

## Security of embedded devices:

In use of embedded devices, we often want an owner to control a device. For example, your smartphone might be the “owner” for your DVD player. Commands from other parties (including previous owners) are to be ignored by the device.

(a) Explain how such a device might be programmed to recognize commands from its owner, and only its owner. Show how this can be done with public-key technology.

*Solution: (There can be more than one solution for this problem)*

*Let the owner be called A and the device C*

- 1. The device should have Internet access and some trusted authorities (CA's) public signature embedded within.*
- 2. When A passes any command, it should send it as follows:  
A ----> C: R, CMD, Signature(hash(CMD)), A's Certificate  
(CMD: command; R is a monotonically increasing nonce)*
- 3. C should verify the certificate using the CA's public key. Given this, the signature of the hash of the command indicates that this message indeed came from A.*
- 4. R ensures that this is not a replay. Since it is monotonically increasing, the device C needs to remember only the last nonce seen (not all of them).*

(b) Explain how such a device might be programmed to recognize commands from its owner, and only its owner. Show how this can be done with secret-key technology.

*Solution: It is possible to solve this using a trusted third party (KDC) but to keep things simple I am assuming that a well-shielded cable can be used to communicate securely with the device in case of performing certain functions that happen only once in a while.*

- 1. The device has the shared key of the first owner installed using the cable.*
- 2. When A passes any command, it should send it as follows:  
A ----> C: R, CMD, MAC(CMD))  
(CMD: command; R is a monotonically increasing nonce)*
- 3. C should verify the MAC of the command which indicates that this message indeed came from A.*
- 4. R ensures that this is not a replay. Since it is monotonically increasing, the device C needs to remember only the last nonce seen (not all of them).*
- 5. Note the commands do not need to go over the secure cable as such. They can happen in an environment where eavesdropping can happen.*

(c) Describe a protocol explaining how “transfer of ownership” can be accomplished, under the public-key framework of your answer in part (a).

*Solution: Assume the new owner is B. To transfer ownership of device C from owner A to owner B, owner B has to send its certificate to A which A should verify. A should then execute the command “CMD*

*= XFER, B, Public-key-of-B.” This causes the device to replace PK of A with PK of B, effectively revoking the old owner and enabling the new owner. The command is*

*protected by integrity as shown in (a).*

**(d)** Describe a protocol explaining how “transfer of ownership” can be accomplished, under the classical (symmetric-key) framework of your answer in part (b).

*Solution:*

*The new owner B sends over a secure link (could be face-to-face or by other means) to the old owner A the message  $(B, H)$  where  $H = \text{hash}(KBC)$  where  $h$  is a secure hash function and  $KBC$  is the symmetric key that the device will eventually receive.*

*The old owner A sends to the device C the command “ $\text{CMD} = \text{XFER1}, B, H$ ”. This tells the device to only accept the next command over a secure link iff the initiating party knows a pre-image of  $H$ . This prevents other parties from becoming the owner. Note, this command does not have to happen over a physically secure link. The new owner B then connects to the device C over a physically secure link (once again, something like a cable or shielded IR) and sends the command “ $\text{XFER2}, KBC$ .” The device verifies that this key is a pre-image of  $H$ . For future commands over physically insecure links, the device only accepts commands MAC’d with this key. This effectively revokes the old owner’s access.*