# COMP2119 Introduction to Data Structures and Algorithms Assignment 1 - Recursion, Mathematical Induction and Algorithm Analysis

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### 1 Asymptotic Bounds

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\begin{array}{l} \Theta(\log n):(d)\,\log n,(j)\,\log\frac{n}{\log n},\\ \Theta(\frac{n}{\log n}):(h)\,\frac{n}{\log n},\\ \Theta(\frac{n}{\log \log n}):(i)\,\frac{n}{\log \log n},\\ \Theta(\frac{n}{\log \log n}):(j)\,\frac{n}{\log \pi},(m)\sqrt{\sum_{i=1}^{n}(i+1)},\\ \Theta(n^{\log \pi}):(e)\,\pi^{\log \pi},(f)\,n^{\log \pi},(k)\,\pi^{\log(2n)},\\ \Theta(n^{\log 2\pi}):(l)\,n^{\log 2\pi},\\ \Theta(n^{\pi}):(a)\,n^{\pi},\\ \Theta(\pi^{n}):(b)\,\pi^{n},\\ \Theta(n^{n}):(n)\,1910n!+316n^{n},(c)\,n^{n} \end{array}
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#### 2 Recurrence Relations

(a) 
$$T(n) = T(n-1) + 3$$
 for  $n > 0$   
 $= T(n-2) + 3 + 3$  for  $n > 1$   
 $= \dots$   
 $\therefore T(n) = T(0) + 3n$   
 $= 3n$   
 $\therefore T(n) = \Theta(n)$ 

(b) Assume that n is a power of 3, i.e.  $n=3^k$  for  $k\in\mathbb{N}$ , and  $\log_3 n=k$ ,  $\therefore T(n)=3T(\frac{n}{3})+n$  for  $n\neq 1$ 

$$= 3(3T(\frac{n}{9}) + \frac{n}{3}) + n \quad \text{for } n >= 9$$

$$= 9T(\frac{n}{9}) + n + n \quad \text{for } n >= 9$$

$$= \dots$$

$$\therefore T(n) = 3^k T(\frac{n}{3^k}) + k * n$$

$$= 3^k T(1) + k * n$$

$$= 0 + k * n$$

$$= kn$$

$$= n * \log_3 n$$

$$\therefore T(n) = \Theta(n \log n)$$

(c) Assume that n is a power of 3, i.e. 
$$n = 3^k$$
 for  $k \in \mathbb{N}$ , and  $\log_3 n = k$ ,

$$T(n) = 4T(\frac{n}{3}) + 1 \text{ for } n >= 3$$

$$= 4(4T(\frac{n}{9}) + 1) + 1 \text{ for } n >= 9$$

$$= 16T(\frac{n}{9}) + 4^{k-1} + \dots + 4^{0} \text{ for } n >= 9$$

$$= 4(4T(\frac{n}{9}) + 4^{k-1} + \dots + 4^{0})$$

(d) Assume that n is a power of 2, i.e. 
$$n = 2^k$$
 for  $k \in \mathbb{N}$ , and  $\log_2 n = k$ ,

$$T(n) = nT(\frac{n}{2}) + n - 1 \quad \text{for } n >= 2$$

$$= n * (\frac{n}{2} * T(\frac{n}{4}) + \frac{n}{2} - 1) + n - 1 \quad \text{for } n >= 4$$

$$= \frac{n^2}{2} * T(\frac{n}{4}) + \frac{n^2}{2} - n + n - 1 \quad \text{for } n >= 4$$

$$= \frac{n^2}{2} * T(\frac{n}{4}) + \frac{n^2}{2} - 1 \quad \text{for } n >= 4$$

$$= \frac{n^2}{2} * (\frac{n}{4} * T(\frac{n}{8}) + \frac{n}{4} - 1) + \frac{n^2}{2} - 1 \quad \text{for } n >= 8$$

$$= \frac{n^3}{2 * 4} * T(\frac{n}{8}) + \frac{n^3}{2 * 4} - \frac{n^2}{2} + \frac{n^2}{2} - 1 \quad \text{for } n >= 8$$

$$= \frac{n^3}{2^1 * 2^2} * T(\frac{n}{2^1 * 2^2}) + \frac{n^3}{2^1 * 2^2} - 1 \quad \text{for } n >= 8$$

$$\therefore T(n) = \frac{n^k}{2^0 * 2^1 * 2^2 * \dots * 2^{k-1}} * T(1) + \frac{n^k}{2^0 * 2^1 * 2^2 * \dots * 2^{k-1}} - 1$$

$$= \frac{n^k}{2^{\frac{(k-1)*k}{2}}} * 2 - 1$$

$$= \frac{n^k}{n^{\frac{k-1}{2}}} * 2 - 1$$

$$= n^{\frac{2k}{2} - \frac{k-1}{2}} * 2 - 1$$

$$= n^{\frac{k+1}{2}} * 2 - 1$$

$$= 2n^{\frac{\log_2 n+1}{2}} - 1$$

$$\therefore T(n) = \Theta(n^{\frac{\log_2 n + 1}{2}})$$

## 3 Mathematical Induction

(a) Let f(n) be the predicate " $1*2^1+2*2^2+3*2^3+...+n*2^n=(n-1)2^{n+1}+2$ " for  $\forall n\in\mathbb{Z}^+$ .

For 
$$n = 1$$
, L.H.S. =  $1 * 2^1$   
= 2  
R.H.S. =  $(1 - 1)2^{1+1} + 2$   
= 2

 $\therefore$  L.H.S. = R.H.S.

 $\therefore f(1)$  is true.

#### Inductive Step:

Assume that f(n) is true when n = k, for some  $k \in \mathbb{Z}^+$ , i.e.  $f(k) = 1 * 2^1 + 2 * 2^2 + 3 * 2^3 + ... + k * 2^k = (k-1)2^{k+1} + 2$ .

Consider the case 
$$n = k + 1$$
, L.H.S.  $= 1 * 2^1 + 2 * 2^2 + 3 * 2^3 + ... + k * 2^k + (k + 1) * 2^{k+1}$   
 $= (k - 1)2^{k+1} + 2 + (k + 1) * 2^{k+1}$  (by induction hypothesis)  
 $= 2^{k+1} * (2k) + 2$   
 $= (k)2^{k+2} + 2$   
 $= (k + 1 - 1)2^{k+1+1} + 2$   
 $= R.H.S.$ 

f(n) is true when n = k + 1.

 $\therefore$  By the principle of mathematical induction, f(n) is true for all  $n \in \mathbb{Z}^+$ .

(b) Let 
$$f(n)$$
 be the predicate " $\frac{1}{\sqrt{1}+\sqrt{2}} + \frac{1}{\sqrt{2}+\sqrt{3}} + \frac{1}{\sqrt{3}+\sqrt{4}} + \dots + \frac{1}{\sqrt{n}+\sqrt{n+1}} = \sqrt{n+1} - 1$  for  $\forall n \in \mathbb{Z}^+$ ".

For 
$$n = 1$$
, L.H.S.  $= \frac{1}{\sqrt{1} + \sqrt{2}}$   
 $= \frac{1}{\sqrt{1} + \sqrt{2}} * \frac{\sqrt{1} - \sqrt{2}}{\sqrt{1} - \sqrt{2}}$   
 $= \frac{\sqrt{1} - \sqrt{2}}{1 - 2}$   
 $= \sqrt{2} - 1$   
R.H.S.  $= \sqrt{1 + 1} - 1$   
 $= \sqrt{2} - 1$ 

 $\therefore$  L.H.S. = R.H.S.

 $\therefore f(1)$  is true.

#### **Inductive Step:**

Assume that f(n) is true when n = k, for some  $k \in \mathbb{Z}^+$ , i.e.  $f(k) = \frac{1}{\sqrt{1} + \sqrt{2}} + \frac{1}{\sqrt{2} + \sqrt{3}} + \frac{1}{\sqrt{3} + \sqrt{4}} + \dots + \frac{1}{\sqrt{k} + \sqrt{k+1}} = \sqrt{k+1} - 1$ .

Consider the case n = k + 1,

L.H.S. 
$$= \frac{1}{\sqrt{1} + \sqrt{2}} + \frac{1}{\sqrt{2} + \sqrt{3}} + \frac{1}{\sqrt{3} + \sqrt{4}} + \dots + \frac{1}{\sqrt{k} + \sqrt{k+1}} + \frac{1}{\sqrt{k+1} + \sqrt{k+2}}$$

$$= \sqrt{k+1} - 1 + \frac{1}{\sqrt{k+1} + \sqrt{k+2}} \quad \text{(by induction hypothesis)}$$

$$= \sqrt{k+1} - 1 + \frac{1}{\sqrt{k+1} + \sqrt{k+2}} * \frac{\sqrt{k+1} - \sqrt{k+2}}{\sqrt{k+1} - \sqrt{k+2}}$$

$$= \sqrt{k+1} - 1 + \frac{\sqrt{k+1} - \sqrt{k+2}}{1-2}$$

$$= \sqrt{k+1} - 1 + \sqrt{k+2} - \sqrt{k+1}$$

$$= \sqrt{k+2} - 1$$

$$= \sqrt{k+1} - 1$$

$$= R.H.S.$$

- f(n) is true when n = k + 1.
- ... By the principle of mathematical induction, f(n) is true for all  $n \in \mathbb{Z}^+$ .

## 4 Algorithm Design