Coordinating Response to Fatal Accidents

An Application of Many-to-One Matching

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What are matching problems?

 Given a graph, G, a "matching" is a pairing up of vertices, each pair connected by an edge. Usually, a matching of minimum cost or maximum cardinality is desired.

$$\begin{split} \sum_{e \in \delta(v)} x_e &= 1 \quad \forall v \in V \\ \sum_{e \in E(S)} x_e &\leq \frac{|S| - 1}{2} \ \forall S \subseteq V, S \text{ odd} \\ x_e &\in \{0, 1\} \quad \forall e \in E \end{split}$$

Where $\delta(v)$ denotes the set of edges incident to node v for $\forall v \in V$

Bipartite Matching

- Here, our graph is divided into two parts, X and Y. The "odd set inequalities" become redundant.
 - This is where most applications lie.

$$\sum_{e \in E(S)} x_e \le \frac{|S| - 1}{2} \ \forall S \subseteq V, S \text{ odd}$$

One-to-One Matching

Standard bipartite matching problem, where each element of X is matched to exactly one element of Y, and every element of Y is matched to exactly one element of X.

$$\sum_{i \in X} x_{ij} = 1 \ \forall j \in Y$$

$$\sum_{j \in Y} x_{ij} = 1 \ \forall i \in X$$

$$x_{ij} \in \{0,1\} \ \ \forall i \in X, j \in Y$$

k-to-One Matching

Each element of X is matched to exactly k elements of Y, and each element of Y is matched to exactly one element of X.

$$\sum_{i \in X} x_{ij} = 1 \ \forall j \in Y$$

$$\sum_{j \in Y} x_{ij} = k \ \forall i \in X$$

$$x_{ij} \in \{0,1\} \ \ \forall i \in X, j \in Y$$

Many-to-One Matching

 Each element of X is matched to at least one element of Y, and each element of Y is matched to exactly one element of X.

$$\sum_{i \in X} x_{ij} = 1 \ \forall j \in Y$$

$$\sum_{j \in Y} x_{ij} \ge 1 \ \forall i \in X$$

$$x_{ij} \in \{0,1\} \ \ \forall i \in X, j \in Y$$

Many-to-Many Matching

 Each element of X is matched to at least one element of Y, and each element of Y is matched to at least one element of X.

$$\sum_{i \in X} x_{ij} \ge 1 \ \forall j \in Y$$

$$\sum_{j \in Y} x_{ij} \ge 1 \ \forall i \in X$$

$$x_{ij} \in \{0,1\} \ \ \forall i \in X, j \in Y$$

Matching Problems are Integer Programs

- In particular, matching problems are binary programs.
- There are polynomial-time algorithms for one-to-one matchings on general graphs.
 - » However, the decision version of k-to-one bipartite matching is NP-Complete.

The College Admissions Problem

- D. Gale and L. S. Shapley
 - » College Admissions and the Stability of Marriage, 1962
- The College Admissions Problem
 - n student applicants
 - m colleges
 - college u_i will accept up to q_i students
- Many-to-One Matching Application
 - Note: Assignment Problem
 - Assign n tasks to m machines each with a specified capacity

The College Admissions Problem

- Preferences
 - » Student applicants rank the colleges in order of preference
 - » Colleges rank the students in order of preference
- Note: For this to work, ignore ties in preference, the "Hungarian admission test" throw away rule, and "couples admission"



- Stable Assignment
 - » No student prefers a college over the one they have been admitted to AND no college prefers a student over the ones admitted
- Optimal Assignment
 - » Stable AND every student is at least "well off" as if the assignment had been different
- "Philosophical Principle"

The College Admissions Problem

The Stable Marriage Problem

- Special case of the College Admissions Problem
 - » Equal number of students and colleges
 - » Each college u_i accepts up to 1 student
 - $\Rightarrow n$ "boys" propose to n "girls"
- Becomes a One-to-One Matching Application
 - » Assignment Problem
 - Assign n tasks to n people

- Consider Preferences
 - "Boys" propose in order of preference
 - "Girls" accept in order of preference
- Is there a stable assignment?



The College Admissions Problem

- Stable Assignments
 - » Gale and Shapley proved that stable marriage assignments exist
 - ⇒ Stable college admission assignments exist
 - The proof "is entirely analogous to the proof given for the marriage problem"

- Optimal Assignments
 - » Optimal College Admission Assignments
 - "deferred acceptance"
- "The main proof is carried out in not in mathematical symbols but in ordinary English... Knowledge of calculus is not presupposed. In fact, one hardly needs to know how to count. Yet any mathematician will immediately recognize the argument as mathematical"

Many-to-One Matching Applications

- Economics
 - "who gets what" especially when "what" is "invisible"
- Dorm room assignments
- Determining the waitlist for organ transplants
 - "pairwise kidney exchange programs"
- Labor Markets
 - » n people applying to m companies with a fixed number of openings
- Coordinating Response to Fatal Accidents

Coordinating Response to Fatal Accidents

- Problem Outline
 - Fatal car accidents require a response from both police and fire department



- Denver Police and Fire
 Departments have separate
 dispatch systems
 - Do not communicate
 - Responders are uninformed of the other entity's likely response time

- Objectives
 - Construct a Many-to-One matching between police and fire stations
 - Generate a list of "Suggested Partnerships"
 - Aid responders by ensuring that they are informed of estimated arrival times of the other entity



Model

Where P := the set of police stations, F := the set of fire stations, C := the set of all car crashes

$$\begin{aligned} &\textit{Minimize} \quad \sum_{i \in P, j \in F} c_{ij} x_{ij} \\ &\textit{Subject To} \quad \sum_{i \in P} x_{ij} = 1 \quad \forall j \in F \\ & \quad \sum_{i \in P} x_{ij} \geq 1 \quad \forall i \in P \\ & \quad x_{ij} \in \{0,1\} \text{ for } \forall i \in P, \forall j \in F \\ &\textit{Where } c_{ij} = \sum_{c \in C} |dist(c,i) - dist(c,j)| \end{aligned}$$

Data

- What we had
 - » Idea + Model Formulation
 - » Denver Open Data Catalog
 - Many data sets
 - Large data sets
 - Data sets with errors
 - » Common Sense
 - What data will work and what will hurt our computers?

- What we needed
 - » Usable data sets for:

P:= the set of police stations

F:= the set of fire stations

C:= the set of all car crashes

- Data to give to Python and AMPL
 - Unique Identifiers
 - Latitude and Longitude

Data

Raw Data Sets (Denver Open Data Catalog)

» Police Stations

- Original: 29 data points
- Includes all police "places"
- Addresses/PO Boxes given

» Fire Stations

- Original: 39 data points
- Addresses given
- Station ID's given

» Traffic Accidents

- Original: 154,967 data points
 - » Five years of all reported traffic incidents
- Latitude and Longitude given
 - » Bonus!
- Preassigned ID's given
 - » Assigned by date/time for all reported traffic incidents

Data

Cleaning Up the Data

- Police and Fire Stations
 - » Only include deployable stations
 - i.e. exclude police bicycle impound and wildfire response stations
 - » PO Box → Physical Address
 - Generate Unique Identifiers

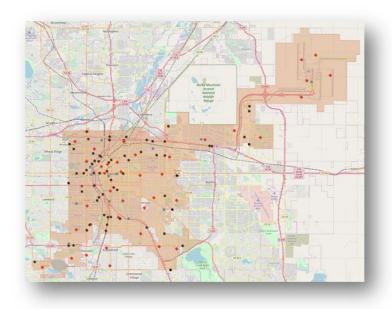
Traffic Accidents

- » Filter data to "Fatal Accidents"
 - » $154,967 \rightarrow 330$ data points
- » Choose a usable time frame
 - Recall: We don't want to hurt our computers
 - 5 years → 15 Months
 - \rightarrow 330 \rightarrow 74 data points

Data

GIS: Filling in the Gaps

- Identify errors in the data
 - Points in "Nowhere Island"?
 - Lat/Lon Missing
 - Points in China?
 - Lat/Lon Inputted Wrong
- Generate latitude and longitude
 - ARCGIS Geometry Toolbox



Finalized Data Sets

- \star *P*:= the set of police stations 14 Data Points
- F:= the set of fire stations 35 Data Points
- C:= the set of all car crashes 70 Data Points

set FIRE: # fire stations

Computer Code - AMPL

```
set POLICE; # police stations
set ACCIDENT; # potentially fatal accidents
param fx {FIRE}; # fire station locations x
param fy {FIRE}; # fire station locations y
param px {POLICE}; # police station locations x
param py {POLICE}; # police station locations y
param ax {ACCIDENT}; # potentially fatal accident
locations x
param ay {ACCIDENT}; # potentially fatal accident
locations y
param incompatibility {i in POLICE, j in FIRE} =
  sum {k in ACCIDENT} <<0; -1, 1>>(<<0; -1, 1>>(px[i] -
ax[k])
              + << 0; -1, 1>> (py[i] - ay[k])
              - << 0; -1, 1>> (fx[i] - ax[k])
              - << 0; -1, 1>> (fy[i] - ay[k])); # how poorly i and
j would work together
```

```
var match {i in POLICE, j in FIRE} binary; # whether i
and j get matched
minimize Total Incompatibility:
  sum {i in POLICE, j in FIRE} incompatibility[i,j] *
match[i,j];
subject to Fire Match{j in FIRE}:
  sum {i in POLICE} match[i,j] == 1;
subject to Police_Match{i in POLICE}:
  sum \{i \text{ in FIRE}\}\ \text{match}[i,i] >= 1:
```

Computer Code - Python

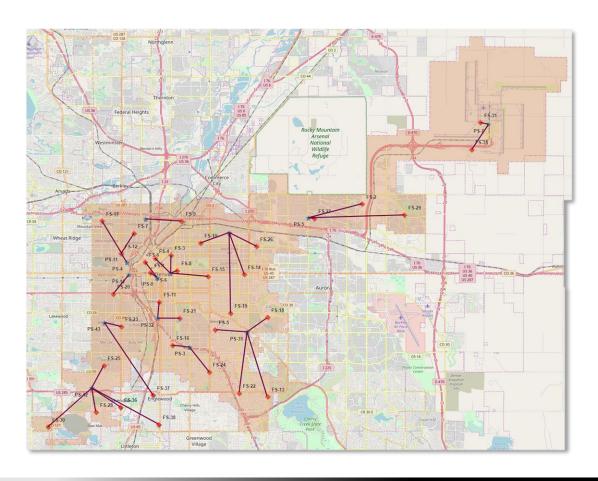
```
with open('fire stations lat and lon.csv', 'r') as fireFile:
fireCSV = csv.reader(fireFile, delimiter=',')
  for i, row in enumerate(fireCSV):
     if i is 0:
       for i in range(len(row)):
          if row[j] == 'FIRE_STATION_NUM':
             namej = j
          if row[i] == 'REAL LAT':
             latj = j
          if row[j] == 'REAL_LON':
             lonj = j
     else:
       fireDat.append([])
       fireDat[i-1].append(row[namej])
       fireDat[i-1].append(int(row[latj])*(10**(-6)))
       fireDat[i-1].append(int(row[lonj])*(10**(-6)))
```

```
with open('Accidents fatal.csv', 'r') as accFile:
  accCSV = csv.reader(accFile, delimiter=',')
  for i, row in enumerate(accCSV):
     if i is 0:
       for j in range(len(row)):
          if row[j] == 'CRASH_NUM':
             namej = j
          if row[j] == 'REAL LAT':
             latj = j
          if row[j] == 'REAL_LON':
             loni = i
     else:
       accDat.append([])
       accDat[i-1].append(row[namej])
       accDat[i-1].append(int(row[latj])*(10**(-6)))
       accDat[i-1].append(int(row[lonj])*(10**(-6)))
```

accDat = []

Results

(PS-0,FS-1)	(PS-4,FS-12)
(PS-0,FS-6)	(PS-42,FS-38)
(PS-1,FS-9)	(PS-42,FS-36)
(PS-11,FS-7)	(PS-42,FS-25)
(PS-11,FS-17)	(PS-42,FS-28)
(PS-12,FS-20)	(PS-42,FS-30)
(PS-2,FS-26)	(PS-43,FS-37)
(PS-2,FS-14)	(PS-43,FS-23)
(PS-2,FS-19)	(PS-5,FS-29)
(PS-2,FS-10)	(PS-5,FS-2)
(PS-3,FS-24)	(PS-5,FS-27)
(PS-3,FS-16)	(PS-6,FS-15)
(PS-31,FS-22)	(PS-6,FS-8)
(PS-31,FS-13)	(PS-6,FS-3)
(PS-31,FS-18)	(PS-6,FS-4)
` · · · '	` · · ·
(PS-31,FS-5)	(PS-7,FS-31)
(PS-32,FS-21)	(PS-7,FS-35)
(PS-32,FS-11)	



Policy Recommendations

- Implementation
 - » Input the set of "Suggested Partnerships" into the separate dispatch systems





- » When a potentially fatal accident happens, the system "pings" the dispatcher with the partnership
- Dispatcher relays the information to responders

Looking Ahead

- Things to Change
 - » Capacities
 - Static and Event Dependent
 - » Dispatcher Needs
 - Human Application
 - » Inverse?
 - When do we NOT need a partnership
 - » Accountability and resource conservation

- More In-Depth Analysis
 - » Patrol Routs
 - GPS Real Time Locators
 - » Drive Time
 - » Spatial Statistics
 - Clustering and weighting accident data
 - Stability
 - Compare results for different timeframes

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