

Review

1. Logical Equivalence

$A \equiv B$ means A and B has the same truth value for every truth assignment.

Prove: $A \equiv B$

- By definition, $A^{-1}(T) = B^{-1}(T)$. i.e. $A^{-1}(T) \subseteq B^{-1}(T)$ and $B^{-1}(T) \subseteq A^{-1}(T)$
- $A \leftrightarrow B \equiv T$, $A \rightarrow B \equiv B \rightarrow A \equiv T$, $A \Rightarrow B$ and $B \Rightarrow A$
- Laws. e.g. $A \rightarrow B \equiv \neg A \vee B$

2. Tautological Implications

$A \Rightarrow B$ means every truth assignment that causes A to be true causes B to be true.

Prove: $A \Rightarrow B$

- By definition, $A^{-1}(T) \subseteq B^{-1}(T)$ or $B^{-1}(F) \subseteq A^{-1}(F)$
- $A \rightarrow B \equiv T$ or $\neg B \rightarrow \neg A \equiv T$
- $A \wedge \neg B \equiv F$
- Laws. e.g. $P \wedge Q \Rightarrow P$

3. Propositional function, Predicate logic

(i) Translation

\forall uses \rightarrow , \exists uses \wedge . Why?

- On empty domain, $\forall x P(x) \equiv T$, $\exists x P(x) \equiv F$
- If $\forall x (P(x) \wedge \dots)$, if x not in domain, making $P(x)$ false, the whole formula is directly F .
- If $\exists x (P(x) \rightarrow \dots)$, if x not in domain, making $P(x)$ false, the whole formula is directly T .

(ii) Proof of \equiv and \Rightarrow

4. Graph definitions

Homework 10

1.

$$(a) \forall x(P(x) \rightarrow \exists y(P(y) \wedge \neg E(x, y) \wedge L(x, y)))$$

$$(b) \exists x(P(x) \wedge \forall y(P(y) \wedge \neg E(x, y) \rightarrow L(y, x)))$$

5. Proof: $\exists x(P(x) \vee Q(x)) \equiv \exists xP(x) \vee \exists xQ(x)$.

(i) Show that $\exists x(P(x) \vee Q(x)) \Rightarrow \exists xP(x) \vee \exists xQ(x)$.

Suppose that $\exists x(P(x) \vee Q(x))$ is T in an interpretation I .

- There is an x_0 s.t. $P(x_0) \vee Q(x_0)$ is T in I .
- At least one of $P(x_0)$ and $Q(x_0)$ is T in I . WLOG, suppose it's $P(x_0)$.
- $\exists xP(x)$ is T in I .
- $\exists xP(x) \vee \exists xQ(x)$ is T in I .

(ii) Show that $\exists x(P(x) \vee Q(x)) \Leftarrow \exists xP(x) \vee \exists xQ(x)$.

Suppose that $\exists xP(x) \vee \exists xQ(x)$ is T in an interpretation I .

- At least one of $\exists xP(x)$ and $\exists xQ(x)$ is T in I . WLOG, suppose it's $\exists xP(x)$.
- There is an x_0 s.t. $P(x_0)$ is T in I
- $P(x_0) \vee Q(x_0)$ is T in I .
- $\exists x(P(x) \vee Q(x))$ is T in I .

6. Proof: $\forall x(P(x) \rightarrow Q(x)) \Rightarrow \forall xP(x) \rightarrow \forall xQ(x)$

Suppose that $\forall xP(x) \rightarrow \forall xQ(x)$ is F in an interpretation I .

- $\forall xP(x)$ is T and $\forall xQ(x)$ is F in I .
- $P(x)$ is T for every x in I and there is an x_0 s.t. $Q(x_0)$ is F in I .
- $P(x_0) \rightarrow Q(x_0)$ is F in I .
- $\forall x(P(x) \rightarrow Q(x))$ is F in I .