

**Question: Show**  $\log(n!) = \Theta(n \log n)$

(a) Upper bound: find a constant  $c_1$ , such that for large enough  $n$ ,  $\log(n!) \leq c_1 n \log n$ , or  $n! \leq c_1 n^n$ . Because  $n! = 1 \times 2 \times 3 \dots \times n \leq (n \times n \times n \dots \times n) = n^n$ . By letting  $c_1 = 1$ , we have  $n! \leq c_1 n^n$  thus  $\log(n!) = O(n \log n)$

(b) Lower bound:  $n! = 1 \times 2 \times 3 \dots \times n$ . By ignoring the first half of the series and keeping only the larger  $\lceil n/2 \rceil$  items, we have

$$n! \geq \lceil n/2 \rceil \times \lceil n/2 \rceil + 1 \times \lceil n/2 \rceil + 2 \dots \times n \geq \lceil n/2 \rceil \times \lceil n/2 \rceil \dots \times \lceil n/2 \rceil \geq \lceil n/2 \rceil^{\lceil n/2 \rceil}$$

Therefore  $\log(n!) \geq \log((n/2)^{n/2}) = n/2 \log n/2$ . We can ignore the floor and ceiling function in asymptotic notations, so  $\log(n!) = \Omega(n \log n)$

**Question: Show**  $\sum_{i=1}^n (1/i) = \Theta(\log n)$

(a) Upper bound:

$$\sum_{i=1}^n (1/i) \leq \sum_{i=1}^n (1/i_{p2})$$

where  $i_{p2}$  is the largest power of 2 that's less than or equal to  $i$ .

So:

$$\begin{aligned} \sum_{i=1}^n (1/i) &\leq (1/1 + 1/2 + 1/2 + 1/4 + 1/4 + 1/4 + 1/4 + 1/8 + \dots) \\ &= 1 + \sum_{i=1}^2 (1/2) + \sum_{i=1}^4 (1/4) + \sum_{i=1}^8 (1/8) + \dots + \sum_{i=1}^{2^k} (1/2^k) = k. \end{aligned}$$

Because there are only  $n$  items in the original series, the constraint on  $k$  would be  $1+2+4+8+\dots+2^{k-1} \leq (n-1)$ . Using geometric series lemma,  $2^k \leq n$ , so  $k = O(\log n)$ , meaning

$$\sum_{i=1}^n (1/i) = O(k) = O(\log n).$$

(b) Similarly for lower bound:

$$\sum_{i=1}^n (1/i) \leq \sum_{i=1}^n (1/i_{p2})$$

where  $i_{p2}$  is the smallest power of 2 that's greater than  $i$ .

So:

$$\begin{aligned} \sum_{i=1}^n (1/i) &\geq (1/1 + 1/2 + 1/4 + 1/4 + 1/8 + 1/8 + 1/8 + 1/8 + \dots) \\ &= 1 + \sum_{i=1}^1 (1/2) + \sum_{i=1}^2 (1/4) + \sum_{i=1}^4 (1/8) + \dots + \sum_{i=1}^{2^k} (1/2^{k+1}) = 1 + k/2. \end{aligned}$$

Again, we have  $1 + 1 + 2 + 4 + \dots + 2^k \geq n$ , therefore  $k \geq \log n - 1$ , and  $\sum_{i=1}^n (1/i) = \Omega(k) = \Omega(\log n)$ .

**Question 0.1 from DPV**

- (a)  $f(n) = \Theta(g(n))$
- (b)  $f(n) = O(g(n))$
- (c)  $f(n) = \Theta(g(n))$
- (d)  $f(n) = \Theta(g(n))$
- (e)  $f(n) = \Theta(g(n))$
- (f)  $f(n) = \Theta(g(n))$
- (g)  $f(n) = \Omega(g(n))$
- (h)  $f(n) = \Omega(g(n))$
- (i)  $f(n) = \Omega(g(n))$
- (j)  $f(n) = \Omega(g(n))$
- (k)  $f(n) = \Omega(g(n))$