# Problem 1: Where do you belong, Krylov?

#### Solution

 $\mathbf{a}$ 

$$Aq^{1} = A \frac{b}{\|b\|} \in \text{span}(\{b, Ab\}) = \mathcal{K}_{2} = \text{span}(\{q^{1}, Aq^{1}\}) = \text{span}(\{q^{1}, q^{2}\})$$

$$v^{i+1} = Aq^{i} - \sum_{k=1}^{i} \langle Aq^{i}, q^{k} \rangle q^{k}$$

Note that the inner products are scalars and denote them as  $\alpha_k$ :

$$= Aq^i - \sum_{k=1}^i \alpha_k q^k$$

$$q^{i+1} = \frac{v^{i+1}}{\|v^{i+1}\|} = \frac{Aq^i - \sum_{k=1}^i \alpha_k q^k}{\|Aq^i - \sum_{k=1}^i \alpha_k q^k\|}$$

Again, note that the denominator norm is a scalar; we will call it  $\gamma_i$ . Substitute and re-arrange:

$$Aq^{i} = \gamma_{i}q^{i+1} + \sum_{k=1}^{i} \alpha_{k}q^{k}$$

$$\implies Aq^i \in \operatorname{span}(\{q_1,\ldots,q^{i+1}\}) = \mathcal{K}_{i+1}$$

Then, we note that for  $i \geq m$ ,  $\mathcal{K}_i = \mathcal{K}_m$ . So:

$$j = \min(i+1, m)$$

b

Taken from Numerical Linear Algebra by Trefethen and Bau.

Consider the Hessenberg decomposition:

$$A = PHP^*$$

where P is a unitary matrix (i.e.  $PP^* = P^*P = PP^{-1} = I$ ). Note that orthonormal matrices are unitary, and restrict our case to state we wish to perform a Hessenberg decomposition via orthonormal matrix Q. Then:

$$A = QHQ^* \implies AQ = QH$$

Let  $Q_n$  be the first n columns of Q and  $H_n$  to be the  $(n+1) \times n$  upper-left section of H (which is also Hessenberg!). Then:

$$AQ_n = Q_{n+1}H_n$$

The n-th column of the equation is:

$$Aq_n = h1_nq_1 + \ldots + h_{nn}q_n + h_{n+1} \, {}_nq_{n+1}$$

Solving for  $q_{n+1}$  will require a technique to ensure it's orthonormality – if we use Modified Gram-Schmidt, we get:

## **Algorithm 1** Algorithm to solve the above equation with an orthonormal $q_{n+1}$

```
b \in \mathbb{R}^n, \ q_1 \leftarrow rac{b}{\|b\|}
for n=1,2,3\ldots do
v \leftarrow Aq_n
for j=1,\ldots,n do
h_{jn} \leftarrow q_j^*v
v \leftarrow v - h_{jn}q_j
end for
h_{n+1,n} \leftarrow \|v\|
q_{n+1} \leftarrow rac{v}{h_{n+1,n}}
end for
```

...which is identical to the provided iterate equations.

 $\mathbf{c}$ 

See the code section at the end for the implementation. The result is:

# Problem 2: Ask Gram, Schmidt, or Givens for QR

#### Solution

 $\mathbf{a}$ 

Taken from "Linear Algebra Done Right" by Sheldon Axler.

Suppose  $v_1, \ldots, v_m$  is a lin. ind. list of vectors in V. Let  $e_1 = \frac{v_1}{\|v_1\|}$ . For  $j = 2, \ldots, m$ , define  $e_j$  inductively by:

$$e_j = \frac{v_j - \langle v_j, e_1 \rangle e_1 - \dots - \langle v_j, e_{j-1} \rangle e_{j-1}}{\|v_j - \langle v_j, e_1 \rangle e_1 - \dots - \langle v_j, e_{j-1} \rangle e_{j-1}\|}$$

(i.e. the Gram-Schmidt process definition). Then,  $e_1, \ldots, e_m$  is an orthonormal list of vectors in V s.t.:

$$\operatorname{span}(v_1, \dots, v_j) = \operatorname{span}(e_1, \dots, e_j) \forall j = 1, \dots, m$$

Note that for j = 1, span $(v_1) = \text{span}(e_1)$  because  $v_1$  is a positive mulitple of  $e_1$ . Suppose 1 < j < m and it has been verified that:

$$span(v_1, ..., v_{i-1}) = span(e_1, ..., e_{i-1})$$

Note that  $v_j \notin \text{span}(v_1, \dots, v_{j-1})$  because  $v_1 \dots, v_m$  are linearly independent. Thus,  $v_j \notin \text{span}(e_1, \dots, e_{j-1})$ . As such, we are not dividing by zero in the definition of  $e_j$ . We can also see that  $||e_j|| = 1$  by its definition.

Let  $k \in [1, j)$ . Then:

$$\langle e_j, e_k \rangle = \left\langle \frac{v_j - \langle v_j, e_1 \rangle e_1 - \dots - \langle v_j, e_{j-1} \rangle e_{j-1}}{v_j - \langle v_j, e_1 \rangle e_1 - \dots - \langle v_j, e_{j-1} \rangle e_{j-1}}, e_k \right\rangle$$

$$= \frac{\langle v_j, e_k \rangle - \langle v_j, e_k \rangle}{v_j - \langle v_j, e_1 \rangle e_1 - \dots - \langle v_j, e_{j-1} \rangle e_{j-1}}$$

$$= 0$$

Thus  $e_1, \ldots, e_j$  is an orthonormal list.

From the definition of  $e_j$ , we see that  $v_j \in \text{span}(e_1, \dots, e_j)$ . Combined with the equivalency of the spans provided above, we know:

$$\operatorname{span}(v_1,\ldots,v_j)\subset\operatorname{span}(e_1,\ldots,e_j)$$

Both these lists are lin. ind., thus both subspaces have dimension j and are equal.

b

The inner loop requires:

$$(4m-1)(1-1)+(4m-1)+(2-1)+\ldots+(4m-1)(n-1)=(4m-1)(0+1+\ldots+n-1)=(4m-1)(\frac{1}{2}n(n+1)-1)$$

The outer portion requires (4m-1)\*n operations. Summing:

$$= (4m+1)(\frac{1}{2}n(n+1)-1) + (4m+1)(n) = (4m+1)(\frac{1}{2}n(n)-1+n)$$

Simplifying and discarding lower order terms we find:

$$=2mn^2$$

### Algorithm 2 Classical Gram-Schmidt, annotated with FLOP counts

```
for j=1,\ldots,n do v \leftarrow a^j for i=1,\ldots,j-1 do R_{ij} \leftarrow \langle v,q^i \rangle \{m \text{ multiplications, } (m-1) \text{ additions } \Longrightarrow 2m-1\} v \leftarrow v - R_{ij}q^i \{m \text{ multiplications, } m \text{ subtractions } \Longrightarrow 2m\} end for {Total cost is } q^j \leftarrow \frac{v}{\|v\|} \{m \text{ multiplications, } m-1 \text{ additions, } 1 \text{ for sqrt } \Longrightarrow 2m\} R_{jj} \leftarrow \langle a^j,q^j \rangle \{m \text{ multiplications, } (m-1) \text{ additions } \Longrightarrow 2m-1\} end for Q=(q^1|\ldots|q^n)
```

 $\mathbf{c}$ 

An upper Hessenberg matrix of size  $n \times n$  will require n-1 operations to zero out the non-zero elements below the diagonal, which will result in a QR decomposition. Givens matrices take the form:

$$G_n = \begin{bmatrix} 1 & \cdots & 0 & \cdots & 0 & \cdots & 0 \\ \vdots & \ddots & \vdots & & \vdots & & \vdots \\ 0 & \cdots & c & \cdots & -s & \cdots & 0 \\ \vdots & & \vdots & \ddots & \vdots & & \vdots \\ 0 & \cdots & s & \cdots & c & \cdots & \vdots \\ \vdots & & \vdots & & \vdots & \ddots & \vdots \\ 0 & \cdots & 0 & \cdots & 0 & \cdots & 1 \end{bmatrix}$$

The location of the c and s terms correspond to zeroing element i and j; that is, to zero element  $A_{i,j}$ , we set  $s = G_{ji} = -Gij$  and  $G_{k,k} = c$  for k = i, j (and 1 for  $k \neq i, j$ ). The resulting algorithm is as follows:

### Algorithm 3 QR Factorization of an Upper Hessenberg Matrix using Givens rotations

```
Input: A^1 \in \mathbb{R}^{n \times n}, A_{ij} = 0 \forall i > j+1 {This could be a parallelized "loop"} for j=1,\ldots,n-1 do i \leftarrow j+1  r^j \leftarrow \sqrt{\left[A^j_{jj}\right]^2 + \left[A^j_{ij}\right]^2}  c \leftarrow A^j_{jj}/r s \leftarrow -A^j_{ij}/r G^j_{ji} \leftarrow s G^j_{ij} \leftarrow -s G^j_{ij} \leftarrow c A^{j+1} \leftarrow G^j A^j end for Q = \prod_{j=1}^{n-1} G^j R = A^{n-1}
```

## Code

#### Problem 1

```
#!/usr/bin/env python3
import numpy as np
     float\_formatter = "{:.2 f}".format
     np.set_printoptions(formatter={"float_kind": float_formatter})
     \begin{array}{lll} \textbf{def} & \texttt{is\_upper\_hessenberg} \ (A: \ \texttt{np.ndarray}) \ \rightarrow \ \textbf{bool:} \\ & \texttt{return} & \texttt{np.isclose} \ (\texttt{np.tril} \ (A, \ k=-2), \ 0.0 \ , \ \ \texttt{atol=1e-3).all} \ () \end{array}
10
     def is-lower_hessenberg(A: np.ndarray) -> bool:
    return np.isclose(np.triu(A, k=2), 0.0, atol=1e-3).all()
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     # Adapted from: https://relate.cs.illinois.edu/course/cs450-s18/file-version/587 fb0505883e0622b9b144b532b5e9a4e7a3684/demos/upload/04-eigenvalues/Arnoldi%20iteration.html def arnoldi_iteration(A, b):
          arnoldi_iteration(A, b):

n = A.shape[0]

H = np.zeros((n, n))

Q = np.zeros((n, n))

# Normalize the input vector

Q[:, 0] = b / np.linalg.norm(b) # Use it as the first Krylov vector

for k in range(n):

u = A @ Q[:, k]

for j in range(k + 1):

qj = Q[:, j]

H[j, k] = qj @ u

u -= H[j, k] * qj

if k + 1 < n:

H[k + 1, k] = np.linalg.norm(u)

Q[:, k + 1] = u / H[k + 1, k]
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35
            return Q, H
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     39
40
            42
43
44
45
            print("Q:")
print(q)
print("H:")
print(h)
46
47
48
            uh = is_upper_hessenberg(h)
lh = is_lower_hessenberg(h)
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51
52
53
54
55
            print("Wow its upper Hessenberg, whoda thunk it")
if lh:
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57
58
                           "Wow it's tridiagonal, its cuz the matrix A is Hermitian (or Symmetric cuz its real)"
```

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