# Improved Fast Direct Methods for Gaussian Processes

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#### February 11, 2022

#### Introduction

For Gaussian Process, the posterior update is given by

$$\begin{split} & \mu_{t+1}(\mathbf{x}) = \mathbf{k}_t(\mathbf{x})^\top \big(\mathbf{K}_t + \eta^2 \mathbf{I}_t\big)^{-1} \mathbf{y}_t, \\ & \sigma_{t+1}(\mathbf{x}) = k(\mathbf{x}, \mathbf{x}) - \mathbf{k}_t(\mathbf{x})^\top \big(\mathbf{K}_t + \eta^2 \mathbf{I}_t\big)^{-1} \mathbf{k}_t(\mathbf{x}). \end{split}$$

Also, the log likelihood used to select hyperparameters is given as:

$$\log p(\mathbf{y}|\mathbf{X},\theta) = -\frac{1}{2}\mathbf{y}^{\intercal}\big(\mathbf{K} + \eta^2\mathbf{I}\big)^{-1}\mathbf{y} - \frac{1}{2}\log|\mathbf{K} + \eta^2\mathbf{I}| - \frac{n}{2}\log 2\pi$$

For large n, the computational cost of inverting  $\widehat{\mathbf{K}}_{XX} = \mathbf{K}_t + \eta^2 \mathbf{I}_t$  and finding  $|\widehat{\mathbf{K}}_{XX}|$  is expensive; the traditional/direct methods for dense systems scales quadratically:

- The time complexity is  $O(n^3)$ .
- The space complexity is  $O(n^2)$ , needed to store the lower triangular factor L along with  $\widehat{\mathbf{K}}_{XX}$ .

Traditionally, Cholesky decomposition is used, factoring the positive-semidefinite matrix  $\widehat{\mathbf{K}}_{XX}$  into  $L \times L^{\top}$ .

#### Related Work

There are generally 2 approaches to address the scaling issue:

- 1. Exact solutions Find the exact solution to the matrix inversion.
  - 1. LU decomposition
  - 2. Cholesky decomposition
    - Implementation: Scikit-Learn (via Numpy, Blas/Lapack)
    - Implementation: GPy (via Scipy, Blas/Lapack)
    - Implementation: GPflow (via Tensorflow)
  - 3. Conjugate gradients method (CG)
    - 1. Exact Gaussian Processes on a Million Data Points
      - Author's Implementation (GPyTorch): Exact GPs with GPU Acceleration
  - 4. Hierarchical Off-Diagonal Low-Rank (HODLR) matrix factorization
    - 1. Fast Direct Methods for Gaussian Processes
      - Author's Implementation: George
- 2. Approximate solutions Trade-off accuracy to improve computational cost.
  - 1. Mixture-of-experts
  - 2. Sampling m Inducing points
  - 3. Random feature expansions
- 3. Numerical issues in maximum likelihood parameter estimation for Gaussian process interpolation

#### Exact Gaussian Processes on a Million Data Points

**Key Idea**: Instead of solving for  $\widehat{\mathbf{K}}_{XX}^{-1}$  to find  $\mathbf{x}$ , we find  $\mathbf{x}$  directly:

$$\mathbf{x}^* = \arg\min_{\mathbf{x}} \left(\frac{1}{2}\mathbf{x}^{\top}\widehat{\mathbf{K}}_{XX}\mathbf{x} - \mathbf{x}^{\top}\mathbf{y}\right)$$

- Unique minimizer  ${f x}$  exist due to positive definite  $\widehat{{f K}}_{XX}$
- Iterative algorithm with matrix-vector multiplication
- Iterations can be speed up with preconditioning
- Able to decide tolerence of solution  $\epsilon = ||\widehat{\mathbf{K}}_{XX}\mathbf{x}^* \mathbf{y}||/||\mathbf{y}||$
- Similar to gradient descent, the difference is that CG enforces that the search direction  $\mathbf{p}_i$  is conjugate to each other.

Note: In the absence of round-off error, it converge to exact solution after n steps.

#### **Process**

Without forming the kernel matrix directly, using partitioned kernel MVM,

- Time complexity of  $O(n^3)$ ??
- Space complexity of O(n), assuming p = n partitions.

Note: Time complexity of the process is not discussed in the paper, which needs more thought.

#### MVM-based GP inference

#### **CG** Computation

#### Sampling m Inducing points

The 2-most popular methods that use this technique:

- 1. Sparse and Variational Gaussian Process (SVGP) [1]
- 2. Sparse Gaussian Process Regression (SGPR) [2]

#### **Process**

With m-inducing points,

- Time complexity of  $O(nm^2)$
- Space complexity of O(nm)

#### Fast Direct Methods for Gaussian Processes

**Key Idea**: Assume that  $\widehat{\mathbf{K}}_{XX}$  is Hierarchical Off-Diagonal Low-Rank (HODLR) matrix, and can be factored

$$\widehat{\mathbf{K}}_{XX} = K^{(0)} = \begin{bmatrix} K_1^{(1)} & U_1^{(1)} V_1^{(1)}^\top \\ {V_1^{(1)} U_1^{(1)}}^\top & K_2^{(1)} \end{bmatrix},$$

where the diagonal blocks are recursive HODLR decompositions

$$K_1^{(1)} = \begin{bmatrix} K_1^{(2)} & U_1^{(2)} V_1^{(2)^\top} \\ V_1^{(2)} U_1^{(2)^\top} & K_2^{(2)} \end{bmatrix}, \quad K_2^{(1)} = \begin{bmatrix} K_3^{(2)} & U_2^{(2)} V_2^{(2)^\top} \\ V_2^{(2)} U_2^{(2)^\top} & K_4^{(2)} \end{bmatrix},$$

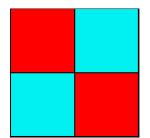
and that  $U_i^{(j)}, V_i^{(j)}$  are  $\frac{n}{2^j} \times r$  matrices. We note that the off-diagonal matrices are low-rank matrices and can be approximated by  $U_i^{(j)}, V_i^{(j)}$  up to a factorization precision  $\epsilon$ . The decompositions can be performed recursively up to the  $\kappa$ -level where  $\kappa \sim \log n$ ; i.e.  $\widehat{\mathbf{K}}_{XX} = K_{\kappa}K_{\kappa-1}\cdots K_0$ , see below.

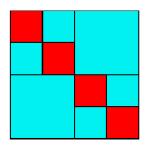
#### Process

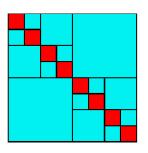
The overall runtime to obtain the matrix inverse of  $\widehat{\mathbf{K}}_{XX}$  and its determinant  $|\widehat{\mathbf{K}}_{XX}|$  is  $O(n\log^2 n)$ :

- The factorization of an  $n \times n$ ,  $\kappa$ -level (where  $\kappa \sim \log n$ ) HODLR matrix takes  $O(n \log^2 n)$ .
- Given a HODLR-type factorization, finding the inverse takes  $O(n \log n)$ .
- Given a HODLR-type factorization, finding the determinant takes  $O(n \log n)$ .

The overall space required scales to the number of dense matrices  $n/2^{\kappa} \times n/2^{\kappa}$ :







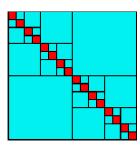


Figure 1: Pictorial representation of the HOLDR decompositons.

Note: Original paper did not discuss space complexity, need to put more thought and experiments; It is also not practical to subdivide completely, the authors implementation subdivide until a certain specified size.

Matrix Inversion. The summarized process to obtain the inverse (details are in the paper):

- 1. Computing the low-rank factorization of all off-diagonal block
- 2. Using these low-rank approximations to recursively factor the matrix into roughly  $O(\log n)$  pieces.
- 3. Sherman-Morrison-Woodbury formula can be applied to the  $O(\log n)$  factors to find the inverse.

**Determinant Computation.** Using the HOLDR factorization, the determinant can be computed directly:

$$|\widehat{\mathbf{K}}_{XX}| = |K_{\kappa}| \times |K_{\kappa-1}| \times \dots \times |K_0|$$

#### Metrics Used

In this section, we review the various metrics used by the various papers, in a attempt to find the best for us:

- 1. Exact Gaussian Processes on a Million Data Points: Dataset is split into 4:2:3 training, validating and testing sets. Validation is used to tune the CG parameters, the metrics are computed on test set.
  - 1. Accuracy Root-Mean-Square Error on the test set, and is heavily used to discuss the performance of the method, on 12 UCI datasets.

$$RMSE = \sqrt{\frac{1}{n}\Sigma_{i=1}^n \Big(y_i^* - y_i)^2}$$

- 2. Accuracy Negative log-likelihood of the exact GP is shown along with the approximate methods, but it is not used to discuss at all.
- 3. Compute Time Training Time, Testing Time
- 2. Fast Direct Methods for Gaussian Processes:
  - 1. Accuracy Relative  $l_2$  precision in the solution to a test problem Cx = b, where b was generated a priori from the known vector x, where C is synthetically created from a linear combination of:
    - 1. Gaussian covariance
    - 2. Multiquadric covariance matrice
    - 3. Exponential covariance
  - 2. Accuracy Difference of the Root-Mean-Square Errors of the exact method vs HOLDR on a synthetic function  $\sin(2x) + \frac{1}{8}e^x + \epsilon$ .
  - 3. Compute Time Assembly Time, Factor Time, Solve Time and Determinant Compute Time.

#### Some thoughts

- 1. RMSE on test set must be used for accuracy comparison.
- 2. We can test synthetic accuracy using Cx = b.
- 3. We need to get a notion of the computational relative floating point error for the operations that we care about.
- 4. What is there a theoretical error that we can achieve?

## Motivation

Our key motivations for improving Holdr is as follows:

- 1. Holdr is fast; The runtime complexity is very good when compared to the next state-of-the-art.
- 2. Holdr is space efficient; The space complexity is very good when compared to the next state-of-the-art.
- 3. Holdr is not accurate; it relies on the assumption that  $\mathbf{K}_{XX}$  resembles a HODLR matrix; failing which, the accuracy of the matrix inversion and determiniant computations is affected in unpredictable ways. The accuracy of the computation improved by setting the factorization precision  $\epsilon$  to a smaller number, but the computation can still fail unpredictably.
- 4.  $\widehat{\mathbf{K}}_{XX}$  can be conditioned, into an equlivant hierarchical matrix  $\widehat{\mathbf{K}}_{X'X'}$  by performing some reordering operations on  $X \to X'$ , improving the matrix structure.

Experiments here are done on the Synthetic Function, sin(x).

#### Holdr is fast

The main advantage of Holdr when compared to other methods is that it is very fast (also on scaling). Even with a small  $\epsilon = 0.000001$ , Holdr is competitive on time.

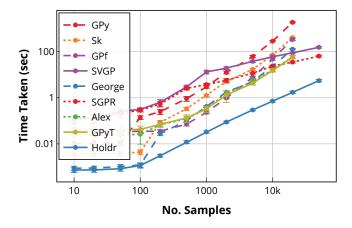


Figure 2: Take taken, Holdr with  $\epsilon = 0.000001$  when compared with other methods.

Note: Some issues with GPyT and Alex, Alex here is using CPU.

#### Holdr is space efficient

The next key advantage of Holdr when compared to other methods is that it is very space efficient (also on scaling). Placeholder, incomplete.

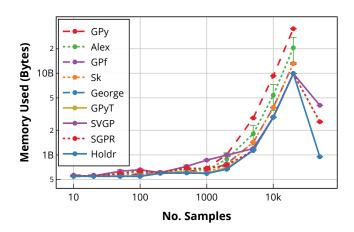


Figure 3: Memory used by the computation, Holdr with  $\epsilon = 0.000001$  when compared with other methods.

Note: Some issues with the chart and experiments, which will be fixed; trust only the last datapoint.

#### *Holdr* is not accurate

From Inconsistent results with HODLRS olver, we see that there are accuracy issues in its use. The accuracy of the computation improved by setting the factorization precision  $\epsilon$  to a smaller number, but the computation can still fail unpredictably.

We can run Holdr with different factorization precisions  $\epsilon \in \{1e^{-i}: i \in [1,6]\}$  to determine its effects on accuracy and runtime complexity. We let the accurate log likelihood of the X to be  $\ell^*$ , computed using Cholesky decomposition. We compute the log likelihood  $\ell_{\text{HODLR}}$  using the HODLR matrix factorization with the various  $\epsilon$ .

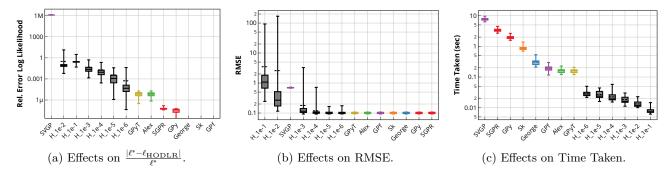


Figure 4: Holdr with different factorization precisions.

We see clearly from 4a that with a small  $\epsilon$ , it improves the average error but at an expense of high deviation. In other words, factoring with a small factorization precision yields better results on the average but fail unpredictably increasingly.

This issue was discussed in passing in the paper, see below; The authors concluded that it is not an issue and presented a suggestion to use kd-tree sort to condition the matrix  $\widehat{\mathbf{K}}_{XX}$ . However, the authors reversed this viewpoint in Issue 128.

We note that the authors claimed that when the data points at which the kernel to be evaluated at are not approximately uniformly distributed, the performance of the factorization may suffer, but only slightly. A higher level of compression could be obtained in the off-diagonal blocks if the hierarchical tree structure is constructed based on spatial considerations instead of point count, as is the case with some kd-tree implementations.

# $\widehat{\mathbf{K}}_{XX}$ can be conditioned

# **Proof of Concept**

The key insight is that the closer the kernel matrix is to an hierarchical matrix, the better the accuracy. Since HOLDR is fast, we want to pay some computational cost to condition the matrix. We see that we can perform some reordering operations on  $X = [x_0, \cdots, x_n] \to X'$  to massage<sup>1</sup> the kernel matrix  $\widehat{\mathbf{K}}_{XX}$  into an equlivant hierarchical matrix  $\widehat{\mathbf{K}}_{X'X'}$ . Here, equlivant means that the reordering operation does not change values of the posterior updates and the marginal likelihood <sup>2</sup>.

$$\widehat{\mathbf{K}}_{XX} = \begin{bmatrix} x_0 & x_1 & x_2 & x_3 & & x_0 & x_2 & x_1 & x_3 \\ 1 & 0 & 4 & 0 \\ 0 & 1 & 0 & 2 \\ 4 & 0 & 1 & 0 \\ 0 & 2 & 0 & 1 \end{bmatrix} \begin{bmatrix} x_0 \\ x_1 \\ x_2 \\ x_3 \end{bmatrix} \rightarrow \begin{bmatrix} 1 & 4 & 0 & 0 \\ 4 & 1 & 0 & 0 \\ 0 & 0 & 1 & 2 \\ 0 & 0 & 2 & 1 \end{bmatrix} \begin{bmatrix} x_0 \\ x_2 \\ x_1 \\ x_3 \end{bmatrix} = \widehat{\mathbf{K}}_{X'X'}$$

In this proof of concept, we will find the hierarchical matrix  $\widehat{\mathbf{K}}_{X'X'}$  and its corresponding reordered X' for which allows us to achieve a good accuracy of the log likelihood  $\log p(\mathbf{y}|\mathbf{X},\theta)$ .

We formulate the problem as a optimization problem and use Genetic Algorithm (GA) to find X', with each *i*-iteration candidate solution to be  $X_i$ . We run the GA with 1000 iterations and population size of 20:

- Mutation operation: Swapping of 2 points  $x_i, x_j$  where i, j are indices from different groups.
- Fitness score: Absolute difference of the accurate log likelihood and the current log likelihood  $|\ell^* \ell^{(i)}|$

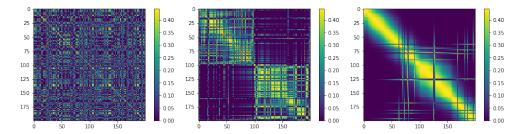


Figure 5: Heatmap of the  $\widehat{\mathbf{K}}_{XX}$  matrices, showing from  $X_0 \to X_{20} \to X_{1000}$ .

From above, we see that the concept works, choosing a solution  $X_i$  that is incrementally closer to the better solution than  $X_{i-1}$ . We summarize the results as follows:

$$\ell^{(0)} = 23.52 \rightarrow \ell^{(1000)} = 31.34 \sim \ell^* = 31.34$$

Hence, we are confident that it would be possible to perform a reordering operation on X such that its kernel matrix is a hierarchical matrix.

# Genetic Algorithm

We have used the *optimal* metric in the Proof of Concept which will not be practical; we are after all aiming to achieve that as a result of our optimization, so using it in our optimization would be cheating, and not to mention not cheap computationally. Here, we will use the same settings as in the proof of concept, with the exception of the fitness function.

#### Sum of off-diagonal matrix

$$\widehat{\mathbf{K}}_{XX} = K^{(0)} = \begin{bmatrix} K_1^{(1)} & U_1^{(1)} V_1^{(1)}^\top \\ {V_1^{(1)} U_1^{(1)}}^\top & K_2^{(1)} \end{bmatrix},$$

Here, we are using the sum over the cells in the off-diagonal matrix  $V_1^{(1)}U_1^{(1)^{\top}}$ . We aim to find a reordering such that the sum over the off-diagonal matrix is minimum.

<sup>&</sup>lt;sup>1</sup>Do we need to show that such set of operations exist?

 $<sup>^{2}</sup>$ Do we need a proof?

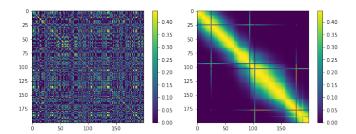


Figure 6: Heatmap of the  $\widehat{\mathbf{K}}_{XX}$  matrices - left is the kernel matrix of  $X_0$ , right is after.

$$\ell^{(0)} = 23.52 \rightarrow \ell^{(\text{GA-Sum})} = 31.44, \frac{|\ell^* - \ell^{(\text{GA-Sum})}|}{\ell^*} = 0.307\%$$

### Rank of off-diagonal matrix

Similar to the sum as above, instead we use the rank of the off-diagonal matrix  $V_1^{(1)}U_1^{(1)^{\top}}$ . This is from the definition of hierarchical matrix.

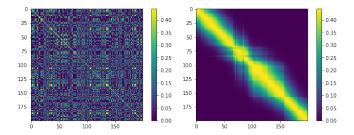


Figure 7: Heatmap of the  $\widehat{\mathbf{K}}_{XX}$  matrices - left is the kernel matrix of  $X_0$ , right is after.

$$\ell^{(0)} = 23.52 \rightarrow \ell^{(\text{GA-Rank})} = 31.44, \frac{|\ell^* - \ell^{(\text{GA-Rank})}|}{\ell^*} = 0.307\%$$

We see that there is not much difference between using the rank and the sum of cells, with the exception of the increased computation needed to calculate the rank.

# Sorting

#### **KD-Tree**

As a simplified example, we build a KD-Tree on X and query the tree on the closest n points to  $x_0$ .

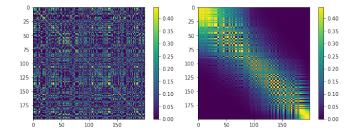


Figure 8: Heatmap of the  $\widehat{\mathbf{K}}_{XX}$  matrices - left is the kernel matrix of  $X_0$ , right is after.

$$\ell^{(0)} = 23.52 \rightarrow \ell^{(\text{Sort-KD})} = 29.14, \frac{|\ell^* - \ell^{(\text{Sort-KD})}|}{\ell^*} = 7.03\%$$

We observe that this method works marginally; However, KD-Trees are unsuitable for finding nearest neighbours in high dimensions due to curse of dimensionality. We expect the performance of this method to rapidly degrade in higher dimensionlity.

# Clustering

## Vanilla K-Means

Here, we perform vanilla K-means on X to find 2 clusters. We do not tweak/restrict the distance metric and the cluster size. Then, we sort the points based on the point's cluster identity.

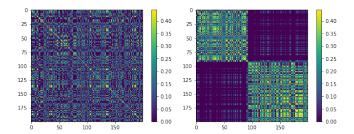


Figure 9: Heatmap of the  $\widehat{\mathbf{K}}_{XX}$  matrices - left is the kernel matrix of  $X_0$ , right is after.

$$\ell^{(0)} = 23.52 \rightarrow \ell^{(\text{Cluster-Vanilla})} = 31.36, \frac{|\ell^* - \ell^{(\text{Cluster-Vanilla})}|}{\ell^*} = 0.0618\%$$

Since we did not make any restrictions; we can expect that the cluster might be imbalanced.

## Sized K-Means

We take into account the equal sized constraint, as given in Same-size k-Means Variation. The implementation used to test is given in https://github.com/ndanielsen/Same-Size-K-Means.

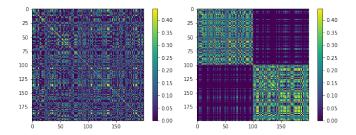


Figure 10: Heatmap of the  $\widehat{\mathbf{K}}_{XX}$  matrices - left is the kernel matrix of  $X_0$ , right is after.

$$\ell^{(0)} = 23.52 \rightarrow \ell^{(\text{Cluster-Sized})} = 31.35, \frac{|\ell^* - \ell^{(\text{Cluster-Sized})}|}{\ell^*} = 0.0171\%$$

#### Distanced K-Means

We take into account the kernel distance  $k(x_i, x_j)$ , replacing the default euclidean distance. Implementation used is from pyclustering.

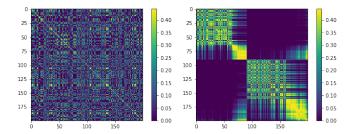


Figure 11: Heatmap of the  $\widehat{\mathbf{K}}_{XX}$  matrices - left is the kernel matrix of  $X_0$ , right is after.

$$\ell^{(0)} = 23.52 \rightarrow \ell^{(\text{Cluster-Distanced})} = 31.29, \frac{|\ell^* - \ell^{(\text{Cluster-Distanced})}|}{\ell^*} = 0.169\%$$

We expect this method to perform better, also the matrix looks a little odd.

# Graph

We can model the problem as a graph problem:

- 1. Vertex: Each  $x_i$
- 2. Edge:  $x_i \to x_j$ , where weight is given by  $k(x_i, x_j)$

Then, we are trying to find the a balanced, minimum cut of the graph. This problem is known also:

- 1. Planted Bisection Problem
- 2. Stochastic block model with two communities
- 3. Minimum bisection problem

#### Related resources:

- 1. On Minimum Bisection and Related Cut Problems in Trees and Tree-Like Graphs
- 2. Exact Recovery in the Stochastic Block Model
- 3. Community Detection in Networks: Algorithms, Complexity, and Information Limits
- 4. MIT mathematics of data science fall 2015
- 5. Minimum Cut Graphs
- 6. An efficient heuristic procedure for partitioning graphs
- 7. A Linear-Time Heuristic for Improving Network Partitions

- 1. Implementation is available here kshitij1489/Graph-Partitioning, but it is in Cpp.
- 8. From Louvain to Leiden: guaranteeing well-connected communities

### Kernighan-Lin Graph Bisection

Partition a graph into two blocks using the Kernighan–Lin algorithm; see An efficient heuristic procedure for partitioning graphs.

This algorithm partitions a network into two sets by iteratively swapping pairs of nodes to reduce the edge cut between the two sets. The pairs are chosen according to a modified form of Kernighan-Lin, which moves node individually, alternating between sides to keep the bisection balanced.

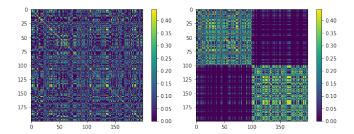


Figure 12: Heatmap of the  $\widehat{\mathbf{K}}_{XX}$  matrices - left is the kernel matrix of  $X_0$ , right is after.

$$\ell^{(0)} = 23.52 \rightarrow \ell^{(\text{Graph-Kernighan})} = 31.38, \frac{|\ell^* - \ell^{(\text{Graph-Kernighan})}|}{\ell^*} = 0.108\%$$

### Louvain Community Detection

Compute the partition of the graph nodes which maximises the modularity using the Louvain heuristices. This is the partition of highest modularity, i.e. the highest partition of the dendrogram generated by the Louvain algorithm. Implementation is from python-louvain, also see Fast unfolding of communities in large networks.

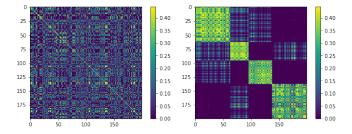


Figure 13: Heatmap of the  $\widehat{\mathbf{K}}_{XX}$  matrices - left is the kernel matrix of  $X_0$ , right is after.

$$\ell^{(0)} = 23.52 \rightarrow \ell^{(\text{Graph-Louvain})} = 29.44, \frac{|\ell^* - \ell^{(\text{Graph-Louvain})}|}{\ell^*} = 6.065\%$$

#### Leiden Community Detection

leidenalg - yet to try

# Linear Algebra

#### Principal Component Analysis

We sort the points using the transformed coordinate using PCA on X; we consider only the principal component. In this version, we use the default empirical sample covariance  $X^{\top}X$ ; it would be ideal to use  $\widehat{\mathbf{K}}_{XX}$ , which should improve the performance. However, that would require us to implement PCA.

The general steps of PCA is as follows:

1. Standardize the dataset.

- 2. Calculate the covariance matrix for the features in the dataset.
- 3. Calculate the eigenvalues and eigenvectors for the covariance matrix.
- 4. Sort eigenvalues and their corresponding eigenvectors.
- 5. Pick k eigenvalues and form a matrix of eigenvectors.
- 6. Transform the original matrix.

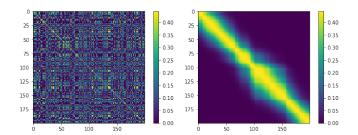


Figure 14: Heatmap of the  $\widehat{\mathbf{K}}_{XX}$  matrices - left is the kernel matrix of  $X_0$ , right is after.

$$\ell^{(0)} = 23.52 \rightarrow \ell^{(\text{La-Pca})} = 31.44, \frac{|\ell^* - \ell^{(\text{La-Pca})}|}{\ell^*} = 0.308\%$$

### Kernel Principal Component Analysis

Placeholder

# Difference Sorting

https://www.youtube.com/watch?v=vE6Ds4jwxb4

# Experiments

### **Dataset**

We used a split the dataset into a training set (80%) and testing set (20%).

For all synthetic experiments, we add  $z_t \sim \mathcal{N}(0, 0.1^2)$  noise to  $y_t = f(\mathbf{x}_t) + z_t$  to simulate real-world noise.

#### [Synthetic] Function

We sample from the 1-Dimensional function  $f(x) = \sin(x)$ , where  $x \in [-10, 10]$ .

#### [Synthetic] GP Kernel

Here, we test our algorithm on synthetic data by sampling functions from GP. We use an RBF kernel with corresponding dimensional lengthscale and scale parameters set to  $\sigma_{\rm opt}^2 = 0.5$  and  $l_{\rm opt} = 0.1$ .

#### [Synthetic] HPOlib2

Commonly used synthetic functions f(x) are from HPOlib2 [3], such as

HPOlib2	d
Forrester	1
Levy	1
Bohachevsky	2
Branin	2
Camelback	2
Hartmann6	6

#### [Real] UCI Data

Curated and processed UCI Dataset [4] from https://github.com/treforevans/uci\_datasets [5].

UCI Dataset	Cite	n	d
3droad	[6]	434874	3
kin40k		40000	8
protein	[7]	45730	9
houseelectric		2049280	11
bike		17379	17
elevators		16599	18
keggdirected	[8]	48827	20
pol		15000	26
keggundirected	[8]	63608	27
buzz	[9]	583250	77
song	[10]	515345	90
slice	[11]	53500	385

#### Methods

All methods use the RBF kernel with  $\sigma_0^2$  set to the population variance of the training sample and  $l_0=1.0$ . The model's noise variance is set  $\eta^2=0.01$  to account for noisy observations.

- Stochastic Methods. Here we optimize the model parameters using the Adam Optimizer [12] which is designed for stochastic objective functions. We use the implementation of Adam from [13].
  - SVGP, implementation from [14]
  - SGPR, implementation from [14]
- Traditional, cholesky decomposition methods
  - $-\ GPy$  popular GP implementation from Sheffield University [14]
  - -Sk scikit-learn implementation [15]
  - George Cholesky implementation inside 'Fast Direct Methods for Gaussian Processes' [16]
  - *GPf* scikit-learn implementation [17]
  - *GPyT GPyT* orch implementation [18]
- Conjugate gradients method (CG)
  - Alex Author's implementation for 'Exact Gaussian Processes on a Million Data Points' [19]
- Hierarchical Off-Diagonal Low-Rank (HODLR) matrix factorization method
  - Holdr Author's implementation for 'Fast Direct Methods for Gaussian Processes' [20]
- Our methods:
  - H KM Holdr with K-Means
  - H KM Holdr with K-Means with equal sized clusters
  - $H\_GA$  Holdr with Genetic Algorithm
  - $H_PCA$  Holdr with PCA
  - $H_KPCA$  Holdr with KPCA
  - H\_Grf Holdr with Kernighan-Lin Graph Bisection

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