Pipelined RISC-V Processor

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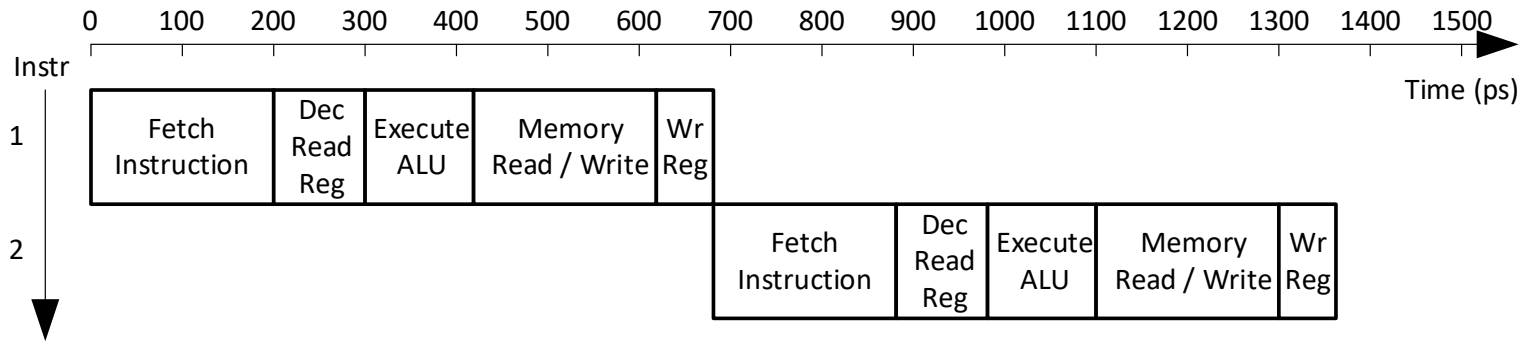
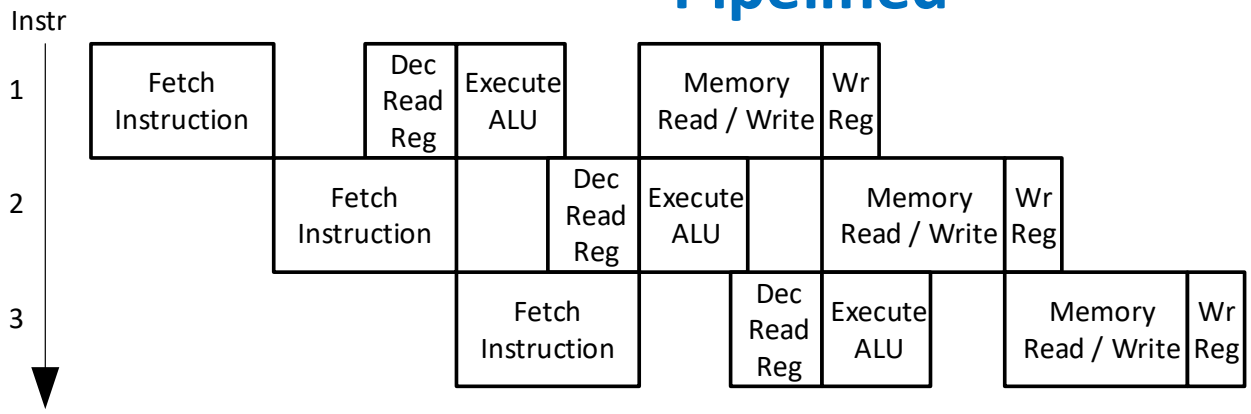
- **Temporal parallelism**
- Divide single-cycle processor into **5 stages**:
  - Fetch
  - Decode
  - Execute
  - Memory
  - Writeback
- Add **pipeline registers** between stages

Single-Cycle vs. Pipelined Processor

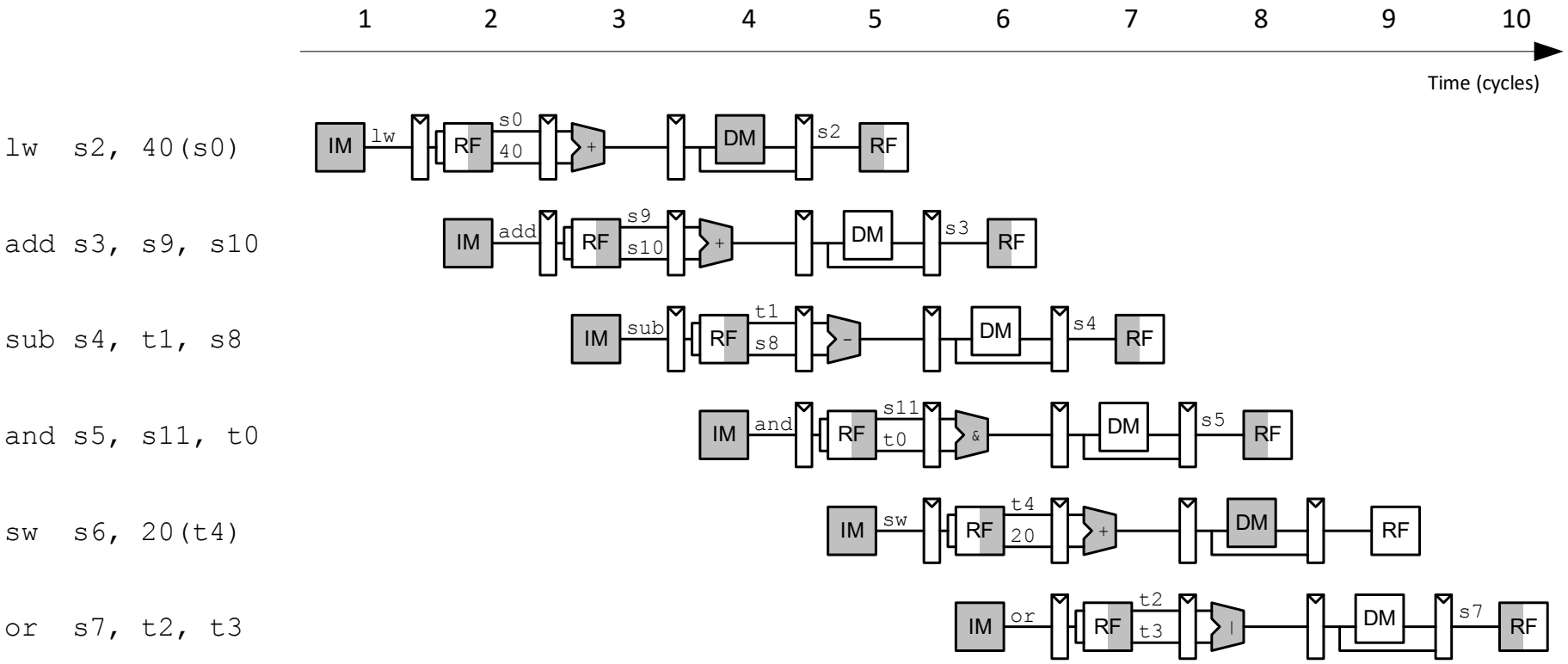
Single-Cycle



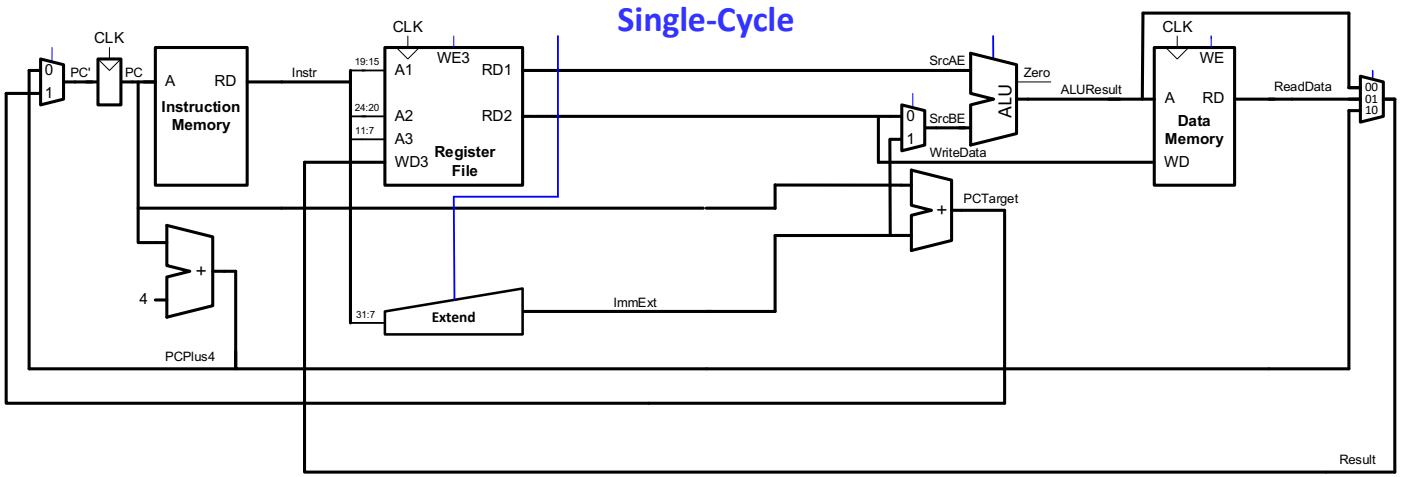
Pipelined

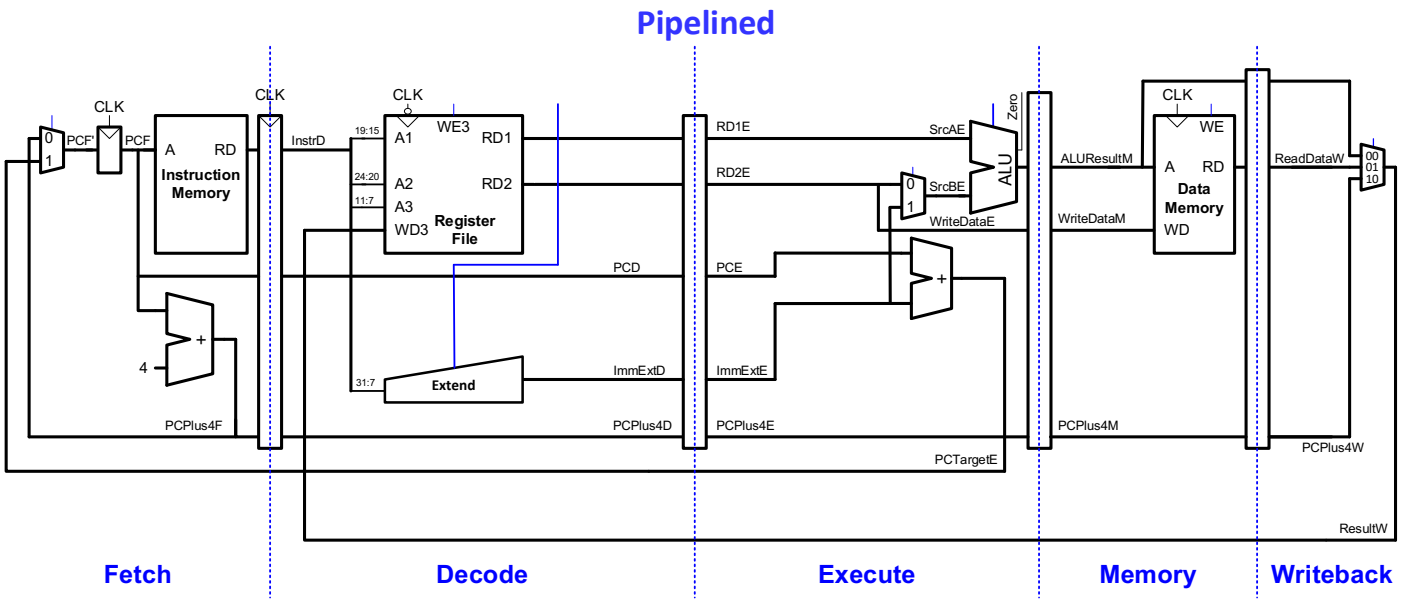


Pipelined Processor Abstraction



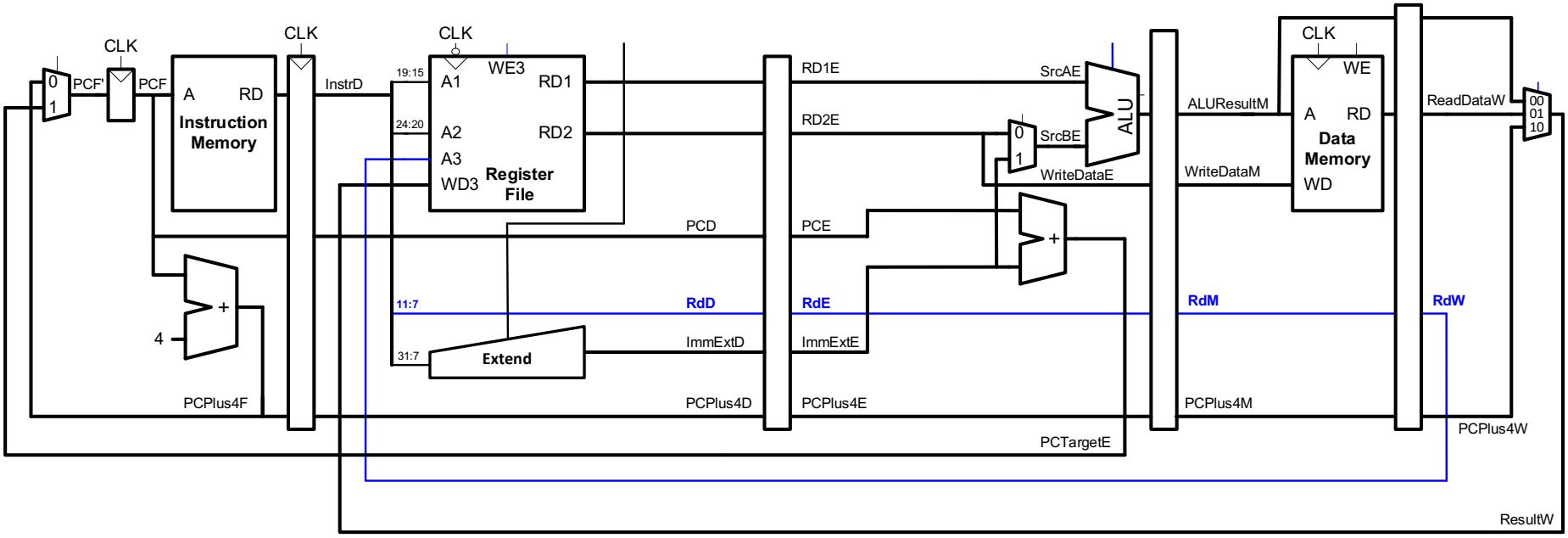
Single-Cycle & Pipelined Datapaths





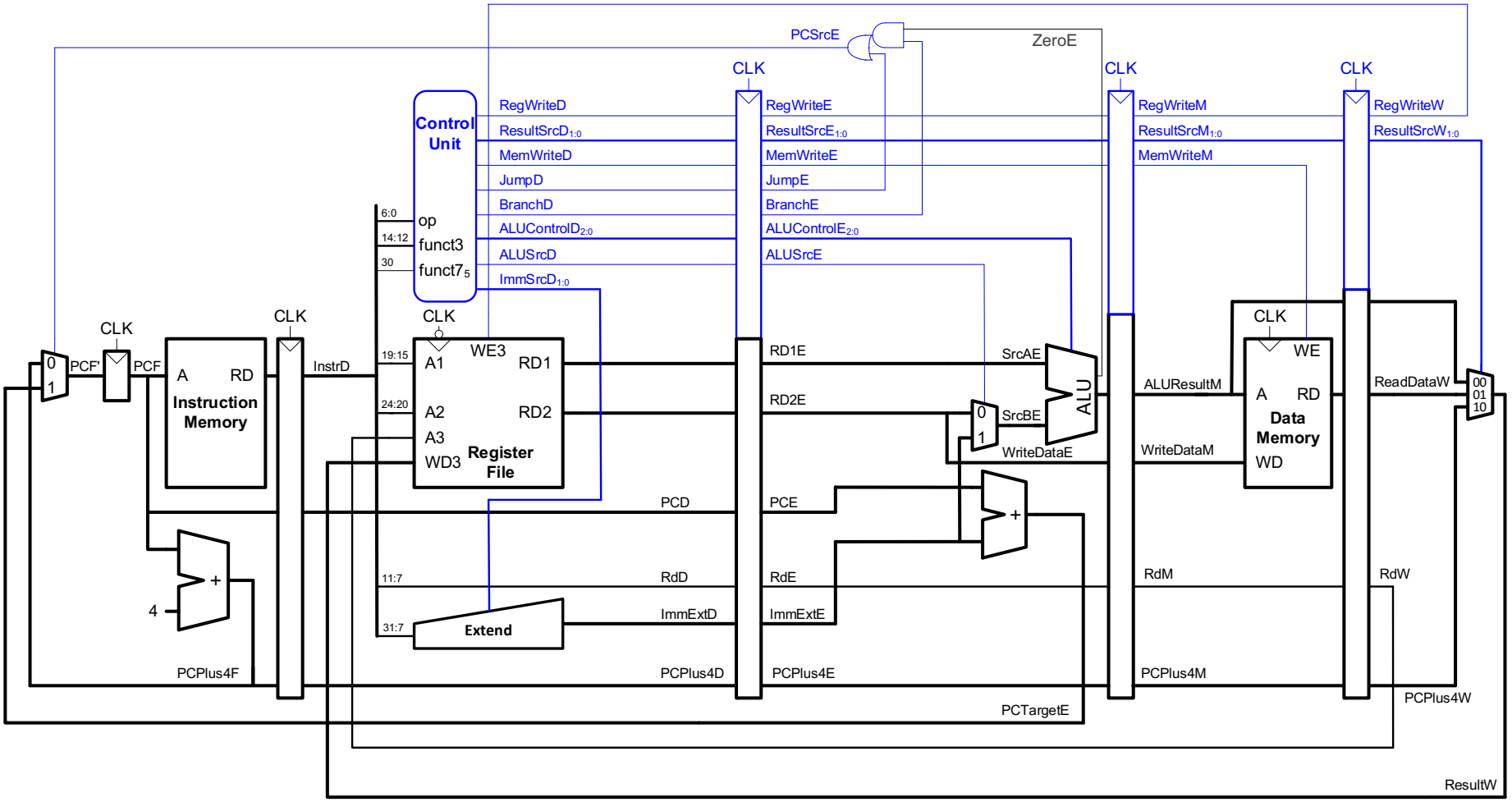
Signals in Pipelined Processor are appended with first letter of stage (i.e., PC**F**, PC**D**, PC**E**).

Corrected Pipelined Datapath



- ***Rd*** must arrive at same time as ***Result***
- Register file written on **falling edge** of *CLK*

Pipelined Processor with Control



- **Same control unit** as single-cycle processor
- **Control signals travel with** the instruction (drop off when used)

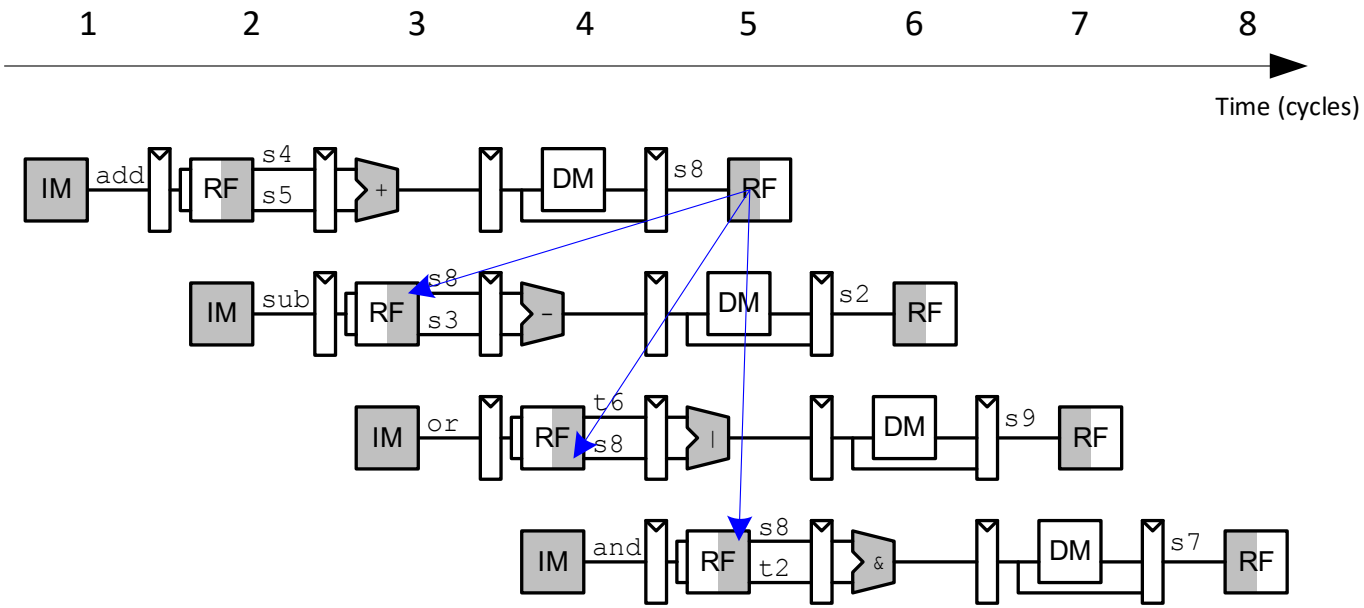
Pipelined Processor Hazards

Pipelined Hazards

- When an instruction depends on result from instruction that hasn't completed
- Types:
  - **Data hazard:** register value not yet written back to register file
  - **Control hazard:** next instruction not decided yet (caused by branch)

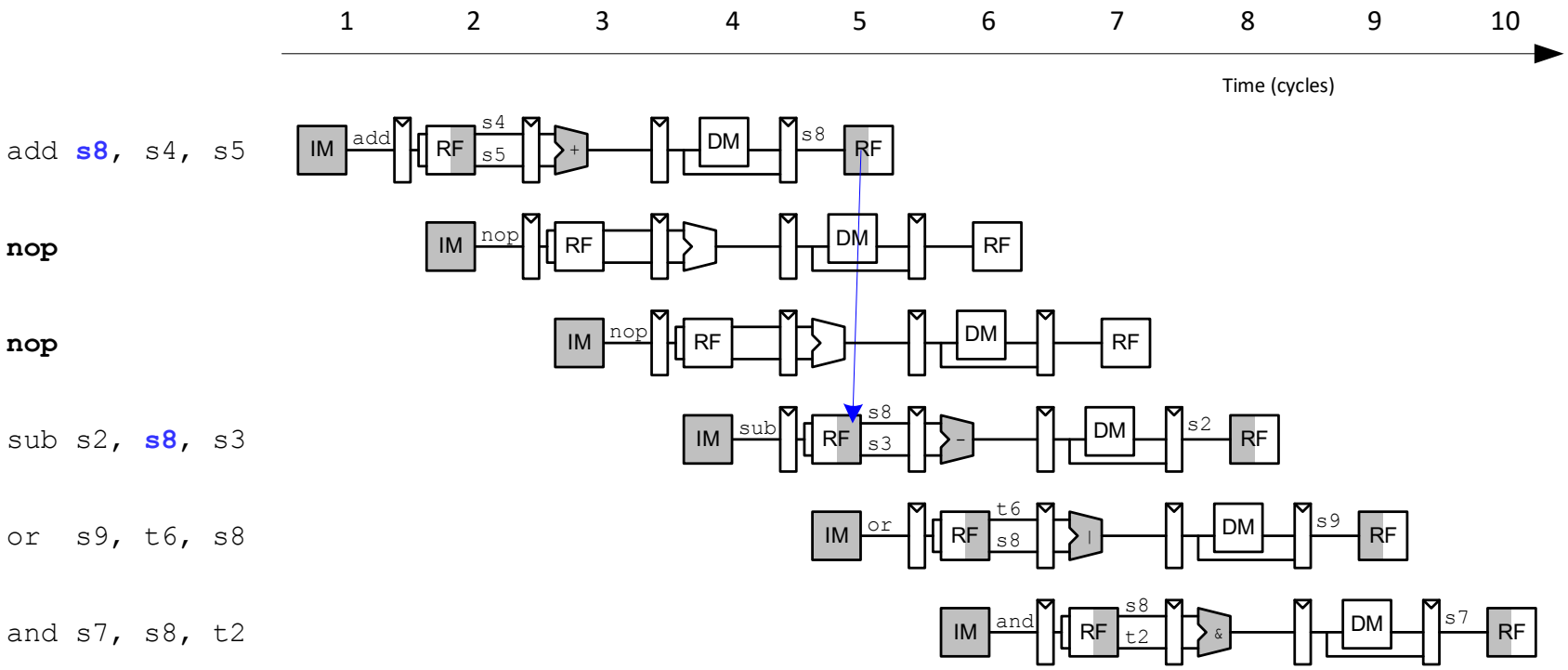
Data Hazard

```
add s8, s4, s5
sub s2, s8, s3
or  s9, t6, s8
and s7, s8, t2
```



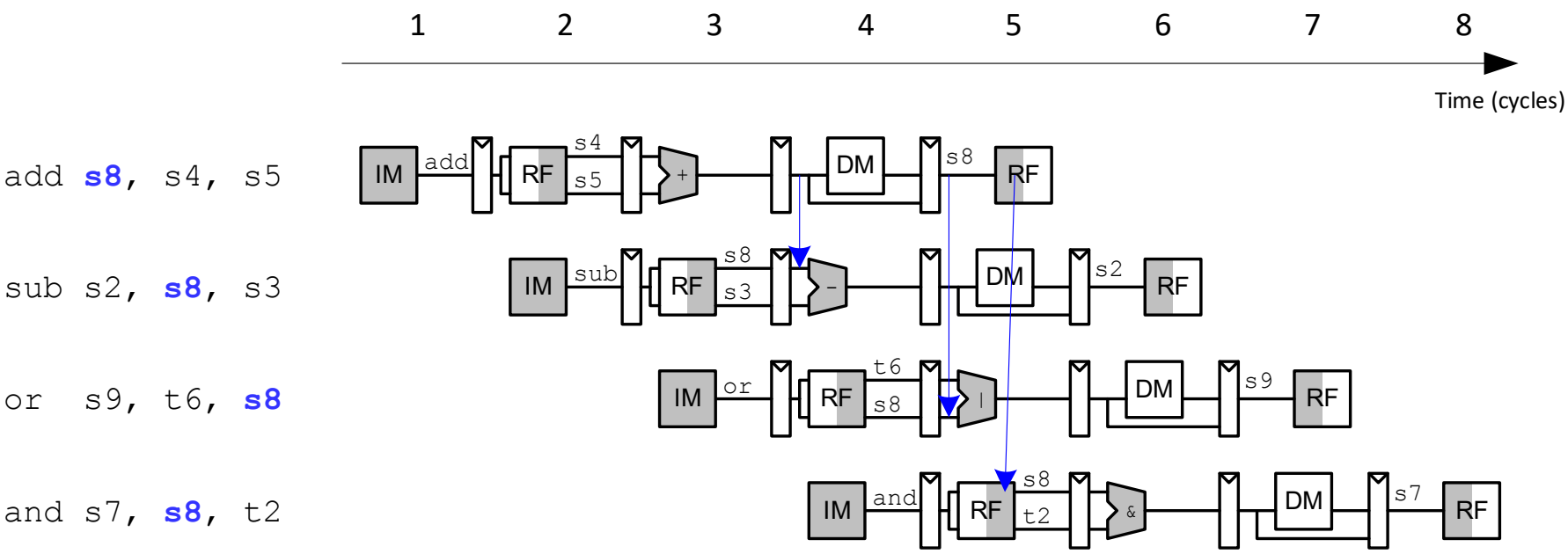
Handling Data Hazards

- **Insert** enough **nops** for result to be ready
- Or move independent useful instructions forward



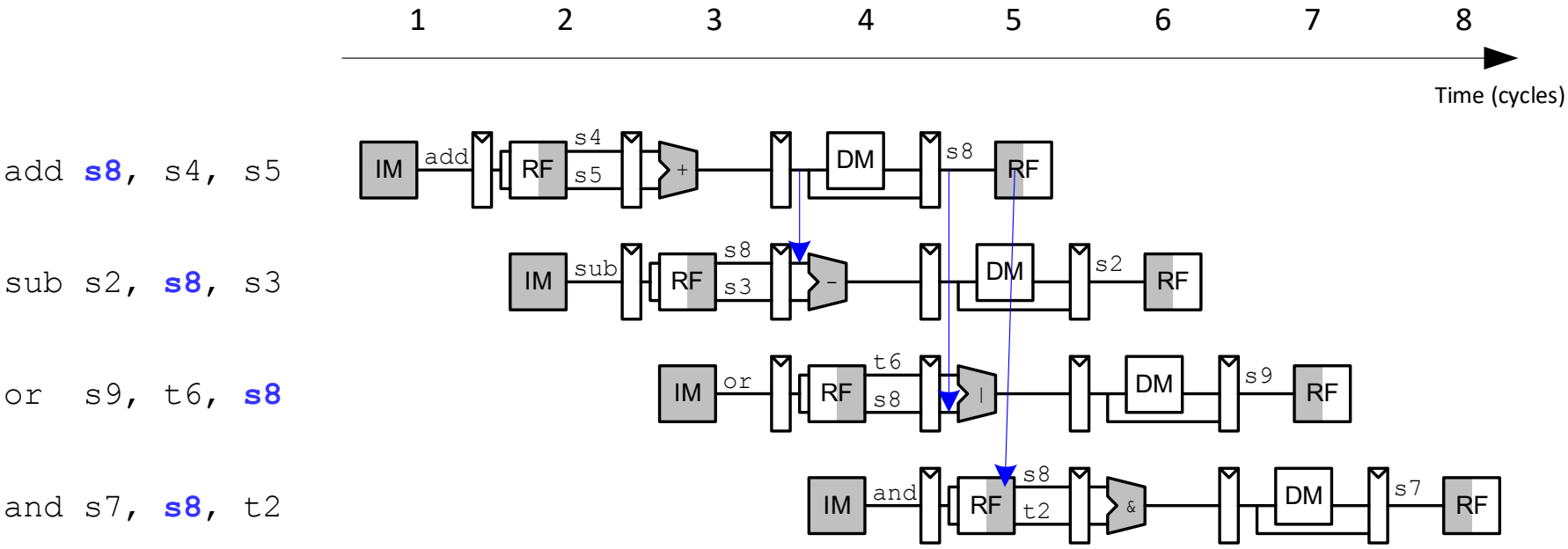
Data Forwarding

- Data is **available on internal busses** before it is written back to the register file (RF).
- **Forward data** from internal busses **to Execute stage**.

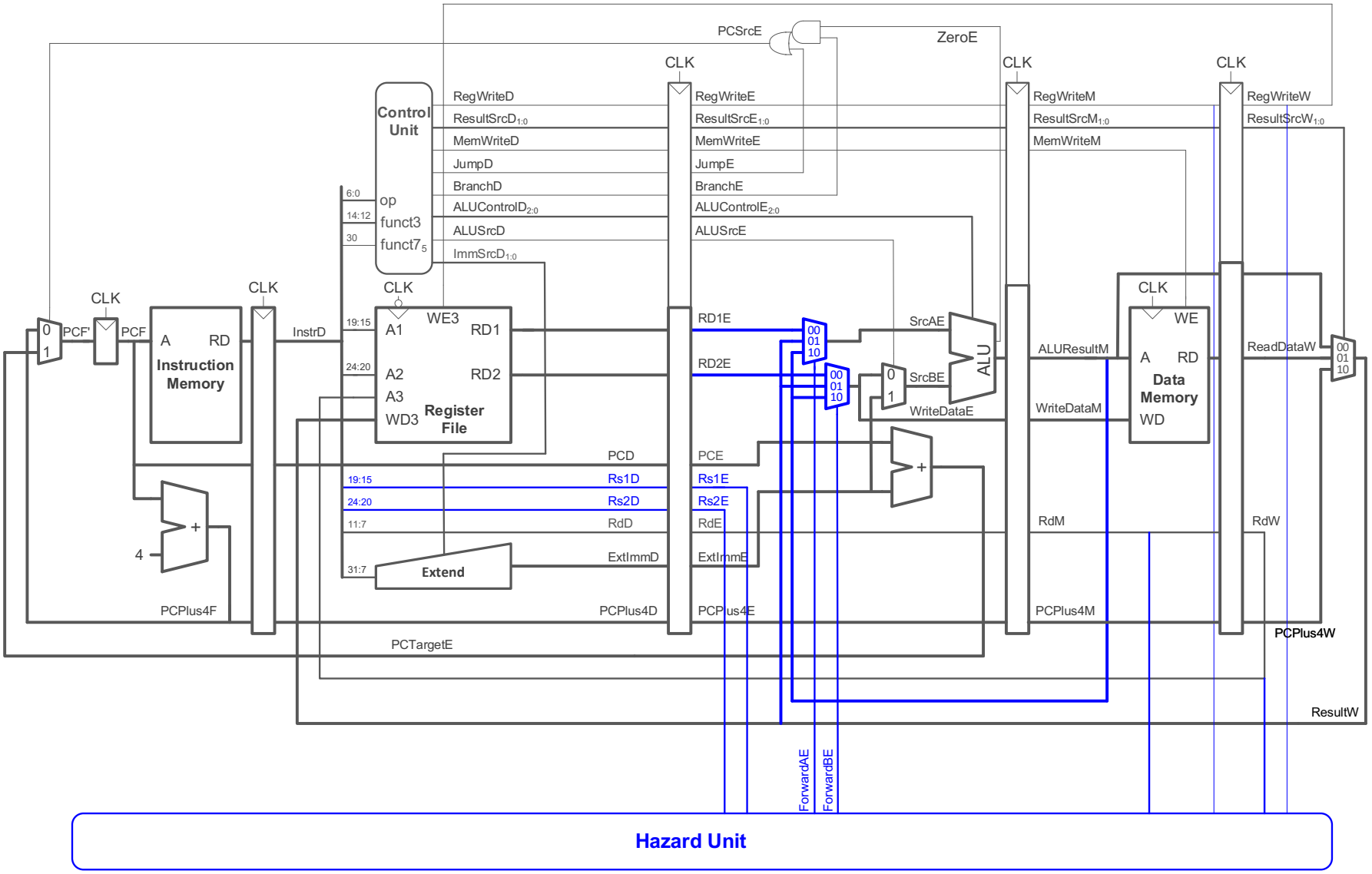


Data Forwarding

- Check if source register **in Execute stage** **matches** destination register of instruction **in Memory or Writeback stage**.
- If so, forward result.



Data Forwarding: Hazard Unit



Data Forwarding

- **Case 1:** **Execute** stage *Rs1* or *Rs2* matches **Memory** stage *Rd*?
  Forward from Memory stage
- **Case 2:** **Execute** stage *Rs1* or *Rs2* matches **Writeback** stage *Rd*?
  Forward from Writeback stage
- **Case 3:** Otherwise use value read from register file (as usual)

Equations for *Rs1*:

if **((*Rs1E* == *RdM*) AND *RegWriteM*)** **// Case 1**
  ForwardAE = 10
else if **((*Rs1E* == *RdW*) AND *RegWriteW*)** **// Case 2**
  ForwardAE = 01
else *ForwardAE* = 00 **// Case 3**

ForwardBE *equations are similar (replace* ***Rs1E*** *with* ***Rs2E****)*

Data Forwarding

- **Case 1:** **Execute** stage *Rs1* or *Rs2* matches **Memory** stage *Rd*?
  Forward from Memory stage
- **Case 2:** **Execute** stage *Rs1* or *Rs2* matches **Writeback** stage *Rd*?
  Forward from Writeback stage
- **Case 3:** Otherwise use value read from register file (as usual)

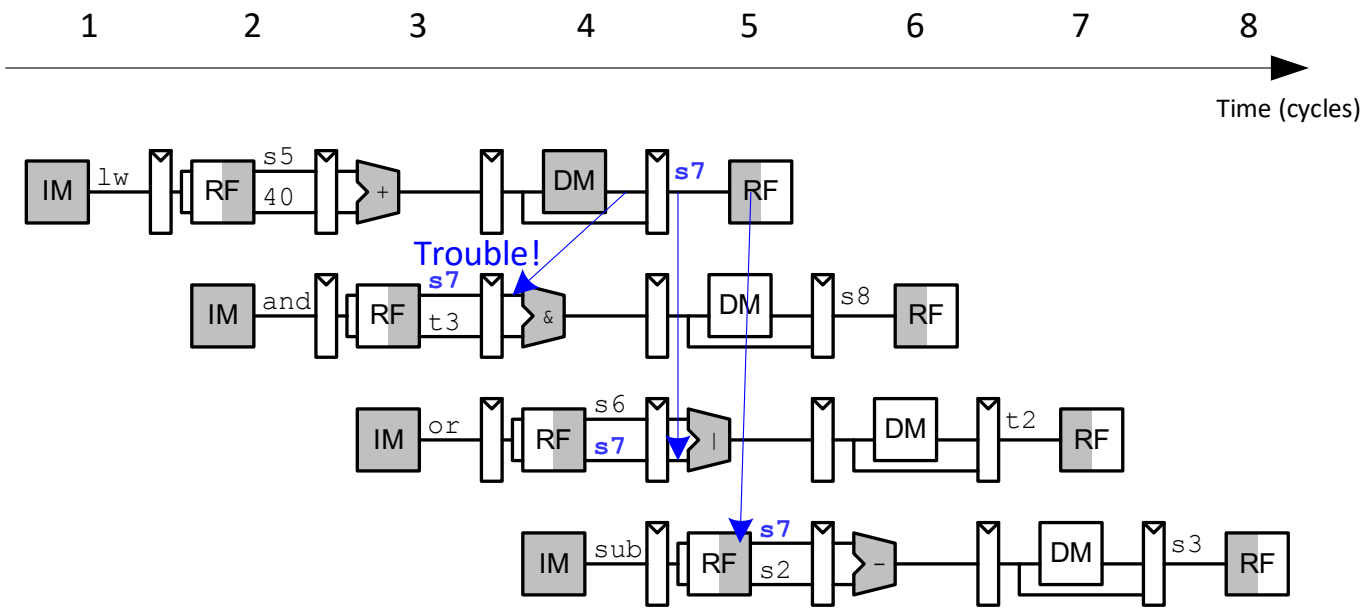
Equations for *Rs1*:

if **((*Rs1E* == *RdM*) AND *RegWriteM*) AND (Rs1E != 0) // Case 1**
  ForwardAE = 10
else if **((*Rs1E* == *RdW*) AND *RegWriteW*) AND (Rs1E != 0) // Case 2**
  ForwardAE = 01
else
  ForwardAE = 00 **// Case 3**

ForwardBE *equations are similar (replace* ***Rs1E*** *with* ***Rs2E****)*

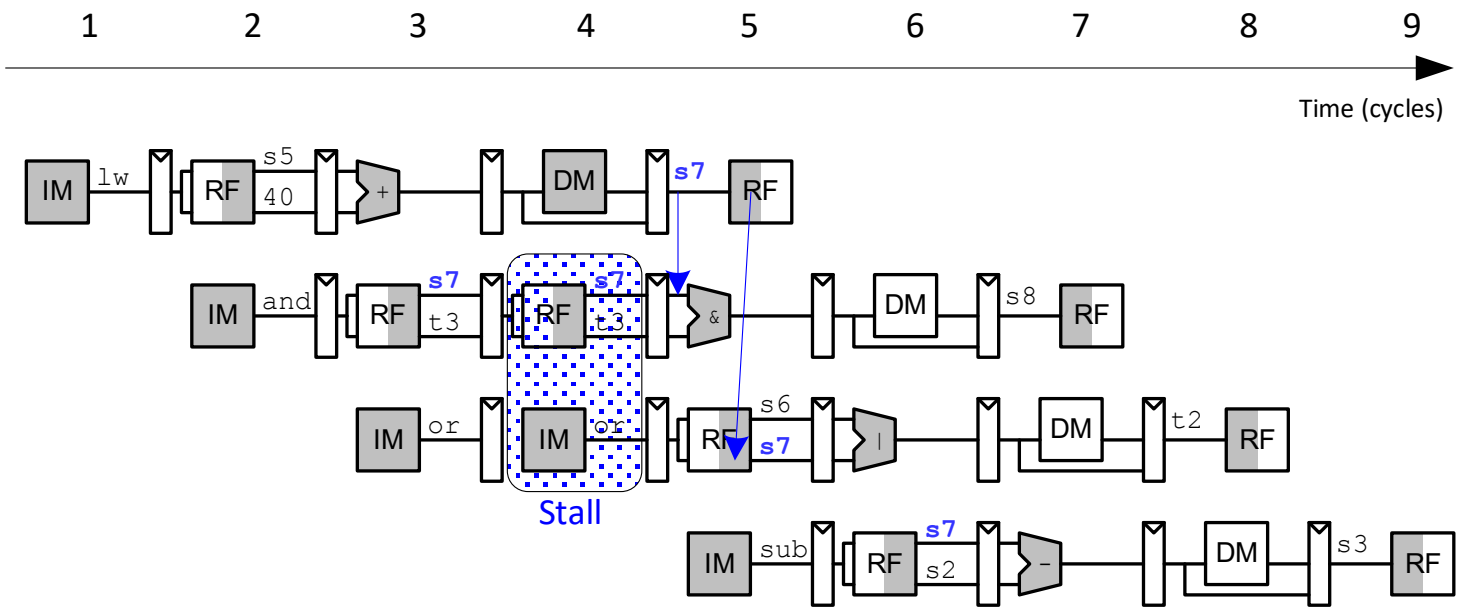
Data Hazard due to lw Dependency

```
lw   s7, 40(s5)
and s8, s7, t3
or   t2, s6, s7
sub s3, s7, s2
```



Stalling to solve `lw` Data Dependency

```
lw   s7, 40(s5)
and s8, s7, t3
or   t2, s6, s7
sub s3, s7, s2
```



Stalling Logic

- Is either **source register in the Decode stage** the same as the **destination register in the Execute stage**?

  AND

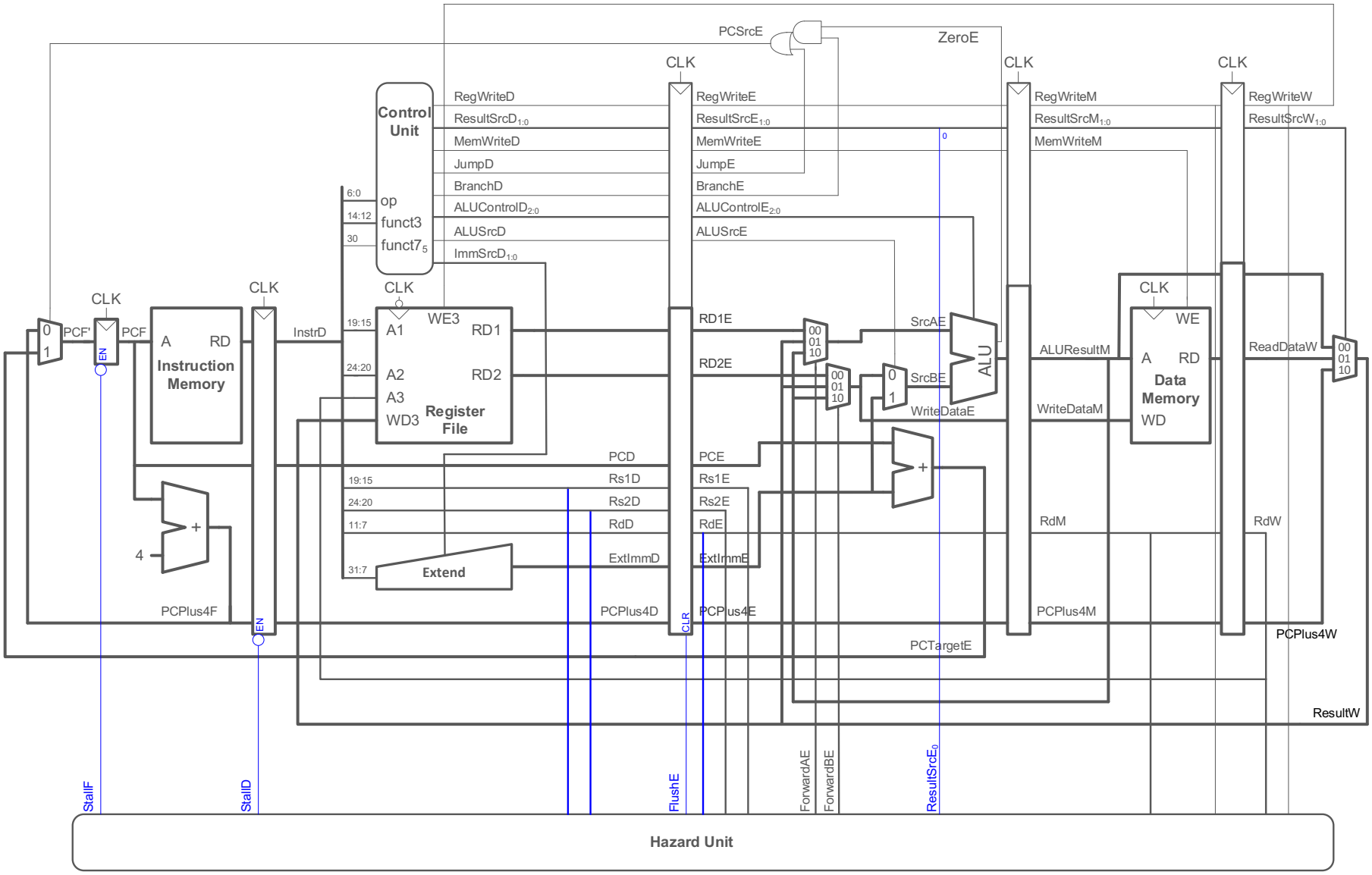
- Is the instruction in the **Execute stage a `lw`**?

lwStall = ((*Rs1D == RdE*) **OR** (*Rs2D == RdE*)) **AND** $ResultSrcE_0$

StallF = ***StallD*** = ***FlushE*** = ***lwStall***

(Stall the Fetch and Decode stages, and flush the Execute stage.)

Stalling Hardware

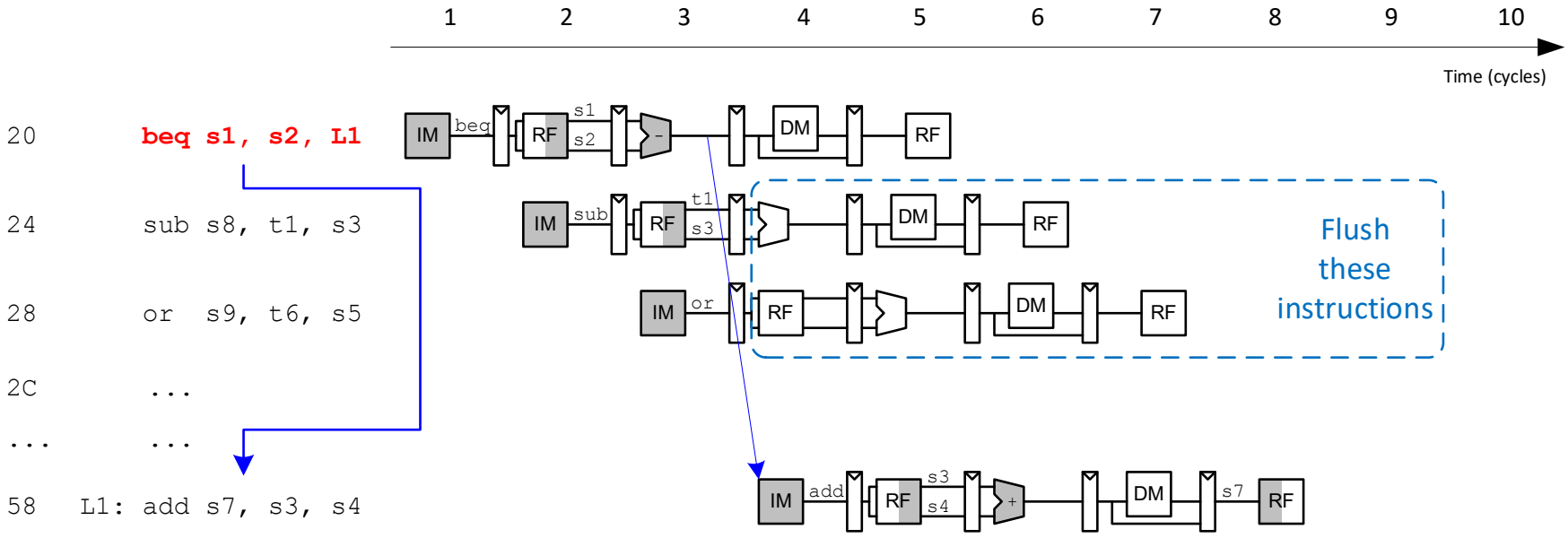


Pipelined Processor Control Hazards

Control Hazards

- **`beq`:**
  - Branch **not determined until the Execute stage** of pipeline
  - **Instructions** after branch **fetched** before branch occurs
  - These **2 instructions must be flushed** if branch happens

Control Hazards



Branch misprediction penalty:

The number of instructions flushed when a branch is taken (in this case, 2 instructions)

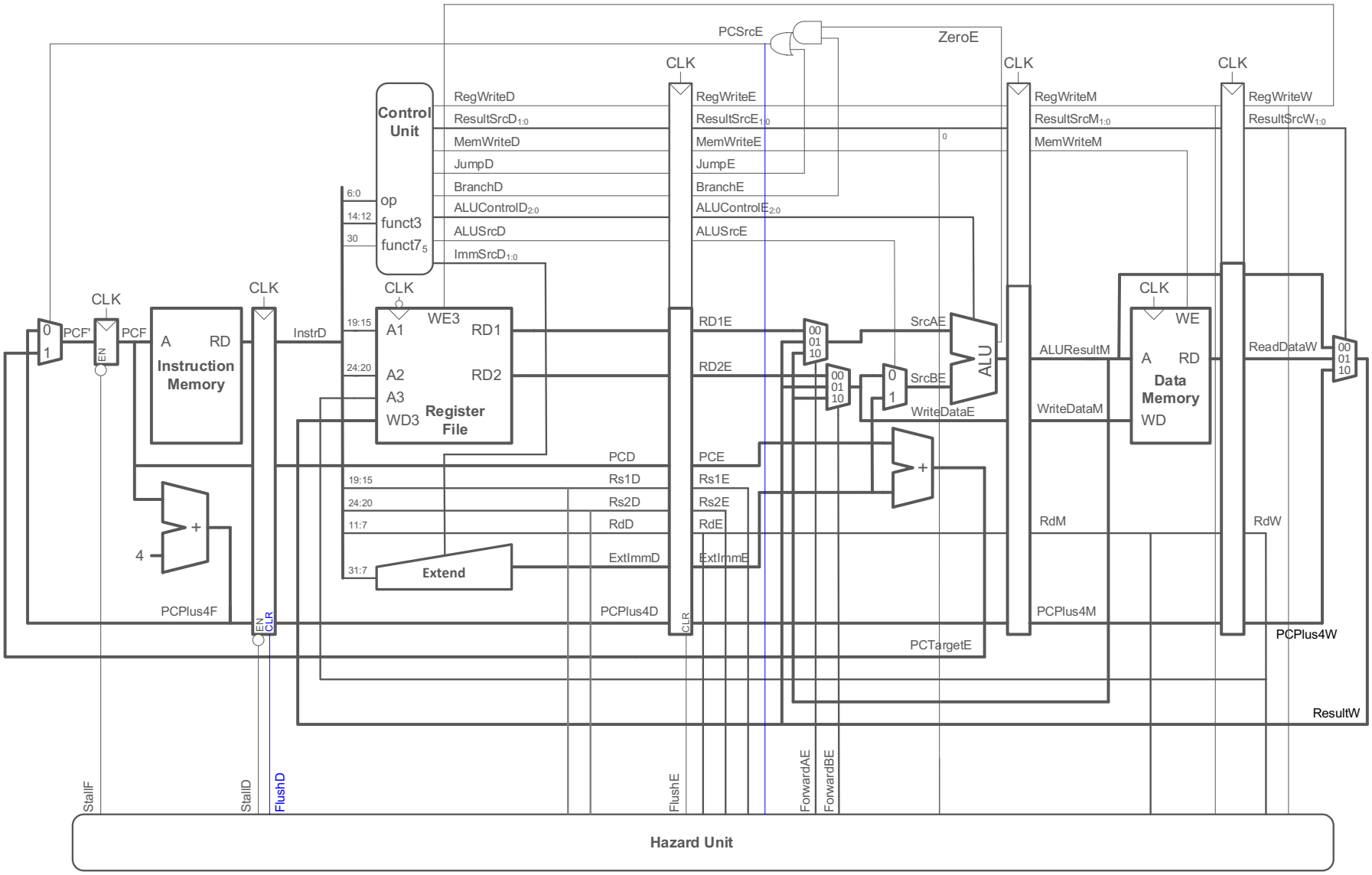
Control Hazards: Flushing Logic

- If branch is taken in execute stage, need to flush the instructions in the Fetch and Decode stages
  - Do this by clearing Decode and Execute Pipeline registers using *FlushD* and *FlushE*
- **Equations:**

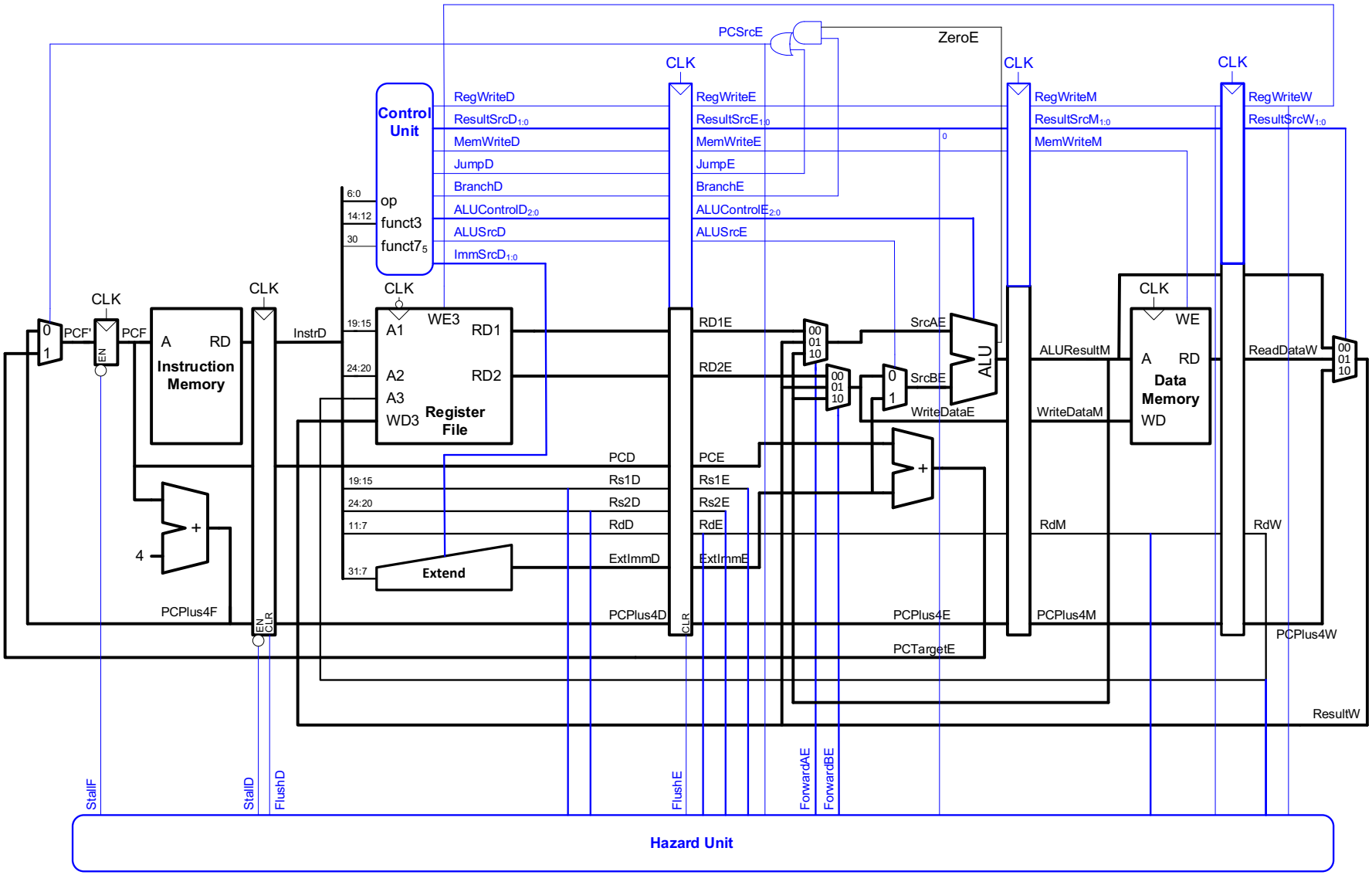
  FlushD = PCSrcE

  FlushE = lwStall OR *PCSrcE*

Control Hazards: Flushing Hardware



RISC-V Pipelined Processor with Hazard Unit



Summary of Hazard Logic

Data hazard logic (shown for SrcA of ALU):

if ((*Rs1E* == *RdM*) AND *RegWriteM*) AND (Rs1E != 0) // Case 1
ForwardAE = 10
else if ((*Rs1E* == *RdW*) AND *RegWriteW*) AND (Rs1E != 0) // Case 2
ForwardAE = 01
else ***ForwardAE*** = 00 // Case 3

Load word stall logic:

lwStall = ((*Rs1D* == *RdE*) OR (*Rs2D* == *RdE*)) AND $ResultSrcE_0$
StallF = ***StallD*** = *lwStall*

Control hazard flush:

FlushD = *PCSrcE*
FlushE = *lwStall* OR *PCSrcE*