proof. We may assume that $\min\{i(M_1), i(M_2), j\} > 0$, otherwise it would evidently be true, since each edge-cut has size at least 0.

Now suppose that a smallest independent edge-cut in M has size m and it holds that $m < \min\{i(M_1), i(M_2), j\}$.

Since $m < j$, severing only the edges created by the partial junction cannot disconnect the graph, as it would not cut all such edges. Also, since $\min\{i(M_1), i(M_2), j\} > 0$, both M_1 and M_2 are connected. Therefore, the cut needs to contain some edges from M_1 or M_2 , suppose it contains edges from M_1 . However, since these edges would be an independent edge-cut in M_1 , it would mean that $m \geq i(M_1)$, leading to a contradiction. \square

Now we present a fundamental theorem in our theoretical framework that addresses the problem introduced in our work: under what conditions does a partial junction of two potentially cyclically k -connected multipoles preserve the property of being potentially cyclically k -connected? This theorem not only serves as a building block for more complex constructions, but also offers insights into the structural relationships between cyclic connectivity and partial junctions, with regard to sizes of minimal independent edge-cuts in multipoles.

Theorem 14. *Let M_1 and M_2 be potentially cyclically k -connected multipoles. Let M be a result of a partial junction $M_1 * M_2$ of size j . Let s_i denote the number of remaining semiedges after the partial junction in M_i , meaning $s_i = |S(M_i)| - j$. If all of the following conditions hold:*

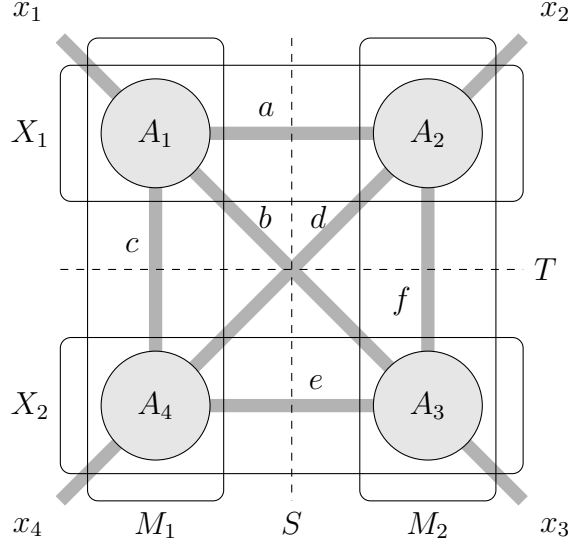
$$(i) \quad i(M_1) + i(M_2) \geq k,$$

$$(ii) \quad s_1 + s_2 \geq k,$$

$$(iii) \quad i(M_1) + s_2 \geq k \text{ and } s_1 + i(M_2) \geq k,$$

then M is potentially cyclically k -connected and $i(M) \geq \min\{i(M_1), i(M_2), j\}$.

Proof. Let us denote the j -cut obtained after the partial junction $M_1 * M_2$ by S . Suppose to the contrary that M contains a cyclic l -pole X_1 , such that $l < k$ and l is the smallest possible. Let T be the minimal edge-cut in M yielding X_1 and denote the second component by X_2 . Let A_1, A_2, A_3 and A_4 denote the induced multipoles of G by $V(X_1) \cap V(M_1), V(X_1) \cap V(M_2), V(X_2) \cap V(M_2)$ and $V(X_2) \cap V(M_1)$ respectively. Let a, b, c, d, e and f denote the number of edges connecting vertices from A_1 and A_2, A_1 and A_3, A_1 and A_4, A_2 and A_4, A_3 and A_4, A_2 and A_3 respectively. For each

Figure 3.5: Intersecting cuts S and T in the multipole M

$i \in \{1, 2, 3, 4\}$, let x_i be the number of semiedges in M incident with vertices in A_i as illustrated in Figure 3.5. We see that $b + c + d + f + x_1 + x_2 = l < k$.

We first prove that A_1 and A_2 are non-empty. Suppose that one is empty, say A_1 . Then M_2 would contain the whole cyclic l -pole X_1 , contradicting the fact that M_2 is potentially cyclically k -connected.

Next, we prove that A_3 and A_4 cannot be empty at the same time, meaning X_2 cannot be empty. Suppose that both were empty, meaning $X_1 = M$. This would mean that $|S(M)| = |S(X_1)| = l < k$, contradicting the fact that $s_1 + s_2 \geq k$.

The cut T separating X_1 from X_2 is independent by Lemma 3. Now we shall explore A_3 and A_4 . If both are non-empty, c and f represent independent edge-cuts in M_1 and M_2 respectively, meaning $c \geq i(M_1)$ and $f \geq i(M_2)$. However $i(M_1) + i(M_2) \geq k$ by (i), therefore $c + f \geq k$, contradicting the assumption that $b + c + d + f + x_1 + x_2 < k$.

Now suppose one of A_3 and A_4 is empty, say A_3 . In this case c represents an independent edge-cut in M_1 , thus $c \geq i(M_1)$. Also, we see that in this case $x_3 = 0$ and thus $x_2 = s_2$. Since $i(M_1) + s_2 \geq k$ by (iii), we have $c + x_2 \geq k$, contradicting the assumption that $b + c + d + f + x_1 + x_2 < k$.

As all possible cases lead to contradictions, we conclude that M is potentially cyclically k -connected.

The fact that $i(M) \geq \min\{i(M_1), i(M_2), j\}$ follows from Lemma 10. \square

This shows that the size of a smallest independent edge-cut in a multipole is a valuable property when analyzing cyclic edge-connectivity of graphs constructed from these multipoles. Specifically, if we perform a partial junction between two multipoles that are potentially cyclically k -connected, and if the sum of the sizes of their smallest independent edge-cuts is sufficiently large, then the resulting multipole will also remain

potentially cyclically k -connected (it is also necessary for the multipoles to satisfy certain conditions relating the number of semiedges to the size of their smallest independent edge-cuts, but all of these conditions concern independent edge-cuts).

Although this Theorem proves useful, as shown in the examples section (Section 3.4), in many cases we encounter difficulties related to the size of smallest independent edge-cuts. In particular examples, the multipoles that serve as building blocks are obtained by removing vertices from graphs with specific properties, which can result in multipoles having relatively small independent edge-cuts. However, not all independent edge-cuts are problematic in terms of undermining the desired cyclic edge-connectivity of the resulting graph. To clarify this idea, we introduce a new definition and notation, which will allow us to handle smaller independent edge-cuts in relation to the number of semiedges and a partial junction with another multipole.

Definition 12. Let M be a multipole and $X \subseteq S(M)$. Let \mathcal{F} denote the set of all independent edge-cuts in M . For any cut $F \in \mathcal{F}$, let M_1 and M_2 denote the two components obtained by severing F in M . Then the *independent cut-semiedge index* of M with X , denoted by $I_S(M, X)$ is defined as

$$I_S(M, X) = \min_{F \in \mathcal{F}} \left\{ |F| + \min \left\{ \left| (S(M_1) \cap S(M)) - X \right|, \left| (S(M_2) \cap S(M)) - X \right| \right\} \right\}.$$

In other words, for each independent edge-cut F , we:

1. count the size of the cut;
2. for each resulting component, count its semiedges that were original semiedges of M , excluding those in X ;
3. take the smaller of these semiedge counts and add it to the size of the cut.

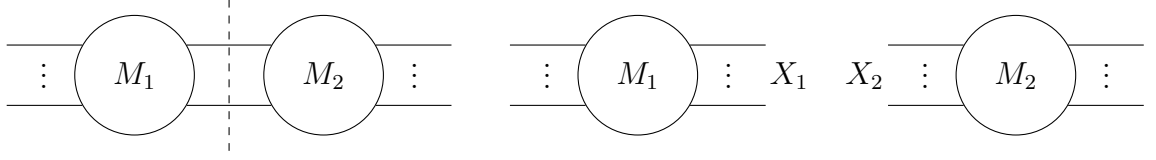
The $I_S(M, X)$ is then the minimum of these sums over all possible independent edge-cuts.

This notion allows smaller independent edge-cuts, if they are dividing semiedges into two components. Using this we can prove a similar theorem to Theorem 14, but with using independent cut-semiedge indices.

Theorem 15. Let M_1 and M_2 be potentially cyclically k -connected multipoles. Let M be a result of a partial junction $M_1 * M_2$ of size j , such that by X_i we denote the set of semiedges of M_i used in the partial junction, for both $i \in \{1, 2\}$. Let s_i denote the number of remaining semiedges after the partial junction in M_i , meaning $s_i = |S(M_i)| - j$. If all of the following conditions hold:

$$(i) \quad I_S(M_1, X_1) + I_S(M_2, X_2) \geq k,$$

$$(ii) \quad s_1 + s_2 \geq k,$$



(a) Partial junction (marked by the dashed line) of M_1 and M_2 (b) Sets of semiedges X_1 and X_2 in regard to the partial junction

Figure 3.6: Sets of semiedges X_1 and X_2

(iii) $I_S(M_1, X_1) + s_2 \geq k$ and $s_1 + I_S(M_2, X_2) \geq k$,

then M is potentially cyclically k -connected.

An illustration of which semiedges are in sets X_1 and X_2 is depicted in Figure 3.6. In the first illustration is illustrated the junction of M_1 and M_2 , marked by the dashed line, and in the second illustration are these two multipoles with marked X_1 and X_2 in regard to that specific junction.

Proof. Let us denote the j -cut obtained after the partial junction $M_1 * M_2$ by S . Suppose to the contrary that M contains a cyclic l -pole X_1 , such that $l < k$ and l is the smallest possible. Let T be the minimal edge-cut in M yielding X_1 and denote the second component by X_2 . We maintain the notation used in the previous theorem, an illustration is depicted in Figure 3.5.

Proofs of the statements that A_1 and A_2 are non-empty, A_3 and A_4 cannot be empty at the same time and that the cut T is independent are the same as in Theorem 14.

Now we shall explore A_3 and A_4 . If both are non-empty, then c and f represent independent edge-cuts in M_1 and M_2 respectively, while x_1 and x_2 represent semiedges from M_1 and M_2 that are not in X_1 and X_2 respectively, meaning $c + x_1 \geq I_S(M_1, X_1)$ and $f + x_2 \geq I_S(M_2, X_2)$. However, it holds that $I_S(M_1, X_1) + I_S(M_2, X_2) \geq k$, therefore $c + f + x_1 + x_2 \geq k$, contradicting the assumption that $b + c + d + f + x_1 + x_2 < k$.

Now suppose one of A_3 and A_4 is empty, say A_3 . In this case c represents an independent edge-cut in M_1 , while x_1 represents the semiedges in M_1 that are not in X_1 , thus $c + x_1 \geq I_S(M_1, X_1)$. Also, we see that in this case $x_3 = 0$ and thus $x_2 = s_2$. Since $I_S(M_1, X_1) + s_2 \geq k$, we have $c + x_1 + x_2 \geq k$, contradicting the assumption that $b + c + d + f + x_1 + x_2 < k$.

As all possible cases lead to contradictions, we conclude that M is potentially cyclically k -connected. \square

We have demonstrated theorems regarding partial junctions of multipoles that result in larger multipoles. Our goal, however, is to create graphs. We demonstrate our main result regarding junctions in the following theorem, where the conditions for multipoles are similar to previous theorems, but the result is a graph.

Theorem 16. *Let M_1 and M_2 be potentially cyclically k -connected multipoles. Let G be the result of a junction $M_1 * M_2$. If $i(M_1) + i(M_2) \geq k$, then G is cyclically k -connected.*

Proof. Suppose to the contrary that G is not cyclically k -connected. Then, by Theorem 3, G contains an independent edge-cut T of size $l < k$, which is cycle-separating, and it separates G into two cyclic l -poles X_1 and X_2 . We maintain the notation used in previous theorems, an illustration is depicted in Figure 3.3.

We first prove that all A_1, \dots, A_4 are non-empty. Suppose that one is empty, say A_1 . Then M_2 would contain the whole cyclic l -pole X_1 , contradicting the assumption that M_2 is potentially cyclically k -connected.

Since T is independent and all A_1, \dots, A_4 are non-empty, it holds that c and f represent independent edge-cuts in M_1 and M_2 respectively. This means that $c \geq i(M_1)$ and $f \geq i(M_2)$. However, it holds that $i(M_1) + i(M_2) \geq k$, therefore $c + f \geq k$, contradicting the assumption that $b + c + d + f < k$. \square

Now we will look at theorems regarding modifications of multipoles and specific junctions with multipoles, that could be problematic in our previous results. First, we show that when subdividing a graph with known cyclic edge-connectivity, we can easily find a lower bound for the size of a smallest independent edge-cut in the resulting multipole. We also show that subdivisions do not change whether a multipole is potentially cyclically k -connected.

Lemma 11. *Let G be a cyclically k -connected graph. Let M be a multipole obtained from G by subdividing some edges and adding a semiedge to each new vertex created by these subdivisions. Then $i(M) \geq k$.*

Proof. Let G' be a graph obtained from M by removing all semiedges. It holds, that the cyclic connectivity of a graph is invariant under subdivisions [25], thus G' is cyclically k -connected. By Theorem 3 it thus holds that $i(G') \geq k$.

Moreover, by adding semiedges to all vertices of degree two (created by the subdivisions), obtaining M , we do not obtain a smaller independent edge-cut. This cut would be present in G' as well and would be independent, contradicting our earlier conclusion.

Therefore, $i(M) \geq k$. \square

As mentioned in the proof, cyclic connectivity of a graph is invariant under edge subdivisions. We now prove a similar result for multipoles: when we subdivide edges in a potentially cyclically k -connected multipole, it remains potentially cyclically k -connected as long as the resulting multipole has at least k semiedges. This minimum requirement of k semiedges is necessary, because when subdividing a potentially cyclically k -connected graph (a special case of a multipole), we need at least k subdivisions so that it will remain potentially cyclically k -connected, otherwise it would have fewer than k semiedges.

Theorem 17. *Let M be a potentially cyclically k -connected multipole. Let M' be a multipole obtained from M by subdividing some edges and adding a semiedge to each new vertex created by these subdivisions. If $|S(M')| \geq k$, then M' is potentially cyclically k -connected.*

Proof. If M is a graph, then by Lemma 11, $i(M') \geq k$. This means that each independent edge-cut in M' needs to cut at least k edges, thus yields multipoles with at least k semiedges. This, together with the fact that M' has at least k semiedges, shows that M' is potentially cyclically k -connected.

Now assume M is not a graph. Suppose to the contrary that M' is not potentially cyclically k -connected, meaning it contains a cyclic l -pole X' for $l < k$. We show that $X = M[V(X') \cap V(M)]$ is a cyclic multipole in M with at most l semiedges.

We can obtain M from M' by reversing the process that created M' from M : removing the semiedges from the subdivision vertices and smoothing all vertices of degree 2. Therefore, X has at most as many semiedges as X' , since we remove all semiedges from the subdivision vertices, and smoothing vertices of degree 2 does not change the number of semiedges in a multipole, as follows from the definition.

By removing a semiedge from a cyclic multipole, the multipole remains cyclic as no cycle can pass through a semiedge. Moreover, smoothing vertices of degree 2 preserves cyclicity: if a cycle contains edges vx and xw where x has degree 2, smoothing creates a new edge vw , ensuring the multipole remains cyclic. Therefore, X is a cyclic multipole with at most $l < k$ semiedges in M , contradicting our assumption that M is potentially cyclically k -connected. \square

Sometimes though, we encounter a problem when we want to obtain a potentially cyclically k -connected multipole by a subdivision. Consider a case where we want a multipole with exactly k semiedges, starting from a multipole that is not a graph. This presents a problem because the original multipole would necessarily have fewer than k semiedges, since subdivisions increase the number of semiedges. This problem is best visible in Section 3.4, in Example 3, where we aim to create a cyclically 5-connected graph. If we were to create the middle vertex by subdividing an edge in a multipole formed by the partial junction $M_1 * M_2$, this partial junction would create a cyclic 4-pole, which cannot be potentially cyclically 5-connected and thus this theorem could not be used.

For this reason we prove a theorem specific for this case, where we tackle this problem of not having potentially cyclically k -connected multipole in some iteration of creating cyclically k -connected graph.

Lemma 12. *Let M be a potentially cyclically k -connected multipole with at least $k + 1$ semiedges, obtained from a partial junction of multipoles M_1 and M_2 satisfying*

the conditions of Theorem 14. Let N be the result of a partial junction of size 2 between M , and a multipole consisting of a single vertex v with three semiedges, where new edges vv_1 and vv_2 are created, where $v_i \in V(M_i)$ for both $i \in \{1, 2\}$. Then N is potentially cyclically k -connected.

Proof. Since M has at least $k + 1$ semiedges, N has at least k semiedges. We therefore only need to verify that there exists no independent edge-cut yielding a cyclic multipole with fewer than k semiedges.

Suppose to the contrary that such a cut exists, and let X be the resulting cyclic l -pole where $l < k$ and l is minimal. We can make several observations about X :

- (i) X must contain vertex v , as otherwise X would exist in M , contradicting the fact that M is potentially cyclically k -connected;
- (ii) X must contain both v_1 and v_2 . If not, X would have at least two semiedges from v , and by excluding v , we would obtain a cyclic multipole with fewer than l semiedges, contradicting the minimality of l . Note that the multipole would remain cyclic since v with two semiedges cannot be part of any cycle.

Given these observations, we consider all possible configurations of X :

- X contains whole M_1 and M_2 . But then X would equal N , which has at least k semiedges, contradicting our assumption that X has fewer than k semiedges;
- Without loss of generality, if X contains all of M_1 but not all of M_2 , then since X contains w_2 , the independent cut must intersect M_2 . This yields at least $i(M_2)$ semiedges from the cut in M_2 , $s_1 - 1$ from M_1 , and one from v . The total number of semiedges is thus at least $i(M_2) + s_1$, which is at least k , leading to a contradiction;
- If X contains neither M_1 nor M_2 entirely, then X must have at least $i(M_1) + i(M_2)$ semiedges. Since $i(M_1) + i(M_2) \geq k$, X has at least k semiedges, leading to a contradiction with the assumption that X has less than k semiedges.

As all cases lead to contradictions, we conclude that N is potentially cyclically k -connected. \square

In this section, we have developed a theoretical framework for proving cyclic edge-connectivity in graphs constructed from smaller multipoles. Two key properties emerged as particularly significant: first, the requirement that multipoles cannot contain other cyclic multipoles with a small number of semiedges, and second, that the sizes of their smallest independent edge-cuts must be sufficiently large. Beyond establishing these theorems for partial junctions, we have also proven results for more

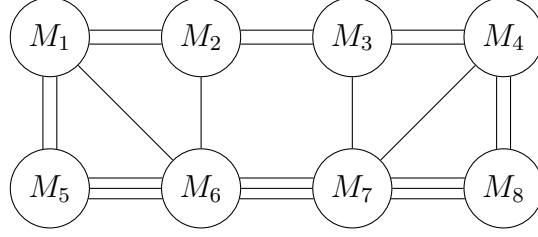


Figure 3.7: Cyclically 5-connected graph constructed from 8 multipoles

specific cases, such as partial junctions with a single vertex and edge subdivisions. The practical application of this theoretical framework is demonstrated in the following section through specific examples, illustrating how these results can be effectively utilized in constructing graphs with desired cyclic edge-connectivity.

3.4 Examples

We begin by examining examples that emerged as challenges during the development of this thesis – specifically, problems that arose when investigating conditions under which graphs resulting from partial junctions of multipoles achieve the desired cyclic edge-connectivity. These examples guided the development of our theoretical framework, helping us generalize from specific cases to theorems applicable to a broader class of graphs.

Following these, we analyze cases from existing literature, where authors have made claims about cyclic connectivity without providing formal proofs. We try and look at these graphs in a perspective of junctions of multipoles and apply our theoretical framework to provide rigorous proofs of these previously unverified claims.

First, however, let us explore examples that demonstrate the sufficient conditions multipoles must satisfy to ensure that graphs constructed through their partial junctions achieve the desired cyclic edge-connectivity.

Example 1. Let M_1, \dots, M_8 be potentially cyclically 5-connected multipoles, such that:

- M_6 and M_7 are 8-poles, and other are 5-poles;
- Let either $i(M_1)$ or $i(M_2)$ be at least 2, and other at least 3.

Let G be the graph depicted in Figure 3.7, resulting from a sequence of partial junctions between M_1, \dots, M_8 . By utilizing the theorems proved in our work, we prove that the result is cyclically 5-connected.

First, take a partial junction $M_1 * M_2$ of size 2. By Theorem 14, the result is potentially cyclically 5-connected and $i(M_1 * M_2) \geq 2$. Similarly, we perform two

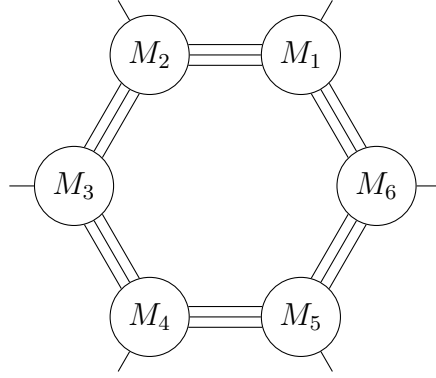


Figure 3.8: Potentially cyclically 6-connected multipole constructed from six 7-poles

more partial junctions of size 2, first with M_3 , and then with M_4 , such that we obtain an 8-pole $((M_1 * M_2) * M_3) * M_4$, let us denote it by N_1 . Again, by Theorem 14, it holds that N_1 is potentially cyclically 5-connected and $i(N_1) \geq 2$.

Now take a partial junction $M_5 * M_6$ of size 3. By Theorem 14, the result is potentially cyclically 5-connected and $i(M_5 * M_6) \geq 3$. Again, we perform two more partial junctions of size 3, first with M_7 , and then with M_8 , such that we obtain an 8-pole $((M_5 * M_6) * M_7) * M_8$, let us denote it by N_2 . Again, by Theorem 14, it holds that N_2 is potentially cyclically 5-connected and $i(N_2) \geq 3$.

Finally, we perform a junction $N_1 * N_2$, resulting in a graph G as depicted in Figure 3.7. By Theorem 16, it holds that G is cyclically 5-connected.

Note that the graph illustrated in Figure 3.7 represents just one possible outcome of the described sequence of partial junctions. Our theorems do not specify on which semiedges the partial junction should be performed, thus the showcased result represents only one of the possible outcomes.

Example 2. Let M_1, \dots, M_6 be potentially cyclically 6-connected 7-poles. Let $i(M_j) \geq 3$ for each $j \in \{1, \dots, 6\}$. We prove that the multipole M depicted in Figure 3.8 resulting from a sequence of partial junctions between M_1, \dots, M_6 is potentially cyclically 6-connected.

First, take a partial junction $M_1 * M_2$ of size 3. By Theorem 14, the result is potentially cyclically 6-connected and $i(M_1 * M_2) \geq 3$. Similarly, we perform a partial junction of the result and M_3 , obtaining a potentially cyclically 6-connected 9-pole $N_1 = (M_1 * M_2) * M_3$, such that $i(N_1) \geq 3$.

We perform the same operations on M_4, M_5 and M_6 , obtaining $N_2 = (M_4 * M_5) * M_6$ with the same properties. We then perform a partial junction $N_1 * N_2$ of size 6, resulting in a 6-pole N . By Theorem 14, it holds that N is potentially cyclically 6-connected.

Again, note that the multipole illustrated in Figure 3.8 represents just one possible outcome of the described sequence of partial junctions.

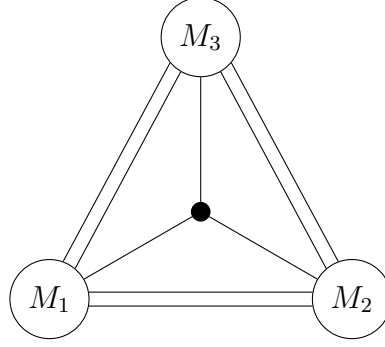
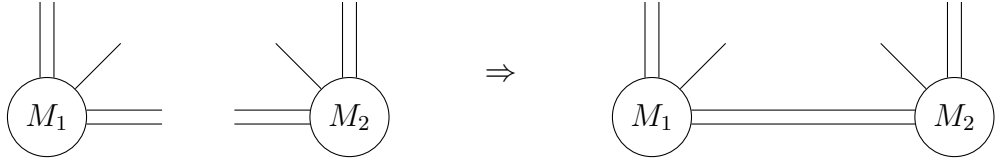


Figure 3.9: Graph from Example 3

Figure 3.10: First step in Example 3 resulting in multipole M_{12}

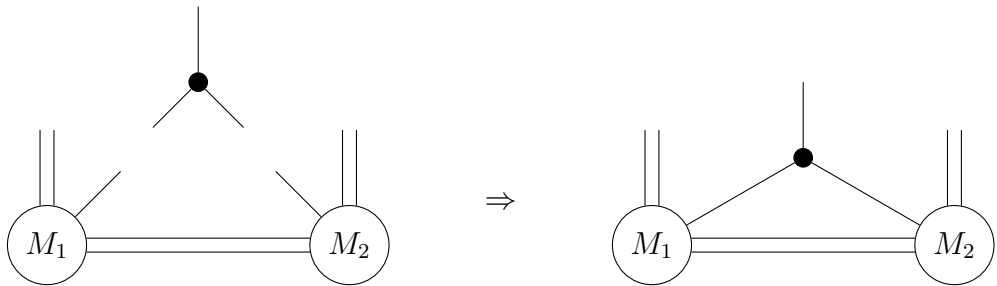
Example 3. Let M_1 , M_2 and M_3 be potentially cyclically 5-connected 5-poles, such that $i(M_1) \geq 2$, $i(M_2) \geq 3$ and $i(M_3) \geq 3$. We prove that a junction of these three multipoles along with one vertex in the middle, as depicted in Figure 3.9, is cyclically 5-connected.

In the first step, we perform a partial junction of size 2 on M_1 and M_2 as illustrated in Figure 3.10, resulting in a multipole M_{12} . By Theorem 14, it is potentially cyclically 5-connected and $i(M_{12}) \geq 2$.

Now we perform a partial junction of size two, between one vertex with three semiedges and M_{12} , resulting in M' , as depicted in Figure 3.11. By Lemma 12, the result is potentially cyclically 5-connected and by Lemma 10 $i(M') \geq 2$.

At last, we perform a junction of M' and M_3 , creating G . By Theorem 16, it holds that G is cyclically 5-connected.

We now examine specific constructions from published literature, where claims about cyclic edge-connectivity were made without formal proofs. For each of these examples, we present the original construction, and provide a formal proof using our theoretical

Figure 3.11: Second step in Example 3 resulting in multipole M'

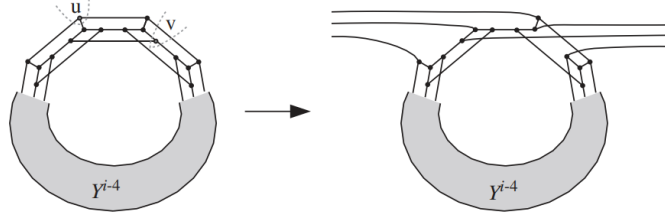


Figure 3.12: [24] Construction of 6-pole F^i from Isaacs snark I_i

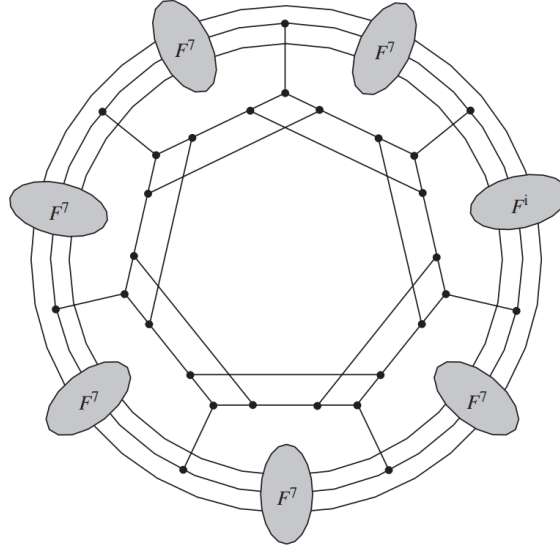


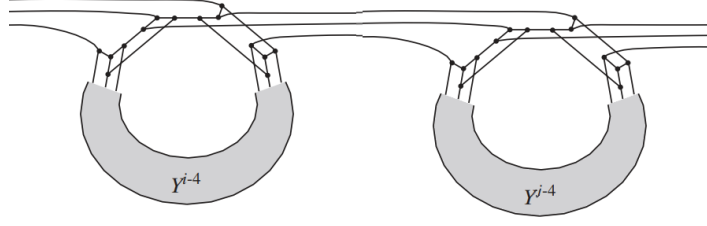
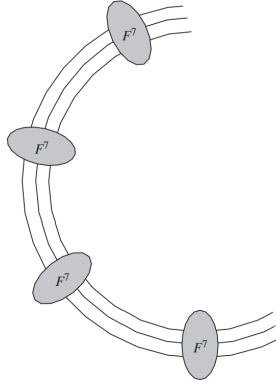
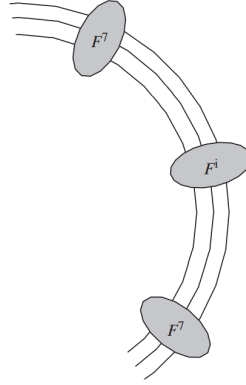
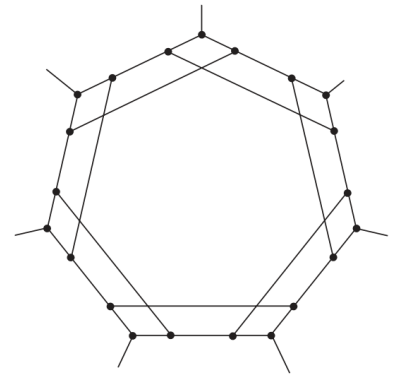
Figure 3.13: [24] Cyclically 6-connected graph $G_{1,i}$

framework, demonstrating how our approach can formalize and verify these previously unproven claims.

Example 4. In 2006, Máčajová and Škoviera presented a sequence of graphs $G_{1,i}$ for odd $i \geq 7$ [24], depicted in Figure 3.13. This graph can be viewed as a junction of several 6-poles F^i , where i is odd and $i \geq 7$, each of which is derived from the Isaacs snark I_i [13] by removing two vertices, as illustrated in Figure 3.12. Specifically, the construction uses six copies of F^7 and one F^i with arbitrary odd $i \geq 7$, along with the structure in the middle. Even though the authors describe this graph as a result of a superposition, we will not dive deep into this theory and present the result as a result of partial junctions.

The authors presented this graph as having cyclic connectivity equal to 6, however they have not provided a formal proof supporting this statement. In this example we prove this statement by utilizing our theoretical framework.

First observation is that each for each odd $i \geq 7$, F_i is potentially cyclically 6-connected: it has six semiedges, and we can verify that it does not contain a cyclic multipole with at most five semiedges. A helpful fact to verify this is that each Isaacs snark I_i with odd $i \geq 7$ is cyclically 6-connected [20]. Secondly, we see that

Figure 3.14: Partial junction $F^i * F^j$ (a) Multipole N_1 (b) Multipole N_2 (c) Multipole I Figure 3.15: Multipoles used in building the graph $G_{1,i}$

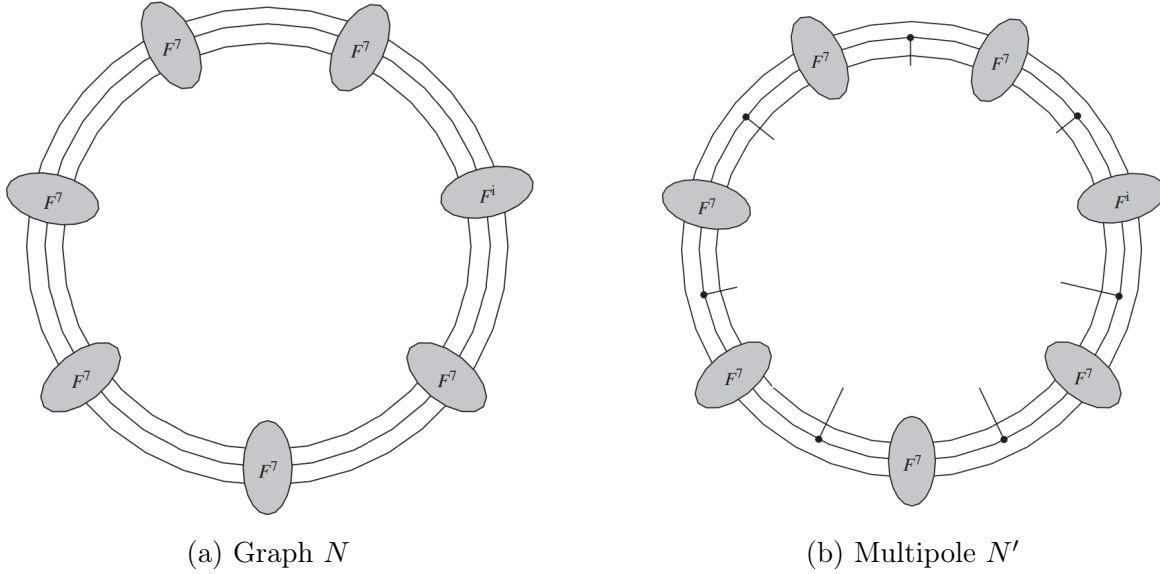
the size of a smallest independent edge-cut in any F_i , $i \geq 7$ is 2. However, when performing a partial junction between any two such multipoles ($F_i * F_j$ as shown in Figure 3.14), by denoting the semiedges used in the partial junction in F_i and F_j by X_i and X_j respectively, we can verify that $I_S(F_i, X_i) \geq 3$ and $I_S(F_j, X_j) \geq 3$. Therefore, by Theorem 15, the result of the partial junction $F_i * F_j$ depicted in Figure 3.14 is potentially cyclically 6-connected. Let us denote the resulting multipole by M_1 . We can verify, that M_1 has no independent 2-edge-cut, thus $i(M_1) \geq 3$, which implies $I_S(M_1, X) \geq 3$ for any $X \subseteq S(M_1)$.

We can now inductively construct multipole N_1 , depicted in Figure 3.15a, through multiple partial junctions with F_7 . By Theorem 15, N_1 will be potentially cyclically 6-connected and it can be verified, that it still holds that $i(N_1) \geq 3$. Similarly, we construct multipole N_2 , however with one F_i with arbitrary odd $i \geq 7$, depicted in Figure 3.15b, with the same properties.

We then perform a junction $N_1 * N_2$, resulting in graph N , as illustrated in Figure 3.16a. By Theorem 16, N is cyclically 6-connected.

Next, we perform seven edge subdivisions to obtain the multipole N' depicted in Figure 3.16b. By Theorem 17, N' is potentially cyclically 6-connected, and by Lemma 11, $i(N') \geq 6$.

Finally, let I be the multipole shown in Figure 3.15c. We can verify that it is potentially cyclically 6-connected since it has 7 semiedges and contains no cyclic

Figure 3.16: N before and after subdivisions

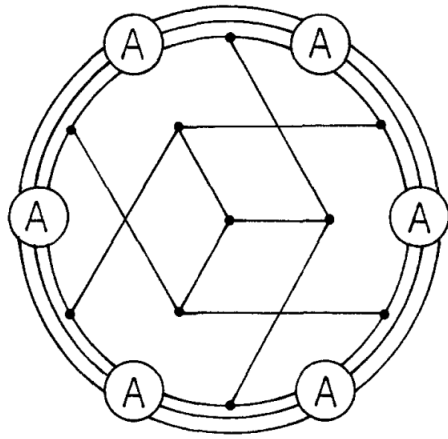
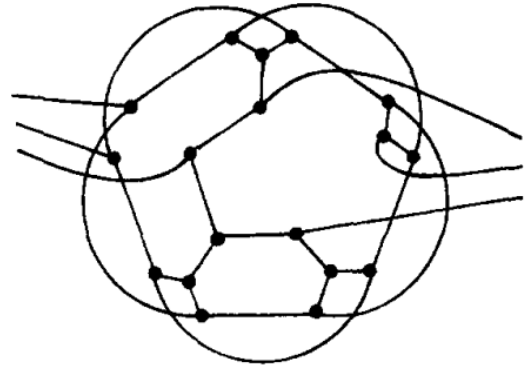
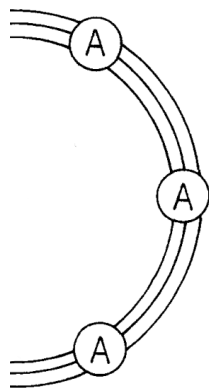
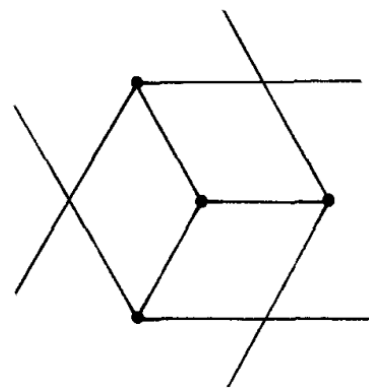
multipole with at most 5 semiedges. Therefore, since $i(N') \geq 6$, by Theorem 16, the final graph $G_{1,i} = N' * I$, depicted at the beginning in Figure 3.13 is cyclically 6-connected.

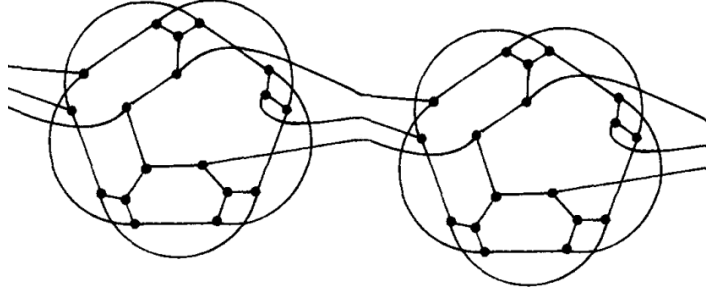
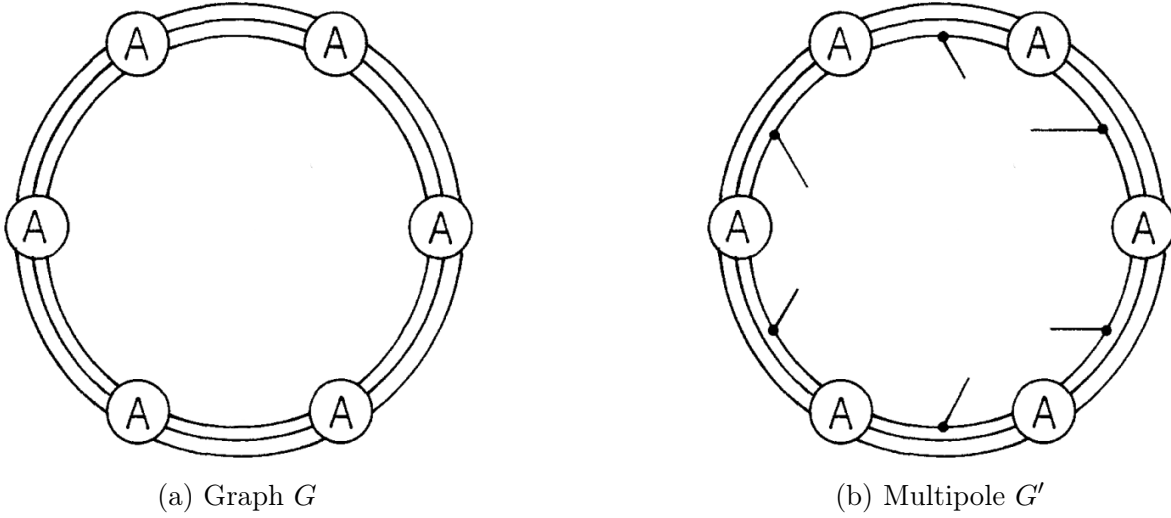
Using our theoretical framework, we have formally proven that $G_{1,i}$ is indeed cyclically 6-connected, as claimed. However, this example has highlighted certain limitations in our framework. Firstly, we lack an efficient method for verifying that a multipole is potentially cyclically k -connected, even when we know the properties of the original graph from which the multipole was derived (in this case by removing a pair of vertices). Secondly, we had to manually verify the size of the smallest independent edge-cut in the multipoles, as the methods in our theoretical framework, specifically Lemma 10, could not establish that this size was at least three. While these limitations required some manual verification steps, the majority of our proof relied on our theoretical framework. These challenges are explored in more detail in Chapter 5.

We now examine a similar construction from another article. While the proof follows a similar approach to the one presented above, some manual verification steps remain necessary.

Example 5. Consider the graph G_{118} presented by Kochol in 1996 [15], depicted in Figure 3.17a. This graph can be viewed as a junction of 6-poles A , created from Isaacs snark I_5 [13] by removing two vertices, along with the construction illustrated in Figure 3.18b. The multipole A is depicted in Figure 3.17b. Even though the author describes this graph as a result of a superposition, as before we will not dive deep into this theory and present the result as a result of partial junctions.

The author presented this graph as having cyclic connectivity equal to 6, however

(a) Graph G_{118} (b) Multipole A Figure 3.17: [15] Graph G_{118} built from multipoles A (a) Multipole N (b) Multipole I Figure 3.18: Multipoles used in building the graph G_{118}

Figure 3.19: Partial junction $A * A$ Figure 3.20: G before and after subdivisions

without a formal proof. In this example we prove this statement by utilizing our theorems.

First observation is that A is potentially cyclically 6-connected. It has six semiedges, and we can verify that it does not contain a cyclic multipole with at most five semiedges. Secondly, we see that the size of a smallest independent edge-cut in A is 2. However, when performing a partial junction between two copies of A ($A * A$ as shown in Figure 3.19), by denoting their semiedges used in the partial junction by X_1 and X_2 respectively, we can verify that $I_S(A, X_i) \geq 3$ for both $i \in \{1, 2\}$. By Theorem 15, the result of the partial junction $A * A$, depicted in Figure 3.19, is potentially cyclically 6-connected. Let us denote this multipole by M_1 . We can verify that M_1 has no independent 2-edge-cut, thus $i(M_1) \geq 3$, which implies $I_S(M_1, X) \geq 3$ for any $X \subseteq S(M_1)$.

Using this result, we can inductively construct two copies of multipole N , depicted in Figure 3.18a, through multiple partial junctions with A . By Theorem 15, N is potentially cyclically 6-connected and it can be verified that it still holds that $i(N) \geq 3$.

We then perform a junction $N * N$, resulting in graph G , as illustrated in Figure 3.20a. By Theorem 16, G is cyclically 6-connected.

Next, we perform six edge subdivisions to obtain the multipole G' depicted in

Figure 3.20b. By Theorem 17, G' is potentially cyclically 6-connected, and by Lemma 11, $i(G') \geq 6$.

Finally, let I be the multipole shown in Figure 3.18b. We see that $i(I) = 1$. However, since it is acyclic, it is potentially cyclically 6-connected. Since $i(G') \geq 6$, by Theorem 16, the final graph $G_{118} = G' * I$, depicted at the beginning in Figure 3.17a is cyclically 6-connected.

Note that this is just one example of constructing such a graph using multipoles denoted by A . Instead of this specific construction, we can use any potentially cyclically 6-connected multipoles instead of A , not necessarily identical, whose smallest independent edge-cut has size at least 3, and the result still holds.

Chapter 4

Cyclic part inflations

The concept of inflations was previously established in Section 2.3, where we presented the necessary definitions and the main theorem. This construction is particularly interesting as it enables the transformation of a graph through an “inflation” process, yielding a larger graph whose cyclic edge-connectivity can be determined from specific properties of the original graph.

This chapter introduces a modification to the mentioned inflation construction. While inflations involve replacing each vertex with a cycle whose length corresponds to the vertex degree, we propose *cyclic part inflations*, where vertices are substituted with cyclic k -parts, where k represents the degree of the respective vertex.

The chapter begins with essential definitions that formalize this operation, followed by our first theorem, which extends and modifies Theorem 11 mentioned in the chapter about inflations. Subsequently, we present a second theorem, which emerged during examining one of the examples mentioned afterwards, when we have proved it specifically for that example, and then made it more abstract to fit a broader class of graphs. At the end, we present illustrative applications of our theorems on specific examples.

Let G be a graph obtained by a sequence of partial junctions between cyclic parts P_1, \dots, P_n . A set of vertices $V \subseteq V(G)$ *corresponds* to a cyclic part P if $V = V(P)$.

We now formally define the concept of cyclic part inflations. Note that the graph from which we create inflations is arbitrary, thus does not have to be cubic.

Definition 13. Let G be an arbitrary graph and let H be a cubic graph resulting from performing partial junctions on cyclic parts P_1, \dots, P_n . Let $\mathcal{V} = \{V_1, \dots, V_n\}$ be a partition of $V(H)$, such that for each i from 1 to n the set V_i corresponds to the cyclic part P_i . Then H is called a *cyclic part inflation* of G , if the graph obtained from H by contracting each set $V_i \in \mathcal{V}$ into a single vertex is isomorphic to G . The set of all cyclic part inflations of G is denoted by $I_P(G)$.

Similarly, if we have a vertex $v \in G$ and some cyclic part inflation $H \in I_P(G)$, we say that the vertex v corresponds to the set of vertices $V \subseteq V(H)$, if V corresponds to

some cyclic part P in H and by contracting we obtain the vertex v in G .

First observation about these cyclic part inflations is in the form of a lemma, and is basically saying that cyclic edge-connectivity of the inflation is bounded from above by the edge-connectivity of the former graph.

Lemma 13. *Let G be a nontrivial graph with $\lambda(G) = k$. Then, for each cyclic part inflation H of G , it holds that $\zeta(H) \leq k$.*

Proof. Let S be a smallest edge-cut in G separating it into multipoles M_1 and M_2 , and let H be a cyclic part inflation of G . In H , the edge-cut S is cycle-separating, since in both M_1 and M_2 , each vertex has been replaced by a cyclic part. Also, since $\lambda(G) = k$, this cut is of size k , therefore the size of a smallest cycle-separating edge-cut in H must be at most k . \square

We now prove our modified version of Theorem 11. Although the proof follows a similar structure to the original theorem, we have made several adjustments to accommodate arbitrary cyclic parts instead of k -cycles. It is worth noting that traditional inflations can be viewed as a special case of cyclic part inflations, since Lemma 5 shows that all k -cycles are cyclic parts. This means our new theorem naturally extends to cover both traditional inflations as well as cyclic part inflations. The proof follows that of Theorem 11 in [12], however several modifications were needed to address cyclic part inflations instead of inflations.

Theorem 18. *Let $k \geq 3$ and let G be a k -connected graph with girth at least k . Then every $H \in I_P(G)$ is cyclically k -edge-connected.*

Proof. Let $H \in I_P(G)$. For each vertex $v \in V(G)$, denote the unique cyclic part in H corresponding to it by P_v . We say that a cycle C in H is *traversal* if there are two distinct vertices v_1 and v_2 in G with $V(P_{v_1}) \cap V(C) \neq \emptyset$, for both $v \in \{v_1, v_2\}$, so it goes through two unique cyclic parts in H . Otherwise, we say that C is *non-traversal*, meaning $C \subseteq P_v$ for some $v \in V(G)$.

Suppose to the contrary that H is not cyclically k -edge-connected. Then H has a cycle-separating edge-cut S with $|S| \leq k - 1$, such that $H - S$ has precisely two cyclic components D_1 and D_2 .

Let D_i^G be the subgraph of G induced by the vertex set

$$\{x \in V(G) \mid V(P_x) \cap V(D_i) \neq \emptyset\},$$

so vertices from G , whose corresponding cyclic parts in H have at least one vertex overlapping with the component D_i . Let us define sets S_V and S_E as

$$\begin{aligned} S_V &= \{x \in V(G) \mid E(P_x) \cap S \neq \emptyset\}, \\ S_E &= \{e \in S \mid \forall x \in V(G) : e \notin E(P_x)\}. \end{aligned}$$

This means that S_V are vertices from G , whose corresponding cyclic parts contain at least one edge from S , and S_E are edges from S which are in no cyclic part, or in other words, connect two vertices in G and therefore connect two cyclic parts in H . It can be seen that $V(D_1^G) \cap V(D_2^G) = S_V$, since the set contains vertices, whose corresponding cyclic parts contain an edge from S , thus are split into the two components.

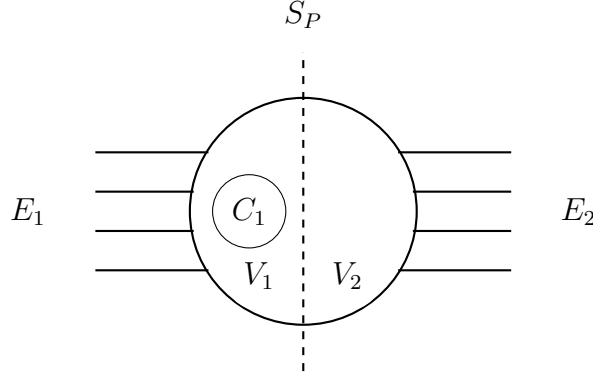
Since S_V contains vertices and S_E contains edges, these sets are naturally disjoint, leading to $|S_V \cup S_E| = |S_V| + |S_E|$. To prove that $|S_V \cup S_E| \leq |S|$, we will show that each element in $S_V \cup S_E$ corresponds to at least one unique representative in S . This correspondence can be understood as an injective function from $S_V \cup S_E$ to S , which we demonstrate through two cases for an element x from $S_V \cup S_E$. When multiple representatives exist for an element, to maintain the formalization of an injective function, we can select just one of the representatives. The crucial point is that these representatives do not overlap:

1. When x is a vertex ($x \in S_V$): The representatives in S are the edges in P_x that appear in the edge-cut S . These representatives cannot overlap with those of any other element from $S_V \cup S_E$: For a different vertex $y \in S_V$, its edges in P_y will be distinct from those in P_x . Similarly, for edges in S_E , the representatives remain distinct, as explained in the next case.
2. When x is an edge ($x \in S_E$): These are edges from S that don't belong to any cyclic part. For clarity, let us denote such edge by e . Its representative in S is simply the edge itself. Again, this representative cannot overlap with those of other elements in $S_V \cup S_E$: for vertices in S_V , their representatives are edges from their corresponding cyclic part, which don't include e , and for other edges in S_E , their representatives are edges distinct from e .

This means, that $|S_V \cup S_E| \leq |S|$, which means after putting everything together that

$$|S_V| + |S_E| = |S_V \cup S_E| \leq |S| \leq k - 1.$$

Suppose that $V(D_i^G) - S_V \neq \emptyset$ for both $i \in \{1, 2\}$, meaning there is a whole cyclic part in both D_1 and D_2 . We prove that by removing vertices S_V and edges S_E from G we obtain a disconnected graph. Let P_1 and P_2 be the whole cyclic parts in D_1 and D_2 respectively. Let v_1 and v_2 be their corresponding vertices in G respectively. Suppose there is a path between v_1 and v_2 in $G - S_V - S_E$. This path goes through edges not in S_E and vertices not in S_V . This would however mean that there is path in $H - S$ between a vertex in P_1 and a vertex in P_2 : it would go through the uncut edges not in S_E and cyclic parts whose corresponding vertices in G are not in S_V , leading to a contradiction.

Figure 4.1: Cyclic part P_v

Now suppose $G - S_V - S_E$ yields two components F_1^G and F_2^G . The number of vertex disjoint paths from a vertex of F_1^G to a vertex of F_2^G in G is at most $|S_V| + |S_E| \leq k - 1$, since each path must either go through a vertex in S_V , or edge in S_E , which contradicts by Menger's Theorem that G is k -connected.

Therefore we may assume without loss of generality that $V(D_1^G) - S_V = \emptyset$, meaning there is no whole cyclic part in D_1 .

Another thing to note, is that in this case, the component D_2 must contain at least one unsplit cyclic part, or in other words $V(D_2^G) - S_V \neq \emptyset$. Suppose that $V(D_2^G) - S_V = \emptyset$, which would mean that each cyclic part has been split by S . However G is k -connected, thus $|V(G)| > k$ and since S contains an edge from each of the cyclic parts, it must hold that $|S| \geq |V(G)| > k$, which contradicts that $|S| \leq k - 1$.

By assumption D_1 contains a cycle, say C_1 . Suppose that this cycle is non-traversalal, thus $C_1 \subseteq P_v$ for some $v \in V(G)$. It must hold that $v \in S_V$, since there is no unsplit cyclic part in D_1 , meaning P_v has been split by S . Let V_1 and V_2 be the sets of vertices from P_v in D_1 and D_2 respectively. Let $S_P = S \cap E(P_v)$ and let E_1 and E_2 be the sets of semiedges in P_v , which have one edge end incident with a vertex in V_1 or V_2 respectively, as illustrated in Figure 4.1. It is evident, that if P_v is a cyclic m -part, $|E_1| + |E_2| = m$. If $|S_P| < |E_2|$, the cut S_P in P_v would yield a cyclic n -pole, where $n = |E_1| + |S_P| < |E_1| + |E_2| = m$, leading to a contradiction by Lemma 6. Therefore $|E_2| \leq |S_P|$.

Let us denote the sets of edges in H , which were created by the junction of semiedges in E_1 or E_2 with some other semiedges by L_1 and L_2 respectively. It is evident that $|L_i| = |E_i|$ for both $i \in \{1, 2\}$. We can now create an edge-cut $S' = (S - S_P) \cup L_2$ in H . This cut is cycle separating as both components now contain a whole cyclic part, and for its size it holds that $|S'| \leq |S| \leq k - 1$. Since both components now contain a whole cyclic part, this contradicts the assumption that G is k -connected, as it was proved before.

This means, that C_1 must be traversalal. Thus C_1 corresponds to a closed trail in

D_1^G , say C_1^G . Since the girth of G is at least k and every closed trail contains a cycle, we have $|V(C_1^G)| \geq k$. It holds that $V(C_1^G) \subseteq V(D_1^G)$, and because each cyclic part in D_1 has been cut by S , we have $V(C_1^G) \subseteq V(D_1^G) \subseteq S_V$. From $|S_V| + |S_E| \leq k - 1$ it leads to a contradiction with the fact that $|V(C_1^G)| \geq k$, since $V(C_1^G) \subseteq V(D_1^G) \subseteq S_V$ and $|S_V| \leq k - 1$. \square

Although this theorem is useful for determining cyclic edge-connectivity when performing junctions of multiple cyclic parts, it has limitations regarding how these junctions can be performed when viewing the resulting graph as a cyclic part inflation. For instance, we cannot apply this theorem to graphs where two cyclic parts are connected by multiple edges. This is because the original graph would then contain vertices connected by multiple edges, resulting in a girth of at most 2 – less than the required minimum of 3.

To overcome this limitation, we present a new theorem that removes the restriction of having at most one edge between a pair of vertices. Instead, it requires a larger number of edges between pairs of vertices, resulting in a larger number of edges connecting cyclic parts in the resulting cyclic part inflation.

Theorem 19. *Let G be an arbitrary graph on v vertices, such that each pair of vertices is connected by at least e edges, and let k be an integer greater than zero, for which it holds that*

$$k \leq 2e(v - 2).$$

Let H be a cyclic part inflation of G . If $\lambda(G) \geq k$, and for each used cyclic part M it holds that $i(M) \geq \frac{k}{2}$, then H is cyclically k -connected.

Proof. We can write the condition inequality equivalently as $v \geq \frac{k+4e}{2e}$.

Suppose to the contrary that H is not cyclically k -connected, thus contains an independent edge-cut S of size $l < k$, such that l is the minimal possible. Let us denote the two cyclic multipoles yielded by this cut by X and $H - X$. By only severing l edges, which were originally from the former graph G , we do not obtain a disconnected graph, since $\lambda(G) \geq k$, therefore S contains some edges from the cyclic parts. Let M be a cyclic part with some of its edges in the cut S . These edges disconnect M , since it is a minimal cut. If it would not disconnect M , we could leave these edges out and obtain a smaller cut, contradicting the assumption about l . Since the cut is independent, it needs to contain at least $\frac{k}{2}$ edges from M .

If S contained edges from another cyclic part, it would have size at least $2\frac{k}{2} = k$, leading to a contradiction with the assumption that $|S| < k$. Thus the remaining edges in S , besides those from M , must be former edges from G .

Note that $\delta(G) \geq k$, since $\delta(G) \geq \lambda(G) \geq k$. Other than the part of M , there must be some whole cyclic parts in X . Otherwise, X would be whole in M , contradicting the fact that M is a cyclic m -part for $m \geq k$ by Lemma 6.

The maximum number of whole cyclic parts in X must be $v - 2$ though. If all $v - 1$ whole cyclic parts were in X , then $H - X$ would be a cyclic l -pole in M , again leading to a contradiction.

Let n be the number of whole cyclic parts in X . To make the graph disconnected, we need to cut all edges connecting these cyclic parts to the ones not in X . Each is connected to another by at least e edges, thus we need to cut at least $ne(v - n - 1)$ edges. The number $(v - n - 1)$ represents the cyclic parts not in X , and ne since each cyclic part in X needs to be cut from the ones in $H - X$ by at least e edges. We prove that $ne(v - n - 1) \geq \frac{k}{2}$ for each n from 1 to $v - 2$.

By equivalent editing the inequality we obtain

$$\begin{aligned} ne(v - n - 1) &\geq \frac{k}{2} \\ -en^2 + nev - ne &\geq \frac{k}{2} \\ -en^2 + ne(v - 1) - \frac{k}{2} &\geq 0. \end{aligned}$$

This is a quadratic function in n , and since the coefficient of n^2 is negative, the parabola opens downward. Therefore, we only need to verify that the inequality holds at both endpoint values.

For $n = 1$ we obtain $e(v - 2) \geq \frac{k}{2}$. For $n = v - 2$ we obtain $(v - 2)e(v - (v - 2) - 1) \geq \frac{k}{2}$, which can be simplified to the form $(v - 2)e \geq \frac{k}{2}$, which is the same inequality as for $n = 1$. thus we need to prove just this one.

We know that $v \geq \frac{k+4e}{2e}$, thus it holds that $e(v - 2) \geq e(\frac{k+4e}{2e} - 2)$ and the sufficient proof will be $e(\frac{k+4e}{2e} - 2) \geq \frac{k}{2}$. By equivalent editing the inequality we obtain

$$\begin{aligned} e\left(\frac{k+4e}{2e} - 2\right) &\geq \frac{k}{2} \\ 2e\left(\frac{k+4e}{2e} - 2\right) &\geq k \\ k + 4e - 4e &\geq k \\ k &\geq k \end{aligned}$$

which is true, thus the inequality holds. This means, that each independent cut in H must contain at least $\frac{k}{2}$ edges from some cyclic part, and at least $\frac{k}{2}$ edges not in any cyclic part (in other words former edges from G), which is in total at least k . This leads to a contradiction with the fact that the size of the cut S is less than k . \square

This theorem holds even for large v and e , and small k , which might lead to an understanding, that this is not really an interesting theorem. For example, if $v = 10$, $e = 3$ and $k = 6$, we have a graph G with 10 vertices, each pair of them connected

by at least three edges, fulfilling the conditions of this theorem, and we prove that a cyclic part inflation of G is cyclically 6-connected. However, the minimal degree of G is at least $(v - 1)e$, since each vertex is connected to each other by at least e edges, specifically in this case it is equal to 27. Thus the cyclic part inflation of G results from a sequence of partial junction of such cyclic parts, that the smallest possible k , for which they can be a cyclic k -part, is 27. And proving if a junction of multiple cyclic 27-parts is cyclically 6-connected might not seem that interesting.

Ideally, the minimum degree of G should be k . That way, the cyclic part inflation will be a junction of multiple cyclic parts, such that at least one of them will be a cyclic k -part. Then, we can prove that the cyclic k -part used in the cyclic part inflation will stay a cyclic k -part even in the whole graph, or in other words, the cyclic edge-connectivity will be equal to k .

Ideally for this, we would need for k , v and e to fulfill a second inequality, where $k \geq (v - 1)e$, or equivalently $\frac{k+e}{e} \geq v$. This way, we know that the minimum possible degree for such v and e will be at most k , thus by adding some edges to the vertices with lower degree we could obtain a graph with minimum degree equal to k , still satisfying the inequality from Theorem 19. If the minimum degree was less than k , it would not satisfy the condition of being k -edge-connected, thus edges need to be added for this condition to be satisfied.

We prove, that for each k such graph with minimal degree equal to k , and fulfilling the inequality from the conditions of Theorem 19 exists, by proving the following lemmas.

Lemma 14. *For each integer $k \geq 2$, there exist integers e and v satisfying*

$$\frac{k+e}{e} \geq v \geq \frac{k+4e}{2e}.$$

Proof. Let $e = 1$ and $v = k + 1$. Substituting these values into the original inequality yields

$$k + 1 \geq k + 1 \geq \frac{k + 4}{2}.$$

The left inequality is immediately satisfied by equality. For the right inequality, we proceed with the following equivalent transformations:

$$k + 1 \geq \frac{k + 4}{2}$$

$$2k + 2 \geq k + 4$$

$$k + 2 \geq 4$$

$$k \geq 2$$

Since this final inequality aligns with our initial assumption that $k \geq 2$, we have demonstrated the existence of suitable values for e and v that satisfy the required inequalities. \square

This demonstrates, that for any given k , we can find values v and e , such that there exists a graph G with v vertices where each pair of vertices is connected by at least e edges, and the minimum degree equal to k . Furthermore, by Theorem 19, any cyclic part inflation of G will be cyclically k -connected if we use cyclic parts whose smallest independent edge-cuts have size at least $\frac{k}{2}$.

However, as previously noted, the primary advantage of Theorem 19 over Theorem 18 is its ability to allow multiple edges between pairs of vertices in the former graph. While our previous proof used $e = 1$, we now extend this result by proving the following lemma that adds the condition for $e \geq 2$.

Lemma 15. *For each integer $k \geq 4, k \neq 5$, there exist integers $e \geq 2$ and v satisfying*

$$\frac{k+e}{e} \geq v \geq \frac{k+4e}{2e}.$$

Proof. We prove this result by considering even and odd values of k separately.

- For even k , let $e = 2$ and $v = \frac{k+2}{2}$. This yields

$$\frac{k+2}{2} \geq \frac{k+2}{2} \geq \frac{k+8}{4}$$

The left inequality holds by equality. For the right inequality, we obtain by equivalent transformations:

$$\begin{aligned} \frac{k+2}{2} &\geq \frac{k+8}{4} \\ 2k+4 &\geq k+8 \\ k+4 &\geq 8 \\ k &\geq 4 \end{aligned}$$

This inequality is satisfied for all even $k \geq 4$.

- For odd k , let $e = 2$ and $v = \frac{k+1}{2}$. This gives us

$$\frac{k+2}{2} \geq \frac{k+1}{2} \geq \frac{k+8}{4}$$

The left inequality is clearly satisfied. For the right inequality:

$$\begin{aligned} \frac{k+2}{2} &\geq \frac{k+8}{4} \\ 2k+2 &\geq k+8 \\ k+2 &\geq 8 \\ k &\geq 6 \end{aligned}$$

This holds for all odd $k \geq 7$, thus covering all odd values $k \geq 4$ except 5.

□

These lemmas demonstrate that Theorem 19 is not limited to the example in the following section, but possibly applies to infinitely many graphs.

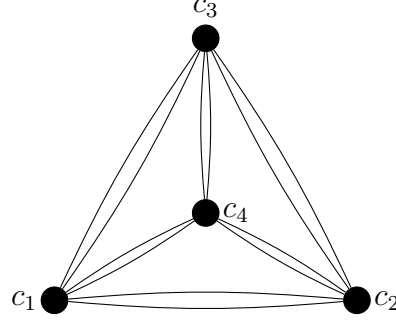
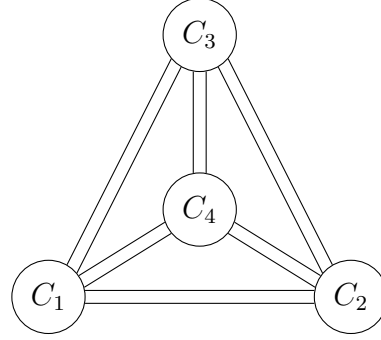
Figure 4.2: Graph G for creating cyclically 6-connected graphs

Figure 4.3: Cyclically 6-edge-connected cyclic part inflation

4.1 Examples

We now proceed to demonstrate practical applications of the theorems presented in this chapter through illustrative examples. Let us begin with an example that showcases the application of Theorem 19.

Example 6. Let G be a graph depicted in Figure 4.2 and let C_1 , C_2 , C_3 and C_4 be cyclic 6-parts, such that $i(C_j) \geq 3$ for each $j \in \{1, 2, 3, 4\}$. Let H be a cyclic part inflation of G , such that each cyclic part C_i corresponds to the vertex c_i in G for each $i \in \{1, 2, 3, 4\}$, as depicted in Figure 4.3.

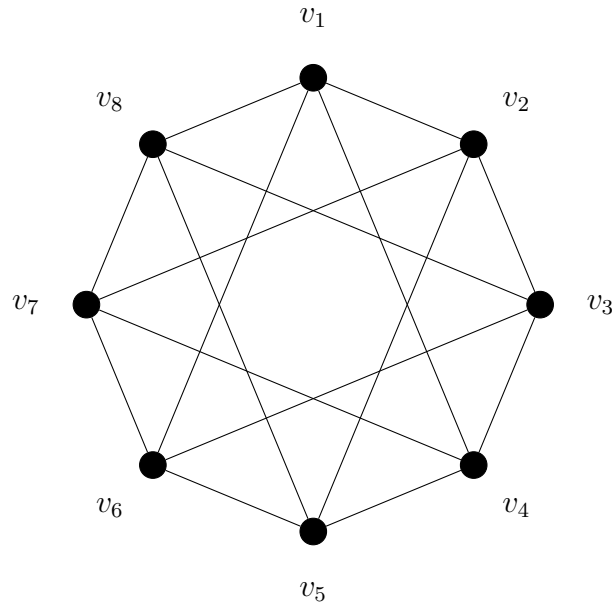
Since the number of vertices (v) is four, each pair of vertices is connected by two edges (e), we can choose $k = 6$ and the inequality

$$k \leq 2e(v - 2)$$

will hold, specifically $6 \leq 4 \cdot 2$. It is evident that $\lambda(G) = 6$ and the condition on the sizes of minimal independent edge-cuts in each cyclic part is fulfilled. This means that by Theorem 19, the graph H is cyclically 6-connected.

This example emerged from a specific challenge encountered during the development of this thesis, specifically when investigating the necessary conditions for a graph H to achieve a cyclic edge-connectivity equal to 6.

Now we present an example showcasing the use of Theorem 18.

Figure 4.4: Graph G in Example 7

Example 7. Let G be the graph depicted in Figure 4.4. It can be verified that $\delta(G) = 4$, $\kappa(G) = 4$ and $g(G) = 4$. This means, that each cyclic part inflation H of G is cyclically 4-connected, and since $\delta(G) = 4$, it holds that $\zeta(H) = 4$.

Chapter 5

Problems for further research

In this chapter, we present problems encountered during our research, including unresolved problems and questions that emerged later. Some of these issues arose from examples in Section 3.4, where we had to manually verify certain properties of multipoles, such as the size of a smallest independent edge-cut, when we could not apply our theoretical framework.

The first problem concerns the independent cut-semiedge index of a multipole created through a partial junction. While we have developed a theoretical framework for establishing lower bounds on the size of smallest independent edge-cuts in multipoles formed by partial junctions, we lack similar results applicable for this index.

Problem 2. Let $M = M_1 * M_2$ be a multipole. Given the values $I_S(M_1, X_1)$ and $I_S(M_2, X_2)$, and considering another subset of semiedges $X \subseteq S(M)$, what is a lower bound of $I_S(M, X)$?

This property presents unique challenges, because it depends not only on the multipole itself, but also on which semiedges will be used in future partial junctions. This makes it difficult to establish a lemma similar to Lemma 10, as simply in independent edge-cuts we do not distinguish between semiedges that will remain, and those used in the partial junction.

The following problem also comes from the exploring of the examples. Specifically, we obtained multipoles from Isaacs snarks by removing a pair of vertices. Despite knowing their cyclic edge-connectivity, we could not formally prove that multipoles obtained by removing these vertices were potentially cyclically k -connected without manual verification. A theoretical framework for deriving these properties from the original graph would be valuable.

Problem 3. Let G be a cyclically k -connected graph, such that $i(G) \geq l$ for some l and k . For a multipole M obtained from G by removing vertices and/or severing edges, what conditions ensure that M is potentially cyclically k -connected, and what determines a lower bound of $i(M)$?

These conditions might be specific rather than general, for instance applying only to the removal of two vertices with a known distance between them.

Other than by a theoretical approach, it would be highly valuable to develop an algorithmic approach for determining whether a multipole is potentially cyclically k -connected.

Problem 4. Develop an efficient algorithm that determines whether a given multipole M is potentially cyclically k -connected for a given integer k .

If M is a graph, the problem is equivalent to the problem of determining whether M is cyclically k -connected, thus the best approach is the algorithm by Dvořák et al. [5]. A solution for multipoles might involve a modification of this algorithm to handle multipoles and search for cyclic multipoles with the smallest number of semiedges, or by performing a junction of the multipole on input with some other multipole, specific for this problem, and running the original algorithm.

Given the fundamental importance of the question whether a multipole is potentially cyclically k -connected in our work, developing an efficient method to determine this property would be particularly valuable.

Regarding inflations, we propose a modified approach, where not every vertex needs to be inflated by a cyclic part – some can remain as single vertices. It would be valuable to develop theorems for this modified approach.

Problem 5. Consider a generalized form of cyclic part inflations, where H is a cyclic part inflation of G , if H results from partial junctions of cyclic parts and multipoles consisting of a single vertex with three semiedges, and G can be obtained by contracting the sets of vertices corresponding to the used cyclic parts in H . Prove analogous results to Theorems 18 and 19 for this modified approach.

Finally, while not formulated as a specific problem, it would be valuable to modify our theoretical framework to include isolated edges as well. Recall that these are edges, where neither end is incident with a vertex. Due to their unique behaviour during junctions, they were excluded from our framework. However, since isolated edges are commonly accepted in the literature, incorporating them would enhance the completeness and broader applicability of our framework.

Conclusion

In this thesis we have made several contributions to the study of cyclic edge-connectivity in cubic graphs, particularly focusing on constructions from smaller building blocks. Our work has extended existing theoretical frameworks and introduced new approaches for analyzing and constructing graphs with specific cyclic edge-connectivity.

Firstly, we have developed a theoretical framework for determining whether a cyclic multipole is a cyclic part without needing to find the original graph. The main result is formulated in Theorem 13, which establishes that a cyclic multipole is a cyclic part if and only if it does not contain a cyclic multipole with a small number of semiedges. This result provides a more direct method for identifying cyclic parts, eliminating the need for constructing the original graph from which the cyclic part could be obtained.

Secondly, we have established conditions for junctions of multiple multipoles to achieve desired cyclic edge-connectivity. Our theoretical framework, particularly through Theorems 14 to 16, provides specific conditions under which these constructions achieve desired cyclic edge-connectivity. These conditions primarily involve the sizes of smallest independent edge-cuts, and the absence of cyclic multipoles with few semiedges. These results are particularly valuable, as they apply to arbitrary values of cyclic edge-connectivity, extending beyond the specific cases studied in previous research. In addition to these theorems, we have proven results regarding modifications of multipoles, for example by edge subdivisions.

Thirdly, we have introduced and analyzed the concept of cyclic part inflations, extending the traditional notion of inflations. Through Theorems 18 and 19, we have established conditions under which these constructions obtain desired cyclic edge-connectivity, providing new methods for constructing cyclically k -connected graphs. This approach offers greater flexibility than traditional inflations, since it allows inflations of the graph not only with k -cycles, but with arbitrary cyclic k -parts.

Our work has also demonstrated practical applications of these theoretical results through several examples, including verifying claims about cyclic edge-connectivity of specific graphs, made in previous publications without formal proofs.

Finally, we have presented open problems that emerged during writing this thesis, suggesting potential directions for future research. These problems emerged primarily while verifying cyclic edge-connectivity in the mentioned examples, where our theoretical

framework lacked methods for certain proofs, requiring manual verification instead.

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