

Exercise 1

We want to implement a gate **mod**[**q**] taking input s of width n and whose action is defined as follows:

$$\mathbf{mod}[\mathbf{q}] = \begin{cases} 0 & \text{if } \text{hw}(s) \equiv 0 \pmod{q}; \\ 1 & \text{otherwise.} \end{cases}$$

The function $\text{hw}(\cdot)$ is the Hamming weight and counts the number of 1s in the input. Such a gate is helpful in an implementation of the recursive Fourier sampling problem (RFS), also called recursive Bernstein-Vazirani problem.

We are given an operator \mathbf{M}_q on $k = \lceil \log_2 q \rceil$ qubits which acts as follows:

$$\mathbf{M}_q|x\rangle = \begin{cases} |x+1 \pmod{q}\rangle & \text{if } x < q; \\ |x\rangle & \text{if } x \geq q. \end{cases}$$

We are interested in the case $q = 3$. We claim that \mathbf{M}_3 with its matrix given below can be used to implement **mod**[**q**] for $q = 3$.

$$\begin{pmatrix} 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix}$$

This matrix looks quite similar to the matrix of the **SWAP** gate. Simply exchange in \mathbf{M}_3 row 0 with row 2 in order to obtain **SWAP**.

- 1.1 An important condition \mathbf{M}_q has to satisfy is $(\mathbf{M}_q)^q = \mathbf{I}$. Show that $(\mathbf{M}_3)^3$ results in the identity.
- 1.2 Implement a gate \mathbf{M}_3 using one **CNOT**⁰ gate and one **SWAP** gate. Recall that **CNOT**⁰ acts when the control qubit is in state $|0\rangle$. Use Qiskit and employ the unitary simulator to obtain the matrix form of your implementation. Is this matrix equal to \mathbf{M}_3 ?
- 1.3 Implement a single-controlled version **cM**₃ of gate \mathbf{M}_3 in Qiskit. Test your implementation and document the test cases.
- 1.4 Implement the **mod**[**q**] gate for $q = 3$ and $n = 4$ input qubits as a gate array in Qiskit.

Start with $k = \lceil \log_2 q \rceil$ ancillary qubits and apply four **cM**₃ gates each controlled by the negation of another single input qubit. The target qubits are the ancillary qubits. Apply then an OR gate with target r exploiting that $a \vee b$ is equivalent to $\neg(\neg a \wedge \neg b)$. Measure the negated r and uncompute.

Test your implementation with all possible 16 values for the input.