

**Describe the data structure(s) you used to store the information in synsets.txt. Why did you make this choice?**

Answer: HashTable.

Using the identity of hashtable's "map keys to values", I was able to use two hashtables to store nouns and synset by their noun ID number, which can be **quickly access** with their ID, for example, get synset using ancestor ID. Moreover, I parse the value(nouns) into strings that separated by comma, which can be easily separated by `String.split()` later in the `printSap` function. There are other data structure can be implemented (like symbol table), because I did not fully utilize hash function from the hashtable due to the fact that I do not really need to.

**Describe the data structure(s) you used to store the information in hypernyms.txt. Why did you make this choice?**

Answer: I did not use any special data structure to store hypernyms. I simply used `Digraph(int)` to store all the hypernyms. Which uses the `Bag<>[]` to store adjacency list of vertex. By using `Bag`, for each vertex, it can store all the edge to it using the id number as reference to vertex, then use the value of hypernyms to add edge to it.

**Describe your algorithm to compute the SAP. What is the worst-case running time as a function of the structure of the graph (height, number of vertices, number of edges, etc)?**

**Best case running time?**

My algorithm assumes there is a root vertex that has no links leaving it (i.e. no ancestors), finds it, and finds the paths from  $v$  and  $w$  to it. As it finds the vertices on the way to the root vertex (using BFS), the algorithm checks to see if the paths from  $v$  to root and  $w$  to root have any shared vertices, Once all the possible scenarios are considered (i.e.  $v$  or  $w$  being the common ancestor), the Ancestor with the smallest distance sum is returned.

Worst case:  $O(E + V + \ln(E))$

Best case:  $O(E + V)$