

COMP2119C
Data Structures & Algorithms
2024-25

Tutorial 1 – Graph Terminology
(Appendix)

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Definition of a Graph

A **graph** $G = (V, E)$ consists of

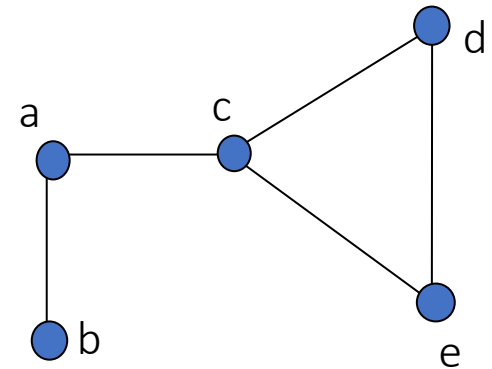
- V , the set of **vertices** (nodes, points), and
- E , the set of **edges** (lines)

Undirected graphs are graphs where edges have no direction (E : **unordered pairs** of V).

Example:

$$V = \{a, b, c, d, e\}$$

$$E = \{\{a, b\}, \{a, c\}, \{c, d\}, \{c, e\}, \{d, e\}\}$$

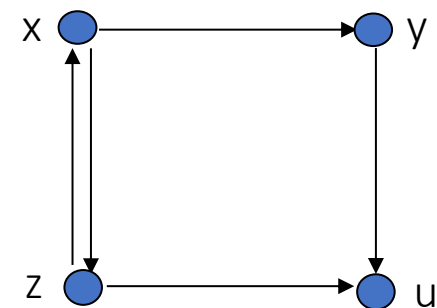


Directed graphs are graphs where edges have directions (E : **ordered pairs** of V).

Example:

$$V = \{u, x, y, z\}$$

$$E = \{(x, y), (x, z), (y, u), (z, x), (z, u)\}$$



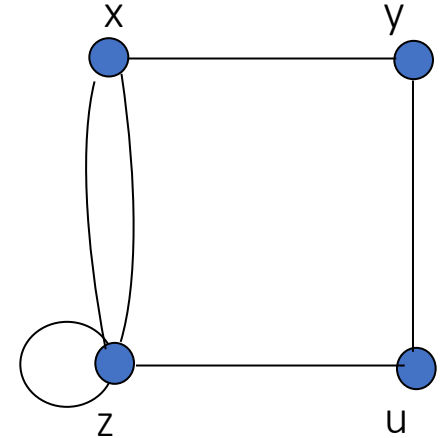
Simple Graphs and Multigraphs

Simple graphs - at most one edge between any pair of vertices, no self-loop. Unless otherwise stated, a “graph” means a simple graph.

Multi-graphs - having **multiple edges** between the same (ordered) pair of vertices, self-loops allowed

If x, y are two vertices, and $e = \{x, y\}$ is in E , then

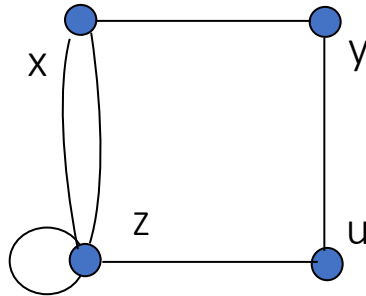
- x and y are **adjacent**
- e is **incident** with x and y
- x and y are the **endpoints** of e



Graph Degree

The **degree** of a vertex of a graph is the number of edges that are incident to the vertex.

Example: $\deg(y)=2$, $\deg(u)=2$, $\deg(x)=3$, $\deg(z)=5$

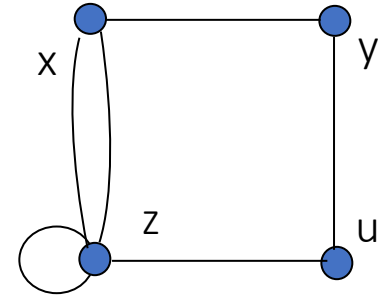


Graph Degree

For an **undirected** graph $G = (V, E)$:

$$\sum_{v \in V} \deg(v) = ? \quad 2|E|$$

Proof: Every edge is counted twice in counting degrees.



e.g. sum of degrees = $3+2+5+2 = 12$, $|E| = 6$

For a **directed** graph $G=(V, E)$,

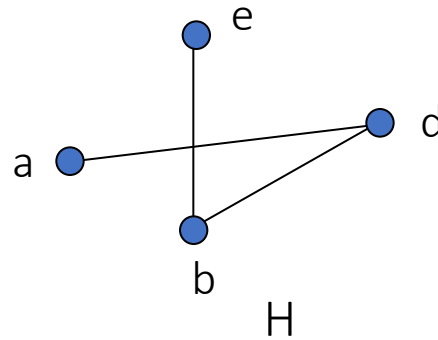
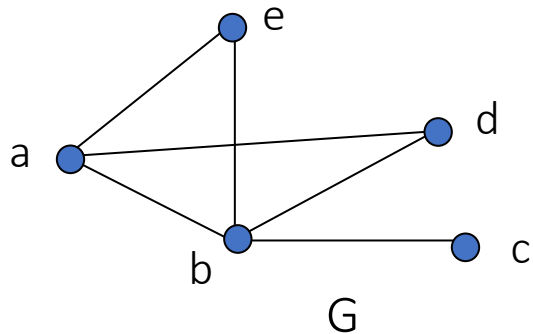
$$\sum_{v \in V} \deg^{-}(v) = \sum_{v \in V} \deg^{+}(v) = |E|$$

in-degree

out-degree

Graph Operations

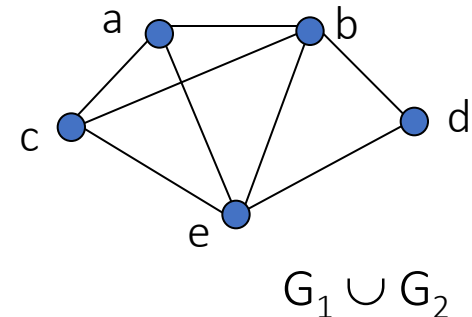
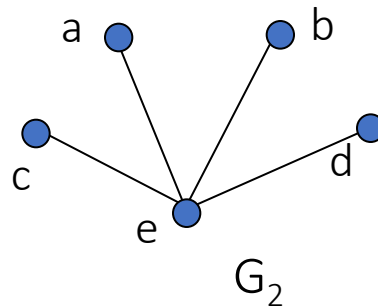
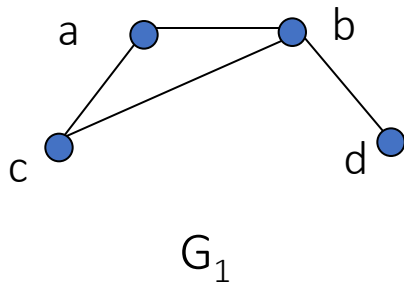
A graph $H = (V', E')$ is a **subgraph** of $G = (V, E)$ if $V' \subseteq V$, $E' \subseteq E$.



Is this a subgraph of G ?

The **union** of two graphs $G_1 = (V_1, E_1)$ and $G_2 = (V_2, E_2)$ is

$$G_1 \cup G_2 = (V_1 \cup V_2, E_1 \cup E_2)$$



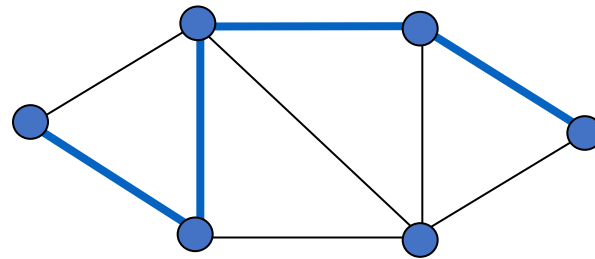
Paths

A **path** of **length n** from u to v in a graph G is a sequence of edges e_1, e_2, \dots, e_n , such that $e_1 = (u, x_1)$, $e_2 = (x_1, x_2)$, $\dots, e_n = (x_{n-1}, v)$.

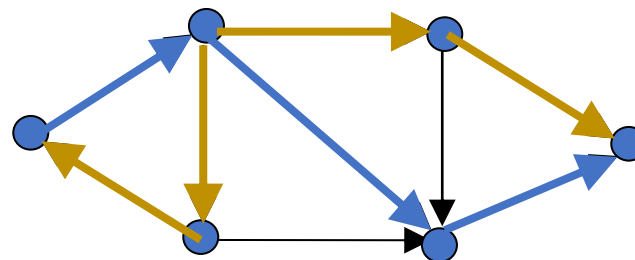
A path is **simple** if it doesn't contain the same vertex more than once.

Example:

➤ Undirected graph:



➤ Directed graph:



— Not a path

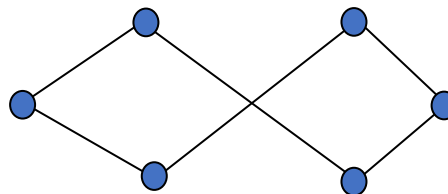
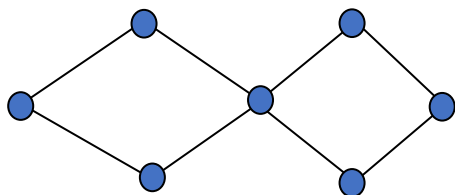
— A path

Circuits

A **circuit** of length n in a graph G is a sequence of edges e_1, e_2, \dots, e_n , such that $e_1 = (x_0, x_1)$, $e_2 = (x_1, x_2)$, $\dots, e_n = (x_{n-1}, x_0)$.

i.e., a circuit is a path that starts and ends at the same vertex.

A **cycle** is a circuit that does not contain the same vertex more than once.



Simple Paths

If there is a path between u and v , then there is a **simple path** between u and v .

Example:

$u \rightarrow a \rightarrow b \rightarrow c \rightarrow d \rightarrow e \rightarrow f \rightarrow a \rightarrow b \rightarrow v$ is a path

$u \rightarrow a \rightarrow b \rightarrow v$ is a simple path

$a \rightarrow b \rightarrow c \rightarrow d \rightarrow e \rightarrow f \rightarrow a$ is a circuit

