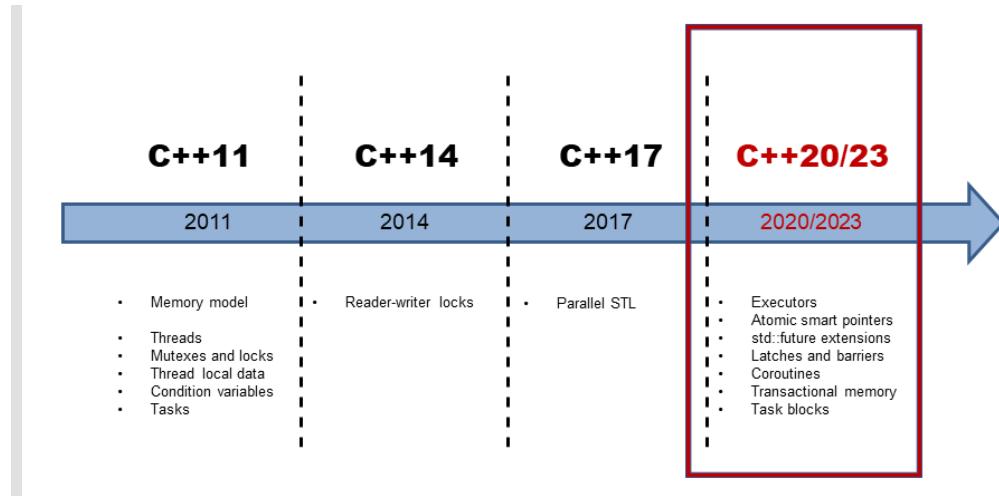


Data Structure and Algorithm in C++



C++ Algorithm

Problem Solving

Abstract data types	Array-based implementation	Link-based implementation
Essential concepts		The ADT bag
<h3><u>Programming vs. problem solving</u></h3> <ul style="list-style-type: none"> Computer scientists/engineers consider a problem as one thing that requires a program as the final solution. Solving a problem requires: <ul style="list-style-type: none"> (Analysis) Gaining an understanding of the problem. (Design) Construct a conceptual solution. (Programming) Implement the solution as a computer program. 		

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Object-oriented Solution

Abstract data types	Array-based implementation	Link-based implementation
Essential concepts		The ADT bag
<h2>Object-oriented solutions</h2> <ul style="list-style-type: none"> An <u>object-oriented (OO) solution</u> is a program that consists of <u>a system of interacting classes of objects</u>. <ul style="list-style-type: none"> Each <u>object</u> has characteristics (<u>attributes</u>) and behaviors (<u>functions</u>) related to the solution. A <u>class</u> is the common type of a set of objects. <u>Object-oriented analysis and design</u> helps us <u>discover</u> and <u>describe</u> these objects and classes (to construct a solution, i.e., a program). 		

Abstract data types	Array-based implementation	Link-based implementation
Essential concepts		The ADT bag
<h2>Object-oriented analysis and design</h2> <ul style="list-style-type: none"> Analysis (in general) is to <u>understand the problem</u> and the <u>requirements</u> of a solution. <ul style="list-style-type: none"> What a solution (program) must do. Not how to implement the solution. The analyzer should generate an accurate understanding of <u>what end users</u> will expect the solution to be. Design (in general) is to <u>describes a solution</u> to a problem. <ul style="list-style-type: none"> To <u>fulfill the requirements</u> discovered during analysis. 		

Abstract data types	Array-based implementation	Link-based implementation
Essential concepts		The ADT bag
<h2>Object-oriented v programming</h2> <ul style="list-style-type: none"> An object-oriented v language enables us to build <u>classes of objects</u>. A class combines: <ul style="list-style-type: none"> <u>Attributes</u> (characteristics) of objects of a single type. <u>Behaviors</u> (operations) Typically: <ul style="list-style-type: none"> Attributes are called data members or <u>member variables</u>. Behaviors are called methods or <u>member functions</u>. 		

Abstract data types	Array-based implementation	Link-based implementation
Essential concepts		The ADT bag
<h2>Object-oriented programming</h2> <ul style="list-style-type: none"> • ✓Encapsulation: <ul style="list-style-type: none"> – A class <u>combines</u> data and operations. – A class <u>hides</u> inner details. • ✓Inheritance: <ul style="list-style-type: none"> – Classes can <u>inherit</u> properties from other classes. – <u>Existing classes can be reused.</u> • Polymorphism: <ul style="list-style-type: none"> – Objects can do appropriate operations at <u>execution time</u>. 		

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Abstract data types	Array-based implementation	Link-based implementation
Essential concepts		The ADT bag
<h2>Operation Contracts (1/3)</h2> <ul style="list-style-type: none"> • An <u>operation contract</u> documents <u>how a module can be used</u> and <u>what limitations it has.</u> • Begin the contract during <u>analysis</u> and finish during <u>design</u>. <ul style="list-style-type: none"> – Used to <u>document code</u>, particularly in <u>header files.</u> ✓ – Does <u>not</u> describe <u>how</u> the module will perform its task. • A module’s operation contract specifies <u>its purpose, assumptions, and input/output.</u> <ul style="list-style-type: none"> – ✓Precondition: Statement of conditions that must exist <u>before</u> a module (function) executes. – ✓Postcondition: Statement of conditions that exist <u>after</u> a module (function) executes. 		

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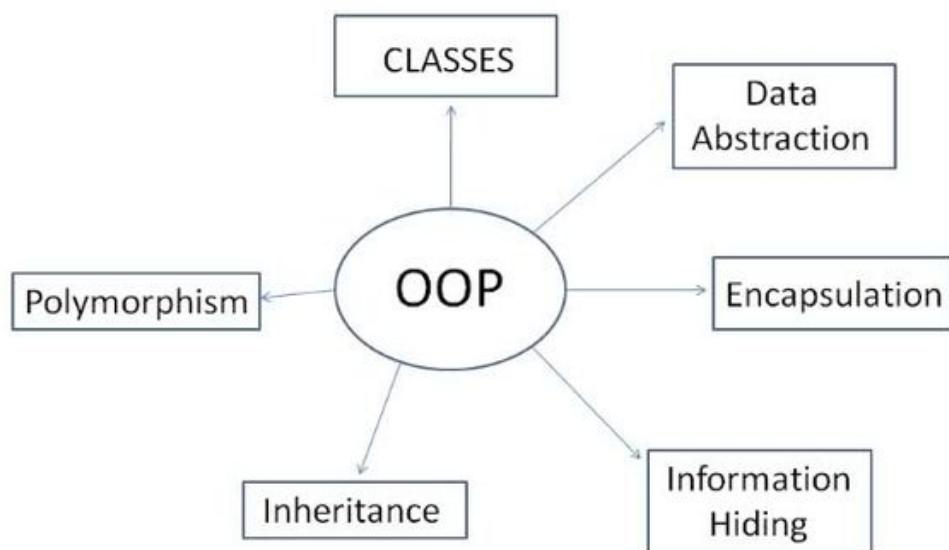
Abstract data types	Array-based implementation	Link-based implementation
Essential concepts		The ADT bag
<h2>Operation Contracts (2/3)</h2> <ul style="list-style-type: none"> • An example: <div style="border: 1px solid black; padding: 10px; width: fit-content;"> purpose // Sorts an array. // Precondition: anArray is an array of num integers; num > 0. // Postcondition: The integers in anArray are sorted. sort(anArray, num) </div> <p>Precondition & postcondition pretty vague. ✓Sort? <u>descending or ascending order?</u></p> 		

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Abstract data types	Array-based implementation	Link-based implementation
Essential concepts		The ADT bag
Operation Contracts (3/3)		
<ul style="list-style-type: none"> Revised specifications: <pre>// Sorts an array into ascending order. // Precondition: anArray is an array of num integers; // 1 <= num <= MAX_ARRAY, where MAX_ARRAY is a global // constant that specifies the maximum size of anArray. // Postcondition: anArray[0] <= anArray[1] <= ... // <= anArray[num-1]. num is unchanged. sort(anArray, num)</pre>		
<ul style="list-style-type: none"> Documentation is very important! <ul style="list-style-type: none"> You, the programmer, may forget everything about your programs. 		

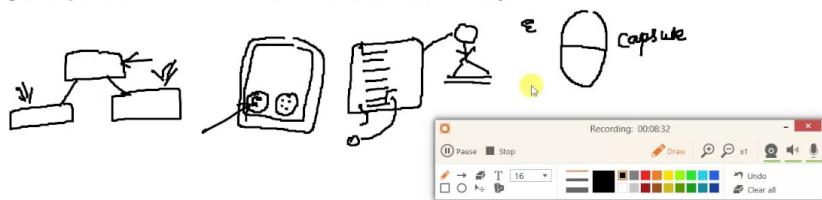
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C++ OOP



Four Pillars of Object Oriented Programming.

- Inheritance:** Process of creating new classes from existing classes.
- Abstraction:** Refers to providing only essential information and hiding the background details from the end user.
- Encapsulation:** Binding together of data and functions that manipulate that data and keep both safe from outside interference and misuse.
- Polymorphism:** Refers to one interface and many forms.



Data Abstraction

What is Data Abstraction?

- A data abstraction is a simplified view of an object that
 - includes only features one is interested in.
 - while hides away the unnecessary details.
- In programming languages, a data abstraction becomes an abstract data type (ADT) or a user-defined type.
- In OOP, it is implemented as a class

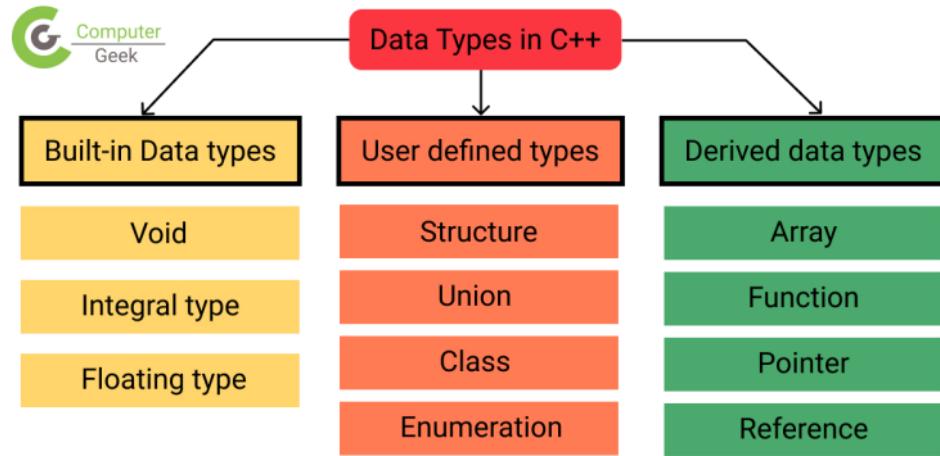


Data Abstraction in C++ (continued)

(11.14)

- Constructors
 - functions to initialize the data members of instances
 - may also allocate storage if part of the object is heap-dynamic
 - can include parameters to provide parameterization of the objects
 - implicitly called when an instance is created
 - » can be explicitly called
 - name is the same as the class name
- Destructors
 - functions to cleanup after an instance is destroyed; usually just to reclaim heap storage
 - implicitly called when the object's lifetime ends
 - » can be explicitly called
 - name is the class name, preceded by a tilde (~)

Struct in C and Class in C++



Structure Types (Using `typedef`)

```

typedef struct car { ← Old Type
    char engine[50];
    char fuel_type[10];
    int fuel_tank_cap;
    int seating_cap;
    float city_mileage;
}car; ← New Type
  
```

`car` becomes a new data type.

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Structure

Structures within Structures

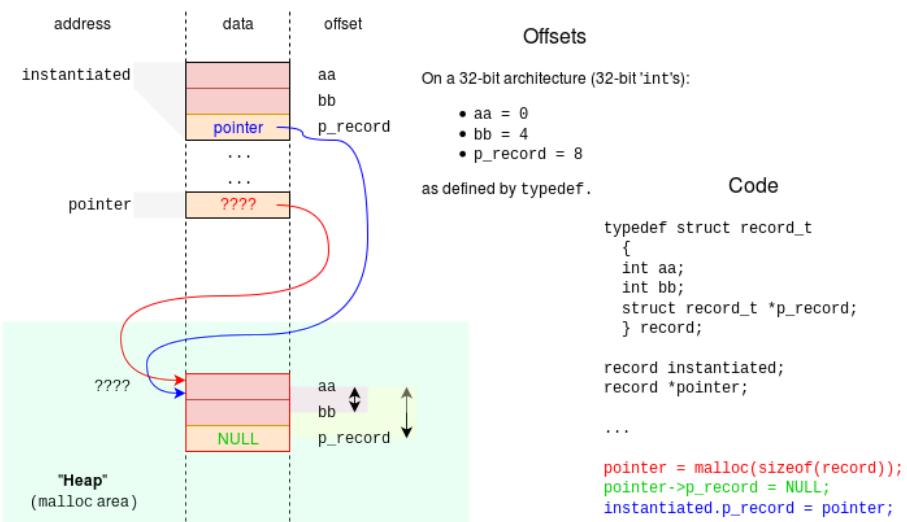
1. C define a variable of structure type as a member of other structure type.
2. The syntax to define the structure within structure is


```

struct <struct_name>{
    <data_type> <variable_name>;
    struct <struct_name>
        { <data_type> <variable_name>;
        .....}<struct_variable>;
    <data_type> <variable_name>;
};
```



Use of struct in C



```
In [ ]: #include <iostream>
using namespace std;

struct MyVector {
    int n;
    int* m;
    void init(int dim);
    void print();
};

void MyVector::init(int dim){
    n = dim;
    m = new int[n];

    for(int i = 0; i < n; i++){
        m[i] = 0;
    }
}

void MyVector::print(){
    cout << "(";
    for(int i = 0; i < n-1; i++){
        cout << m[i] << ", ";
    }
    cout << m[n-1] << ")" << endl;
}

int main(){
    cout << "Welcome to Learning C++ Data Structure and Algorithm" << endl;
    cout << "Topic: OOP - Class" << endl << endl;

    MyVector v;
    v.init(3);
    v.m[0] = 3;
    v.print();

    delete [] v.m;

    return 0;
}
```

Class in C++

```

class myClass {
    public:
        int myNum;
    private:
        string myStr;
    protected:
        float myFloat;
}

public:
    myClass() {
    ...
}
myClass(int x, string y){
    ...
}

private:
    void myMethodOne() {
    ...
}

protected:
    void myMethodTwo() {
    ...
}
};

```

Data and Access Modifiers Constructors Methods

Access Modifiers are used to assign the accessibility to the class members.

A constructor is a member function which initialises objects of a class.

Methods are functions that perform operations on the class data.

enjoyalgorithms.com

```

#include<iostream>
using namespace std;
class student
{
    private :
        int id;
        char name[20];
    public :
        int a;
        Void getdata(void);
        Void display (void)
        {
            cout << id << '\t' << name << endl;
        }
};

```

Data Members

Member Functions

```

keyword
class classname
{
    Access specifiers: //private/public/protected
    Data members/variables; //variables
    Member functions () {} //methods
};

//end of class with a semicolon

```

Struct vs Class

Structure	Class
It is a value type.	It is a reference type.
Its object is created on the stack memory.	Its object is created on the heap memory.
It does not support inheritance.	It supports inheritance.
The member variable of structure cannot be initialized directly.	The member variable of class can be initialized directly.
It can have only parameterized constructor.	It can have all the types of constructor and destructor.

Visibility in Class

Basic concepts
Friends and static members

Constructors and the destructor
Object pointers and the copy constructor

Visibility

- We can/must set visibility of members in a class:
 - Public members can be accessed anywhere.
 - Private members can be accessed only in the class.
 - Protected members will be discussed later in ~~this~~ semester.
- These three keywords are the visibility modifiers.
- By default, all members' visibility level is private.
 - That is why ~~v.~~init(5) generates a compilation error; init() is private and cannot be invoked outside the class (e.g., in the main function).
- By setting visibility, we can hide/open our instance members.
 - Usually all instance variables are private.
 - Let's see how to do this.

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Basic concepts
Friends and static members

Constructors and the destructor
Object pointers and the copy constructor

Encapsulation

- The concepts of packaging (grouping member variables and member functions) and data hiding together form the concept of "encapsulation".
 - Roughly speaking, we pack data (member variables) into a black box and provide only controlled interfaces (member functions) for others to access these data.
 - Others should not even know how those interfaces are implemented.
- For OOP, there are three main characteristics/functionalitys:
 - ~~– Encapsulation.~~
 - Inheritance.
 - Polymorphism.
- The last two will be discussed later in ~~this~~ semester.
Later

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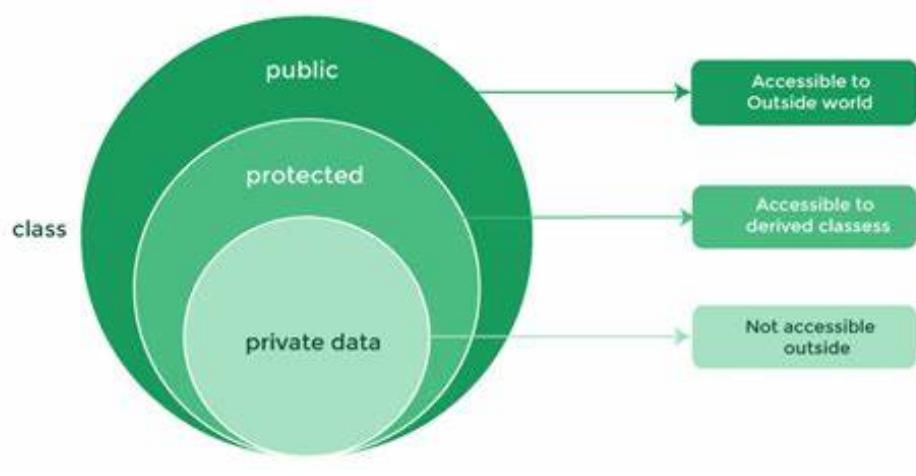
Encapsulation

Access Control

- ▶ We use the **private section** of the struct to **encapsulate** our data type as in our ADTs in C
- ▶ But if the struct is defined in the header file, doesn't this **break encapsulation?**
 - ◆ No! Encapsulation is not about *hiding information* from the user, it's about ensuring that he uses our ADT **in a safe way**, which we achieve by enabling him to use only its public interface
- ▶ However, unlike ADTs in C style, we do not get **compile-time** encapsulation
 - ◆ In other words, changing the private part of a struct will require **recompiling files** that use the ADT

Encapsulation

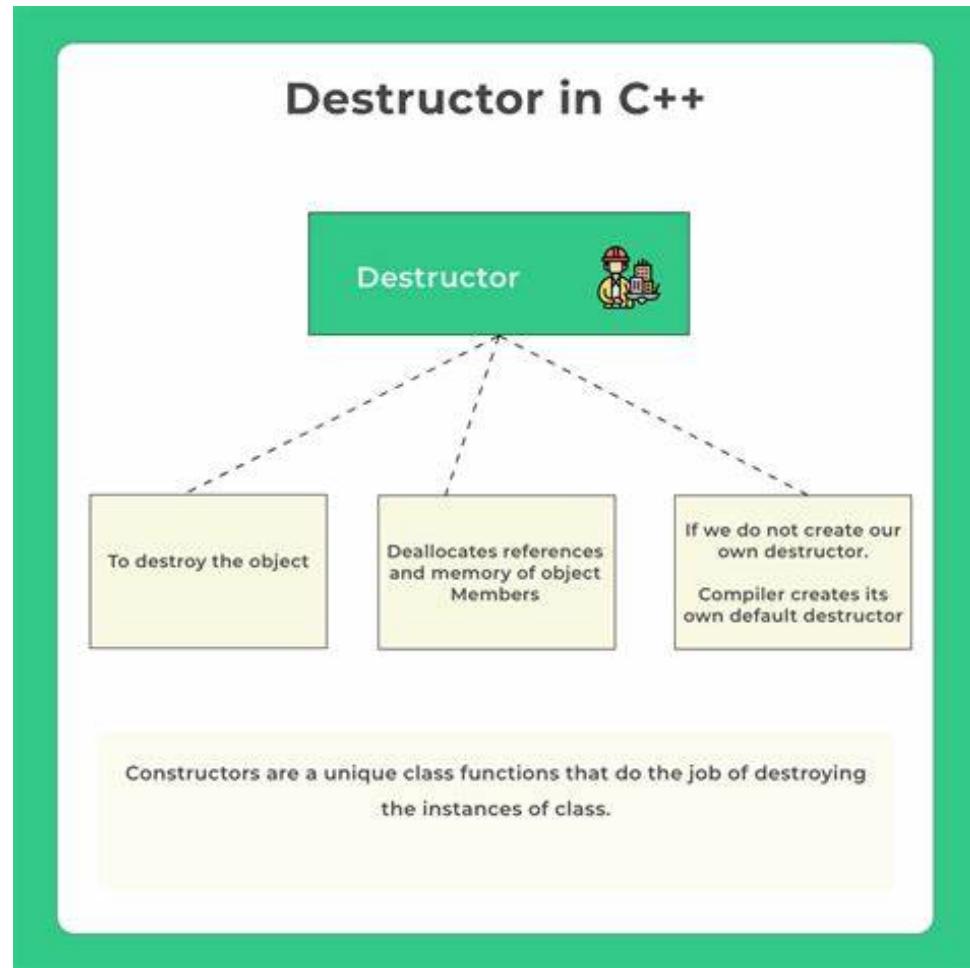
- **Encapsulation** is the **packing of data** and **functions** into a single component. The features of **encapsulation** are supported using classes in most object-oriented programming languages, although other alternatives also exist.
- **Encapsulation** is:
 - A language mechanism for restricting access to some of the object's components. (public, private, protected)
 - A language construct that facilitates the bundling of data with the methods (or other functions) operating on that data.



Constructor and Deconstructor

Class Constructors

- A class constructor is a member function whose **purpose is to initialize the private data members** of a class object
- The name of a constructor is **always** the name of the class, and there is no return type for the constructor
- A class **may have several constructors** with different parameter lists. A constructor with no parameters is the **default** constructor
- A constructor is **implicitly and automatically invoked** when a class object is declared--if there are parameters, their values are listed in parentheses in the declaration



Basic concepts Friends and static members	Constructors and the destructor Object pointers and the copy constructor
--	---

Constructors

- A constructor's name is **the same as** the class.
- It does not return anything, not even **void**.
- You can (and usually will) overload them.
- The constructor with **no parameter** is the **default constructor**.
- If, and only if, a programmer does not define **any constructor**, the **compiler** makes a default one which **does nothing**.
- A constructor may be private.
 - Be invoked only by **other constructors**.

```

class MyVector
{
private:
    int n;
    int* m;
public:
    MyVector() // constructors
    MyVector(int dim);
    MyVector(int dim, int value);
    void print();
};

```

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Basic concepts
Friends and static members
Destructors
Constructors and the destructor
Object pointers and the copy constructor

- A destructor is invoked right before an object is **destroyed**.
 - It must be public and have **no** parameter.
 - The **compiler** provides a default destructor that does nothing.
- To define your own destructor, use **~**.
 - Typically we release dynamically allocated space in a destructor.

```

class MyVector {
private:
    int n;
    int* m;
public:
    MyVector();
    MyVector(int dim, int value = 0);
    ~MyVector();
    void print();
};

MyVector::MyVector(){
    n = 0;
    m = nullptr;
}

MyVector::MyVector(int dim, int value){
    n = dim;
    m = new int[n];
    for(int i = 0; i < n; i++){
        m[i] = value;
    }
}

MyVector::~MyVector(){
    delete [] m;
}

void MyVector::print(){
    cout << "(";
    for(int i = 0; i < n-1; i++){
        cout << m[i] << ", ";
    }
    cout << m[n-1] << ")" << endl;
}

int main()
{
    if (true)
        MyVector v1(1);
        // no memory leak
    return 0;
}

```

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```

In [ ]: #include <iostream>
using namespace std;

class MyVector {
private:
    int n;
    int* m;
public:
    MyVector();
    MyVector(int dim, int value = 0);
    ~MyVector();
    void print();
};

MyVector::MyVector(){
    n = 0;
    m = nullptr;
}

MyVector::MyVector(int dim, int value){
    n = dim;
    m = new int[n];
    for(int i = 0; i < n; i++){
        m[i] = value;
    }
}

MyVector::~MyVector(){
    delete [] m;
}

void MyVector::print(){
    cout << "(";
    for(int i = 0; i < n-1; i++){
        cout << m[i] << ", ";
    }
    cout << m[n-1] << ")" << endl;
}

int main()
{
    if (true)
        MyVector v1(1);
        // no memory leak
    return 0;
}

```

```
cout << "Welcome to Learning C++ Data Structure and Algorithm" << endl;
cout << "Topic: OOP - Class" << endl << endl;

MyVector v1(1);
MyVector v2(3, 8);

v1.print();
v2.print();

return 0;
}
```

```
In [ ]: #include <iostream>
using namespace std;
```

```
class A {
public:
    A() { cout << "A\n"; }
    ~A() { cout << "~A\n"; }
};

class B {
private:
    A a;
public:
    B() { cout << "B\n"; }
    ~B() { cout << "~B\n"; }
};

int main(){
    B b;

    return 0;
}
```

Static Member

Types of variables in C++

```
class GFG {
public:
    static int a; Static Variable
    int b; Instance Variable
public:
    func()
    {
        int c; Local Variable
    };
};
```



Basic concepts Constructors and the destructor
Friends and static members Object pointers and the copy constructor

Static members

- A class contains some instance variables and functions.
 - Each object has its own copy of instance variables and functions.
- A member variable/function may be an attribute/operation of a class.
 - When the attribute/operation is class-specific rather than object-specific.
 - A class-specific attribute/operation should be identical for all objects.
- These variables/functions are called static members.

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Basic concepts Constructors and the destructor
Friends and static members Object pointers and the copy constructor

Static members: an example

- We have to initialize a static variable globally.
- To access static members, use `class name::member name`.

```
int Window::barColor = 0; // default
int Window::getBarColor()
{
    return barColor;
}

void Window::setBarColor(int color)
{
    barColor = color;
}
```

```
int main()
{
    Window w; // not used
    cout << Window::getBarColor();
    cout << "\n";
    Window::setBarColor(1);
    return 0;
}
```

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Constant Object

Motivations and prerequisites
Assignment and self-assignment operators

Comparison and indexing operators
Addition operators

Constant objects

- A constant object cannot invoke a function that modifies its instance variables.
 - In C++, functions that may be invoked by a constant object must be declared as a **constant instance function**.
- For a constant instance function:
 - It can be called by non-constant objects.
 - It cannot modify any instance variable.
- For a non-constant instance function:
 - It cannot be called by constant objects even if no instance variable is modified.

```
class MyVector
{
private:
    int n;
    int* m;
public:
    MyVector();
    MyVector(int dim, int v[]);
    MyVector(const MyVector& v);
    ~MyVector();
    void print() const;
};
```

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Member Initializers

Motivations and prerequisites
Assignment and self-assignment operators

Comparison and indexing operators
Addition operators

✓ Member initializers

- We need a **member initializer**.
 - ✓ A specific operation for initializing an instance variable.
 - Can also be used for initializing non-constant instance variables.
 - ✓ Member initializers are used a lot in general.

```
MyVector::MyVector() : n(0)
{
    m = nullptr;
}
MyVector:: MyVector(int dim, int v[]) : n(dim)
{
    for(int i = 0; i < n; i++)
        m[i] = v[i];
}
MyVector:: MyVector(const MyVector& v) : n(v.n)
{
    m = new double[n];
    for(int i = 0; i < n; i++)
        m[i] = v.m[i];
}
```

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Overloading

Overloading in C++

- ❑ What is overloading
 - Overloading means assigning multiple meanings to a function name or operator symbol
 - It allows multiple definitions of a function with the same name, but different signatures.
- ❑ C++ supports
 - Function overloading
 - Operator overloading

3

Ritika sharma

Function signature and Parameters

MAJU

- **Function signature** is the combination of the **function name** and the **parameter list**.
- **Variables** defined in the **function header** are known as **formal parameters**.
- When a **function** is **invoked**, you pass a value to the parameter. This **value** is referred to as **actual parameter or argument**.

Motivations and prerequisites
Assignment and self-assignment operators

Comparison and indexing operators
Addition operators

Overloading an operator

- An operator is overloaded by “implementing a **special instance function**”.
 - It cannot be implemented as a static function.
- Let op be the operator to be overloaded, the “special instance function” is always named


```
operator op
```

 - The keyword **operator** is used for overloading operators.
- Let’s overload **==** for **MyVector**.

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Motivations and prerequisites
Assignment and self-assignment operators

Comparison and indexing operators
Addition operators

Overloading ==

- To overload **==**, simply do this:

```
class MyVector
{
private:
    int n;
    double* m;
public:
    // others
    bool operator== (const MyVector& v) const
    {
        if(this->n != v.n)
            return false;
        else {
            for(int i = 0; i < n; i++) {
                if(this->m[i] != v.m[i])
                    return false;
            }
        }
        return true;
    }
}
```

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Motivations and prerequisites
Assignment and self-assignment operators

Comparison and indexing operators
Addition operators

Preventing assignments and copying

- In some cases, we **disallow** assignments between objects of a certain class.
 - To do so, overload the assignment operator as a **private** member.
- In some cases, we disallow creating an object by **copying** another object.
 - To do so, implement the copy constructor as a **private** member.
- The copy constructor, assignment operator, and destructor form a group.
 - If there is no pointer, none of them is needed.
 - If there is a pointer, all of them are needed.

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<p>Basic concepts Friends and static members</p>	<p>Constructors and the destructor Object pointers and the copy constructor</p>
--	---

friend for functions and classes

- To “open” private members, another way is to declare “**friends**.”
- One class can allow its friends to access its private members.
- Its friends can be global functions or other classes.
 - Then inside **test()** and member functions of **Test**, those private members of **MyVector** can be accessed.
 - MyVector** cannot access **Test**’s members.

```
class MyVector
{
    // ...
    friend void test();
    friend class Test;
};
```

```
In [ ]: #include <iostream>
#include <cstdlib>
using namespace std;

class MyVector {
private:
    int n;
    double* m;
public:
    MyVector();
    MyVector(int dim, double v[]);
    MyVector(const MyVector& v); //Deep Copy
    ~MyVector();
    void print() const;
    bool operator==(const MyVector& v) const;
    bool operator!=(const MyVector& v) const;
    double operator[](int i) const;
    double& operator[](int i);
    MyVector& operator=(const MyVector& v);
    const MyVector& operator+=(const MyVector& v);
    const MyVector operator+(const MyVector& v);
    const MyVector operator+(double d);
    friend const MyVector operator+(const MyVector& v, double d);
};

MyVector::MyVector() : n(0), m(nullptr) {}

MyVector::MyVector(int dim, double v[]) : n(dim){
    m = new double[n];
    for(int i = 0; i < n; i++){
        m[i] = v[i];
    }
}

//Deep Copy
MyVector::MyVector(const MyVector& v): n(v.n){
    m = new double[n];
    for(int i = 0; i < n; i++){
        m[i] = v.m[i];
    }
}
```

```

MyVector::~MyVector(){
    delete [] m;
}

void MyVector::print() const {
    cout << "(";
    for(int i = 0; i < n - 1; i++){
        cout << m[i] << ", ";
    }
    cout << m[n - 1] << ")";
}

bool MyVector::operator==(const MyVector& v) const {
    if(n != v.n)
        return false;
    else {
        for(int i = 0; i < n; i++){
            if(m[i] != v.m[i])
                return false;
        }
    }
    return true;
}

bool MyVector::operator!=(const MyVector& v) const {
    return (*this == v);
}

double MyVector::operator[](int i) const {
    if (i < 0 || i >= n){
        exit(1); //terminate the program! required <cstdlib>
    }
    return m[i];
}

double& MyVector::operator[](int i){
    if (i < 0 || i >= n){
        exit(1);
    }
    return m[i];
}

MyVector& MyVector::operator=(const MyVector& v){
    if(this != &v){
        if(this->n != v.n){
            delete [] this->m;
            this->n = v.n;
            this->m = new double[this->n];
        }
        for(int i = 0; i < n; i++){
            this->m[i] = v.m[i];
        }
    }
    return *this;
}

const MyVector& MyVector::operator+=(const MyVector& v){
    if(this->n == v.n){
        for(int i = 0; i < n; i++){

```

```

        this->m[i] += v.m[i];
    }
}
return *this;
}

const MyVector MyVector::operator+ (const MyVector& v){
    MyVector sum(*this); //MyVector(const MyVector& v)
    sum += v;
    return sum;
}

const MyVector MyVector::operator+ (double d){
    MyVector sum(*this);
    for(int i = 0; i < n; i++){
        sum[i] += d;
    }
    return sum;
}

const MyVector operator+ (const MyVector& v, double d){
    MyVector sum(v);
    for(int i = 0; i < v.n; i++){
        sum[i] += d;
    }
    return sum;
}

const MyVector operator+ (double d, const MyVector& v){
    return v + d;
}

const MyVector operator+ (const MyVector& v1, const MyVector& v2){
    MyVector sum(v1);
    sum += v2;
    return sum;
}

int main(){
    cout << "Welcome to Learning C++ Data Structure and Algorithm" << endl;
    cout << "Topic: OOP - Class" << endl << endl;

    double d1[5] = {1, 2, 3, 4, 5};
    MyVector a1(5, d1);

    double d2[4] = {1, 2, 3, 4};
    MyVector a2(4, d2);

    MyVector a3(a1);

    cout << (a1 == (a2) ? "Y" : "N");
    cout << endl;
    cout << (a1 == (a3) ? "Y" : "N");
    cout << endl;

    cout << "a1[3]: " << a1[3] << endl;
    a1[1] = 999;
    cout << "a1[1]: " << a1[1] << endl;
}

```

```

cout << "a1: ";
a1.print();
cout << endl;

cout << "a2: ";
a2.print();
cout << endl;

cout << "a3: ";
a3.print();
cout << endl;

a1 += a3;
cout << "a1 + a3: ";
a1.print();
cout << endl;

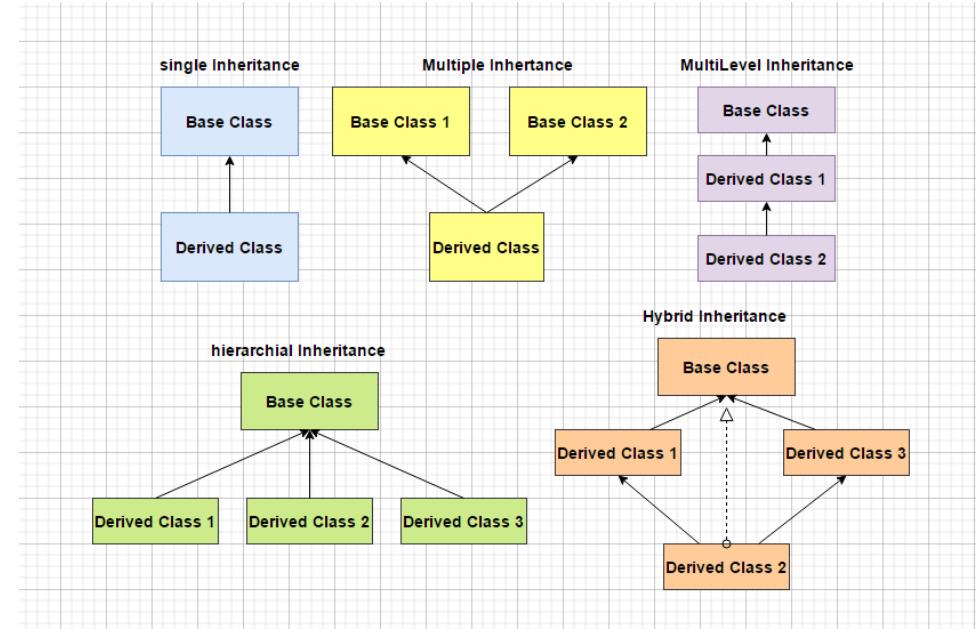
a1 = a1 + 4.2;
cout << "a1 + 4.2: ";
a1.print();
cout << endl;

a2 = 100+ a2;
cout << "100 + a2: ";
a2.print();
cout << endl;

return 0;
}

```

Inheritance



Base-class member-access specifier	Type of inheritance		
	public inheritance	protected inheritance	private inheritance
public	public in derived class. Can be accessed directly by member functions, friend functions and nonmember functions.	protected in derived class. Can be accessed directly by member functions and friend functions.	private in derived class. Can be accessed directly by member functions and friend functions.
protected	protected in derived class. Can be accessed directly by member functions and friend functions.	protected in derived class. Can be accessed directly by member functions and friend functions.	private in derived class. Can be accessed directly by member functions and friend functions.
private	Hidden in derived class. Can be accessed by member functions and friend functions through public or protected member functions of the base class.	Hidden in derived class. Can be accessed by member functions and friend functions through public or protected member functions of the base class.	Hidden in derived class. Can be accessed by member functions and friend functions through public or protected member functions of the base class.

Inheritance An example Polymorphism

Inheritance visibility

- In general, the visibility of a member in a child class depends on:
 - The member visibility by the parent.
 - The inheritance modifier.

Member visibility by the parent	Inheritance modifier		
	public	protected	private
✓ public	public	protected	private
✓ protected	protected	protected	private
private	private	private	private

- If you have no idea, just use public inheritance.

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Inheritance An example Polymorphism

Invoking parent class' constructors

- The parent class' constructor will not be inherited.
- One of them will be invoked before the child class' constructor is invoked.
 - >Create the parent before creating the child!
- ✓ If not specified, the parent's **default** constructor will be invoked.

```
MyVector::MyVector() : n(0), m(nullptr)
{
}

MyVector2D::MyVector2D()
{
  this->n = 2;
  // this->m = nullptr is redundant
}
```

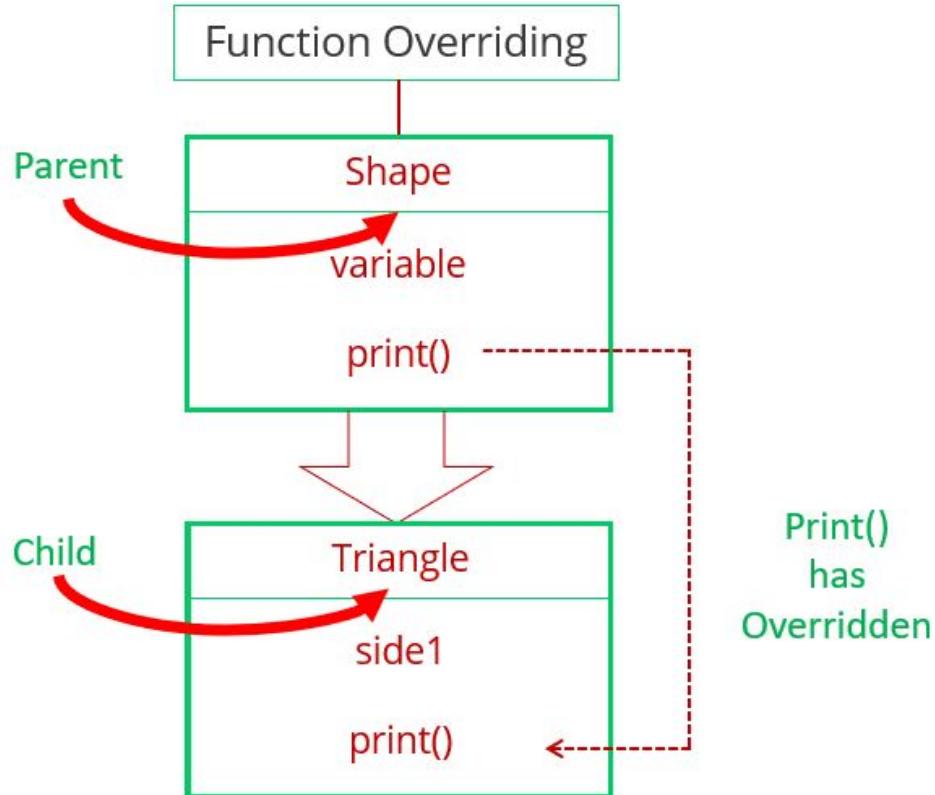
```
int main()
{
  MyVector2D v;
  return 0;
}
```

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Inheritance	An example	Polymorphism
<h2>Invoking copy constructors</h2> <ul style="list-style-type: none"> How about the copy constructor? If we <u>do not define one for the child</u>, the system provides a <u>default</u> one. <u>Before</u> the child's <u>default</u> copy constructor is invoked, the parent's copy constructor will be <u>automatically</u> invoked. 		
<pre>MyVector::MyVector(const MyVector& v) { this->n = v.n; this->m = new double[n]; for(int i = 0; i < n; i++) this->m[i] = v.m[i]; } class MyVector2D : public MyVector { public: MyVector2D(); MyVector2D(double m[]); // no copy constructor };</pre>		
Programming Design – Inheritance and Polymorphism	10 / 62	Ling-Chieh Kung (NTU IM)

Inheritance	An example	Polymorphism
<h2>Invoking parent class' destructor</h2> <ul style="list-style-type: none"> When an <u>object of the child class</u> is to be destroyed: <ul style="list-style-type: none"> First the <u>child's</u> destructor is invoked. <u>Then</u> the parent's destructor is invoked <u>automatically</u>, even if we do not define a destructor for the child. 		
<pre>MyVector::~MyVector() { delete [] m; } class MyVector2D : public MyVector { public: MyVector2D(); MyVector2D(double m[]); // no destructor };</pre>		
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Overriding vs Overloading



Overloading vs. Overriding

- Overriding a base class member function is similar to overloading a function or operator
 - But for overriding, definitions are distinguished by their scopes rather than by their signatures
- C++ can distinguish method definitions according to either static or dynamic type
 - Depends on whether a method is virtual or not
 - Depends on whether called via a reference or pointer vs. directly on an object
 - Depends on whether the call states the scope explicitly (e.g., `Foo::baz();`)

CSE 332: C++ Overloading

```
In [ ]: #include <iostream>
#include <cstdlib>
using namespace std;
```

```
class MyVector {
protected:
    int n;
    double* m;
public:
    MyVector();
```

```

MyVector(int dim, double v[]);
MyVector(const MyVector& v); //Deep Copy
~MyVector();
void print() const;
bool operator==(const MyVector& v) const;
bool operator!=(const MyVector& v) const;
double operator[](int i) const;
double& operator[](int i);
MyVector& operator=(const MyVector& v);
const MyVector& operator+=(const MyVector& v);
const MyVector operator+(const MyVector& v);
const MyVector operator+(double d);
friend const MyVector operator+(const MyVector& v, double d);
};

MyVector::MyVector() : n(0), m(nullptr) {
}

MyVector::MyVector(int dim, double v[]) : n(dim){
    m = new double[n];
    for(int i = 0; i < n; i++){
        m[i] = v[i];
    }
}

//Copy Constructor: Deep Copy
MyVector::MyVector(const MyVector& v): n(v.n){
    m = new double[n];
    for(int i = 0; i < n; i++){
        m[i] = v.m[i];
    }
}

MyVector::~MyVector(){
    delete [] m;
}

void MyVector::print() const {
    cout << "(";
    for(int i = 0; i < n - 1; i++){
        cout << m[i] << ", ";
    }
    cout << m[n - 1] << ")";
}

bool MyVector::operator==(const MyVector& v) const {
    if(n != v.n)
        return false;
    else {
        for(int i = 0; i < n; i++){
            if(m[i] != v.m[i])
                return false;
        }
    }
    return true;
}

bool MyVector::operator!=(const MyVector& v) const {
    return (*this == v);
}

```

```

double MyVector::operator[] (int i) const {
    if (i < 0 || i >=n){
        exit(1); //terminate the program! required <cstdlib>
    }
    return m[i];
}

double& MyVector::operator[] (int i){
    if (i < 0 || i >= n){
        exit(1);
    }
    return m[i];
}

MyVector& MyVector::operator= (const MyVector& v){
    if(this != &v){
        if(this->n != v.n){
            delete [] this->m;
            this->n = v.n;
            this->m = new double[this->n];
        }
        for(int i = 0; i < n; i++){
            this->m[i] = v.m[i];
        }
    }
    return *this;
}

const MyVector& MyVector::operator+= (const MyVector& v){
    if(this->n == v.n){
        for(int i = 0; i < n; i++){
            this->m[i] += v.m[i];
        }
    }
    return *this;
}

const MyVector MyVector::operator+ (const MyVector& v){
    MyVector sum(*this); //MyVector(const MyVector& v)
    sum += v;
    return sum;
}

const MyVector MyVector::operator+ (double d){
    MyVector sum(*this);
    for(int i = 0; i < n; i++){
        sum[i] += d;
    }
    return sum;
}

const MyVector operator+ (const MyVector& v, double d){
    MyVector sum(v);
    for(int i = 0; i < v.n; i++){
        sum[i] += d;
    }
    return sum;
}

const MyVector operator+ (double d, const MyVector& v){
    return v + d;
}

```

```

}

const MyVector operator+ (const MyVector& v1, const MyVector& v2){
    MyVector sum(v1);
    sum += v2;
    return sum;
}

//Inheritance
class MyVector2D : public MyVector {
public:
    MyVector2D();
    MyVector2D(double m[]);
    void setValue(double d1, double d2);
    void print() { MyVector::print(); }
    void print() const {
        cout << "2D: ";
        MyVector::print();
    };
};

MyVector2D::MyVector2D(){
    this->n = 2;
}

MyVector2D::MyVector2D(double m[]) : MyVector(2, m){

}

void MyVector2D::setValue(double d1, double d2){
    if (this->m == nullptr){
        this->m = new double[2];
    }
    this->m[0] = d1;
    this->m[1] = d2;
}

int main(){
    cout << "Welcome to Learning C++ Data Structure and Algorithm" << endl;
    cout << "Topic: OOP - Class" << endl << endl;

    double d1[2] = {1, 2};

    const MyVector2D v(d1);
    v.print();
    cout << endl;

    MyVector2D u;
    u.setValue(3, 4);
    u.print();
    cout << endl;

    return 0;
}

```

Polymorphism

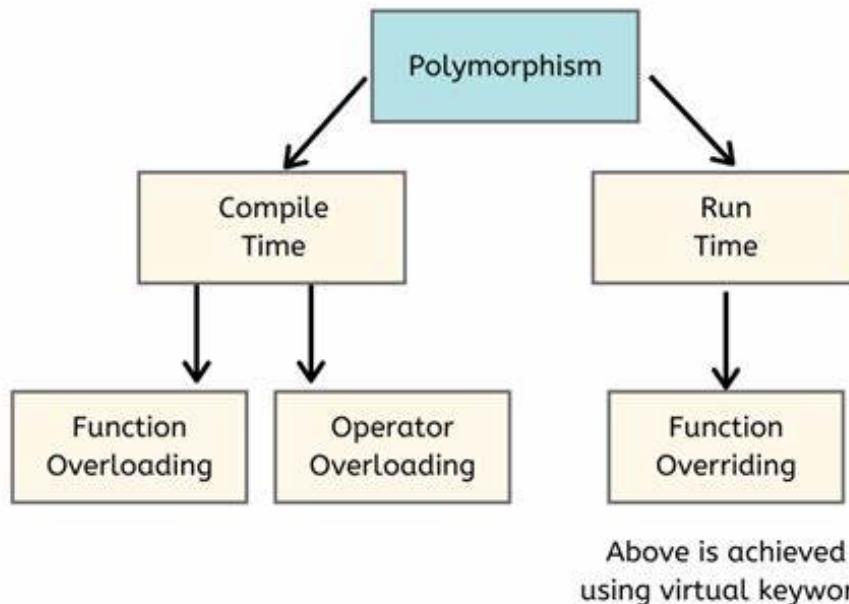
Overview of C++ Polymorphism

- Two main kinds of types in C++: native and user-defined
 - “User” defined types: declared classes, structs, unions
 - including types provided by the C++ standard libraries
 - Native types are “built in” to the C++ language itself: int, long, float, ...
 - A `typedef` creates a new type *name* for another type (type aliasing)
- Public inheritance creates sub-types
 - Inheritance only applies to user-defined classes (and structs)
 - A publicly derived class is-a subtype of its base class
 - Known as “inheritance polymorphism”
- Template parameters also induce a subtype relation
 - Known as “interface polymorphism”
 - We’ll cover how this works in depth, in later sessions
- Liskov Substitution Principle (for both kinds of polymorphism)
 - if S is a subtype of T, then wherever you need a T you can use an S

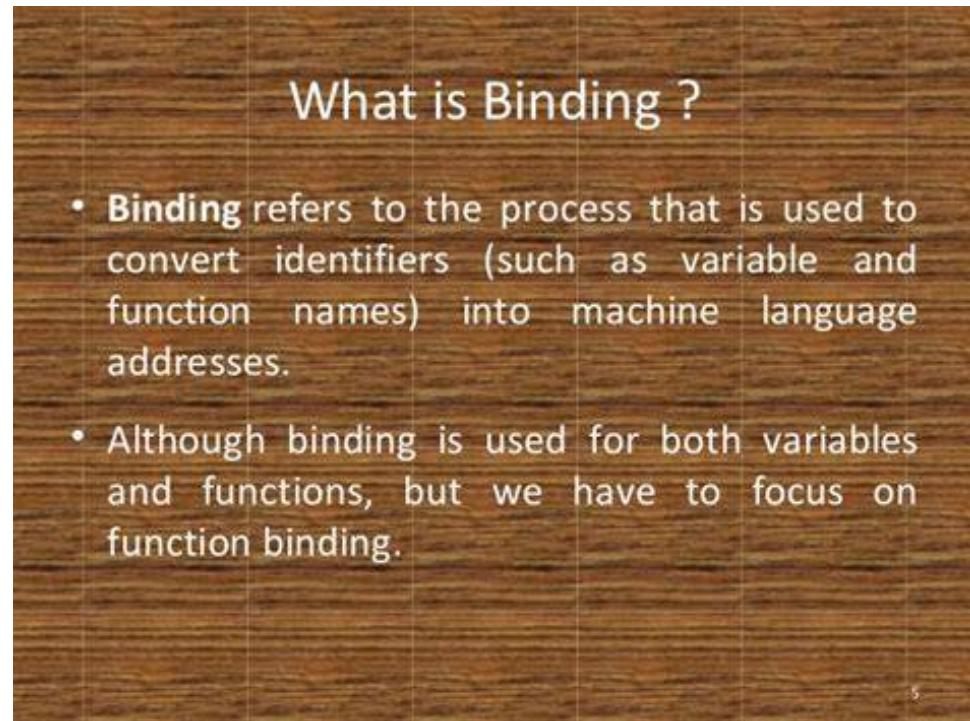
CSE 332: C++ Polymorphism



Polymorphism in C++



Early Binding and Late Binding



Sr. No.	Early Binding	Late Binding
1.	Early binding is the process of linking a function with an object during the compilation process.	Late binding is a run-time polymorphism with method overriding.
2.	Static binding is another name for early binding.	Dynamic binding is another name for late binding.
3.	Early binding happens at compile time.	Late binding happens at run time.
4.	Execution speed is faster in early binding.	Execution speed is lower in late binding.
5.	The class information is used by Early binding to resolve method calls.	The object is used by Late binding to resolve method calls.

Inheritance	An example	Polymorphism
<h3>Early binding vs. late binding</h3> <ul style="list-style-type: none"> When we do A a = b or A* a = &b, we are using polymorphism. For A a = b, the system does early binding: <ul style="list-style-type: none"> a occupies only four bytes for storing i. a does not have a space for storing j. Its type is set to be A at compilation. For A* a = &b, the system does late binding: <ul style="list-style-type: none"> a is just a pointer. It can point to an A object or a B object. Its “type” can be set at the run time. <pre style="border: 1px solid blue; padding: 10px;"> class A { protected: i public: void a() { cout << "a\n"; } void f() { cout << "af\n"; } }; class B : public A { private: j public: void b() { cout << "b\n"; } void f() { cout << "bf\n"; } }; </pre>		

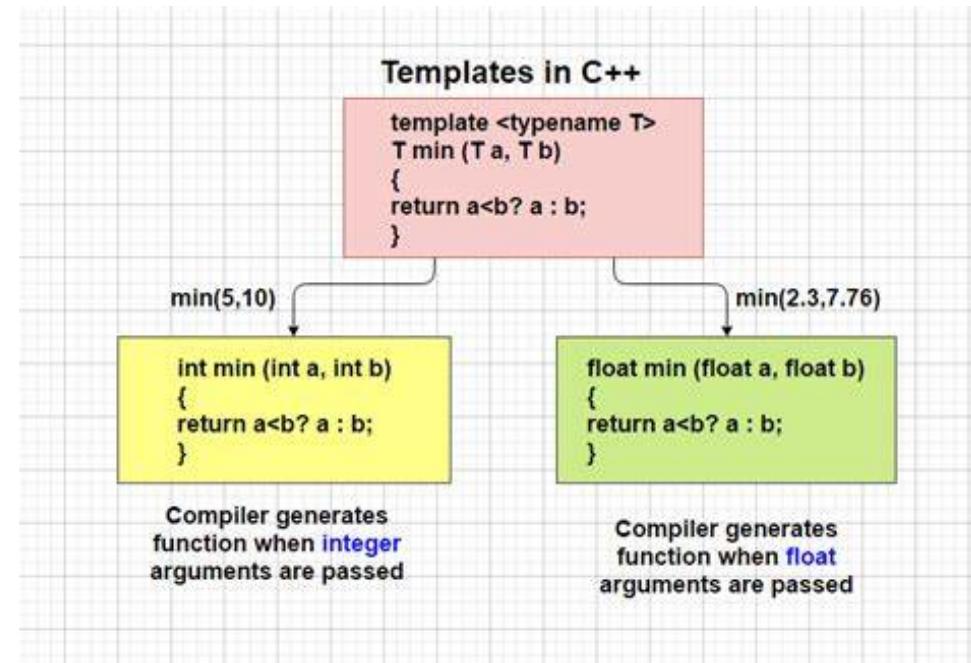
Virtual Function

Inheritance	An example	Polymorphism
	<h2>Virtual functions</h2> <ul style="list-style-type: none"> If we declare a parent's member function to be virtual, its invocation priority will be lower than a child's (if we use late binding). <ul style="list-style-type: none"> A child cannot declare a parent's function as virtual (<u>it is of no use</u>). In summary, we need: <ul style="list-style-type: none"> Late binding + virtual functions. 	
	<pre>class Parent { protected: int x; int y; public: Parent(int a, int b) : x(a), y(b) {} virtual void print() { cout << x << " " << y; } }; class Child : public Parent { protected: int z; public: Child(int a, int b, int c) : Parent(a, b) { z = c; } void print() { cout << z; } };</pre>	

Inheritance	An example	Polymorphism
	<h2>Abstract classes</h2> <ul style="list-style-type: none"> The two virtual functions are different in their natures: <ul style="list-style-type: none"> print() is invoked in the children's implementations. beatMonster() should not be invoked by any one. We may set beatMonster() to be a pure virtual function. <pre>class Character { // ... virtual void beatMonster(int exp) = 0; };</pre> <ul style="list-style-type: none"> Now we do not need to implement it. Moreover, we <u>cannot</u> create Character objects! 	

Inheritance	An example	Polymorphism
	<h2>Summary</h2> <ul style="list-style-type: none"> Polymorphism is a technique to make our program <u>clearer</u>, <u>more flexible</u> and <u>more powerful</u>. <ul style="list-style-type: none"> It is based on <u>inheritance</u>. It is tightly related to <u>function overriding</u>, <u>late binding</u>, and <u>virtual functions</u>. The key action is to “use a parent pointer to point to a child object”. To implement late binding, you need to <ul style="list-style-type: none"> Declare and override virtual functions. Do <u>late binding</u> by using parent pointers to point to child objects. 	

C++ Templates



Templates The standard library <vector> Exception handling

Template declaration

- To declare a type parameter, use the keywords `template` and `typename`.

```
template<typename T>
class TheClassName
{
    // T can be treated as a type inside the class definition block
};
```

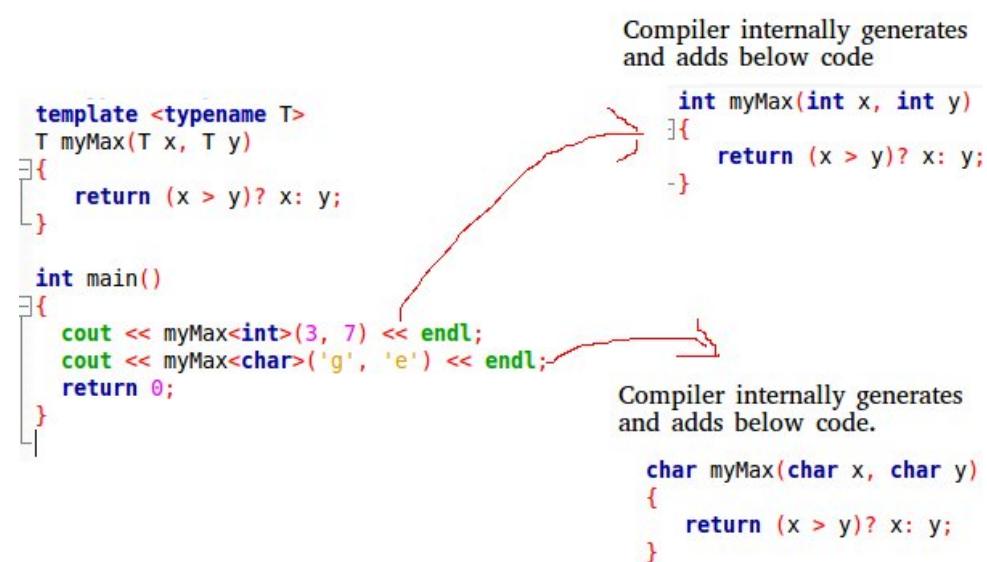
– Some old codes write `class` instead of `typename`. Both are fine.

- We then do this to all member functions:

```
template<typename T>
T TheClassName<T>::f(T t)
{
    // t is a variable whose type is T
};
```

```
template<typename T>
void TheClassName<T>::f(int i)
{
    // follow the rule even if T is not used
};
```

Programming Design – Templates, Vectors, and Exceptions 7 / 55 Ling-Chieh Kung (NTU IM)



In []: `#include <iostream>
#include <cmath>
using namespace std;`

```

template<typename KeyType>
class Character {
protected:
    static const int EXP_LV =100;
    KeyType name;
    int level;
    int exp;
    int power;
    int knowledge;
    int luck;
    void levelUp(int pInc, int kInc, int lInc);
public:
    Character(KeyType n, int lv, int po, int kn, int lu);
    virtual void print();
    virtual void beatMonster(int exp) = 0; //pure virtual function
    KeyType getName();
};

//construct
template<typename KeyType>
Character<KeyType>::Character(KeyType n, int lv, int po, int kn, int lu)
: name(n), level(lv), exp(pow(lv - 1, 2) * EXP_LV),
  power(po), knowledge(kn), luck(lu) {
}

template<typename KeyType>
void Character<KeyType>::print(){
    cout << this->name
        << ": Level " << this->level
        << " (" << this->exp << "/" << pow(this->level, 2) * EXP_LV
        << "), " << this->power << "-" << this->knowledge
        << "-" << this->luck << endl;
}

template<typename KeyType>
KeyType Character<KeyType>::getName(){
    return this->name;
}

template<typename KeyType>
void Character<KeyType>::beatMonster(int exp){
    this->exp += exp;
    while(this->exp >= pow(this->level, 2) * EXP_LV){
        this->levelUp(0, 0, 0);
    }
}

template<typename KeyType>
void Character<KeyType>::levelUp(int pInc, int kInc, int lInc){
    this->level++;
    this->power += pInc;
    this->knowledge += kInc;
    this->luck += lInc;
}

template<typename KeyType>
class Warrior : public Character<KeyType> {

```

```

private:
    static const int PO_LV = 10;
    static const int KN_LV = 5;
    static const int LU_LV = 5;
public:
    Warrior(KeyType n, int lv = 0);
    void print();
    void beatMonster(int exp);
};

template<typename KeyType>
Warrior<KeyType>::Warrior(KeyType n, int lv)
    : Character<KeyType>(n, lv, lv * PO_LV, lv * KN_LV, lv * LU_LV) {}

template<typename KeyType>
void Warrior<KeyType>::print() {
    cout << "Warrior: ";
    Character<KeyType>::print();
}

template<typename KeyType>
void Warrior<KeyType>::beatMonster(int exp){ // function overloading
    this->exp += exp;
    while(this->exp >= pow(this->level, 2) * Character<KeyType>::EXP_LV){
        this->levelUp(PO_LV, KN_LV, LU_LV);
    }
}
}

template<typename KeyType>
class Wizard : public Character<KeyType> {
private:
    static const int PO_LV = 4;
    static const int KN_LV = 9;
    static const int LU_LV = 7;
public:
    Wizard(KeyType n, int lv = 0);
    void print();
    void beatMonster(int exp);
};

template<typename KeyType>
Wizard<KeyType>::Wizard(KeyType n, int lv)
    : Character<KeyType>(n, lv, lv * PO_LV, lv * KN_LV, lv * LU_LV) {}

template<typename KeyType>
void Wizard<KeyType>::print() {
    cout << "Wizard: ";
    Character<KeyType>::print();
}

template<typename KeyType>
void Wizard<KeyType>::beatMonster(int exp){ // function overloading
    this->exp += exp;
    while(this->exp >= pow(this->level, 2) * Character<KeyType>::EXP_LV){
        this->levelUp(PO_LV, KN_LV, LU_LV);
    }
}

```

```
template<typename KeyType>
class Team {
private:
    int memberCount;
    Character<KeyType>* member[10];
public:
    Team();
    ~Team();
    void addWarrior(KeyType name, int lv);
    void addWizard(KeyType name, int lv);
    void memberBeatMonster(KeyType name, int exp);
    void printMember(KeyType name);
};

template<typename KeyType>
Team<KeyType>::Team(){
    memberCount = 0;
    for (int i = 0; i < 10; i++){
        member[i] = nullptr;
    }
}

template<typename KeyType>
Team<KeyType>::~Team(){
    for (int i = 0; i < memberCount; i++){
        delete member[i];
    }
}

template<typename KeyType>
void Team<KeyType>::addWarrior(KeyType name, int lv){
    if (memberCount < 10){
        member[memberCount] = new Warrior<KeyType>(name, lv);
        memberCount++;
    }
}

template<typename KeyType>
void Team<KeyType>::addWizard(KeyType name, int lv){
    if (memberCount < 10){
        member[memberCount] = new Wizard<KeyType>(name, lv);
        memberCount++;
    }
}

template<typename KeyType>
void Team<KeyType>::memberBeatMonster(KeyType name, int exp){
    for(int i = 0; i < memberCount; i++){
        if (member[i]->getName() == name){
            member[i]->beatMonster(exp);
            break;
        }
    }
}

template<typename KeyType>
void Team<KeyType>::printMember(KeyType name){
    for(int i = 0; i < memberCount; i++){
        if(member[i]->getName() == name){
            member[i]->print();
            break;
        }
    }
}
```

```

        }
    }

}

int main(){
    cout << "Welcome to Learning C++ Data Structure and Algorithm" << endl;
    cout << "Topic: OOP - Class" << endl << endl;

    Team<string> t1;
    t1.addWarrior("Alice", 1);
    t1.memberBeatMonster("Alice", 10000);
    t1.addWizard("Bob", 2);
    t1.printMember("Alice");

    Team<int> t2;
    t2.addWarrior(1, 1);
    t2.memberBeatMonster(1, 10000);
    t2.addWizard(2, 2);
    t2.printMember(1);

    return 0;
}

```

Abstract Data Type

Abstract data types	Array-based implementation	Link-based implementation
Essential concepts		The ADT bag
Abstract data types	<p>After one obtains basic programming ability, she/he may (should) study the role of data abstraction (and algorithms) in problem solving.</p> <ul style="list-style-type: none"> – Apply data abstraction to increase modularity. – Conceptually, build a “wall” between a program and its data structures. <p>She/he should design abstract data types (ADTs) first before doing implementations (i.e., building data structures).</p> <ul style="list-style-type: none"> – Data abstraction focuses on what the operations do to data, not on their implementation. <p>A data structure is an implementation of an ADT.</p> <ul style="list-style-type: none"> – It is defined within a programming language to store a collection of data. – E.g., array-based or link-based implementation of ADT bag. 	

Data Structures – The Data Structure “Bag” 12 / 88 Ling-Chieh Kung (NTU IM)

Abstract data types	Array-based implementation	Link-based implementation
Essential concepts		The ADT bag
<h2>Abstract data types</h2>		
<ul style="list-style-type: none"> The interface of a class identifies the publicly accessible methods (and data). A complete interface provides methods for accomplishing any reasonable task consistent with the responsibilities of the class. <ul style="list-style-type: none"> Very important. A minimal interface provides only essential methods. <ul style="list-style-type: none"> Easier for understanding and maintenance. Less important than completeness. 		

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ADT Bag

Abstract data types	Array-based implementation	Link-based implementation
Essential concepts		The ADT bag
<h2>The ADT “Bag”</h2>		
<ul style="list-style-type: none"> We want to implement an ADT <u>“Bag”</u> of items whose types are <u>“ItemType”</u>: <ul style="list-style-type: none"> -getCurrentSize(): integer -isEmpty(): boolean -add(newEntry: ItemType): boolean // duplications are allowed - remove(anEntry: ItemType): boolean // remove only one copy - clear(): void - getFrequencyOf(anEntry: ItemType): integer - contains(anEntry: ItemType): boolean - print(): void Specifications indicate what the operations do, not how to implement Let’s begin with a bag of C++ strings. 		

Data Structures – The Data Structure “Bag” 21 / 88 Ling-Chieh Kung (NTU IM)

Abstract data types	Array-based implementation	Link-based implementation
Essential concepts		The ADT bag
<h2>Implementing the ADT “Bag”</h2>		
<ul style="list-style-type: none"> Implementing the ADT as a C++ class provides ways to enforce a wall: <ul style="list-style-type: none"> This prevents one to access the data without using the defined operations. 		

Figure 3-1 (Carrano and Henry, 2013)

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Abstract data types	Array-based implementation	Link-based implementation
Essential concepts		The ADT bag
<h2>Implementing the ADT “Bag”</h2>		
<ul style="list-style-type: none"> The abstract class BagInterface defines the behaviors of the ADT. <ul style="list-style-type: none"> In many cases, no member variable is needed for the abstract class. <pre>class BagInterface { public: virtual int getCurrentSize() const = 0; virtual bool isEmpty() const = 0; virtual bool add(const string& newEntry) = 0; virtual bool remove(const string& anEntry) = 0; virtual void clear() = 0; virtual int getFrequencyOf(const string& anEntry) const = 0; virtual bool contains(const string& anEntry) const = 0; virtual void print() const = 0; };</pre>		

Abstract data types	Array-based implementation	Link-based implementation
A static array implementation		A dynamic array implementation
<h2>Implementing the ADT “Bag”</h2>		
<ul style="list-style-type: none"> Now it is time to choose a data structure. For the ADT bag, we will try two different data structures: <ul style="list-style-type: none"> An array and a linked list. Let's use a fixed-size (static) array first. <ul style="list-style-type: none"> In a fixed-size array, each item occupies one entry of the array. These items are unsorted. A one-dimensional unsorted array will be used. <p>Figure 3-2 (Carrano and Henry, 2013)</p>		

Abstract data types	Array-based implementation	Link-based implementation
A static array implementation		A dynamic array implementation
<h2>Implementing the ADT “Bag”</h2>		
<ul style="list-style-type: none"> To implement the ADT, we write a (concrete) class to inherit BagInterface and implement all the operations. 	<pre>✓class ArrayBag : public BagInterface { private: // ... public: ArrayBag(); int getCurrentSize() const; bool isEmpty() const; bool add(const string& newEntry); bool remove(const string& anEntry); void clear(); bool contains(const string& anEntry) const; int getFrequencyOf(const string& anEntry) const; void print() const; };</pre>	

Abstract data types	Array-based implementation	Link-based implementation
A static array implementation		A dynamic array implementation
<h2>Using a fixed-size array</h2> <ul style="list-style-type: none"> What else should we keep track of when we use a <u>fixed array</u>? <ul style="list-style-type: none"> The <u>maximum size</u> of the array. The <u>number of items</u> currently stored. We add <u>member variable declarations</u> into the private section. Note that only a <u>constant static variable</u> may be initialized in the class definition part. <ul style="list-style-type: none"> A <u>non-constant static variable</u> must be initialized outside it. 		

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Abstract data types	Array-based implementation	Link-based implementation
A static array implementation		A dynamic array implementation
<h2>Implementing member functions</h2> <ul style="list-style-type: none"> Some member functions are pretty straightforward. <ul style="list-style-type: none"> Note that these should be <u>constant</u> member functions. May you add a constructor to set the maximum number of items? 		

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Abstract data types	Array-based implementation	Link-based implementation
A static array implementation		A dynamic array implementation
<h2>The member function <u>remove()</u></h2> <ul style="list-style-type: none"> For removing a given item: <ul style="list-style-type: none"> Remove <u>one copy if it exists in the array</u>. Return true if the <u>item exists</u> or false otherwise. Note that we need to determine <u>whether a given item exists</u>. <ul style="list-style-type: none"> This is exactly the function <code>contains()</code>. Moreover, we need to <code>locate</code> that item. We may enhance <u>modularity</u> (and make the development more efficient) by implementing a member function to be a <u>building block</u>. 		

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Abstract data types	Array-based implementation	Link-based implementation
A static array implementation		
		A dynamic array implementation

The member function getIndexOf()

- Let's implement a member function getIndexOf().
 - getIndexOf(anEntry: ItemType): integer
 - Given an item, return the index of its first copy in the array or -1 otherwise.

Figure 3-4 (Carrano and Henry, 2013)

- Define things **precisely**: Returning an array **index** or a **rank** of an item?

Abstract data types	Array-based implementation	Link-based implementation
A static array implementation		
		A dynamic array implementation

Privatizing getIndexOf()

- The function getIndexOf() should be **private**!
- If it is public, clients will know some **details** of the private array.
 - Data hiding (encapsulation, the “wall”) would be damaged.
- As long as one thing should not be used by (or is useless to) clients, hide it!

```
class ArrayBag : public BagInterface
{
private:
    static const int DEFAULT_CAPACITY = 6;
    string items[DEFAULT_CAPACITY];
    int itemCount;
    int maxItems;
public:
    ✓ int getIndexOf(const string& target) const;
    // those member functions...
};
```

Figure 3-5 (Carrano and Henry, 2013)

Abstract data types	Array-based implementation	Link-based implementation
A static array implementation		
		A dynamic array implementation

The member function remove()

- With getIndexOf(), remove() can be easily implemented.
- The pseudocode:
- How to “remove the item **while ensuring no hole**?”

```
remove(anEntry)
{
    ✓ Search the array items for anEntry
    if anEntry is in the bag at items[index]
    {
        Decrement the counter itemCount
        remove the item while ensuring no hole
        return true
    }
    else
        return false
}
```

Figure 3-6 (Carrano and Henry, 2013)

Abstract data types	Array-based implementation	Link-based implementation
A static array implementation		A dynamic array implementation
Ensuring no hole (idea 2) <i>hole</i>		
Figure 3-6 (Carrano and Henry, 2013)		

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Abstract data types	Array-based implementation	Link-based implementation
A static array implementation		A dynamic array implementation
Remarks		
<ul style="list-style-type: none"> The array-based implementation of the ADT bag has been completed! ✓ Before we start the implementation, always “<u>design</u>” your program first. <ul style="list-style-type: none"> Design the ADT. In particular, design the operations and behaviors. Write an <u>abstract class</u> to specify the operations. Write a <u>concrete class</u> to inherit the abstract class. <u>Implement the class</u>. For a specific function, <u>pseudocodes</u> are helpful. Test your implementation. Do not forget <u>data hiding</u>. <ul style="list-style-type: none"> If something is not useful for clients, hide it. E.g., do not let clients know whether you starts at <code>items[0]</code> or <code>items[1]</code>. E.g., do not let clients know whether you store your items in an array. 		

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Implement by static Array

```
In [ ]: #include <iostream>
#include <string>
using namespace std;

class BagInterface {
public:
    virtual int getCurrentSize() const = 0;
    virtual bool isEmpty() const = 0;
    virtual bool add(const string& newEntry) = 0;
    virtual bool remove(const string& anEntry) = 0;
    virtual void clear() = 0;
    virtual int getFrequencyOf(const string& anEntry) const = 0;
    virtual bool contains(const string& anEntry) const = 0;
    virtual void print() const = 0;
};

class ArrayBag : public BagInterface{
private:
    static const int DEFAULT_CAPACITY = 6;
    string items[DEFAULT_CAPACITY];
    int itemCount;
    int maxItems;
```

```
int getIndexOf(const string& target) const;
public:
    ArrayBag();
    int getCurrentSize() const;
    bool isEmpty() const;
    bool add(const string& newEntry);
    bool remove(const string& anEntry);
    void clear();
    int getFrequencyOf(const string& anEntry) const;
    bool contains(const string& anEntry) const;
    void print() const;
};

ArrayBag::ArrayBag()
: itemCount(0),
  maxItems(ArrayBag::DEFAULT_CAPACITY){}

int ArrayBag::getCurrentSize() const {
    return itemCount;
}

bool ArrayBag::isEmpty() const {
    return itemCount == 0;
}

void ArrayBag::print() const {
    for (int i = 0; i < itemCount; i++){
        cout << items[i] << " ";
    }
    cout << endl;
}

bool ArrayBag::add(const string& newEntry){
    bool hasRoomToAdd = (itemCount < maxItems);
    if (hasRoomToAdd){
        items[itemCount] = newEntry;
        itemCount++;
    }
    return hasRoomToAdd;
}

int ArrayBag::getIndexOf(const string& target) const {
    bool found = false;
    int result = -1;
    int searchIndex = 0;
    while(!found && (searchIndex < itemCount)){
        if (items[searchIndex] == target){
            found = true;
            result = searchIndex;
        }
        else {
            searchIndex++;
        }
    }
    return result;
}

bool ArrayBag::remove(const string& anEntry){
    int locateIndex = getIndexOf(anEntry);
    bool canRemoveItem = (locateIndex > -1);
    if (canRemoveItem){
```

```

        itemCount--;
        items[locateIndex] = items[itemCount];
    }
    return canRemoveItem;
}

bool ArrayBag::contains(const string& anEntry) const {
    return getIndexOf(anEntry) > -1;
}

//void ArrayBag::clear(){
//    while(!isEmpty()){
//        remove(items[0]);
//    }
//}

void ArrayBag::clear(){
    itemCount = 0;
}

int ArrayBag::getFrequencyOf(const string& anEntry) const{
    int frequency = 0;
    int curIndex = 0;
    while(curIndex < itemCount){
        if(items[curIndex] == anEntry) {
            frequency++;
        }
        curIndex++;
    }
    return frequency;
}

int main(){
    ArrayBag bag;
    bag.add("aaa");
    bag.print();
    cout << bag.isEmpty() << endl;
    cout << bag.getCurrentSize() << endl;
    bag.add("bbb");
    bag.print();
    cout << bag.isEmpty() << endl;
    cout << bag.getCurrentSize() << endl;

    return 0;
}
}

```

Implement by Dynamic Memory Allocate

```

In [ ]: #include <iostream>
#include <string>
using namespace std;

class BagInterface {
public:
    virtual int getCurrentSize() const = 0;
    virtual bool isEmpty() const = 0;
    virtual bool add(const string& newEntry) = 0;
    virtual bool remove(const string& anEntry) = 0;
    virtual void clear() = 0;
}

```

```

        virtual int getFrequencyOf(const string& anEntry) const = 0;
        virtual bool contains(const string& anEntry) const = 0;
        virtual void print() const = 0;
    };

    class ArrayBag : public BagInterface{
private:
    static const int DEFAULT_CAPACITY = 6;
    string* items;
    int itemCount;
    int maxItems;
    int getIndexOf(const string& target) const;
public:
    ArrayBag();
    ~ArrayBag();
    int getCurrentSize() const;
    bool isEmpty() const;
    bool add(const string& newEntry);
    bool remove(const string& anEntry);
    void clear();
    int getFrequencyOf(const string& anEntry) const;
    bool contains(const string& anEntry) const;
    void print() const;
};

ArrayBag::ArrayBag():
    itemCount(0),
    maxItems(ArrayBag::DEFAULT_CAPACITY){
    items = new string[DEFAULT_CAPACITY];
}

ArrayBag::~ArrayBag(){
    delete [] items;
}

int ArrayBag::getCurrentSize() const {
    return itemCount;
}

bool ArrayBag::isEmpty() const {
    return itemCount == 0;
}

void ArrayBag::print() const {
    for (int i = 0; i < itemCount; i++){
        cout << items[i] << " ";
    }
    cout << endl;
}

bool ArrayBag::add(const string& newEntry){
    bool hasRoomToAdd = (itemCount < maxItems);
    if (hasRoomToAdd){
        string* oldArray = items;
        items = new string[2 * maxItems];
        for (int index = 0; index < maxItems; index++){
            items[index] = oldArray[index];
        }
        delete [] oldArray;
        maxItems = 2 * maxItems;
    }
}

```

```
        }
        items[itemCount] = newEntry;
        itemCount++;
        return true;
    }

    int ArrayBag::getIndexOf(const string& target) const {
        bool found = false;
        int result = -1;
        int searchIndex = 0;
        while(!found && (searchIndex < itemCount)){
            if (items[searchIndex] == target){
                found = true;
                result = searchIndex;
            }
            else {
                searchIndex++;
            }
        }
        return result;
    }

    bool ArrayBag::remove(const string& anEntry){
        int locateIndex = getIndexOf(anEntry);
        bool canRemoveItem = (locateIndex > -1);
        if (canRemoveItem){
            itemCount--;
            items[locateIndex] = items[itemCount];
        }
        return canRemoveItem;
    }

    bool ArrayBag::contains(const string& anEntry) const {
        return getIndexOf(anEntry) > -1;
    }

    //void ArrayBag::clear(){
    //    while(!isEmpty()){
    //        remove(items[0]);
    //    }
    //}

    void ArrayBag::clear(){
        itemCount = 0;
    }

    int ArrayBag::getFrequencyOf(const string& anEntry) const{
        int frequency = 0;
        int curIndex = 0;
        while(curIndex < itemCount){
            if(items[curIndex] == anEntry) {
                frequency++;
            }
            curIndex++;
        }
        return frequency;
    }

    int main(){
        ArrayBag bag;
```

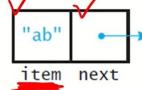
```

bag.add("aaa");
bag.print();
cout << bag.isEmpty() << endl;
cout << bag.getCurrentSize() << endl;
bag.add("bbb");
bag.print();
cout << bag.isEmpty() << endl;
cout << bag.getCurrentSize() << endl;

return 0;
}

```

Implement by Linked List

Abstract data types	Array-based implementation	Link-based implementation
Nodes	Member functions	Constructors and the destructor
Nodes  <ul style="list-style-type: none"> The type of an item is ItemType. <pre> template<typename ItemType> class BagInterface { public: virtual bool add(const ItemType& newEntry) = 0; // all other functions }; </pre>	 <p>Figure 4-1 (Carrano and Henry, 2013)</p> <ul style="list-style-type: none"> For each item, we will associate a pointer pointing to the next item-pointer pair. The combination of the item and pointer is often called a node. <ul style="list-style-type: none"> The pointer points to the next node, not the next item! 	
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Abstract data types	Array-based implementation	Link-based implementation
Nodes	Member functions	Constructors and the destructor
The class Node  <pre> template<class ItemType> class Node { private: ItemType item; Node<ItemType>* next; public: Node(); Node(const ItemType& anItem); Node(const ItemType& anItem, Node<ItemType>* nextNodePtr); void setItem(const ItemType& anItem); void setNext(Node<ItemType>* nextNodePtr); ItemType getItem() const; Node<ItemType>* getNext() const; }; </pre>	 <p>Figure 4-1 (Carrano and Henry, 2013)</p>	
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Abstract data types	Array-based implementation	Link-based implementation
Nodes	Member functions	Constructors and the destructor

The class Node

- The class **Node** has **getters** and **setters** for all instance variables.
 - It does no data hiding at all.
- Sometimes people implement **Node** as a structure.
 - Member variables of a structure are by default public.
 - Though they can be set to be private.
- Sometimes a better way is to set **LinkedBag** a **friend** of **Node**. ✓

Basic concepts	Constructors and the destructor
Friends and static members	Object pointers and the copy constructor

friend for functions and classes

- Declare friends only if data hiding is preserved.
 - Do not set everything public!
 - Use structures rather than classes when nothing should be private (this is recommended but not required).
 - Be careful in offering public member functions (e.g., getters and setters).
- friend** in fact **help you hide data**.
 - If a **private** member should be accessed only by another class/function, we should declare a friend instead of writing a getter/setter.

Abstract data types	Array-based implementation	Link-based implementation
Nodes	Member functions	Constructors and the destructor

The head pointer headPtr

- An object of **LinkedBag** will have an **item counter** itemCount.
 - Its type is **int**.
- An object of **LinkedBag** will have a head pointer headPtr.
 - Its type is **Node< itemType >***. It is a pointer, not a node.
 - headPtr** is a **“static member”**; all other nodes are **“dynamic members.”**

Figure 4-5 (Carrano and Henry, 2013)

- This “static” is different from that “static” for “static member variables.”

Abstract data types	Array-based implementation	Link-based implementation
Nodes	Member functions	Constructors and the destructor
<h2>Copy constructor (shallow copy)</h2>		
<ul style="list-style-type: none"> If we do not define a copy constructor by ourselves, the default copy constructor of LinkedBag will be like: 		
<pre>template<class ItemType> LinkedBag<ItemType>::LinkedBag (const LinkedBag<ItemType>& aBag) { itemCount = aBag->itemCount; headPtr = aBag->headPtr; }</pre>		
Figure 4-8 (a) (Carrano and Henry, 2013)		

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Abstract data types	Array-based implementation	Link-based implementation
Nodes	Member functions	Constructors and the destructor
<h2>Copy constructor (deep copy)</h2>		
<ul style="list-style-type: none"> What we want is deep copy: A new list of items should be dynamically created. 		
Figure 4-8 (b) (Carrano and Henry, 2013)		
<ul style="list-style-type: none"> To do so, we need to define our own copy constructor for LinkedBag. 		

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```
In [ ]: #include <iostream>
#include <string>
#include <vector>
using namespace std;

//Class Node
template<typename ItemType>
class Node{
private:
    ItemType item;
    Node<ItemType*>* next;
public:
    Node();
    Node(const ItemType& anItem);
    Node(const ItemType& anItem, Node<ItemType*>* nextNodePtr);
    void setItem(const ItemType& anItem);
    void setNext(Node<ItemType*>* nextNodePtr);
    ItemType getItem() const;
    Node<ItemType*>* getNext() const;
};

//Constructors of Node
template<typename ItemType>
Node<ItemType>::Node()
    : next(nullptr) {}
```

```

template<typename ItemType>
Node<ItemType>::Node(const ItemType& anItem)
    : item(anItem), next(nullptr) {}

template<typename ItemType>
Node<ItemType>::Node(const ItemType& anItem, Node<ItemType>* nextNodePtr)
    : item(anItem), next(nextNodePtr) {}

// Getters and Setters
template<typename ItemType>
void Node<ItemType>::setItem(const ItemType& anItem){
    item = anItem;
}

template<typename ItemType>
void Node<ItemType>::setNext(Node<ItemType>* nextNodePtr){
    next = nextNodePtr;
}

template<typename ItemType>
ItemType Node<ItemType>::getItem() const {
    return item;
}

template<typename ItemType>
Node<ItemType>* Node<ItemType>::getNext() const {
    return next;
}

//Class BagInterface
template<typename ItemType>
class BagInterface {
public:
    virtual int getCurrentSize() const = 0;
    virtual bool isEmpty() const = 0;
    virtual bool add(const ItemType& newEntry) = 0;
    virtual bool remove(const ItemType& anEntry) = 0;
    virtual void clear() = 0;
    virtual int getFrequencyOf(const ItemType& anEntry) const = 0;
    virtual bool contains(const ItemType& anEntry) const = 0;
    virtual void print() const = 0;
};

//class LinkedBag
template<typename ItemType>
class LinkedBag : public BagInterface<ItemType> {
private:
    Node<ItemType>* headPtr;
    int itemCount;
public:
    LinkedBag(); // constructor
    LinkedBag(const LinkedBag<ItemType>& aBag); // copy constructor
    virtual ~LinkedBag(); // destructor
    int getCurrentSize() const;
    bool isEmpty() const;
    bool add(const ItemType& newEntry);
    vector<ItemType> toVector() const;
    bool remove(const ItemType& anEntry);
}

```

```

        void clear();
        Node<ItemType>* getPointerTo(const ItemType& anEntry) const;
        int getFrequencyOf(const ItemType& anEntry) const;
        bool contains(const ItemType& anEntry) const;
        void print() const;
    };

template<typename ItemType>
LinkedBag<ItemType>:: LinkedBag(){
    headPtr = nullptr;
    itemCount = 0;
}

template<typename ItemType>
LinkedBag<ItemType>::LinkedBag(const LinkedBag<ItemType>& aBag){
    itemCount = aBag->itemCount;
    Node<ItemType>* origChainPtr = aBag->headPtr;
    if(origChainPtr == nullptr){
        headPtr = nullptr;
    }
    else {
        headPtr = new Node<ItemType>();
        headPtr->setItem(origChainPtr->getItem());
        Node<ItemType>* newChainPtr = headPtr;
        while(origChainPtr != nullptr){
            origChainPtr = origChainPtr->getNext();
            ItemType nextItem = origChainPtr->getItem();
            Node<ItemType>* newNodePtr = new Node<ItemType>(nextItem);
            newChainPtr->setNext(newNodePtr);
            newChainPtr = newChainPtr->getNext();
        }
        newChainPtr->setNext(nullptr);
    }
}

template<typename ItemType>
LinkedBag<ItemType>::~LinkedBag(){
    clear();
}

template<typename ItemType>
bool LinkedBag<ItemType>::isEmpty() const{
    return itemCount == 0;
}

template<typename ItemType>
int LinkedBag<ItemType>::getCurrentSize() const {
    return itemCount;
}

//add to the beginning
template<typename ItemType>
bool LinkedBag<ItemType>::add(const ItemType& newEntry){
    Node<ItemType>* newNodePtr = new Node<ItemType>();
    newNodePtr->setItem(newEntry);
    newNodePtr->setNext(headPtr);
    headPtr = newNodePtr;
    itemCount++;
    return true;
}

```

```

template<typename ItemType>
vector<ItemType> LinkedBag<ItemType>::toVector() const{
    vector<ItemType> bagContents;
    Node<ItemType>* curPtr = headPtr;
    while(curPtr != nullptr ){
        bagContents.push_back(curPtr->getItem());
        curPtr = curPtr->getNext();
    }
    return bagContents;
}

template<typename ItemType>
bool LinkedBag<ItemType>::remove(const ItemType& anEntry){
    Node<ItemType>* entryNodePtr = getPointerTo(anEntry);

    bool canRemoveItem = !isEmpty() && (entryNodePtr != nullptr);

    if (canRemoveItem) {
        entryNodePtr->setItem(headPtr->getItem());//Replace
        Node<ItemType>* nodeToDeletePtr = headPtr;
        headPtr = headPtr->getNext();
        delete nodeToDeletePtr;
        nodeToDeletePtr = nullptr;
        itemCount--;
    }
    return canRemoveItem;
}

template<typename ItemType>
Node<ItemType>* LinkedBag<ItemType>::getPointerTo(const ItemType& anEntry) const {
    bool found = false;
    Node<ItemType>* curPtr = headPtr;
    while(!found && (curPtr != nullptr)){
        if(anEntry == curPtr->getItem()){
            found = true;
        }
        else
            curPtr = curPtr->getNext();
    }
    return curPtr;
}

template<typename ItemType>
bool LinkedBag<ItemType>::contains(const ItemType& anEntry) const {
    return (getPointerTo(anEntry) != nullptr);
}

template<typename ItemType>
int LinkedBag<ItemType>::getFrequencyOf(const ItemType& anEntry) const {
    int frequency = 0;
    Node<ItemType>* curPtr = headPtr;
    while(curPtr != nullptr){
        if(anEntry == curPtr->getItem()){
            frequency++;
        }
        curPtr = curPtr->getNext();
    }
    return frequency;
}

```

```

template<typename ItemType>
void LinkedBag<ItemType>::clear(){
    Node<ItemType>* nodeToDeletePtr = headPtr;
    while(headPtr != nullptr){
        headPtr = headPtr->getNext();
        delete nodeToDeletePtr;
        nodeToDeletePtr = headPtr;
    }
    itemCount = 0;
}

void testBag(BagInterface<string>* bagPtr){
    string item[] = {"aa", "bb", "cc"};
    for(int i = 0; i < 3; i++){
        bagPtr->add(item[i]);
    }
    cout << bagPtr->isEmpty();
    cout << bagPtr->getCurrentSize();

    bagPtr->remove("two");
    cout << bagPtr->isEmpty();
    cout << bagPtr->getCurrentSize();
}

int main(){
    pass

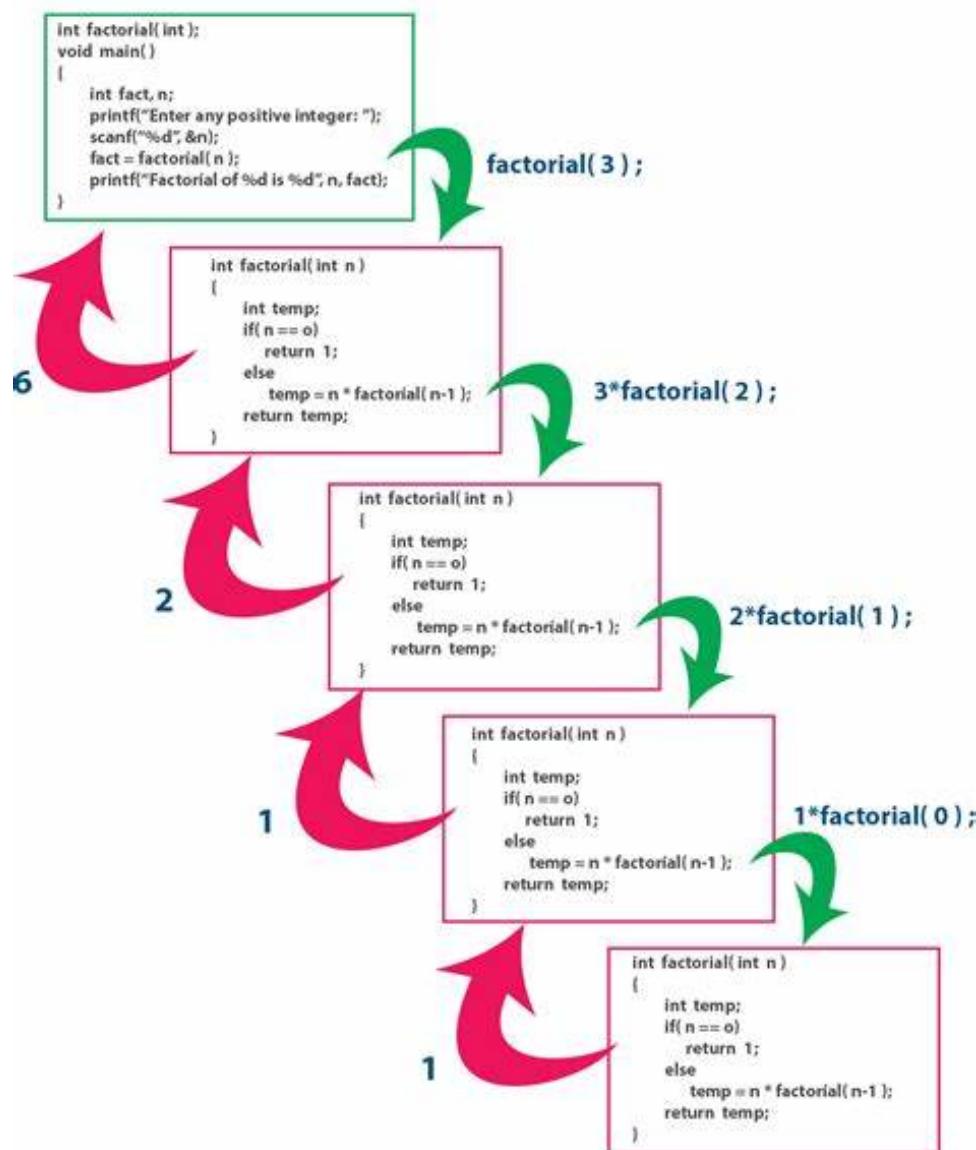
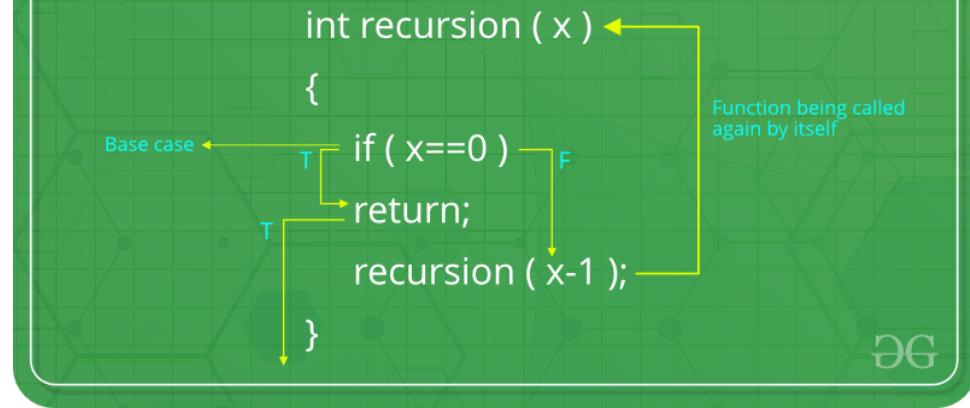
    return 0;
}

```

Recursive Function

	Algorithms and complexity	Recursion	
	Recursive functions		
	遞迴		
	<ul style="list-style-type: none"> • A function is <u>recursive</u> if it invokes itself (directly or indirectly). • The process of using recursive functions is called <u>recursion</u>. • Why recursion? <ul style="list-style-type: none"> – Many problems can be solved by <u>dividing the original problem into one or several smaller pieces of subproblems</u>. – Typically <u>subproblems are quite similar</u> to the original problem. – With recursion, we write one function to solve the problem by <u>using the same function</u> to solve subproblems. 		
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Recursive Functions



```

In [ ]: #include <iostream>
using namespace std;

int fact(int n);

```

```

int main(){

    cout << "fact(5): " << fact(5) << endl;

    return 0;
}

int fact(int n){
    if(n==0)
        return 1;
    else
        return n * fact(n-1);
}

```

Algorithms and complexity

Recursion

Searching and sorting

Example 1: finding the maximum

- Let's try to implement the strategy.
 - First, I know I need to write a function whose header is:
- `double max(double array[], int len);`
- This function returns the maximum in `array` (containing `len` elements).
 - I want this to happen, though at this moment I do not know how.
- Now let's implement it:
 - If the function really works, subtask 1 can be completed by invoking
- `double subMax = max(array, len - 1);`
- Subtask 2 is done by comparing `subMax` and `array[len - 1]`.

Programming Design – Algorithms and Recursion

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Algorithms and complexity

Recursion

Searching and sorting

Some remarks

- There must be a stopping condition in a recursive function. Otherwise, the program will not terminate.
- In many cases, a recursive strategy can also be implemented with loops.
 - E.g., writing a loop for finding a maximum and factorial.
 - But sometimes it is hard to use loops to imitate a recursive function.
- Compared with an equivalent iterative function, a recursive implementation is usually simpler and easier to understand.
- However, it generally uses more memory spaces and is more time-consuming.
 - Invoking functions has some cost.

Programming Design – Algorithms and Recursion

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```

In [ ]: #include <iostream>
using namespace std;

double findMax(double array[], int len);

int main(){
    double a[5] = {5, 7, 2, 4, 3};
    cout << "The max item is " << findMax(a, 5) << endl;

    return 0;
}

```

```

double findMax(double array[], int len){
    if (len == 1){
        return array[0];
    }
    else {
        double subMax = findMax(array, len-1);
        if (array[len-1] > subMax){
            return array[len - 1];
        }
        else {
            return subMax;
        }
    }
}

```

```

In [ ]: #include <iostream>
using namespace std;

int choose(int n, int k);

int main(){
    cout << "Choose(3 from 5 items): " << choose(5, 3) << endl;
    return 0;
}

int choose(int n, int k){
    if(k == 0 || k == n)
        return 1;
    else if(k > n)
        return 0;
    else
        return choose(n-1, k-1) + choose(n-1, k);
}

```

Recursion	Algorithm complexity
Basic concepts	Number of recursive calls

The Hanoi towers

```

solveHT(count, source, destination, spare)
if count is 1
    Move a disk directly from source to destination
else
    solveHT(count - 1, source, spare, destination)
    solveHT(1, source, destination, spare)
    solveHT(count - 1, spare, destination, source)

```

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Recursion	Algorithm complexity
Basic concepts	Number of recursive calls
<h2>The Hanoi towers</h2>	
$N = 1$ (1) : move the disk once from source to destination directly.	1 move
$N = 2$ (1) (1) : move the disk from source to auxiliary, then from auxiliary to destination, then from source to destination.	3 moves
$N = 3$ (1) (1) (1) : move the bottom two disks from source to auxiliary, then move the bottom disk from source to destination, move the auxiliary disk to destination, then move the bottom two disks from auxiliary to destination.	7 moves
$\text{moves}(N)$ equals to the number of nodes in a binary tree: $2^N - 1$.	

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Recursion	Algorithm complexity
Basic concepts	Number of recursive calls
<h2>Recursion and efficiency</h2>	
<ul style="list-style-type: none"> Recursion is a powerful problem-solving technique that often produces very clear solutions to complex problems. However, some recursive solutions are inefficient that they should not be used. Factors that contribute to the inefficiency of some recursive solutions: <ul style="list-style-type: none"> Overhead associated with function calls. Inherent inefficiency of some recursive algorithms. 	

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Find Prime Number

Algorithms and complexity	Recursion	Searching and sorting
<h2>Example: listing all prime numbers</h2>		
<ul style="list-style-type: none"> Given an integer n, let's list all the prime numbers no greater than n. Consider the following (imprecise) algorithm: <ul style="list-style-type: none"> For each number i no greater than n, check whether it is a prime number. To check whether i is a prime number: <ul style="list-style-type: none"> Idea: If any number $j < i$ can divide i, i is not a prime number. Algorithm: For each number $j < i$, check whether j divides i. If there is any j that divides i, report no; otherwise, report yes. Before we write a program, we typically prefer to formalize our algorithm. <ul style="list-style-type: none"> We write pseudocodes, a description of steps in words organized in a program structure. This allows us to ignore the details of implementations. 		

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Algorithms and complexity Recursion Searching and sorting

Example: listing all prime numbers

- One pseudocode for listing all prime numbers no greater than n is:

Given an integer n :
for i from 2 to n
 assume that i is a prime number
 for j from 2 to $i - 1$
 if j divides i
 set i to be a composite number
 if i is still considered as prime
 print i

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```
In [ ]: #include <iostream>
using namespace std;

int main(){
    int num = 0;
    cout << "Learning Algorithm and Data Structure" << endl;
    cout << "Enter your Integer Number: ";
    cin >> num;

    for (int i = 2; i <= num; i++){
        bool isPrime = true;
        for (int j = 2; j < i; j++){
            if ((i % j) == 0){
                isPrime = false;
                break;
            }
        }
        if (isPrime == true) {
            cout << i << " is prime" << endl;
        }
    }
    return 0;
}
```

```
In [ ]: #include <iostream>
using namespace std;

bool isPrime(int x);

int main(){
    int num = 0;
    cout << "Learning Algorithm and Data Structure" << endl;
    cout << "Enter your Integer Number: ";
    cin >> num;

    for (int i = 2; i <= num; i++){
        if (isPrime(i) == true) {
            cout << i << " is prime" << endl;
        }
    }
    return 0;
}
```

```

bool isPrime(int x) {
    for (int i = 2; i < x; i++){
        if ((x % i) == 0){
            return false;
        }
    }
    return true;
}

```

```

In [ ]: #include <iostream>
using namespace std;

bool isPrime(int x);

int main(){
    int num = 0;
    cout << "Learning Algorithm and Data Structure" << endl;
    cout << "Enter your Integer Number: ";
    cin >> num;

    for (int i = 2; i <= num; i++){
        if (isPrime(i) == true) {
            cout << i << " is prime" << endl;
        }
    }
    return 0;
}

bool isPrime(int x) {
    for (int i = 2; i * i < x; i++){
        if ((x % i) == 0){
            return false;
        }
    }
    return true;
}

```

```

In [ ]: #include <iostream>
using namespace std;

const int MAX_LEN = 10000;

void ruleOutPrime(int x, bool isPrime[], int n);

int main(){
    int num = 0;
    cout << "Learning Algorithm and Data Structure" << endl;
    cout << "Enter your Integer Number(2~10000): ";
    cin >> num;

    bool isPrime[MAX_LEN] = {0};

    for (int i = 2; i < num; i++){
        isPrime[i] = true;
    }

    for (int i = 2; i <= num; i++){
        if(isPrime[i] == true){
            cout << i << " is prime" << endl;
            ruleOutPrime(i, isPrime, num);
        }
    }
}

```

```

        }

    }

    return 0;
}

void ruleOutPrime(int x, bool isPrime[], int n){
    for (int i = 1; x * i < n; i++){
        isPrime[x * i] = false;
    }
}

```

Dynamic Programming

```
In [ ]: #include <iostream>
#include <stdexcept>
using namespace std;

int fib(int n);

int main(){
    cout << "fib(5): " << fib(5) << endl;
    return 0;
}

int fib(int n){

    if(n < 0)
        throw logic_error("Error!");
    if(n == 0)
        return 0;
    if (n == 1)
        return 1;
    else
        return fib(n-1) + fib(n-2);
}
```

```
In [ ]: #include <iostream>
#include <stdexcept>
using namespace std;

int fib(int n);

int main(){
    cout << "fib(6): " << fib(6) << endl;
    return 0;
}

int fib(int n){

    int A[20];
```

```

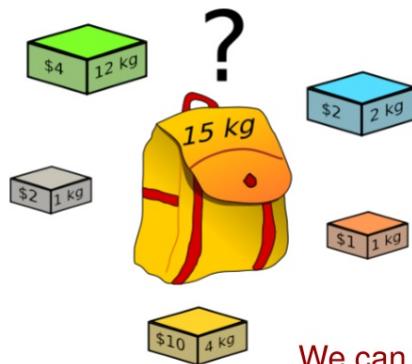
A[0] = 0;
A[1] = 1;
for(int i = 2; i <= n; i++){
    A[i] = A[i-1] + A[i-2];
}
return A[n];
}

```

The Knapsack Problem

Knapsack problem

- Each item has a value and a weight
- **Objective:** maximize value
- **Constraint:** knapsack has a weight limitation



Three versions:

0-1 knapsack problem: take each item or leave it

Fractional knapsack problem: items are divisible

Unbounded knapsack problem: unlimited supplies of each item.

Which one is easiest to solve?

We can solve the fractional knapsack problem using greedy algorithm

Recursion	The knapsack problem	Algorithm complexity
	<h2>The knapsack problem</h2> <ul style="list-style-type: none"> • Problem input: <ul style="list-style-type: none"> – The weight of items: w_1, w_2, \dots, w_n. – The value of items: v_1, v_2, \dots, v_n. – The weight limit of the knapsack B. • Problem formulation: <ul style="list-style-type: none"> – Let $x_i = 1$ if item i is selected and 0 otherwise. – The problem: <div style="border: 1px solid black; padding: 10px;"> $\begin{aligned} \max \quad & \sum_{i=1}^n v_i x_i \\ \text{s.t.} \quad & \sum_{i=1}^n w_i x_i \leq B \\ & x_i \in \{0, 1\} \quad \forall i = 1, \dots, n \end{aligned}$ </div>	Algorithm complexity The maximum flow problem

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Recursion Algorithm complexity

The knapsack problem The maximum flow problem

Box trace

- Example: $w = (8, 3, 4, 5)$ and $B = 16$.

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Recursion Algorithm complexity

The knapsack problem The maximum flow problem

The dynamic programming algorithm

- The idea is great, but a recursive implementation is inefficient (why?).
- By applying “dynamic programming,” we solve this problem with a matrix:

w_i/B	0	1	2	3	4	5	6
2	NS	IMP	S	IMP	IMP	IMP	IMP
3	NS	IMP	NS	S	IMP	S	IMP
4	NS	IMP	NS	NS	S	NS	S
5	NS	IMP	NS	NS	NS	S or NS	NS

- The last cell is what we want. The answer is “Yes, we may fill a knapsack of capacity 6 with the four items.”
- How to determine the items to be selected?

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Recursion Algorithm complexity

The knapsack problem The maximum flow problem

Efficiency

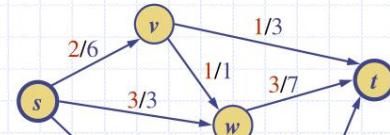
- The complexity of the dynamic programming algorithm: $O(nB)$. (circled with a red X)
- Is this algorithm efficient?
 - Typically yes, but no if the knapsack capacity is really large.
 - This is why it is “pseudo”-polynomial: Its complexity depends on the “value” of inputs, not just the “number” of inputs.

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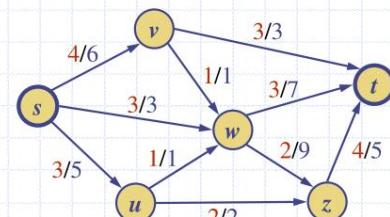
The Max Flow Problem

Maximum Flow

- ◆ A flow for a network N is said to be maximum if its value is the largest of all flows for N
- ◆ The *maximum flow problem* consists of finding a maximum flow for a given network N
- ◆ Applications
 - Hydraulic systems
 - Electrical circuits
 - Traffic movements
 - Freight transportation



Flow of value $8 = 2 + 3 + 3 = 1 + 3 + 4$



Maximum flow of value $10 = 4 + 3 + 3 = 3 + 3 + 4$

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Maximum Flow

5

Recursion

The knapsack problem

Algorithm complexity

The maximum flow problem

The maximum flow problem

- A very natural application is in logistic/shipping/delivery.
 - The source node is a factory. The destination node is a market.
 - An edge's capacity is the maximum number of goods that can be shipped along that edge.
 - It may be easily used to model a problem with multiple factories and multiple markets (how?).
- There are many other applications.
- It may also be defined on undirected networks.
 - Let's consider only directed networks.

Data Structures – Recursion and Algorithm Complexity

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Recursion

The knapsack problem

Algorithm complexity

The maximum flow problem

Efficiency

- The worst-case efficiency of the augmenting path algorithm is $O(mf^*)$.
 - m is the number of edges: each augmenting path's maximum length.
 - f^* is the maximum flow: the maximum number of iterations.
- Alternatively, if you prefer, $O(mU)$.
 - $U = \sum_{(i,j) \in E} u_{ij}$, the total capacity, is an upper bound of f^* .
- The augmenting path algorithm is thus **pseudo-polynomial**.
- There exists polynomial algorithm for the maximum flow problem.

Data Structures – Recursion and Algorithm Complexity

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	Max-flow (Primal)	Min-cut (Dual)
variables	$f_{uv} \forall (u, v) \in E$ [a variable per edge]	$d_{uv} \forall (u, v) \in E$ [a variable per edge] $z_v \forall v \in V \setminus \{s, t\}$ [a variable per non-terminal node]
objective	maximize $\sum_{v \in \{s, v\} \in E} f_{sv}$ [max total flow from source]	minimize $\sum_{(u, v) \in E} c_{uv} d_{uv}$ [min total capacity of edges in cut]
constraints	subject to $\begin{aligned} f_{uv} &\leq c_{uv} & \forall (u, v) \in E \\ \sum_u f_{uv} - \sum_w f_{wu} &= 0 & v \in V \setminus \{s, t\} \end{aligned}$ [a constraint per edge and a constraint per non-terminal node]	subject to $\begin{aligned} d_{uv} - z_u + z_v &\geq 0 & \forall (u, v) \in E, u \neq s, v \neq t \\ d_{sv} + z_v &\geq 1 & \forall (s, v) \in E \\ d_{ut} - z_u &\geq 0 & \forall (u, t) \in E \end{aligned}$ [a constraint per edge]
sign constraints	$f_{uv} \geq 0 \quad \forall (u, v) \in E$	$d_{uv} \geq 0 \quad \forall (u, v) \in E$ $z_v \in \mathbb{R} \quad \forall v \in V \setminus \{s, t\}$

C++ Performance Evaluation - Space Complexity and Time Complexity

What does a BETTER Algorithm mean ?

- Faster ? (Less execution time) – Time Complexity
- Less Memory ? – Space Complexity
- Easy to read ?
- Less Line of Code ?
- Less Hw/Sw needs ?

*Note: Algorithm Analysis **does not give you accurate/exact values**(time, space etc), however it gives **estimates** which can be used to study the **behavior** of the algorithm.

What is Asymptotic Algorithm Analysis ?

- Definition: In mathematical analysis, **asymptotic** analysis of algorithm is a method of defining the **mathematical boundaries** of its **run-time** performance.
- Using the asymptotic analysis, we can easily estimate about the average case, best case and worst case scenario of an algorithm.

Simple words: It is used to mathematically calculate the running time of any operation inside an algorithm.

Asymptotic Algorithm analysis is to estimate the time complexity function for arbitrarily large input.

Time Complexity : is a computational way to show how(**behavior**) runtime of a program increases as the size of its input increases.

What is Space complexity?

- **Definition:** The space complexity of an algorithm or a computer program is the **amount of memory space** required to solve an instance of the computational problem as a function of the size of the input.
- **Simple words :** It is the memory required by an algorithm to execute a program and produce output.
- Similar to time complexity, Space complexity is often expressed asymptotically in big O notation, such as $O(n)$, $O(n\log(n))$, $O(n^2)$ etc., where **n** is the input size in units of bits needed to represent the input.

- ❖ We want to define time taken by an algorithm without depending on the implementation details.
- ❖ because
- ❖ A given algorithm will take different amounts of time on the same inputs depending on such factors as:
 - Processor speed;
 - Instruction set,
 - Disk speed,
 - Brand of compiler and etc.
- ❖ The way around is to estimate efficiency of each algorithm asymptotically.
 - **Time Complexity:** Running time of the program as a function of the size of input
 - **Space Complexity:** Amount of computer memory required during the program execution, as a function of the input size

Mr. Nitin M Shivale (Asst. Prof JSPM'S BSIOTR)

Space Complexity

- ❑ When memory was expensive we focused on making programs as **space** efficient as possible and developed schemes to make memory appear larger than it really was (virtual memory and memory paging schemes)
- ❑ Space complexity is still important in the field of embedded computing (hand held computer based equipment like cell phones, palm devices, etc)

Algorithms and complexity Recursion Searching and sorting

Complexity

- Running time may be affected by the hardware, number of programs running at the same time, etc.
 - The number of basic operations is a better measurement.
 - Basic operations include simple arithmetic, comparisons, etc.
- Convince yourself that algorithm 2 does fewer basic operations.
- The calculation of complexity needs training.
 - This will be formally introduced in Discrete Mathematics, Data Structures, and/or Algorithms.

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Complexity Terminology of graphs The “big O” notation Graph algorithms

Time complexity: example

- Consider the previous example.
- Let's count the number of basic operations algorithm 1.
- For the first part of algorithm 1, we have $5mn + 10m + 2$ basic operations.

	Decl.	Assi.	Arith.	Comp.
(1)	m	m	0	0
(2)	1	$m + 1$	m	m
(3)	m	m	0	0
(4)	m	$m(n + 1)$	mn	mn
(5)	0	mn	mn	0
(6)	0	m	0	0

```

int rowSum[MAX_ROW_CNT] = {0}; // (1)
for(int i = 0; i < m; i++) // (2)
{
    int aRowSum = 0; // (3)
    for(int j = 0; j < n; j++) // (4) ✓
        aRowSum += A[i][j]; // (5)
    rowSum[i] = aRowSum; // (6)
}

// the remaining are skipped
  
```

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Complexity Terminology of graphs The “big O” notation Graph algorithms

Time complexity: example

- Let's analyze algorithm 2.
- The bottleneck is the two nested loops.
- The complexity is roughly mn :
 - This is how the execution time would grow as the input size increases.
- To formalize the above idea, let's introduce the “big O” notation.

```

int maxRowSum(int A[] [MAX_COL_CNT],
              int m, int n)
{
    int maxRowSumValue = 0;
    int maxRowNumber = 0;
    for(int i = 0; i < m; i++)
    {
        int aRowSum = 0;
        for(int j = 0; j < n; j++)
            aRowSum += A[i][j]; MN

        if(aRowSum > maxRowSumValue)
        {
            maxRowSumValue = aRowSum;
            maxRowNumber = i + 1;
        }
    }
    return maxRowNumber;
}
  
```

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Big "O" Notation

Complexity
Terminology of graphs

The “big O” notation

- Mathematically, let $f(n) \geq 0$ and $g(n) \geq 0$ be two functions defined for $n \in \mathbb{N}$. We say

$$f(n) \in O(g(n))$$

if and only if there exists a positive number c and a number N such that

$$f(n) \leq cg(n)$$

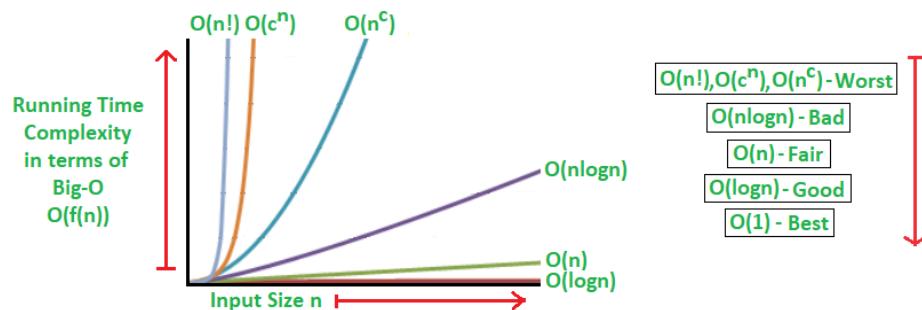
for all $n \geq N$.

- Intuitively, that means **when n is large enough, $g(n)$ will dominate $f(n)$** .
- If $f(n)$ is the number of operations that an algorithm takes to complete a task, we say **the algorithm's time complexity** is $g(n)$.
 - We write $f(n) \in O(g(n))$, but some people write $f(n) = O(g(n))$.

Complexity
Terminology of graphs

Worst-case time complexity

- In many cases, the number of operations of running an algorithm depends on not only the **number of input values** but also **contents of input values**.
- People talk about two kinds of time complexity:
 - Average-case time complexity**: the **expected** number of operations required for a randomly drawn input. The probability distribution matters.
 - Worst-case time complexity**: the **maximum possible** number of operations required for a randomly drawn input.
- The “big O” notation typically deals with worst-case complexity.



Container	Insertion	Access	Erase	Find	Persistent Iterators
vector / string	Back: $O(1)$ or $O(n)$ Other: $O(n)$	$O(1)$	Back: $O(1)$ Other: $O(n)$	Sorted: $O(\log n)$ Other: $O(n)$	No
deque	Back/Front: $O(1)$ Other: $O(n)$	$O(1)$	Back/Front: $O(1)$ Other: $O(n)$	Sorted: $O(\log n)$ Other: $O(n)$	Pointers only
list / forward_list	Back/Front: $O(1)$ With iterator: $O(1)$ Index: $O(n)$	Back/Front: $O(1)$ With iterator: $O(1)$ Index: $O(n)$	Back/Front: $O(1)$ With iterator: $O(1)$ Index: $O(n)$	$O(n)$	Yes
set / map	$O(\log n)$	-	$O(\log n)$	$O(\log n)$	Yes
unordered_set / unordered_map	$O(1)$ or $O(n)$	$O(1)$ or $O(n)$	$O(1)$ or $O(n)$	$O(1)$ or $O(n)$	Pointers only
priority_queue	$O(\log n)$	$O(1)$	$O(\log n)$	-	-

Data Structures

Data Structure	Time Complexity								Space Complexity	
	Average				Worst					
	Indexing	Search	Insertion	Deletion	Indexing	Search	Insertion	Deletion		
Basic Array	O(1)	O(n)	-	-	O(1)	O(n)	-	-	O(n)	
Dynamic Array	O(1)	O(n)	O(n)	O(n)	O(1)	O(n)	O(n)	O(n)	O(n)	
Singly-Linked List	O(n)	O(n)	O(1)	O(1)	O(n)	O(n)	O(1)	O(1)	O(n)	
Doubly-Linked List	O(n)	O(n)	O(1)	O(1)	O(n)	O(n)	O(1)	O(1)	O(n)	
Skip List	O(log(n))	O(log(n))	O(log(n))	O(log(n))	O(n)	O(n)	O(n)	O(n)	O(n log(n))	
Hash Table	-	O(1)	O(1)	O(1)	-	O(n)	O(n)	O(n)	O(n)	
Binary Search Tree	O(log(n))	O(log(n))	O(log(n))	O(log(n))	O(n)	O(n)	O(n)	O(n)	O(n)	
Cartesian Tree	-	O(log(n))	O(log(n))	O(log(n))	-	O(n)	O(n)	O(n)	O(n)	
B-Tree	O(log(n))	O(log(n))	O(log(n))	O(log(n))	O(log(n))	O(log(n))	O(log(n))	O(n)	O(n)	
Red-Black Tree	O(log(n))	O(log(n))	O(log(n))	O(log(n))	O(log(n))	O(log(n))	O(log(n))	O(n)	O(n)	
Splay Tree	-	O(log(n))	O(log(n))	O(log(n))	-	O(log(n))	O(log(n))	O(log(n))	O(n)	
AVL Tree	O(log(n))	O(log(n))	O(log(n))	O(log(n))	O(log(n))	O(log(n))	O(log(n))	O(n)	O(n)	

C++ Type of Data Structure

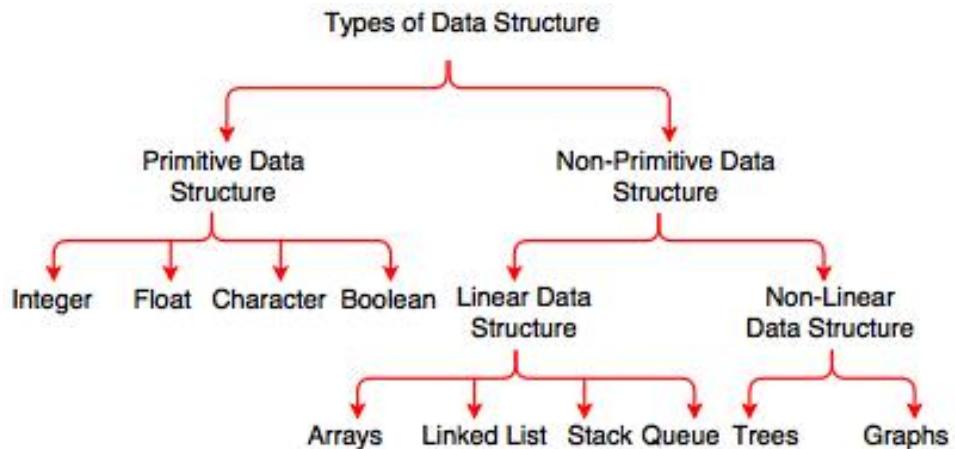
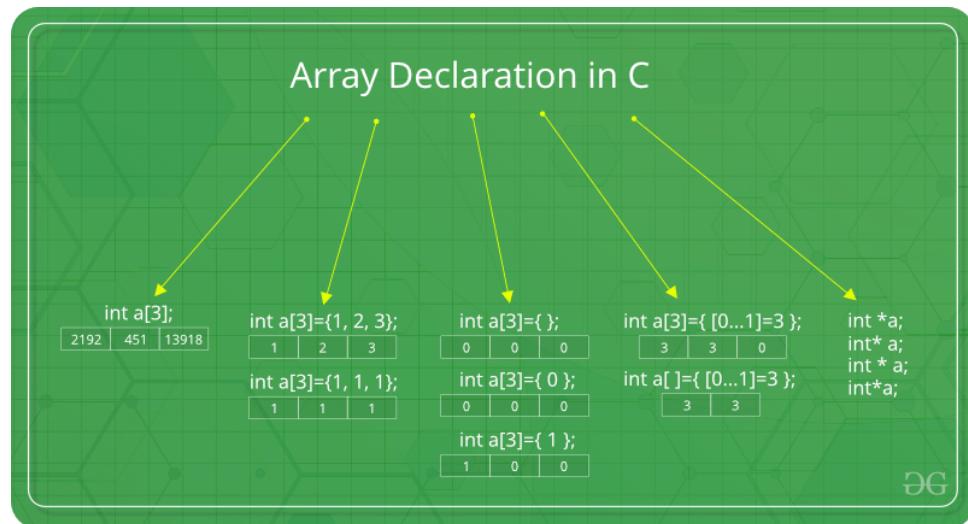
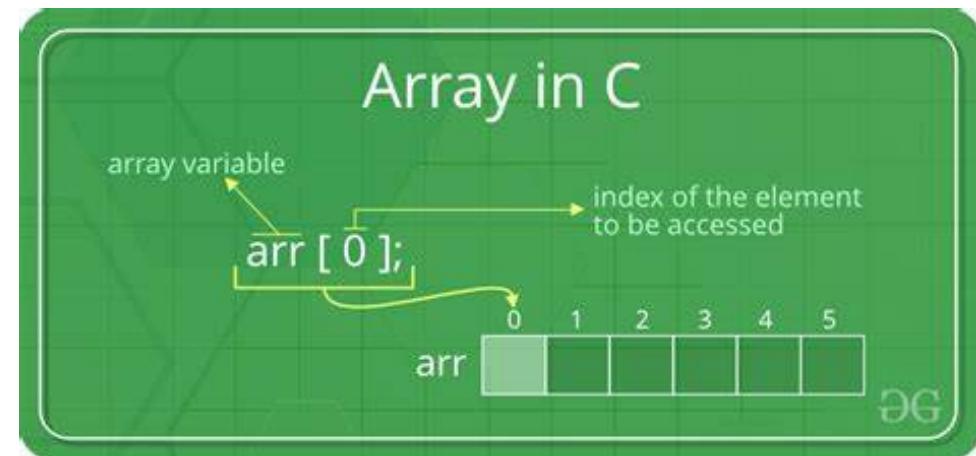


Fig. Types of Data Structure

Linear Data Structure

Array



Operations in Array

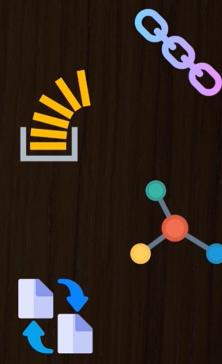
language built-in array	container library array
<pre>#include <iostream> using namespace std; int main() { int myarray[3] = {10,20,30}; for (int i=0; i<3; ++i) ++myarray[i]; for (int elem : myarray) cout << elem << '\n'; }</pre>	<pre>#include <iostream> #include <array> using namespace std; int main() { array<int,3> myarray {10,20,30}; for (int i=0; i<myarray.size(); ++i) ++myarray[i]; for (int elem : myarray) cout << elem << '\n'; }</pre>

Linked List

Some Applications of Linked List Data Structure

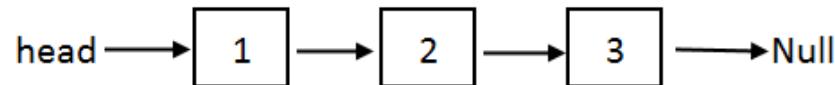
Some Applications of Linked List are as follows –

- Linked Lists can be used to implement **Stacks** , **Queues**.
- Linked Lists can also be used to implement **Graphs**. (Adjacency list representation of Graph).
- Implementing Hash Tables :- Each Bucket of the hash table can itself be a linked list. (Open chain hashing).
- **Undo** functionality in Photoshop or Word . Linked list of states

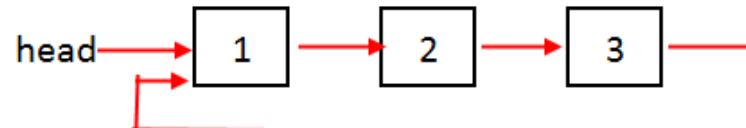


[/simplesnippets](#) [/simplesnippets](#) [/simplesnippets](#) [/simplesnippet](#) [www.simplesnippets.tech](#)

CLL

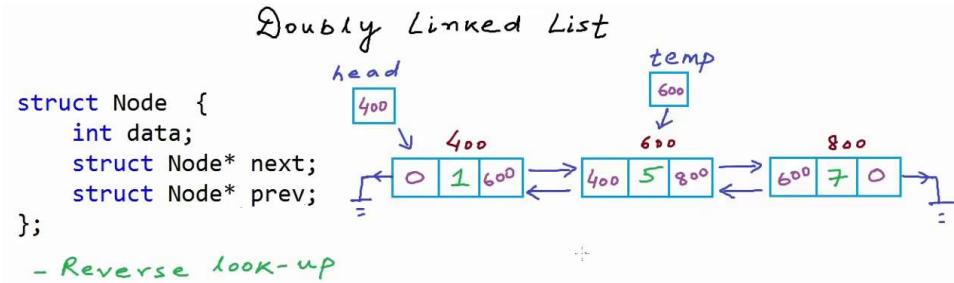


Singly Linked List

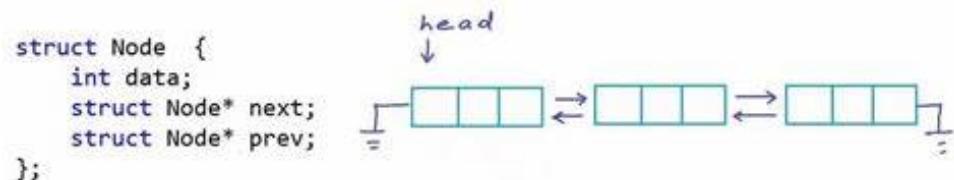


Circular Linked List

DLL



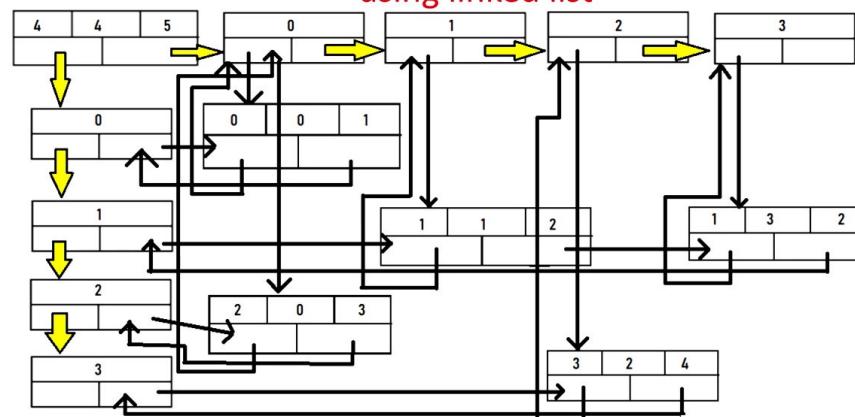
mycodeschool.com

Doubly Linked List - Implementation

mycodeschool.com

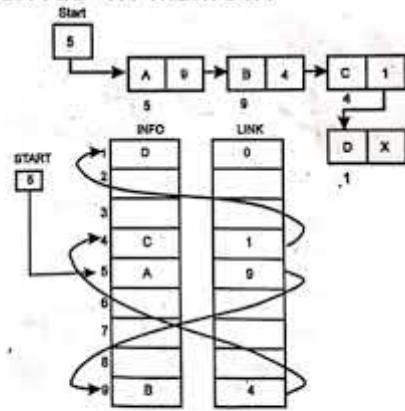
Matrix

This is the final representation of sparse matrix using linked list

**Linked List in Memory**

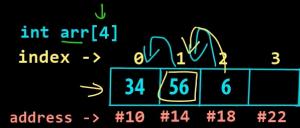
LINKED LIST REPRESENTED IN MEMORY

- Linked lists can be represented in memory by using two arrays respectively known as INFO and LINK, such that INFO[K] and LINK[K] contains information of element and next node address respectively.
- The list also requires a variable 'Name' or 'Start', which contains address of first node. Pointer field of last node denoted by NULL which indicates the end of list.
- Consider a linked list given on the side :
- Figure shows linked list. It indicates that the node of a list need not occupy adjacent elements in the array INFO and LINK.



3) Insertion & deletion of elements (Cost)

ARRAYS



1) Insertion/Deletion at the Beginning -

$O(n) = n$ Linear

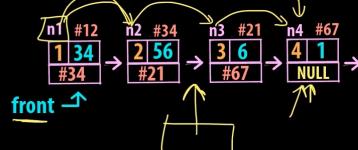
2) Insertion/Deletion at the End -

$O(n) = 1$ Constant

3) Insertion/Deletion at random nth position -

$\rightarrow O(n) = n$ Linear

LINKED LIST



1) Insertion/Deletion at the Beginning -

$O(1) = 1 \rightarrow \text{constant}$

2) Insertion/Deletion at the End -

$O(n) = n \rightarrow \text{linear}$

3) Insertion/Deletion at random nth position -

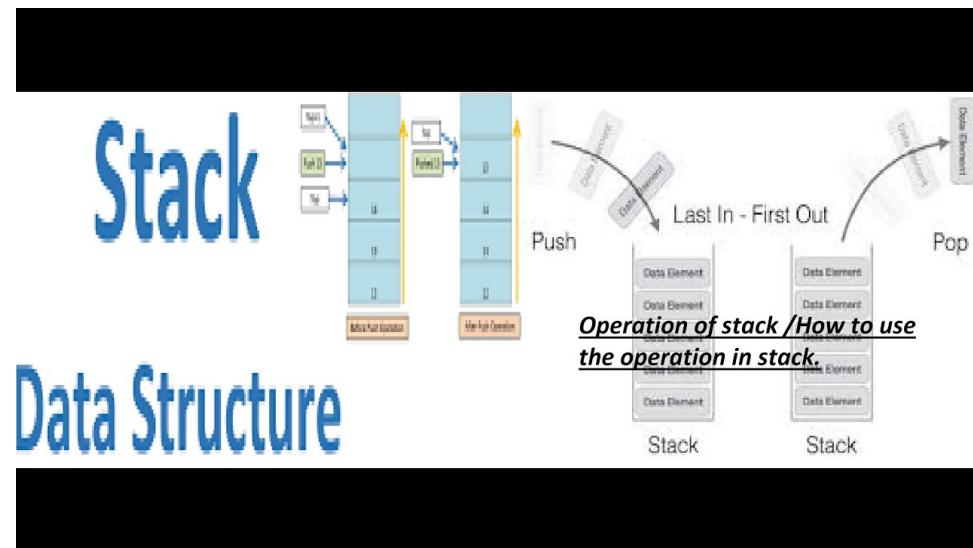
$O(n) = n \rightarrow \text{linear}$

Stack

Some Applications of Stack Data Structure

- Balancing of symbols
- Infix to Postfix /Prefix conversion
- Redo-undo features at many places like editors, photoshop.
- Forward and backward feature in web browsers
- Used in many algorithms like Tower of Hanoi, tree traversals, stock span problem, histogram problem.
- Other applications can be Backtracking, Knight tour problem, rat in a maze, N queen problem and sudoku solver
- In Graph Algorithms like Topological Sorting and Strongly Connected Components

```
{
    cout<<"Hello";
    cout<<"World";
}
```



Stack Application

- Function recursion
- System stack
 - Used in the run time to process **recursive function calls**
 - Store the **return addresses** of previous outer procedures

```

main PROC
    .
    call Sub1
    exit
main ENDP

Sub1 PROC
    .
    call Sub2
    ret
Sub1 ENDP

Sub2 PROC
    .
    call Sub3
    ret
Sub2 ENDP

Sub3 PROC
    .
    ret
Sub3 ENDP
  
```

By the time **Sub3** is called, the stack contains all three return addresses:

System Stack
Return address to Sub2
Return address to Sub1
Return address to main

[Introduction](#) [Applications of stacks](#) [Implementation of a stack](#)

Stacks: formal definition in words

Data: A finite number of objects, not necessarily distinct, having the same data type and **ordered** by when they were added.

Operations	Descriptions
isEmpty()	Task: Check whether this stack is empty. Input: None. Output: True if the stack is empty; otherwise false.
push(newEntry)	Task: Add newEntry to the top of this stack. Input: newEntry. Output: True if the operation is successful; otherwise false.
pop()	Task: Remove the top (the lastly added item) of this stack. Input: None. Output: True if the operation is successful; otherwise false.
peek()	Task: Return the top of this stack. Do not modify the stack. Input: None. Output: The top of the stack.

Data Structures – Stacks

8 / 78 Ling-Chieh Kung (NTU IM)

[Introduction](#) [Applications of stacks](#) [Implementation of a stack](#)

An UML diagram for the class Stack

- Many people describe an information system by drawing a [UML diagram](#).
– UML = unified modeling language.
- A [class](#) in a UML diagram is depicted as follows:

```

classDiagram
    class Stack {
        +isEmpty(): boolean
        +push(newEntry: ItemType): boolean
        +pop(): boolean
        +peek(): ItemType
    }
  
```

Figure 6-2
(Carrano and Henry, 2013)

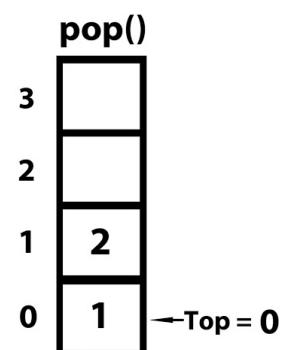
Data Structures – Stacks

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Stack: Array Implementation

```

class StackArrayImpl{
    public:
        int pop(){
            if(top == -1){
                cout << "Stack is Empty" << endl;
                return -1;
            }
            return stackArr[top--];
        }
}
  
```



NOOBCODER.COM
LIKE, COMMENT, SUBSCRIBE, SHARE

```

In [ ]: #include <iostream>
#include <stdexcept>
using namespace std;

const int MAX_STACK = 999;

template<typename ItemType>
  
```

```

class StackInterface {
public:
    virtual bool isEmpty() const = 0; //pure virtual function
    virtual bool push(const ItemType& newEntry) = 0; //pure virtual function
    virtual bool pop() = 0; //pure virtual function
    virtual ItemType peek() const = 0; //pure virtual function
};

template<typename ItemType>
class ArrayStack : public StackInterface<ItemType>{
private:
    ItemType items[MAX_STACK];
    int top;
public:
    ArrayStack(); //Default constructor
    bool isEmpty() const;
    bool push(const ItemType& newEntry);
    bool pop();
    ItemType peek() const;
};

template<typename ItemType>
ArrayStack<ItemType>::ArrayStack() : top(-1){}

template<typename ItemType>
bool ArrayStack<ItemType>::isEmpty() const{
    return top < 0;
}

template<typename ItemType>
bool ArrayStack<ItemType>::push(const ItemType& newEntry) {
    bool result = false;
    if(this->top < MAX_STACK - 1){
        top++;
        this->items[top] = newEntry;
        result = true;
    }
    return result;
}

template<typename ItemType>
bool ArrayStack<ItemType>::pop(){
    bool result = false;
    if(!this->isEmpty()){
        top--;
        return true;
    }
    return result;
}

template<typename ItemType>
ItemType ArrayStack<ItemType>::peek() const{
    if(!this->isEmpty()){
        return items[top];
    }
    else
        throw logic_error("...");
}

```

```

int main(){
    StackInterface<int>* as = new ArrayStack<int>();
    as->push(4);
    as->push(11);
    as->pop();
    try {
        cout << as->peek() << endl;
        as->pop();
        cout << as->peek() << endl;
    }
    catch (logic_error e){
        cout << "empty!" << endl;
    }

    return 0;
}

```

Stack: Linked List Implementation

Class StackLinkedListImpl{

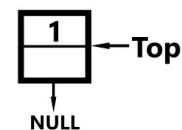
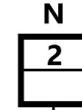
public:

```

void push(int data){
    Node *n = new Node(data);
    if(top == NULL)
        top = n;
    else{
        n->next = top;
        top = n;
    }
}

```

push(2)



NOOBCODER.COM

LIKE, COMMENT, SUBSCRIBE, SHARE

```

In [ ]: #include <iostream>
#include <stdexcept>
using namespace std;

template<typename ItemType>
class Node {
private:
    ItemType item;
    Node<ItemType>* next;
public:
    Node();
    Node(const ItemType& anItem);
    Node(const ItemType& anItem, Node<ItemType>* nextNodePtr);
    void setItem(const ItemType& anItem);
    void setNext(Node<ItemType>* nextNodePtr);
    ItemType getItem() const;
    Node<ItemType>* getNext() const;
};

template<typename ItemType>
Node<ItemType>::Node() : next(nullptr) {}


```

```

template<typename ItemType>
Node<ItemType>:: Node(const ItemType& anItem)
    : item(anItem), next(nullptr) {}

template<typename ItemType>
Node<ItemType>::Node(const ItemType& anItem, Node<ItemType>* nextNodePtr)
    : item(anItem), next(nextNodePtr) {}

template<typename ItemType>
void Node<ItemType>::setItem(const ItemType& anItem){
    item = anItem;
}

template<typename ItemType>
void Node<ItemType>::setNext(Node<ItemType>* nextNodePtr){
    next = nextNodePtr;
}

template<typename ItemType>
ItemType Node<ItemType>::getItem() const {
    return item;
}

template<typename ItemType>
Node<ItemType>* Node<ItemType>::getNext() const{
    return next;
}

template<typename ItemType>
class stackInterface {
public:
    virtual bool isEmpty() const = 0;
    virtual bool push(const ItemType& newItem) = 0;
    virtual bool pop() = 0;
    virtual ItemType peek() const = 0;
};

template<typename ItemType>
class LinkedStack : public stackInterface<ItemType>{
private:
    Node<ItemType>* topPtr;
public:
    //Constructor and Destructor
    LinkedStack(); // Default constructor
    LinkedStack(const LinkedStack<ItemType>& aStack); //Copy constructor
    virtual ~LinkedStack(); // Destructor
    // Stack Operations: Interface
    bool isEmpty() const;
    bool push(const ItemType& newItem);
    bool pop();
    ItemType peek() const;
};

template<typename ItemType>
LinkedStack<ItemType>::LinkedStack() : topPtr(nullptr) {}

template<typename ItemType>
LinkedStack<ItemType>::~LinkedStack() {

```

```

        while(!isEmpty()) { // Pop until stack is empty
            pop();
        }
    }

template<typename ItemType>
LinkedStack<ItemType>::LinkedStack(const LinkedStack<ItemType>& aStack){
    //Point to nodes in the original chain
    Node<ItemType>* origChainPtr = aStack.topPtr;

    if(origChainPtr == nullptr){
        this->topPtr = nullptr; // Original Stack is empty
    }
    else {
        //Copy the first node
        topPtr = new Node<ItemType>();
        topPtr->setItem(origChainPtr->getItem());
        Node<ItemType>* newChainPtr = topPtr;

        //Copy remaining nodes
        while(origChainPtr->getNext() != nullptr){
            origChainPtr = origChainPtr->getNext();
            ItemType nextItem = origChainPtr->getItem();
            Node<ItemType>* newNodePtr = new Node<ItemType>(nextItem);
            newChainPtr->setNext(newNodePtr);
            newChainPtr = newChainPtr->getNext();
        }
        newChainPtr->setNext(nullptr); //Mark the bottom of stack
    }
}

template<typename ItemType>
bool LinkedStack<ItemType>::isEmpty() const {
    return topPtr == nullptr;
}

template<typename ItemType>
bool LinkedStack<ItemType>::push(const ItemType& newItem){
    Node<ItemType>* newNodePtr = new Node<ItemType>(newItem, topPtr);
    topPtr = newNodePtr;
    return true;
}

template<typename ItemType>
bool LinkedStack<ItemType>::pop(){
    bool result = false;
    if(!isEmpty()){
        Node<ItemType>* nodeToDeletePtr = topPtr;
        topPtr = topPtr->getNext();
        delete nodeToDeletePtr;
        result = true;
    }
    return result;
}

template<typename ItemType>
ItemType LinkedStack<ItemType>::peek() const {
    if(!this->isEmpty()){
        return topPtr->getItem();
    }
}

```

```

    else {
        throw logic_error("...");
    }

}

int main(){

    LinkedStack<int> ls;
    ls.push(4);
    ls.push(11);
    LinkedStack<int> lsCopy(ls);
    lsCopy.pop();
    try {
        cout << lsCopy.peek() << endl;
        cout << ls.peek() << endl;
    }
    catch(logic_error e) {
        cout << "empty!" << endl;
    }

    return 0;
}

```

```

In [ ]: #include<iostream>
using namespace std;

class Node{
public:
    int data;
    Node* next;

    Node();
    Node(int d);
};

Node::Node(){
    data = 0;
    next = nullptr;
}

Node::Node(int d){
    data = d;
    next = nullptr;
}

class Stack{
public:
    Node* top;

    Stack();
    bool isEmpty();
    bool checkIfNodeExist(Node* newNode);
    void push(Node* newNode);
    Node* pop();
    Node* peek();
    int count();
    void display();
};

```

```

Stack::Stack(){
    top = nullptr;
}

bool Stack::isEmpty(){
    if(top==nullptr)
        return true;
    else
        return false;
}

bool Stack::checkIfNodeExist(Node* newNode){
    Node* temp = top;
    while(temp!=nullptr){
        if(temp->data==newNode->data)
            return true;
        else {
            temp = temp->next;
        }
    }
    return false;
}

void Stack::push(Node* newNode){
    if(top==nullptr){
        top = newNode;
        cout << "Pushed First Node Successfully" << endl << endl;
    }
    else if(checkIfNodeExist(newNode)){
        cout << "Node with value already exist" << endl;
        cout << "Please enter different key and value" << endl << endl;
    }
    else {
        newNode->next = top;
        top = newNode;
        cout << "Pushed Successfully" << endl;
    }
}

Node* Stack::pop(){
    if(isEmpty()){
        return nullptr;
    }
    else {
        Node* temp = top;
        top = top->next;
        return temp;
    }
}

Node* Stack::peek(){
    if(isEmpty()){
        return nullptr;
    }
    else {
        return top;
    }
}

int Stack::count(){
    int count = 0;
}

```

```

Node* ptr = top;
while(ptr!=nullptr){
    count++;
    ptr = ptr->next;
}
return count;
}

void Stack::display(){
    if(isEmpty()){
        cout << "Stack is empty";
    }
    else {
        cout << "Stack: " << endl;
        Node* ptr = top;
        while(ptr!=nullptr){
            cout << "Value: " << ptr->data << endl;
            ptr = ptr->next;
        }
        cout << endl;
    }
}

int main(){
    Stack s;
    int option, data;

    do {
        cout << "What operation do you want to perform?" << endl;
        cout << "Select Option Number or 0 to Exit" << endl << endl;
        cout << "1. Push" << endl;
        cout << "2. Pop" << endl;
        cout << "3. Peek" << endl;
        cout << "4. isEmpty" << endl;
        cout << "5. Count" << endl;
        cout << "6. Display" << endl;
        cout << "9. Clear Screen" << endl;
        cout << "0. Exit" << endl << endl;

        cout << "Enter Your Option Number: ";
        cin >> option;

        Node* newNode = new Node();

        switch(option){
            case 0:
                break;
            case 1:
                cout << "Enter Value of Node" << endl;
                cin >> data;
                newNode->data = data;
                s.push(newNode);
                break;
            case 2:
                cout << "Pop Function Called." << endl;
                newNode = s.pop();
                if(newNode==nullptr){
                    cout << "Stack Underflow" << endl << endl;
                }
        }
    } while(option != 0);
}

```

```

        }
    else{
        cout << "Pop Value: " << newNode->data << endl << endl;
    }

    delete newNode;

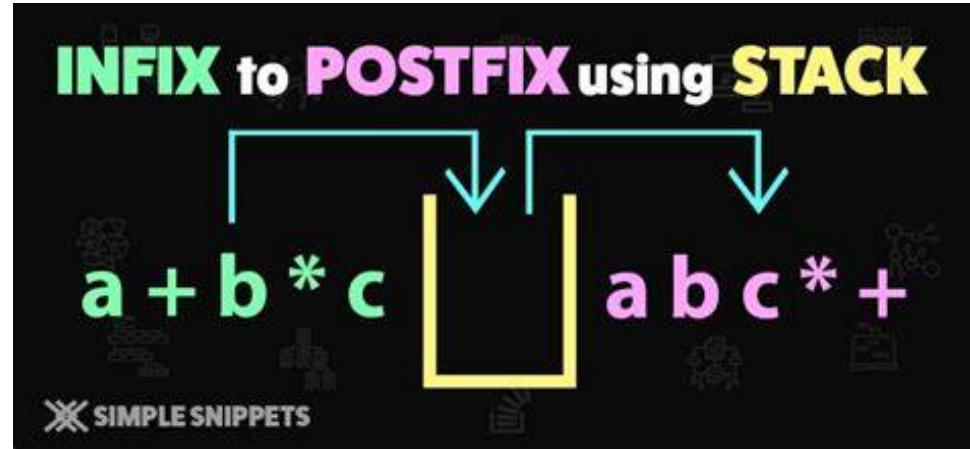
    break;
case 3:
    cout << "Peek Function Called." << endl;
    newNode = s.peek();
    cout << "The Top Value of Stack: " << newNode->data << endl << endl;
    break;
case 4:
    if(s.isEmpty()){
        cout << "Stack is Empty" << endl;
    }
    else {
        cout << "Stack is not Empty" << endl;
    }
    break;
case 5:
    cout << "Count Function Called." << endl;
    cout << "No of Node in the Stack: " << s.count() << endl << endl;
    break;
case 6:
    s.display();
    cout << endl;
    break;
case 9:
    system("cls");
    break;
default:
    cout << "Please Enter the Proper Option Number" << endl;
    break;
}

} while(option!=0);

return 0;
}

```

Infix Prefix & Postfix Expressions Using Stack



Pseudocode of Infix to Postfix conversion using STACK Data Structure

Rules for infix to Postfix using STACK DS

- Scan Expression from Left to Right
- PRINT OPERANDS as the arrive.
- If OPERATOR arrives & Stack IS EMPTY, PUSH this OPERATOR onto the Stack.
- If incoming OPERATOR has HIGHER precedence than the TOP of the Stack, PUSH it on stack.
- If incoming OPERATOR has LOWER precedence than the TOP of the Stack, then POP and PRINT the TOP. Then test the incoming against the NEW TOP of stack.
- If incoming OPERATOR has EQUAL precedence with TOP of Stack, use ASSOCIATIVITY Rules.
- For ASSOCIATIVITY of LEFT to RIGHT - POP and PRINT the TOP of stack, then PUSH the incoming
- For ASSOCIATIVITY of RIGHT to LEFT - PUSH incoming OPERATOR on stack.
- At the end of Expression, POP & PRINT all from the stack.
- If incoming SYMBOL is '(' PUSH it onto Stack.
- If incoming SYMBOL is ')' POP the stack and PRINT OPERATORS till ')' is found.
- If TOP of stack is ')' PUSH OPERATOR on Stack

Pseudocode -

```

FUNCTION InfixToPostfix(stack,infix)
string postfix.
LOOP i=0 to i<infix.length
1. IF infix[i] > OPERAND then postfix+=infix[i]
2. ELSE IF infix[i] > '(' then PUSH to stack
3. ELSE IF infix[i] > ')' then POP & PRINT stack
    till stack gets empty OR ')' is found. POP that ')'
4. ELSE IF infix[i] > OPERATOR (+,-,*,/,^) then
    4.1. IF stack > EMPTY then PUSH infix[i] on stack
    4.2. ELSE IF stack > NOT EMPTY
        4.1.1 IF precedence(infix[i]) > precedence(stack.top)
            then -> PUSH infix[i] on Stack
        4.1.2 ELSE IF (precedence(infix[i]) == precedence(stack.top))
            && ( infix[i]=='^' )
            then -> PUSH infix[i] on Stack
        4.1.3 ELSE
            4.1.3.1 WHILE stack > (NOT EMPTY)
                && precedence(infix[i])<=precedence(stack.top)
                then POP and PRINT stack
            4.1.3.2 PUSH current OPERATOR on Stack
    END FOR LOOP

```

Q1) Convert INFIX to POSTFIX -

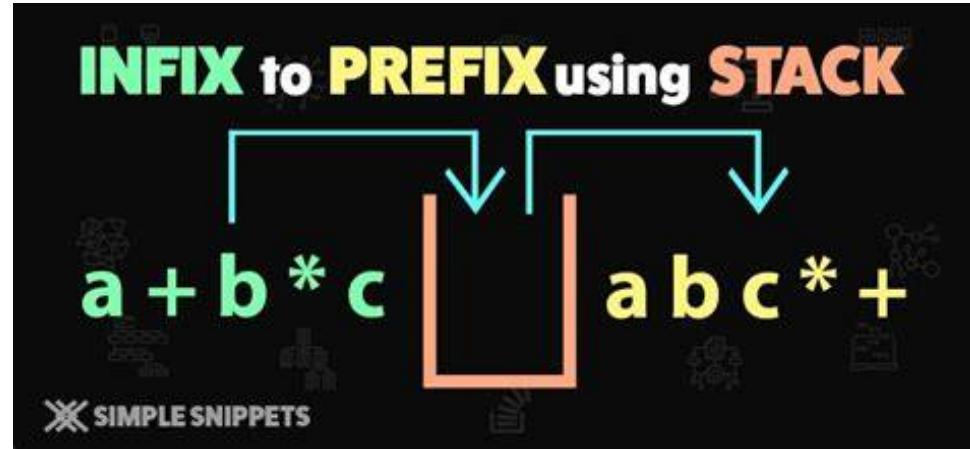
((a + b - c) * d ^ e ^ f) / g



Manual Conversion -
ab+c- def^*g/
Postfix -

PRIORITY TABLE

Priority level	Operators
1.	Unary +,Unary -,NOT
2.	^
3.	/, *
4.	+, -
5.	<, >, <=, >=
6.	==, !=
7.	&&
8.	



Pseudocode of Infix to PREFIX conversion using STACK Data Structure

Rules for INFIX to PREFIX using Stack DS -

- 1) REVERSE infix expression & swap '(' to ')' & ')' to '('
- 2) Scan Expression from Left to Right
- 3) Print OPERANDs as they arrive
- 4) If OPERATOR arrives & Stack is empty, PUSH to stack
- 5) If incoming OPERATOR has HIGHER precedence than the TOP of the Stack, PUSH it on stack
- 6) If incoming OPERATOR has EQUAL precedence with TOP of Stack && incoming OPERATOR is '^', POP & PRINT TOP of Stack. Then test the incoming OPERATOR against the NEW TOP of stack.
- 7) If incoming OPERATOR has EQUAL precedence with TOP of Stack, PUSH it on Stack.
- 8) If incoming OPERATOR has LOWER precedence than the TOP of the Stack, then POP and PRINT the TOP of Stack. Then test the incoming OPERATOR against the NEW TOP of stack.
- 9) At the end of Expression, POP & PRINT all OPERATORS from the stack
- 10) If incoming SYMBOL is ')' PUSH it onto Stack.
- 11) If incoming SYMBOL is ')' POP the stack & PRINT OPERATORS till ')' is found or Stack Empty. POP out that ')' from stack
- 12) If TOP of stack is ')' PUSH OPERATOR on Stack
- 13) At the end REVERSE output string again.

Pseudocode - Infix to Prefix

```

FUNCTION InfixToPrefix(stack, infix)
    infix = reverse(infix)
    LOOP i=0 to i<infix.length
        IF infix[i] is OPERAND -> prefix+=infix[i]
        ELSE IF infix[i] is '(' -> Stack.push(infix[i])
        ELSE IF infix[i] is ')' -> POP & PRINT Stack values
            till ')' is found & stack NOT EMPTY .POP that ')'
        ELSE IF infix[i] is A OPERATOR(+,-,*,/,\^) ->
            IF Stack IS EMPTY -> PUSH OPERATOR on stack
            ELSE ->
                IF precedence(infix[i])>precedence(stack.top)
                    -> PUSH infix[i] on Stack
                ELSE IF precedence(infix[i])==precedence(stack.top)
                    -> POP & PRINT TOP of Stack till this condition is true.
                    -> PUSH infix[i] on Stack
                ELSE IF precedence(infix[i])<precedence(stack.top)
                    -> POP & PRINT till Stack NOT EMPTY &&
                        precedence(infix[i])<precedence(stack.top)
                    -> PUSH infix[i] onto Stack
            END LOOP
            POP & PRINT Remaining Elements of Stack
            prefix = reverse(prefix) & RETURN
        
```

Q1) Convert INFIX to PREFIX -

$((a+b-c)^*d^e^f)/g$

Stack S

Manual Conversion -
 $/ \rightarrow +abc^d^e^f/g$
Using Stack -

Rules & Pseudocode of POSTFIX to INFIX Conversion using Stack DS

Rules for POSTFIX to INFIX using Stack -

- 1) Scan POSTFIX expression from LEFT to RIGHT
- 2) If the incoming symbol is a OPERAND, PUSH it onto the Stack
- 3) If the incoming symbol is a OPERATOR, POP 2 OPERANDs from the Stack, ADD this incoming OPERATOR in between the 2 OPERANDs, ADD '(' & ')' to the whole expression & PUSH this whole new expression string back into the Stack.
- 4) At the end POP & PRINT the full INFIX expression from the Stack.

$\Rightarrow (((a+b)-c)*(d^(e^f)))/g$

Pseudocode -

```

Function PostfixToInfix(string postfix)
    1. stack
    2. LOOP: i=0 to postfix.length
        2.1. IF postfix[i] is OPERAND ->
            2.1.1 s.push(postfix[i])
        2.2. ELSE IF postfix[i] is OPERATOR ->
            2.2.1 op1 = s.top()
            2.2.2 s.pop()
            2.2.3 op2 = s.top()
            2.2.4 s.pop()
            2.2.5 exp = '(' + op2 + postfix[i] + op1 + ')'
            -> 2.2.6 s.push(exp)
    END LOOP
    RETURN s.top
    
```

Q1) Convert POSTFIX to INFIX -

$a|b|+|c|-|d|e|f|^|\wedge|*|g|/$

$OP_1 = g$
 $OP_2 = (((a+b)-c)*(d^(e^f)))/g$

Stack S

Manual Conversion -
 $((a+b)-c)*(d^(e^f))/g$
Using Stack -

$((((a+b)-c)*(d^(e^f)))/g)$

Rules & Pseudocode of PREFIX to INFIX Conversion using Stack DS

Rules for PREFIX to INFIX using Stack -

- 1) REVERSE the PREFIX expression
OR
- 1) Scan PREFIX expression from RIGHT to LEFT
- 2) If the incoming symbol is a OPERAND, PUSH it onto the Stack
- 3) If the incoming symbol is a OPERATOR, POP 2 OPERANDs from the Stack, ADD this incoming OPERATOR in between the 2 OPERANDs, ADD '()' to the whole expression & PUSH this whole new expression string back into the Stack.
- 4) At the end POP & PRINT the full INFIX expression from the Stack.

Pseudocode -

```
Function PrefixToInfix(string prefix)
1. stack s
2. LOOP: i=prefix.length-1 to 0 i=0 . .
    2.1 IF prefix[i] is OPERAND ->
        2.1.1 s.push(prefix[i])
    => 2.2 ELSE IF prefix[i] is OPERATOR ->
        2.2.1 op1 = s.top()
        2.2.2 s.pop()
        2.2.3 op2 = s.top()
        2.2.4 s.pop()
        2.2.5 exp = '('+op1+prefix[i]+op2+')'
        2.2.6 s.push(exp)
    END LOOP
3. RETURN s.top
```

exp =

Q1) Convert PREFIX to INFIX -

$\frac{a+b-c}{d*(e^f)}$

$op_1 = ((a+b)-c)$

$op_2 = *$

$exp = ((a+b)-c)$

$op_1 = ((a+b)-c)$

$op_2 = ^$

$exp = ((a+b)-c)^{(d*(e^f))}$

$op_1 = ((a+b)-c)$

$op_2 = /$

$exp = (((a+b)-c)*(d^(e^f)))/g$

$op_1 = (((a+b)-c)*(d^(e^f)))$

$op_2 = /$

$exp = (((a+b)-c)*(d^(e^f)))/g$

$op_1 = (((a+b)-c)*(d^(e^f)))$

$op_2 = /$

$exp = (((a+b)-c)*(d^(e^f)))/g$

$op_1 = (((a+b)-c)*(d^(e^f)))$

$op_2 = /$

$exp = (((a+b)-c)*(d^(e^f)))/g$

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$op_1 = (((a+b)-c)*(d^(e^f)))$

$op_2 = /$

$exp = (((a+b)-c)*(d^(e^f)))/g$

$op_1 = (((a+b)-c)*(d^(e^f)))$

$op_2 = /$

```

        return true;
    else
        return false;
}

bool isOperand(char c){
    if((c>='A' && c<='Z') || (c>='a' && c<='z'))
        return true;
    else
        return false;
}

int precedence(char c){
    if(c=='^')
        return 3;
    else if(c=='*' || c=='/')
        return 2;
    else if(c=='+' || c=='-')
        return 1;
    else
        return -1;
}

string infixToPostfix(string infix){
    stack<char> s;
    string postfix;

    for(int i=0; i<infix.length(); i++){
        // check if operand
        if((infix[i]>='A' && infix[i]<='Z')
           || (infix[i]>='a' && infix[i]<='z')){
            postfix += infix[i];
        }
        else if(infix[i]=='('){
            s.push(infix[i]);
        }
        else if(infix[i]==')'){
            while((s.top()!='(') && (!s.empty())){
                postfix += s.top();
                s.pop();
            }
            if(s.top()=='('){
                s.pop();
            }
        }
        else if(isOperator(infix[i])){
            if(s.empty()){
                s.push(infix[i]);
            }
            else{
                if(precedence(infix[i])>precedence(s.top())){
                    s.push(infix[i]);
                }
                else if((precedence(infix[i])==precedence(s.top())) && (infix[i]==
                    s.top())){
                    s.push(infix[i]);
                }
                else{
                    while(!s.empty() && (precedence(infix[i])<=precedence(s.top()))
                        postfix += s.top();
                    }
                    s.push(infix[i]);
                }
            }
        }
    }
}

```

```

        s.pop();
    }
    s.push(infix[i]);
}
}

while(!s.empty()){
    postfix += s.top();
    s.pop();
}

return postfix;
}

string infixToPrefix(string infix){
    stack<char> s;
    string prefix;

    reverse(infix.begin(), infix.end()); // include<algorithm>

    for(int i=0; i<infix.length(); i++){
        if(infix[i]=='(')
            infix[i] = ')';
        else if(infix[i]==')')
            infix[i] = '(';
    }

    cout << "\nReverse Infix: " << infix << endl;

    for(int i=0; i<infix.length(); i++){
        if((infix[i]>='A' && infix[i]<='Z') || (infix[i]>='a' && infix[i]<='z')){
            prefix += infix[i];
        }
        else if(infix[i]=='('){
            s.push(infix[i]);
        }
        else if(infix[i]==')'){
            while((s.top()!='(') && (!s.empty())){
                prefix += s.top();
                s.pop();
            }
            if(s.top()=='('){
                s.pop();
            }
        }
        else if(isOperator(infix[i])){
            if(s.empty()){
                s.push(infix[i]);
            }
            else {
                if(precedence(infix[i]) > precedence(s.top())){
                    s.push(infix[i]);
                }
                else if((precedence(infix[i])==precedence(s.top())) && (infix[i]==
                    while((precedence(infix[i])==precedence(s.top())) &&(infix[i]==
                        prefix += s.top();
                        s.pop();
                }
            }
        }
    }
}

```

```

        s.push(infix[i]);
    }
    else if(precedence(infix[i])==precedence(s.top())){
        s.push(infix[i]);
    }
    else {
        while((!s.empty()) && (precedence(infix[i])<precedence(s.top())))
            prefix += s.top();
        s.pop();
    }
    s.push(infix[i]);
}
}

while((!s.empty())){
    prefix += s.top();
    s.pop();
}

reverse(prefix.begin(), prefix.end());

return prefix;
}

string postToInfix(string postfix){
stack<string> s;

for(int i=0; i<postfix.length(); i++){
    if(isOperand(postfix[i])){
        string op(1, postfix[i]);
        s.push(op);
    }
    else {
        string op1 = s.top();
        s.pop();
        string op2 = s.top();
        s.pop();
        s.push('(' + op2 + postfix[i] + op1 + ')');
    }
}

return s.top();
}

string preToInfix(string prefix){
stack<string> s;

for(int i=prefix.length()-1; i>=0; i--){
    if(isOperand(prefix[i])){
        string op(1, prefix[i]);
        s.push(op);
    }
    else {
        string op1 = s.top();
        s.pop();
        string op2 = s.top();
        s.pop();
        s.push('(' + op1 + prefix[i] + op2 + ')');
    }
}
}

```

```

    }

    return s.top();
}

string postToPrefix(string postfix){
    stack<string> s;

    for(int i=0; i<postfix.length(); i++){
        if(isOperand(postfix[i])){
            string op(1, postfix[i]);
            s.push(op);
        }
        else {
            string op1 = s.top();
            s.pop();
            string op2 = s.top();
            s.pop();
            s.push(postfix[i] + op2 + op1);
        }
    }

    return s.top();
}

string preToPostfix(string prefix){
    stack<string> s;

    for(int i=prefix.length()-1; i>=0; i--){
        if(isOperand(prefix[i])){
            string op(1, prefix[i]);
            s.push(op);
        }
        else {
            string op1 = s.top();
            s.pop();
            string op2 = s.top();
            s.pop();
            s.push(op1 + op2 + prefix[i]);
        }
    }

    return s.top();
}

int main(){

    string infix_exp, postfix_exp, prefix_exp;
    string post_in_exp, pre_in_exp, post_pre_exp, pre_post_exp;
    cout << "Enter a Infix Expression: " << endl;
    cin >> infix_exp;

    cout << "\nInfix Express: " << infix_exp << endl;
    postfix_exp = infixToPostfix(infix_exp);
    cout << "\nPostfix Expression: " << postfix_exp << endl;
    prefix_exp = infixToPrefix(infix_exp);
}

```

```

cout << "\nPrefix Expression: " << prefix_exp << endl;
cout << endl;

post_in_exp = postToInfix(postfix_exp);
cout << "\nPostfix to Infix: " << post_in_exp << endl;
cout << endl;

pre_in_exp = preToInfix(prefix_exp);
cout << "\nPrefix to Infix: " << pre_in_exp << endl;
cout << endl;

post_pre_exp = postToPrefix(postfix_exp);
cout << "\nPostfix to Prefix: " << post_pre_exp << endl;
cout << endl;

pre_post_exp = preToPostfix(prefix_exp);
cout << "\nPrefix to Postfix: " << pre_post_exp << endl;
cout << endl;

return 0;
}

```

Queue

Some Applications of Queue Data Structure

Queue is used when things but have to be processed in **First In First Out** order. Like –

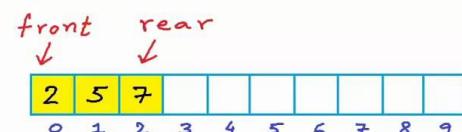
- CPU scheduling, Disk Scheduling.
- Handling of interrupts in **real-time** systems. The interrupts are handled in the same order as they arrive, First come first served.
- In real life, Call Center phone systems will use Queues, to hold people calling them in an order, until a service representative is free.
- When data is transferred **asynchronously** between two processes. Queue is used for synchronization.



```

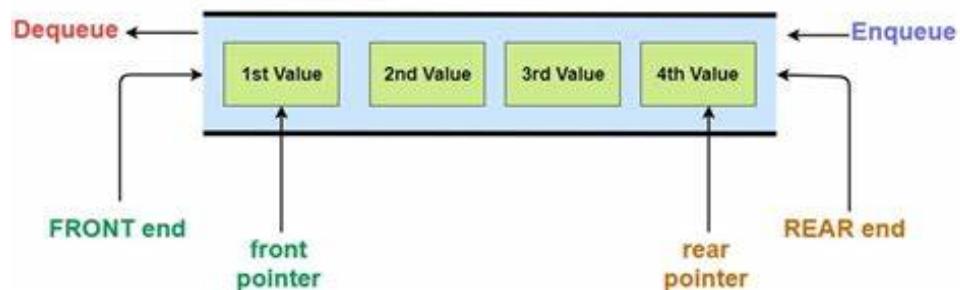
Dequeue()
{
    if IsEmpty()
        return
    else if front == rear
        front = rear = -1
    else
        front = front + 1
}

```



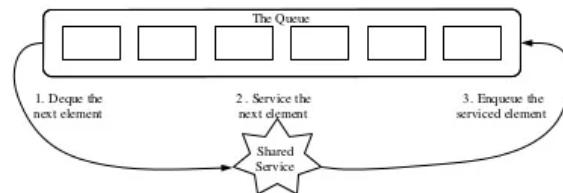
Enqueue(2)
Enqueue(5)
Enqueue(7)

Queue is a linear data structure which operates in a LIFO(Last In Last Out) or FIFO(First In First Out) pattern.



Application: Round Robin Schedulers

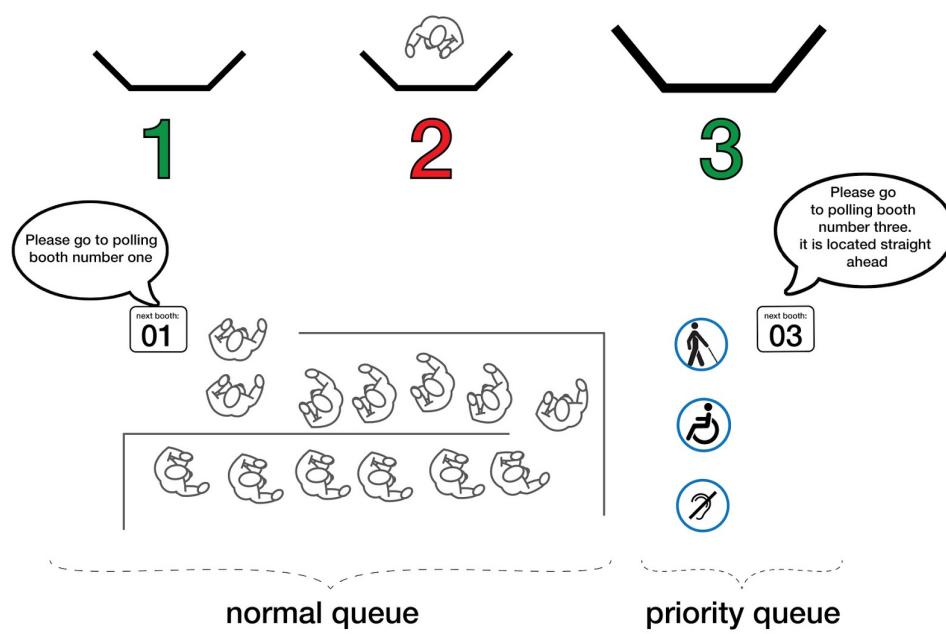
- We can implement a round robin scheduler using a queue, Q , by repeatedly performing the following steps:
 - $e = Q.dequeue()$
 - Service element e
 - $Q.enqueue(e)$



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Priority Queue

<code>public class IndexMinPQ<Item extends Comparable<Item>></code>	
<code>IndexMinPQ(int maxN)</code>	<i>create a priority queue of capacity maxN with possible indices between 0 and maxN-1</i>
<code>void insert(int k, Item item)</code>	<i>insert item; associate it with k</i>
<code>void change(int k, Item item)</code>	<i>change the item associated with k to item</i>
<code>boolean contains(int k)</code>	<i>is k associated with some item?</i>
<code>void delete(int k)</code>	<i>remove k and its associated item</i>
<code>Item min()</code>	<i>return a minimal item</i>
<code>int minIndex()</code>	<i>return a minimal item's index</i>
<code>int delMin()</code>	<i>remove a minimal item and return its index</i>
<code>boolean isEmpty()</code>	<i>is the priority queue empty?</i>
<code>int size()</code>	<i>number of items in the priority queue</i>



Queue Summary

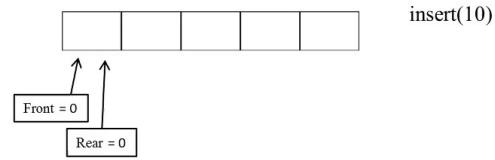
- Queue Operation Complexity for Different Implementations

	Array Fixed-Size	Array Expandable (doubling strategy)	List Singly-Linked
dequeue()	O(1)	O(1)	O(1)
enqueue(o)	O(1)	O(n) Worst Case O(1) Best Case O(1) Average Case	O(1)
front()	O(1)	O(1)	O(1)
Size(), isEmpty()	O(1)	O(1)	O(1)

Implementation of Queue using Array

```
int arr[5] int rear = -1; int front = -1;
```

- Initially when the queue is empty, the value of both front and rear will be -1
- For insertion, the value of rear is incremented by 1 and the element is inserted at the new rear position



```
In [ ]: #include <iostream>
#include <string>
using namespace std;

const int MAX_QUEUE = 5;

template<typename ItemType>
class Queue {
private:
    ItemType items[MAX_QUEUE];
    int front;
    int rear;
public:
    Queue(); //default constructor
    ~Queue();
    bool isEmpty() const;
    bool isFull() const;
    void enqueue(const ItemType& newEntry);
    ItemType dequeue();
    int size();
    void display();
};

template<typename ItemType>
Queue<ItemType>::Queue() {
    front = -1;
    rear = -1;

    for (int i = 0; i < MAX_QUEUE; i++){
        items[i] = -1;
    }
}

template<typename ItemType>
Queue<ItemType>::~Queue() {
    delete [] items;
}

template<typename ItemType>
bool Queue<ItemType>::isEmpty() const{
    return ((front < 0) && (rear < 0));
};

template<typename ItemType>
bool Queue<ItemType>::isFull() const {
    return ((this->front == 0) && (this->rear == (MAX_QUEUE - 1)));
};

template<typename ItemType>
void Queue<ItemType>::enqueue(const ItemType& newEntry){
    if(this->rear == MAX_QUEUE - 1){
        cout << "Can not Enqueue." << endl;
        return;
    }
    if(this->isEmpty()){
        this->front = 0;
        this->rear = 0;
        this->items[rear] = newEntry;
    }
    else{
        this->rear++;
        this->items[rear] = newEntry;
    }
}
```

```

    }
    else {
        rear++;
        this->items[rear] = newEntry;
    }
}

template<typename ItemType>
ItemType Queue<ItemType>::dequeue(){
    int result = 0;

    if(this->isEmpty()){
        return -1;
    }
    if(this->front == this->rear){
        result = this->items[front];
        this->items[front] = -1;
        rear = -1;
        front = -1;
        return result;
    }
    else {
        result = this->items[front];
        this->items[front] = -1;
        front++;
        return result;
    }
};

template<typename ItemType>
int Queue<ItemType>::size(){
    if(!this->isEmpty())
        return (this->rear - this->front + 1);
    else
        return 0;
}

template<typename ItemType>
void Queue<ItemType>::display(){
    cout << "The elements in Queue are as below: " << endl;
    for (int i = 0; i < MAX_QUEUE; i++){
        cout << items[i] << " ";
    }
    cout << endl << endl << endl;
}

int main(){
    Queue<int> q1;//Call default constructor

    int option, item;

    do {
        cout << "1. Enqueue" << endl;
        cout << "2. Dequeue" << endl;
        cout << "3. isEmpty" << endl;
        cout << "4. isFull" << endl;
        cout << "5. Count" << endl;
        cout << "6. Display" << endl;
        cout << "7. Clear Screen" << endl;
}

```

```
cout << "0. Exit" << endl;
cout << endl;

cout << "Select Your Option Number: ";
cin >> option;

switch(option){
    cout << endl;
    case 0:
        break;
    case 1:
        cout << "Enqueue your number(>=0): ";
        cin >> item;
        q1.enqueue(item);
        cout << endl << endl;
        break;
    case 2:
        cout << "Dequeue..." << endl;
        item = q1.dequeue();
        if (item == -1){
            cout << "Queue is empty" << endl << endl;
        }
        else {
            cout << "Dequeued element: " << item << endl << endl;
        }
        break;
    case 3:
        if(q1.isEmpty())
            cout << "Queue is empty" << endl << endl;
        else
            cout << "Queue is not empty" << endl << endl;
        break;
    case 4:
        if(q1.isFull())
            cout << "Queue is full" << endl << endl;
        else
            cout << "Queue is not full" << endl << endl;
        break;
    case 5:
        cout << "Count of queue: " << q1.size() << endl << endl;
        break;
    case 6:
        q1.display();
        break;
    case 7:
        system("cls");
        break;
    default:
        cout << "Please enter the proper option" << endl;
        break;
}

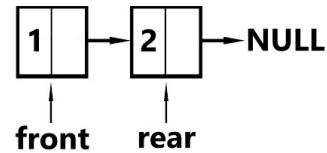
} while (option != 0);

return 0;
}
```

Queue: LinkedList Implementation

```
class QueueLinkedListImpl{
public:
    int dequeue(){
        ➔ if(front == NULL){
            cout << "Empty Queue" << endl;
            return -1;
        }
        Node *temp = front;
        int tempData = front->data;
        if(front == rear)
            front = rear = NULL;
        else
            front = front->next;

        delete temp;
        return tempData;
    }
};
```



ORDER.C

```
In [ ]: #include <iostream>
using namespace std;

class Node{
public:
    int data;
    Node* next;

    Node();
    Node(int d);
};

Node::Node(){
    data = 0;
    next = nullptr;
}

Node::Node(int d){
    data = d;
    next = nullptr;
}

class Queue{
public:
    Node* front;
    Node* rear;

    Queue();
    bool isEmpty();
    bool checkIfNodeExist(Node* newNode);
    void enqueue(Node* newNode);
    Node* dequeue();
    int count();
    void display();
};

Queue::Queue() : front(nullptr), rear(nullptr) {}

bool Queue::isEmpty(){
    if(front==nullptr && rear==nullptr)
        return true;
    else
```

```
        return false;
    }

bool Queue::checkIfNodeExist(Node* newNode){
    Node* ptr = front;
    while(ptr!=nullptr){
        if(ptr->data==newNode->data)
            return true;
        else
            ptr = ptr->next;
    }
    return false;
}

void Queue::enqueue(Node* newNode){
    if(checkIfNodeExist(newNode)){
        cout << "Node with value: " << newNode->data << " already exist" << endl;
        cout << "Enter different value" << endl << endl;
    }
    if(isEmpty()){
        front = newNode;
        rear = newNode;
        cout << "Enqueue first node successfully." << endl << endl;
    }
    else {
        rear->next = newNode;
        rear = newNode;
        cout << "Enqueue successfully" << endl << endl;
    }
}

Node* Queue::dequeue(){
    if(isEmpty()){
        cout << "Queue is empty" << endl << endl;
    }
    else if(front==rear){
        cout << "Dequeue Value: " << front->data << endl << endl;
        front = nullptr;
        rear = nullptr;
    }
    else {
        cout << "Dequeue Value: " << front->data << endl << endl;
        front = front->next;
    }
    return front;
}

int Queue::count(){
    if(isEmpty())
        return 0;
    else {
        int count = 0;
        Node* ptr = front;
        while(ptr!=nullptr){
            count++;
            ptr = ptr->next;
        }
        return count;
    }
}
```

```

void Queue::display(){
    if(isEmpty()){
        cout << "Queue is empty" << endl << endl;
    }
    else {
        cout << "Queue: " << endl;
        Node* ptr = front;
        while(ptr!=nullptr){
            cout << ptr->data << " -> ";
            ptr = ptr->next;
        }
    }
    cout << endl << endl;
}

int main(){
    Queue q;
    int option, data;

    do {
        cout << "What operation do you want to perform?" << endl;
        cout << "Select Option Number or 0 to Exit" << endl << endl;
        cout << "1. Enqueue" << endl;
        cout << "2. Dequeue" << endl;
        cout << "3. isEmpty" << endl;
        cout << "4. Count" << endl;
        cout << "8. Display" << endl;
        cout << "9. Clear Screen" << endl;
        cout << "0. Exit" << endl << endl;

        cout << "Enter Your Option Number: ";
        cout << endl;
        cin >> option;

        Node* newNode = new Node();

        switch(option){
            case 0:
                break;
            case 1:
                cout << "Enqueue Function Called." << endl;
                cout << "Enter Value of Node: ";
                cin >> data;
                newNode->data = data;
                q.enqueue(newNode);
                cout << endl;
                break;
            case 2:
                cout << "Dequeue Function Called." << endl;
                newNode = q.dequeue();
                delete newNode;
                cout << endl;
                break;
            case 3:
                if(q.isEmpty()){
                    cout << "Queue is Empty" << endl;
                }
                else {
                    cout << "Queue is not Empty" << endl;
                }
        }
    }
}

```

```

        break;
    case 4:
        cout << "Count Function Called." << endl;
        cout << "No of Node in the Queue: " << q.count() << endl << endl;
        break;
    case 8:
        q.display();
        cout << endl;
        break;
    case 9:
        system("cls");
        break;
    default:
        cout << "Please Enter the Proper Option Number" << endl;
        break;
    }
}

} while(option!=0);

return 0;
}

```

Circular Queue

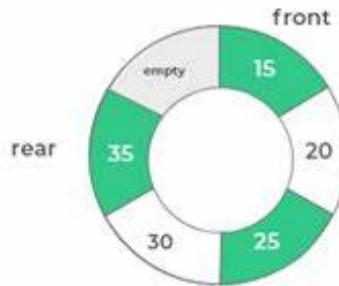
CIRCULAR QUEUE DATA STRUCTURE

**THEORY
&
C++
PROGRAM**

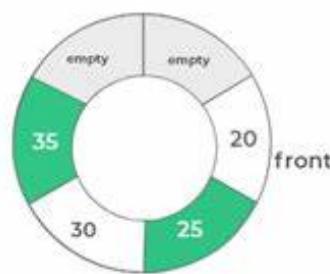
- Tanmay Sakpal

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Circular Queue using Array in C



Enqueue (50)

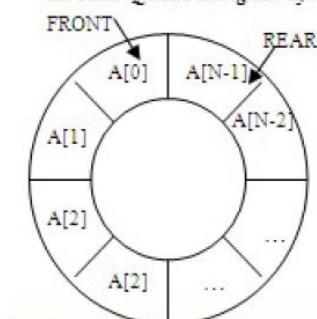


Dequeue()

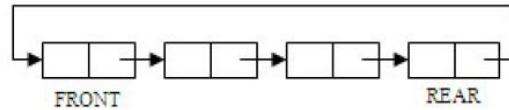
Enqueue() $\text{rear} = (\text{rear} + 1) \% \text{SIZE};$

Dequeue() $\text{front} = (\text{front} + 1) \% \text{SIZE};$

Circular Queue using Arrays



Circular Queue using Linked list



```
In [ ]: #include <iostream>
#include <string>
#include <cmath>
using namespace std;

const int MAX_QUEUE = 5;

template<typename ItemType>
class CircularQueue {
private:
    ItemType items[MAX_QUEUE];
    int front;
    int rear;
    int itemCount;
public:
    CircularQueue(); //default constructor
```

```
~CircularQueue();
bool isEmpty() const;
bool isFull() const;
void enqueue(const ItemType& newEntry);
ItemType dequeue();
int size();
void display();
};

template<typename ItemType>
CircularQueue<ItemType>::CircularQueue() {
    front = -1;
    rear = -1;
    itemCount = 0;

    for (int i = 0; i < MAX_QUEUE; i++){
        items[i] = -1;
    }
}

template<typename ItemType>
CircularQueue<ItemType>::~CircularQueue() {
    delete [] this->items;
}

template<typename ItemType>
bool CircularQueue<ItemType>::isEmpty() const{
    return ((front < 0) && (rear < 0));
}

template<typename ItemType>
bool CircularQueue<ItemType>::isFull() const {
    return (((this->rear + 1) % MAX_QUEUE) == this->front);
}

template<typename ItemType>
void CircularQueue<ItemType>::enqueue(const ItemType& newEntry){
    if(this->isFull()){
        cout << "Queue is full" << endl << endl;
        return;
    }
    if(this->isEmpty()){
        this->front = 0;
        this->rear = 0;
        this->items[rear] = newEntry;
    }
    else {
        rear = (rear + 1) % MAX_QUEUE;
        this->items[rear] = newEntry;
    }
    itemCount++;
}

template<typename ItemType>
ItemType CircularQueue<ItemType>::dequeue(){
    int result = 0;

    if(this->isEmpty()){
        return -1;
    }
}
```

```

if(this->front == this->rear){
    result = this->items[front];
    this->items[front] = -1;
    rear = -1;
    front = -1;
    itemCount--;
    return result;
}
else {
    result = this->items[front];
    this->items[front] = -1;
    front = (front + 1) % MAX_QUEUE;
    itemCount--;
    return result;
}
};

template<typename ItemType>
int CircularQueue<ItemType>::size(){
    if(!this->isEmpty())
        return itemCount;
    else
        return 0;
}

template<typename ItemType>
void CircularQueue<ItemType>::display(){
    cout << "The elements in Queue are as below: " << endl;
    for (int i = 0; i < MAX_QUEUE; i++){
        cout << items[i] << " ";
    }
    cout << endl << endl << endl;
}

int main(){

    CircularQueue<int> q1;//Call default constructor

    int option, item;

    do {
        cout << "1. Enqueue" << endl;
        cout << "2. Dequeue" << endl;
        cout << "3. isEmpty" << endl;
        cout << "4. isFull" << endl;
        cout << "5. Count" << endl;
        cout << "6. Display" << endl;
        cout << "7. Clear Screen" << endl;
        cout << "0. Exit" << endl;
        cout << endl;

        cout << "Select Your Option Number: ";
        cin >> option;

        switch(option){
            cout << endl;
            case 0:
                break;
            case 1:
                cout << "Enqueue your number(>=0): ";

```

```

        cin >> item;
        q1.enqueue(item);
        cout << endl << endl;
        break;
    case 2:
        cout << "Dequeue..." << endl;
        item = q1.dequeue();
        if (item == -1){
            cout << "Queue is empty" << endl << endl;
        }
        else {
            cout << "Dequeued element: " << item << endl << endl;
        }
        break;
    case 3:
        if(q1.isEmpty())
            cout << "Queue is empty" << endl << endl;
        else
            cout << "Queue is not empty" << endl << endl;
        break;
    case 4:
        if(q1.isFull())
            cout << "Queue is full" << endl << endl;
        else
            cout << "Queue is not full" << endl << endl;
        break;
    case 5:
        cout << "Count of queue: " << q1.size() << endl << endl;
        break;
    case 6:
        q1.display();
        break;
    case 7:
        system("cls");
        break;
    default:
        cout << "Please enter the proper option" << endl;
        break;
    }

} while (option != 0);

return 0;
}

```

Non-Linear Data Structure

Tree

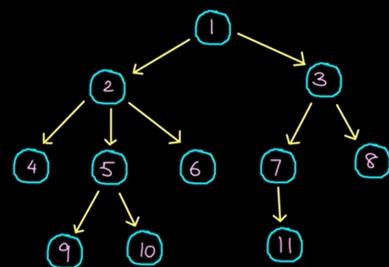
Introduction to TREE data structure

TREE - A tree is a non linear data structure that simulates a hierachial tree structure with a root value & subtrees of children with parent node, represented as set of linked nodes.

Important TREE terms -

1. **Root** - Root is a special node in a tree. The entire tree is referenced through it. It does not have a parent.
2. **Parent Node** - Parent node is an immediate predecessor of a node.
3. **Child Node** - All immediate successors of a node are its children.
4. **Siblings** - Nodes with the same parent are called Siblings.
5. **Leaf** - Last node in the tree. There is no node after this node.
6. **Edge** - Edge is a connection between one node to another. It is a line between two nodes or a node and a leaf.
7. **Path** - Path is a number of successive edges from source node to destination node.
8. **Degree of Node** - Degree of a node is equal to number of children, a node have.

Logical Tree Diagram



SIMPLE SNIPPETS Check video description for CODE & THEORY on our official website

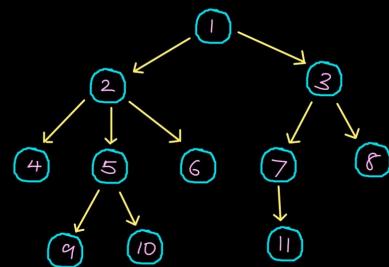
Introduction to TREE data structure

TREE - A tree is a non linear data structure that simulates a hierachial tree structure with a root value & subtrees of children with parent node, represented as set of linked nodes.

Important TREE terms / properties -

- 1) Tree can be termed as a RECURSIVE data structure.
- 2) In a valid tree for N Nodes we have N-1 Edges/Links.
- 3) Depth of Node - Depth of a node represents the number of edges from the tree's root node to the node.
- 4) Height of Node - Height of a node is the number of edges on the longest path between that node & a leaf.
- 5) Height of Tree - Height of tree is the height of its root node.

Logical Tree Diagram



SIMPLE SNIPPETS

Introduction to TREE data structure

TREE - A tree is a non linear data structure that simulates a hierachial tree structure with a root value & subtrees of children with parent node, represented as set of linked nodes.

Types of Trees -

1. **General Tree**
2. **Binary Tree**
3. **Binary Search Tree**
4. **AVL Tree**
5. **Spanning Tree**
6. **B-Tree**
7. **B+ Tree**
8. **Heap**

Applications of Tree Data Structure -

1. Store hierarchical data, like folder structure, organization structure data.
2. Binary Search Tree is a tree that allows fast search, insert, delete on a sorted data.
3. Heap is a tree data structure which is implemented using arrays and used to implement priority queues.
4. B-Tree and B+ Tree are used to implement indexing in databases.
5. Used to store router-tables in routers.
6. Used by compilers to build syntax trees.
7. Used to implement expression parsers and expression solvers.

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Binary Tree

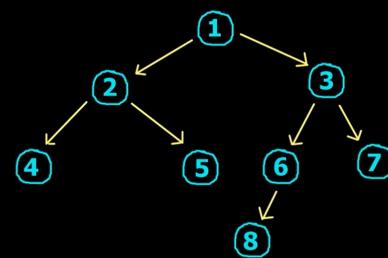
Binary Tree Data Structure

Binary Tree - A binary tree is a tree data structure in which each node has at most two children, which are referred to as the left child (LC) and the right child (RC).

Important Binary Tree Terms & Properties -

- 1) A binary tree is called **STRICT/PROPER** binary tree, when each node has 2 or 0 children.
- 2) A binary tree is called **COMPLETE** binary tree if all levels except the last are completely filled and all nodes are as left as possible.
- 3) A binary tree is called **PERFECT** binary tree if all levels are completely filled with 2 children each.
- 4) Max number of nodes at level ' h ' = 2^h
- 5) For a binary tree, maximum no of nodes with height ' h ' = $2^0 + 2^1 + \dots + 2^h$
 $= 2^{(h+1)} - 1$

Binary Tree Diagram

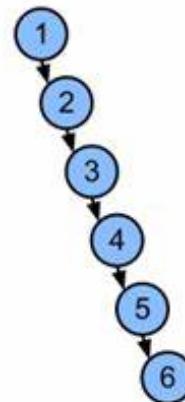


SIMPLE SNIPPETS

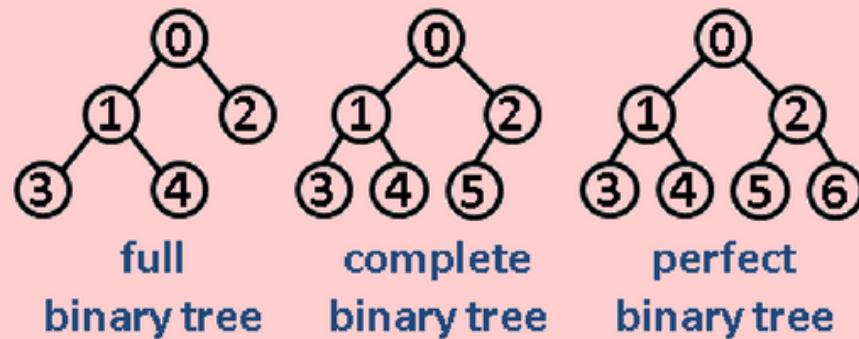
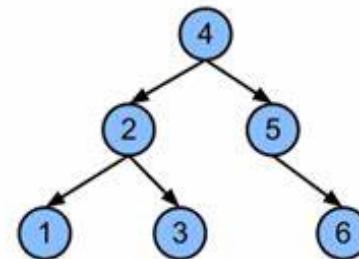


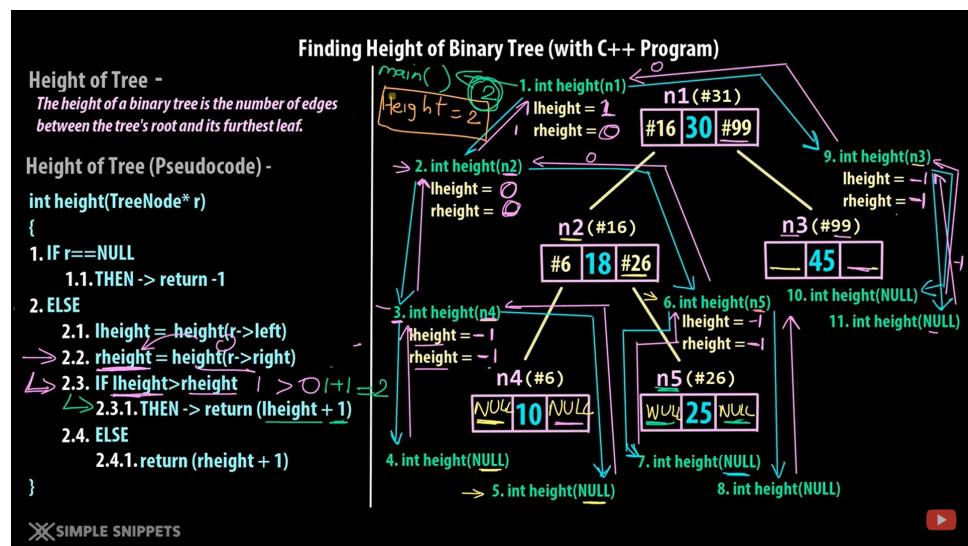
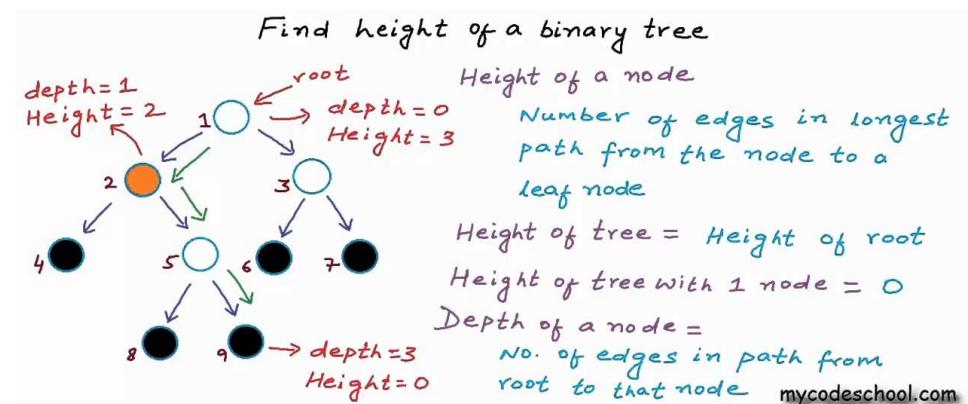
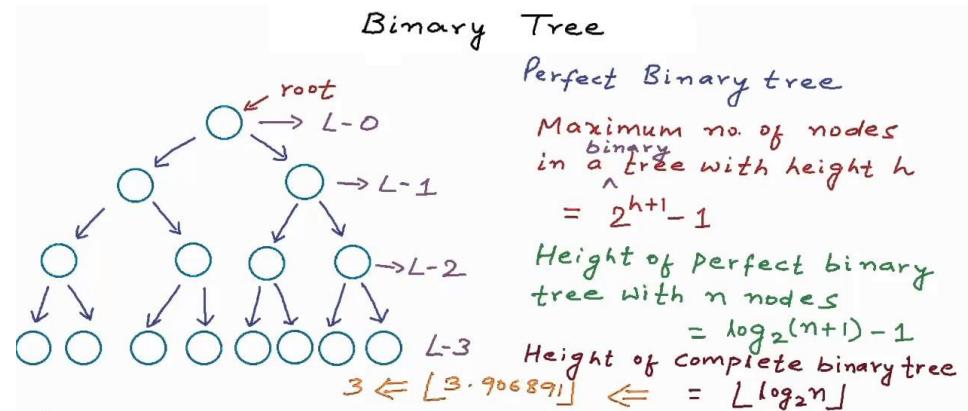
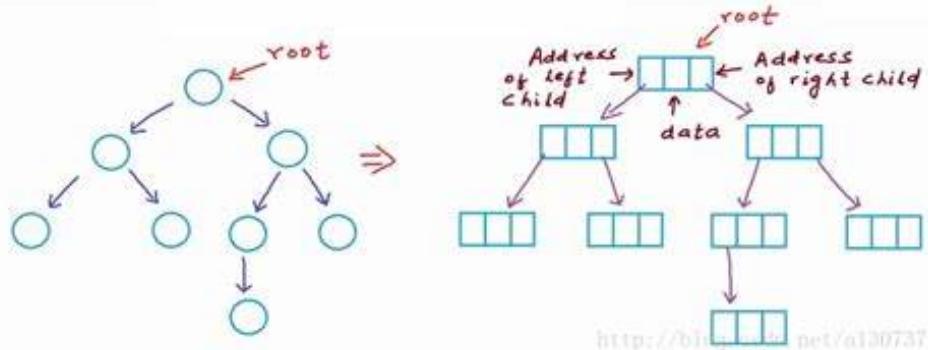
Balanced binary tree

Non-balanced



Balanced





(Each node have 0 or 2 children)

Binary tree & its types

- Full/Proper/strict
- Complete Binary tree
- Perfect Binary tree
- Degenerate Binary tree

(except last level)

all internal nodes have 2 children & all leaves are at same level

Diagram of a binary tree:

Max. nodes	Min. nodes
$2^{h+1} - 1$	$h + 1$
$2^{h+1} - 1$	2^{h+1}
$2^{h+1} - 1$	2^h

Minheight	Max.height
$\lceil \log_2(n+1) \rceil - 1$	$n-1$
$\lceil \log_2(n+1) \rceil - 1$	$\frac{n-1}{2}$
$\lceil \log_2(n+1) \rceil - 1$	$\log n$

BT Implementation

Binary Tree Data Structure

Binary Tree - A binary tree is a tree data structure in which each node has at most two children, which are referred to as the left child (LC) and the right child (RC).

1. Binary Tree as dynamic nodes in memory

Node (`*left, data , *right`)

```

    graph TD
      n1["n1(#10)  
#23 | 5 | #27"] --> n2["n2(#23)  
#73 | 99 | #66"]
      n1 --> n3["n3(#27)  
#13 | 66 | #71"]
      n2 --> n4["n4(#73)  
8 | ]"]
      n2 --> n5["n5(#66)  
23 | ]"]
      n3 --> n6["n6(#13)  
7 | ]"]
      n3 --> n7["n7(#71)  
45 | ]"]
  
```

2. Binary Tree as conventional arrays

arr = [4 9 3 1 5 7 2]

For node at index i
Left child index = $2i+1$ | Right child index = $2i+2$

BT Traversal

Binary Tree Data Structure (Traversal Techniques)

Tree Traversal: Tree traversal (also known as tree search and walking the tree) refers to the process of visiting (checking and/or updating) each node in a tree data structure, exactly once. Such traversals are classified by the order in which the nodes are visited.

Tree Traversal

↓

Depth-first search/traversal (DFS)

(These searches are referred to as depth-first search (DFS), since the search tree is deepened as much as possible on each child before going to the next sibling.)

- 1. Pre-order (NLR)
- 2. In-order (LNR)
- 3. Post-order (LRN)

↓

Breadth-first search/traversal (BFS)

Trees can also be traversed in level-order, where we visit every node on a level before going to a lower level. This search is referred to as breadth-first search (BFS), as the search tree is broadened as much as possible on each depth before going to the next depth.

Binary Tree Data Structure (Traversal Techniques)

→ Depth-first search/Traversal of binary tree (DFS)

1. Pre-order (NLR) (node - left - right)

- Access the data part of the current node.
- Traverse the left subtree by recursively calling the pre-order function.
- Traverse the right subtree by recursively calling the pre-order function.

```

graph TD
    Root[30] --> L18[18]
    Root --> R43[43]
    L18 --> L10[10]
    L18 --> R25[25]
    R43 --> L32[32]
    R43 --> R48[48]
    L10 --- X
    R25 --- X
    L32 --- X
    R48 --- X
  
```

30, 18, 10, 25, 43, 32, 48

2. In-order (LNR) (left - node - right)

- Traverse the left subtree by recursively calling the in-order function.
- Access the data part of the current node.
- Traverse the right subtree by recursively calling the in-order function.

In BST in-order traversal retrieves the keys in ascending sorted order.

```

graph TD
    Root[30] --> L18[18]
    Root --> R43[43]
    L18 --> L10[10]
    L18 --> R25[25]
    R43 --> L32[32]
    R43 --> R48[48]
    L10 --- X
    R25 --- X
    L32 --- X
    R48 --- X
  
```

10, 18, 25, 30, 32, 43, 48

3. Post-order (LRN) (left - right - node)

- Traverse the left subtree by recursively calling the post-order function.
- Traverse the right subtree by recursively calling the post-order function.
- Access the data part of the current node.

```

graph TD
    Root[30] --> L18[18]
    Root --> R43[43]
    L18 --> L10[10]
    L18 --> R25[25]
    R43 --> L32[32]
    R43 --> R48[48]
    L10 --- X
    R25 --- X
    L32 --- X
    R48 --- X
  
```

10, 25, 18, 32, 48, 43, 30

SIMPLE SNIPPETS

Binary Tree Data Structure (Traversal Techniques)

Tree Traversal : Tree traversal (also known as tree search and walking the tree) refers to the process of visiting (checking and/or updating) each node in a tree data structure, exactly once. Such traversals are classified by the order in which the nodes are visited.

DFS -	
1. Pre-order (NLR) - 30, 18, 10, 25, 43, 32, 48	
2. In-order (LNR) - 10, 18, 25, 30, 32, 43, 48	
3. Post-order (LRN) - 10, 25, 18, 32, 48, 43, 30	
BFS -	30, 18, 43, 10, 25, 32, 48

```

graph TD
    Root[30] --> L18[18]
    Root --> R43[43]
    L18 --> L10[10]
    L18 --> R25[25]
    R43 --> L32[32]
    R43 --> R48[48]
    L10 --- X
    R25 --- X
    L32 --- X
    R48 --- X
  
```

SIMPLE SNIPPETS

(BFS) Level Order Tree Traversal Technique(Working + C++ Program)

Print Level Order (Pseudocode) -

```

void printLevelOrderBFS(TreeNode* r)
{
    1. h = height(r) // calculate height of tree
    2. FOR i=0 to i<=h
        2.1.printGivenLevel(r,i)
}
  
```

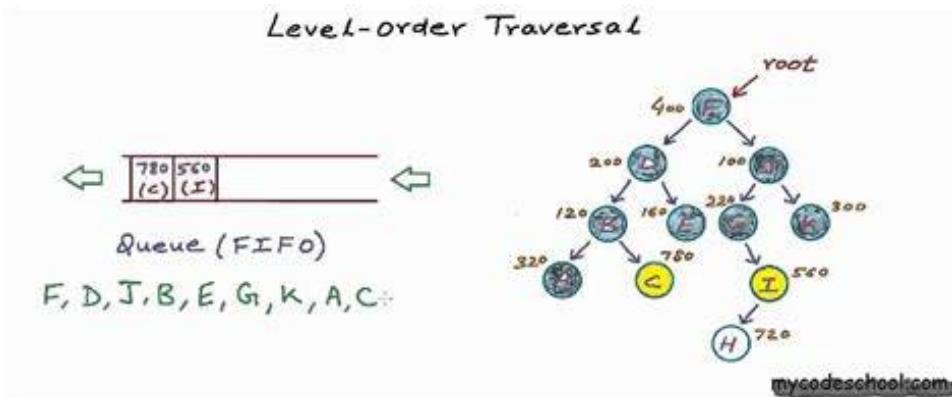
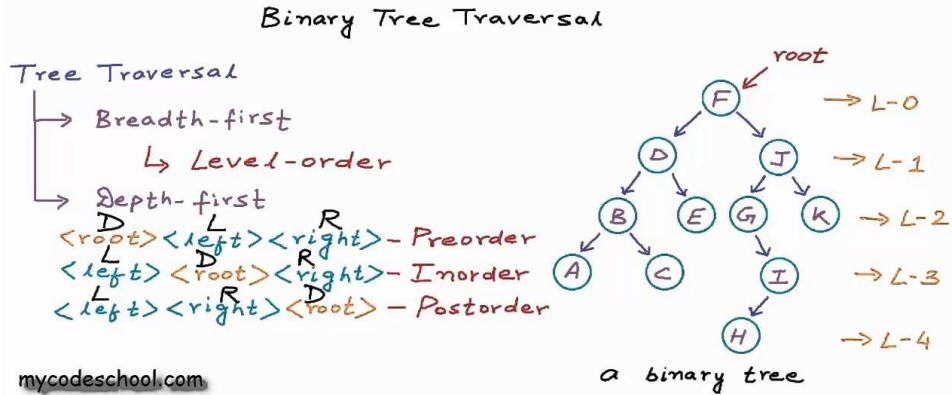
Print Given Level (Pseudocode) -

```

void printGivenLevel(TreeNode* r, int level)
{
    1. IF r==NULL
        1.1.THEN > return
    2. ELSE IF level==0
        2.1.PRINT (r->value)
    3. ELSE
        3.1. printGivenLevel(r->left, level-1)
        3.2. printGivenLevel(r->right, level-1)
}
  
```

Output - 30 18 43 10 25
n1(#31) #16 30 #99
n2(#16) #6 18 #26
n3(#99) NULL 45 NULL
n4(#6) 10
n5(#26) 25

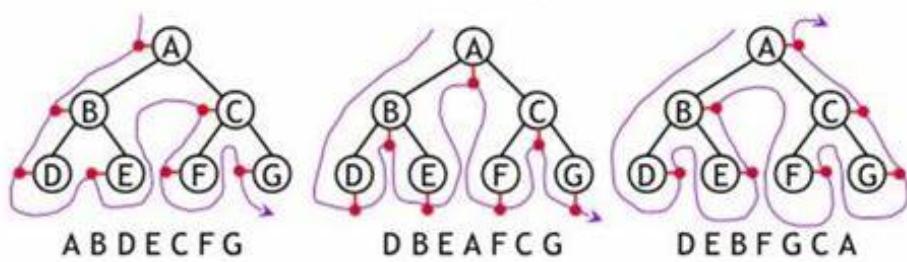
SIMPLE SNIPPETS



The order in which the nodes are visited during a tree traversal can be easily determined by imagining there is a “flag” attached to each node, as follows:



To traverse the tree, collect the flags:



BST vs BT

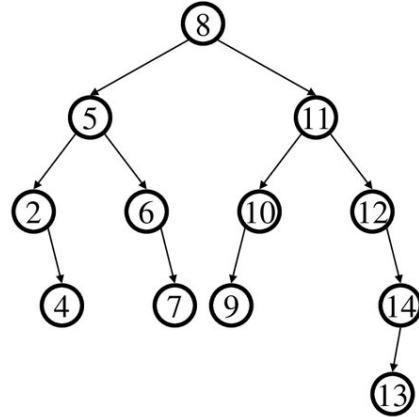
Binary Search Tree Dictionary Data Structure

Binary tree property

- ▶ each node has ≤ 2 children
- ▶ result:
 - ▶ storage is small
 - ▶ operations are simple
 - ▶ average depth is small

Search tree property

- ▶ all keys in left subtree smaller than root's key
- ▶ all keys in right subtree larger than root's key
- ▶ result:
 - ▶ easy to find any given key
 - ▶ Insert/delete by changing links



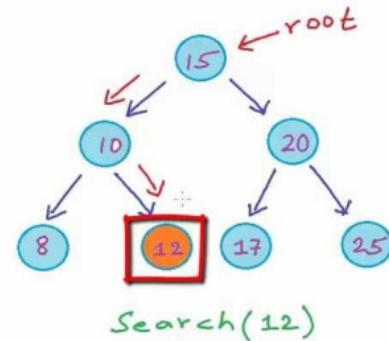
Binary Search Trees

Binary Search Tree/BST

Binary Search Tree

Binary Search Tree(BST)

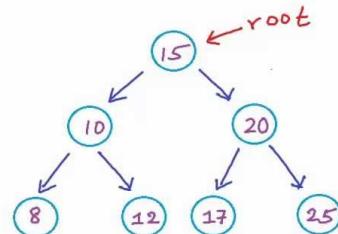
↳ a binary tree in which for each node, value of all the nodes in left subtree is lesser or equal and value of all the nodes in right subtree is greater.



```

struct BstNode {
    int data;
    BstNode* left;
    BstNode* right;
};

int FindMin(BstNode* root) {
    if(root == NULL) {
        cout<<"Error: Tree is empty\n";
        return -1;
    }
    else if(root->left == NULL) {
        return root->data;
    }
    // Search in left subtree.
    return FindMin(root->left);
}
  
```



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Binary Search Tree - Implementation in C/C++

```
BstNode* Insert(BstNode* root, int data) {
    if(root == NULL) { // empty tree
        root = GetNewNode(data);
    }
    else if(data <= root->data) {
        root->left = Insert(root->left, data);
    }
    else {
        root->right = Insert(root->right, data);
    }
    return root;
}
```

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BST implementation in C/C++ - Part II

```
BstNode* Insert(BstNode* root, int data) {
    if(root == NULL) { // empty tree
        root = GetNewNode(data);
    }
    else if(data <= root->data) {
        root->left = Insert(root->left, data);
    }
    else {
        root->right = Insert(root->right, data);
    }
    return root;
}
```

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Binary SEARCH Tree Data Structure

Binary search tree - *BST is a binary tree data structure, in which the values in the left subtrees of every node are smaller and the values in the right subtrees of every node are larger.*

<p>$\downarrow \quad x = 65$</p> <p>Array</p> <p>$x = 19$</p> <p>$arr = [25 23 14 36 29 19 52]$</p> <p>$x = 19$</p> <p>$sorted_arr = [14 19 23 25 29 36 52]$</p> <p>$\downarrow \quad x = 52$</p>	<p>Values/Keys = {25, 23, 14, 36, 29, 19, 52}</p>	<p style="text-align: center;">→ Time Complexity (Big O)</p> <table border="1" style="width: 100%; border-collapse: collapse;"> <thead> <tr> <th>Search/access/update</th> <th>Insert /Add/</th> <th>Delete</th> </tr> </thead> <tbody> <tr> <td>$O(N)$ linear search</td> <td>$O(N)$ Start</td> <td>$O(N)$ Random</td> </tr> <tr> <td>$O(1)$ random access</td> <td>$O(1)$ End</td> <td>$O(1)$ Start</td> </tr> <tr> <td>$O(\log n)$ binary search</td> <td>$O(N)$ Random</td> <td>$O(N)$ End</td> </tr> </tbody> </table> <p>LinkedList</p> <p>$x = 52$</p> <p>$n1 \rightarrow [25 45]$ $n2 \rightarrow [23 67]$ $n3 \rightarrow [14 21]$ $n4 \rightarrow [36 35]$ $n5 \rightarrow [29 41]$ $n6 \rightarrow [19 99]$ $n7 \rightarrow [52 null]$</p> <p>$\rightarrow (\#45)$ $\rightarrow (\#67)$ $\rightarrow (\#21)$ $\rightarrow (\#35)$ $\rightarrow (\#41)$ $\rightarrow (\#99)$</p>	Search/access/update	Insert /Add/	Delete	$O(N)$ linear search	$O(N)$ Start	$O(N)$ Random	$O(1)$ random access	$O(1)$ End	$O(1)$ Start	$O(\log n)$ binary search	$O(N)$ Random	$O(N)$ End
Search/access/update	Insert /Add/	Delete												
$O(N)$ linear search	$O(N)$ Start	$O(N)$ Random												
$O(1)$ random access	$O(1)$ End	$O(1)$ Start												
$O(\log n)$ binary search	$O(N)$ Random	$O(N)$ End												

SIMPLE SNIPPETS

Insert Node - Recursive

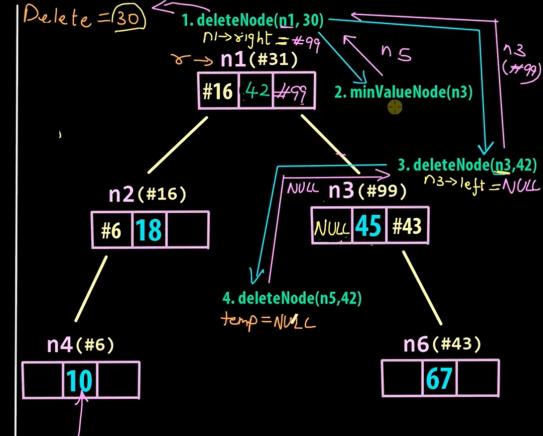
```

Delete Node(Pseudocode) -
TreeNode* deleteNode(TreeNode*r, int v) {
    1. IF r==NULL THEN -> // Base condition
        1.1 return r
    2. ELSE IF v < r->value THEN -> // if value smaller go left sub tree
        2.1 r->left = deleteNode(r->left, v)
    3. ELSE IF v > r->value THEN -> // if value larger go right sub tree
        3.1 r->right = deleteNode(r->right, v)
    4. ELSE // if value matches
        4.1 IF r->left==NULL THEN -> // node with only right child
            4.1.1 temp = r->right //OR no child
            4.1.2 delete r
            4.1.3 return temp
        4.2 ELSE IF r->right==NULL THEN -> // node with only left child
            4.2.1 temp = r->left
            4.2.2 delete r
            4.2.3 return temp
        4.3 ELSE // node with TWO children
            4.3.1 temp = minValueNode(r->right)
            4.3.2 r->value = temp->value
            4.3.3 r->right = deleteNode(r->right, temp->value)
    5. return r
}

```

SIMPLE SNIPPETS

Delete Node Operation in Binary Search Tree (with C++ Code)



Binary SEARCH Tree Data Structure - INSERTION (RECURSIVE)

BST INSERT(Recursive) PSEUDOCODE -

```

TreeNode* insertRecursive(TreeNode*r, TreeNode *new_node)
{
    1. IF r == NULL THEN ->
        1.1. r=new_node # 77
        1.2. (return r)
    2. IF new_node->value IS LESS THAN r->value THEN ->
        2.1. r->left = insertRecursive(r->left, new_node)
    3. ELSE IF new_node->value IS GREATER THAN r->value THEN ->
        3.1. r->right = insertRecursive(r->right, new_node)
    4. ELSE
        4.1. PRINT ("NO DUPLICATES ALLOWED")
        4.2. return r
    5. return r
}

```

Main Function -

```

BST obj; // create object of BST class
TreeNode *new_node = new TreeNode(); // create node in HEAP memory
new_node->value = val; // take from user
obj.root= obj.insertRecursive(obj.root,new_node); // insert node
LHS = RHS

```

left [] pointer value [] right pointer

→ 1. obj.root= obj.IR() new_node; → → null

56

RAM STACK HEAP

#77	Y	
*new_node		#77
obj*root	NULL	(#77)
		56
*new_node		
#77		



Binary SEARCH Tree Data Structure - INSERTION (RECURSIVE)

BST INSERT(Recursive) PSEUDOCODE -

```

TreeNode* insertRecursive(TreeNode*r, TreeNode *new_node)
{
    1. IF r == NULL THEN ->
        1.1. r=new_node #90
        1.2. (return r)
    2. IF new_node->value IS LESS THAN r->value THEN ->
        2.1. r->left = insertRecursive(r->left, new_node)
    3. ELSE IF new_node->value IS GREATER THAN r->value THEN ->
        3.1. r->right = insertRecursive(r->right, new_node)
    4. ELSE
        4.1. PRINT ("NO DUPLICATES ALLOWED")
        4.2. return r
    5. return r
}

```

Main Function -

```

BST obj; // create object of BST class
TreeNode *new_node = new TreeNode(); // create node in HEAP memory
new_node->value = val; // take from user
obj.root= obj.insertRecursive(obj.root,new_node); // insert node

```

left [] pointer value [] right pointer

① 2.1. 1. obj.root= obj.IR() new_node; Y = #77

② 2. r->left = IR(r->left, new_node); #90 n2 (#90)

RAM STACK HEAP

#90	Y	
*new_node		#77
obj*root	#77	(#77)
		56
*new_node		
#90		
		35



Binary SEARCH Tree Data Structure - INSERTION (RECURSIVE)

BST INSERT(Recursive) PSEUDOCODE -

```

TreeNode* insertRecursive(TreeNode* r, TreeNode *new_node)
{
    1. IF r == NULL THEN -->
        ↳ 1.1. r=new_node
        ↳ 1.2. return r
    2. IF new_node->value IS LESS THAN r->value THEN -->
        ↳ 2.1. r->left = insertRecursive(r->left, new_node)
    3. ELSE IF new_node->value IS GREATER THAN r->value THEN -->
        ↳ 3.1. r->right = insertRecursive(r->right, new_node)
    4. ELSE
        4.1. PRINT ("NO DUPLICATES ALLOWED")
        4.2. return r
    }
}

```

Main Function -

```

BST obj;           // create object of BST class
TreeNode *new_node = new TreeNode(); // create node in HEAP memory
new_node->value = val; // take from user
obj.root=obj.insertRecursive(obj.root,new_node); // insert node

```

left pointer value pointer

RAM **STACK** **HEAP**

obj *root	(#77)	n1
	#77	n2
		n3 (#12)
		n4 (#12)
		n5
		n6 (#12)
		n7
		40

SIMPLE SNIPPETS

Delete Node

Delete Node(Pseudocode) -

```

TreeNode* deleteNode(TreeNode* r, int v) {
    1. IF r==NULL THEN -> // Base condition
        1.1 return r
    2. ELSE IF v < r->value THEN -> // if value smaller go left sub tree
        2.1 r->left = deleteNode(r->left, v)
    3. ELSE IF v > r->value THEN -> // if value larger go right sub tree
        3.1 r->right = deleteNode(r->right, v)
    4. ELSE // if value matches
        4.1 IF r->left==NULL THEN -> // node with only right child
            4.1.1 temp = r->right      //OR no child
            4.1.2 delete r
            4.1.3 return temp
        4.2 ELSE IF r->right==NULL THEN -> // node with only left child
            4.2.1 temp = r->left
            4.2.2 delete r
            4.2.3 return temp
        4.3 ELSE // node with two children
            4.3.1 temp = minValueNode(r->right)
            4.3.2 r->value = temp->value
            4.3.3 r->right = deleteNode(r->right, temp->value)
    }
    5. return r
}

```

Delete Node Operation in Binary Search Tree (with C++ Code)

SIMPLE SNIPPETS

AVL Tree

AVL tree is a self-balancing Binary Search Tree (BST) where the difference between heights of left & right subtrees cannot be more than one for all nodes. This difference is called the Balance Factor(BF). BF={-1,0,1}

Binary Tree

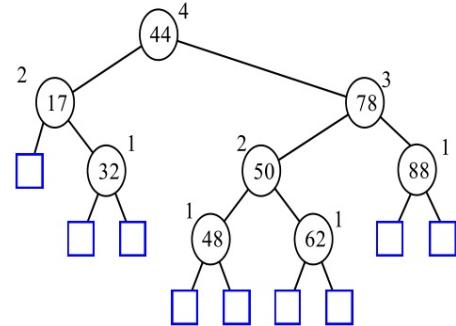
Binary SEARCH Tree

AVL Tree

SIMPLE SNIPPETS

AVL Tree Definition

- AVL trees are balanced.
- An AVL Tree is a **binary search tree** such that for every internal node v of T , the *heights of the children of v can differ by at most 1*.



An example of an AVL tree where the heights are shown next to the nodes:

The AVL Tree Data Structure

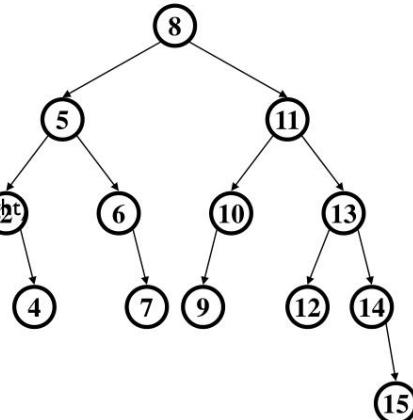
An AVL tree is a BST

In addition: Balance property:
balance of every node is
between -1 and 1

balance(node) = $\text{height}(\text{node.left}) - \text{height}(\text{node.right})$

Result: Worst-case depth is $O(\log n)$

How are we going to maintain this? Worry
about that later...



Q: Why does this work? Why are you allowed to change the structure of a tree like this?

Short answer: In the left tree we know that $D < B < E < A < C$ and this remains true for the right subtree, so we didn't break the BST invariant and, therefore, this is a valid transformation.

To maintain the balance in AVL tree, re-balancing is done using ROTATIONS. Lets Insert - 10,20,30 in AVL Tree

1. Order of Insertion -30, 20, 10 BST	AVL	3. Order of Insertion -30, 10, 20 L-R Imbalance LEFT \rightarrow R-L Rotation
2. Order of Insertion -10, 20, 30 BST	AVL	4. Order of Insertion -10, 30, 20 R-L Imbalance RIGHT \rightarrow L-R Rotation

SIMPLE SNIPPETS

To maintain the balance in AVL tree, re-balancing is done using ROTATIONS

LEFT LEFT Case/imbalance T1, T2, T3 and T4 are subtrees. Right(z) Rotation	LEFT RIGHT case/imbalance Left(y) Rotation \rightarrow Right(z) Rotation T1, T2, T3 and T4 are subtrees.
 Right(z) Rotation T1, T2, T3 and T4 are subtrees.	RIGHT LEFT case/imbalance Right(y) Rotation \rightarrow Left(z) Rotation T1, T2, T3 and T4 are subtrees."/> T1, T2, T3 and T4 are subtrees.

SIMPLE SNIPPETS

Binary Search Tree - Insertion (Recursive)

```
TreeNode* insertRecursive(TreeNode *r, TreeNode *new_node)
{
    1. IF r == NULL THEN -->
        1.1. r=>new_node
        1.2. return r
    2. IF new_node->value IS LESS THAN r->value THEN -->
        2.1. r->left = insertRecursive(r->left, new_node)
    3. ELSE IF new_node->value IS GREATER THAN r->value THEN -->
        3.1. r->right = insertRecursive(r->right, new_node)
    4. ELSE
        4.1. PRINT ("NO DUPLICATES ALLOWED")
        4.2. return r
    5. return r
}
```

AVL - Insertion (Recursive)

```
TreeNode* insertRecursive(TreeNode *r, TreeNode *new_node)
1. IF r == NULL THEN -->
    1.1. r=>new_node
    1.2. return r
2. IF new_node->value LESS THAN r->value THEN -->
    2.1. r->left = insert(r->left,new_node)
3. ELSE IF new_node->value GREATER THAN r->value THEN -->
    3.2. r->right = insert(r->right,new_node)
4. ELSE
    4.1. PRINT("No duplicate values")
    4.2. return r
5. bf = getBalanceFactor(r)
6. IF bf > 1 AND new_node->value < r->left->value
    6.1. return rightRotate(r)
7. IF bf < -1 && new_node->value > r->right->value
    7.1. return leftRotate(r)
8. IF bf > 1 && new_node->value > r->left->value
    8.1. r->left = leftRotate(r->left)
    8.2. return rightRotate(r)
9. IF bf < -1 && new_node->value < r->right->value
    9.1. r->right = rightRotate(r->right);
    9.2. return leftRotate(r)
10. return r
```

SIMPLE SNIPPETS

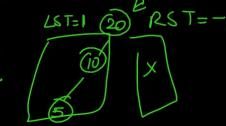
**Binary Search Tree Delete Node (Recursive) -**

```
TreeNode* deleteNode(TreeNode* r, int v) {
    1. IF r==NULL THEN --> // Base condition
        1.1 return r
    2. ELSE IF v < r->value THEN --> // if value smaller go left sub tree
        2.1 r->left = deleteNode(r->left, v)
    3. ELSE IF v > r->value THEN --> // if value larger go right sub tree
        3.1 r->right = deleteNode(r->right, v)
    4. ELSE // if value matches
        4.1 IF r->left==NULL THEN --> // node with only right child
            4.1.1 temp = r->right //OR no child
            4.1.2 delete r
            4.1.3 return temp
        4.2 ELSE IF r->right==NULL THEN --> // node with only left child
            4.2.1 temp = r->left
            4.2.2 delete r
            4.2.3 return temp
        4.3 ELSE // node with TWO children
            4.3.1 temp = minValueNode(r->right)
            4.3.2 r->value = temp->value
            4.3.3 r->right = deleteNode(r->right, temp->value)
    } 5. return r
}
```

SIMPLE SNIPPETS

AVL Tree Delete Node (Recursive) -

```
TreeNode* deleteNode(TreeNode* r, int v) {
    //First 4 Steps same as BST DELETION
    2. 
    5. bf = getBalanceFactor(r) // 2
    / 6. IF bf == 2 AND getBalanceFactor(r->left) >= 0
    / 6.1. return rightRotate(r)
    7. ELSE IF bf == 2 AND getBalanceFactor(r->right) == -1
    7.1. r->left = leftRotate(r->left)
    7.2. return rightRotate(r)
    8. ELSE IF bf == -2 AND getBalanceFactor(r->right) <= -1
    8.1. return leftRotate(r)
    9. ELSE IF bf == -2 && getBalanceFactor(r->right) == 1
    9.1. r->right = rightRotate(r->right);
    9.2. return leftRotate(r)
    10. return r
}
```



```
In [ ]: #include <iostream>
#define SPACE 5
using namespace std;

class TreeNode{
public:
    int value;
    TreeNode* left;
    TreeNode* right;

    TreeNode();
    TreeNode(int v);
};

TreeNode::TreeNode() : value(0), left(nullptr), right(nullptr){}

TreeNode::TreeNode(int v){
    value = v;
    left = nullptr;
    right = nullptr;
}
```

```
TreeNode::TreeNode() : value(0), left(nullptr), right(nullptr){}
```

```
TreeNode::TreeNode(int v){
    value = v;
    left = nullptr;
    right = nullptr;
}
```

```
class BST{
public:
    TreeNode* root;
```

```

BST();

bool isEmpty();
void insertNode(TreeNode* newNode, int v);
TreeNode* searchNode(int v);
int height(TreeNode* root);
TreeNode* minValueNode(TreeNode* root);
TreeNode* deleteNode(TreeNode* root, int v);
void print2D(TreeNode* root, int space);
void printPreorder(TreeNode* root);
void printInorder(TreeNode* root);
void printPostorder(TreeNode* root);
void printLevelOrder(TreeNode* root);
void printGivenLevel(TreeNode* root, int level);
};

BST::BST() : root(nullptr){}

bool BST::isEmpty(){
    if(root==nullptr)
        return true;
    else
        return false;
}

void BST::insertNode(TreeNode* newNode, int v){
    newNode->value = v;
    if(root==nullptr){
        root = newNode;
        cout << "Value Inserted as Root Node!" << endl;
    }
    else {
        TreeNode* ptr = root;
        while(ptr!=nullptr){
            if(newNode->value == ptr->value){
                cout << "Value Already exist. Insert Another Value." << endl;
                return;
            }
            else if((newNode->value < ptr->value) && (ptr->left==nullptr)){
                ptr->left = newNode;
                cout << "Value Inserted in the left" << endl;
                break;
            }
            else if(newNode->value < ptr->value){
                ptr = ptr->left;
            }
            else if((newNode->value > ptr->value) && (ptr->right==nullptr)){
                ptr->right = newNode;
                cout << "Value Inserted in the right" << endl;
                break;
            }
            else {
                ptr = ptr->right;
            }
        }
    }
}

TreeNode* BST::searchNode(int v){

```

```

if(root==nullptr)
    return root;
TreeNode* ptr = root;
while (ptr!=nullptr){
    if(ptr->value == v)
        return ptr;
    else if(ptr->value > v){
        ptr = ptr->left;
    }
    else {
        ptr = ptr->right;
    }
}
return nullptr;
}

int BST::height(TreeNode* root){
    if(root==nullptr)
        return -1;
    else {
        int leftHeight = height(root->left);
        int rightHeight = height(root->right);

        if(leftHeight > rightHeight)
            return (leftHeight + 1);
        else
            return (rightHeight + 1);
    }
}

TreeNode* BST::minValueNode(TreeNode* root){
    TreeNode* ptr = root;
    while(ptr->left!=nullptr){
        ptr = ptr->left;
    }
    return ptr;
}

TreeNode* BST::deleteNode(TreeNode* root, int v){

    if(root==nullptr)
        return root;
    TreeNode* ptr = root;
    if(ptr->value > v){
        ptr->left = deleteNode(ptr->left, v);
    }
    else if(ptr->value < v){
        ptr->right = deleteNode(ptr->right, v);
    }
    else{ // find the node with v
        if(ptr->left==nullptr){ // maybe have right child or no right child when pt
            TreeNode* temp = ptr->right;
            delete ptr;
            return temp;
        }
        else if(ptr->right==nullptr){ // only left child
            TreeNode* temp = ptr->left;
            delete ptr;
            return temp;
        }
        else { // have two children
    }
}

```

```

        TreeNode* temp = minValueNode(ptr->right);
        ptr->value = temp->value;
        ptr->right = deleteNode(ptr->right, temp->value);
    }
}
return ptr;
}

void BST::print2D(TreeNode* root, int space){
    if(root == nullptr)
        return;
    space += SPACE; // increase distance between level
    print2D(root->right, space);
    cout << endl;
    for(int i=SPACE; i<space; i++){
        cout << " ";
    }
    cout << root->value << endl;
    print2D(root->left, space);
}

void BST::printPreorder(TreeNode* root){
    if(root==nullptr){
        return;
    }
    cout << root->value << " ";
    printPreorder(root->left);
    printPreorder(root->right);
}

void BST::printInorder(TreeNode* root){
    if(root==nullptr){
        return;
    }
    printInorder(root->left);
    cout << root->value << " ";
    printInorder(root->right);
}

void BST::printPostorder(TreeNode* root){
    if(root==nullptr){
        return;
    }
    printPostorder(root->left);
    printPostorder(root->right);
    cout << root->value << " ";
}

void BST::printGivenLevel(TreeNode* root, int level){
    if(root==nullptr){
        return;
    }
    else if (level==0)
        cout << root->value << " ";
    else {
        printGivenLevel(root->left, level - 1);
        printGivenLevel(root->right, level - 1);
    }
}

void BST::printLevelOrder(TreeNode* root){

```

```

        int h = height(root);
        for(int i=0; i<=h; i++){
            printGivenLevel(root, i);
        }
    }

class AVLTree : public BST {
public:

    AVLTree();

    int getBalanceFactor(TreeNode* ptrNode);
    TreeNode* insertRecursive(TreeNode* root, TreeNode* newNode, int v);
    TreeNode* rotateRecursive(TreeNode* root);
    TreeNode* insertAndRotate(TreeNode* root, TreeNode* newNode, int v);
    TreeNode* deleteAndRotate(TreeNode* root, int v);
    TreeNode* rightRotate(TreeNode* ptrNode);
    TreeNode* leftRotate(TreeNode* ptrNode);
    TreeNode* recursiveSearch(TreeNode* root, int v);
};

AVLTree::AVLTree(){
    root = nullptr;
}

TreeNode* AVLTree::insertRecursive(TreeNode* root, TreeNode* newNode, int v){
    newNode->value = v;
    if(root==nullptr){
        root = newNode;
        cout << "Value inserted successfully" << endl;
        return root;
    }
    if(newNode->value < root->value){
        root->left = insertAndRotate(root->left, newNode, v);
    }
    else if(newNode->value > root->value){
        root->right = insertAndRotate(root->right, newNode, v);
    }
    else {
        cout << "No duplicate value allowed" << endl;
        return root;
    }
    return root;
}

int AVLTree::getBalanceFactor(TreeNode* root){
    if (root==nullptr)
        return -1;
    return (height(root->left) - height(root->right));
}

TreeNode* AVLTree::rightRotate(TreeNode* root){
    TreeNode* ptrNode = root->left;
    ptrNode->right = root;
    root->left = nullptr;

    return ptrNode;
}

TreeNode* AVLTree::leftRotate(TreeNode* root){
}

```

```

TreeNode* ptrNode = root->right;
ptrNode->left = root;
root->right = nullptr;
return ptrNode;
}

/*
TreeNode* AVLTree::insertAndRotate(TreeNode* root, TreeNode* newNode, int v){
    newNode->value = v;
    if(root==nullptr){
        root = newNode;
        cout << "Value inserted successfully" << endl;
        return root;
    }
    if(newNode->value < root->value){
        root->left = insertAndRotate(root->left, newNode, v);
    }
    else if(newNode->value > root->value){
        root->right = insertAndRotate(root->right, newNode, v);
    }
    else {
        cout << "No duplicate value allowed" << endl;
        return root;
    }

    // Transform to AVL Tree
    int balanceFactor = getBalanceFactor(root);

    if((balanceFactor>1) && (newNode->value < root->left->value))
        return rightRotate(root);

    if((balanceFactor>1) && (newNode->value > root->left->value)){
        root->left = leftRotate(root->left);
        return rightRotate(root);
    }

    if((balanceFactor<-1) && (newNode->value > root->right->value))
        return leftRotate(root);

    if((balanceFactor<-1) && (newNode->value < root->right->value)){
        root->right = rightRotate(root->right);
        return leftRotate(root);
    }
    return root;
}
*/
}

TreeNode* AVLTree::rotateRecursive(TreeNode* root){
    int balanceFactor = getBalanceFactor(root);

    if((balanceFactor>1) && (getBalanceFactor(root->left)>=0))
        return rightRotate(root);

    if((balanceFactor>1) && (getBalanceFactor(root->left)==-1)){
        root->left = leftRotate(root->left);
        return rightRotate(root);
    }

    if((balanceFactor<-1) && (getBalanceFactor(root->right)<=0))
        return leftRotate(root);
}

```

```

if((balanceFactor<-1) && (getBalanceFactor(root->right)==1)){
    root->right = rightRotate(root->right);
    return leftRotate(root);
}
return root;
}

TreeNode* AVLTree::insertAndRotate(TreeNode* root, TreeNode* newNode, int v){

    root = insertRecursive(root, newNode, v);
    root = rotateRecursive(root);

    return root;
}

TreeNode* AVLTree::deleteAndRotate(TreeNode* root, int v){

    root = deleteNode(root, v);
    root = rotateRecursive(root);

    return root;
}

TreeNode* AVLTree::recusiveSearch(TreeNode* root, int v){
    if(root==nullptr || root->value==v)
        return root;
    else if(v < root->value)
        return recusiveSearch(root->left, v);
    else
        return recusiveSearch(root->right, v);
}

int main(){
    int option, value, height;
    int space = 2;
    BST obj;
    AVLTree balanceObj;

    do{
        cout << "Welcome to learning Data Structure and Algorithm." << endl << endl;

        cout << "Select Option Number or Enter 0 to exit." << endl;
        cout << "1. Insert Node" << endl;
        cout << "2. Search Node" << endl;
        cout << "3. Delete Node" << endl;
        cout << "4. Height of Tree" << endl;
        cout << "5. Print Values in 2D" << endl;
        cout << "6. Insert Node to AVL Tree" << endl;
        cout << "7. Delete Node in AVL Tree" << endl;
        cout << "8. Search Node in AVL Tree" << endl;
        cout << "9. Print AVL Tree" << endl;
        cout << "10. Clear Screen" << endl;
        cout << "0. Exit Program" << endl;
        cout << endl;

        cin >> option;

        TreeNode* newNode = new TreeNode(); // in heap memory

        switch(option){

```

```

case 0:
    break;
case 1:
    cout << "Insert" << endl;
    cout << "Enter Value of Tree Node to Insert in BST: ";
    cin >> value;
    obj.insertNode(newNode, value);
    cout << endl << endl;
    break;
case 2:
    cout << "Search" << endl;
    cout << "Enter Value of Tree Node to Search: ";
    cin >> value;
    newNode = obj.searchNode(value);
    if(newNode!=nullptr){
        cout << "Value " << value << " Found" << endl;
    }
    else {
        cout << "Value " << value << " Not Found" << endl;
    }
    break;
case 3:
    cout << "Delete" << endl;
    cout << "Enter Value of Tree Node to Delete: ";
    cin >> value;
    if(obj.searchNode(value)!=nullptr){
        obj.deleteNode(obj.root, value);
        cout << "Value Deleted" << endl << endl;
    }

    else {
        cout << "Value " << value << " not Found" << endl << endl;
    }

    break;
case 4:
    cout << "Tree Height" << endl;
    height = obj.height(obj.root);
    cout << "Height: " << height << endl;
    cout << endl << endl;
    break;
case 5:
    cout << "2D Print" << endl;
    obj.print2D(obj.root, space);
    cout << endl << endl;
    cout << "Preorder: " << endl;
    obj.printPreorder(obj.root);
    cout << endl << endl;
    cout << "Inorder: " << endl;
    obj.printInorder(obj.root);
    cout << endl << endl;
    cout << "Postorder: " << endl;
    obj.printPostorder(obj.root);
    cout << endl << endl;
    cout << "Levelorder: " << endl;
    obj.printLevelOrder(obj.root);
    cout << endl << endl;
    break;
case 6:
    cout << "Insert Node to AVL Tree" << endl;
    cout << "Enter Value: ";

```

```
        cin >> value;
        balanceObj.root = balanceObj.insertAndRotate(balanceObj.root, newNode);
        cout << endl << endl;
        break;
    case 7:
        cout << "Delete Node in AVL Tree" << endl;
        cout << "Enter Value: ";
        cin >> value;
        balanceObj.root = balanceObj.deleteAndRotate(balanceObj.root, value);
        cout << endl << endl;
        break;
    case 8:
        cout << "Search Node in AVL Tree" << endl;
        cout << "Enter Value to Search: ";
        cin >> value;
        newNode = balanceObj.searchNode(value);
        if(newNode!=nullptr){
            cout << "Value " << value << " Found" << endl;
        }
        else {
            cout << "Value " << value << " Not Found" << endl;
        }
        break;
    case 9:
        cout << "Print AVL Tree" << endl;
        balanceObj.print2D(balanceObj.root, space);
        cout << endl << endl;
        break;
    case 10:
        system("cls");
        break;
    default:
        cout << "Please Enter the Proper Option." << endl;
        break;
    }
} while (option !=0);

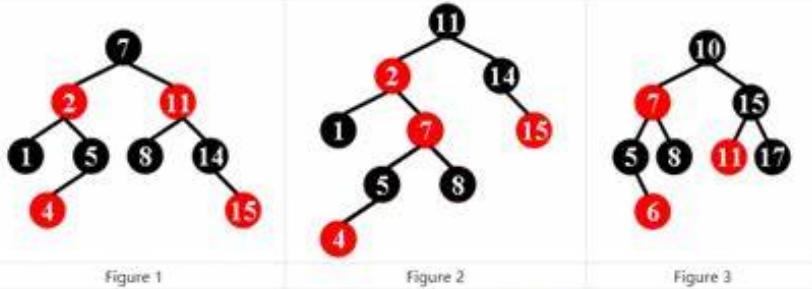
return 0;
}
```

Red Black Tree

There is a kind of balanced binary search tree named **red-black tree** in the data structure. It has the following 5 properties:

- (1) Every node is either red or black.
- (2) The root is black.
- (3) Every leaf (NULL) is black.
- (4) If a node is red, then both its children are black.
- (5) For each node, all simple paths from the node to descendant leaves contain the same number of black nodes.

For example, the tree in Figure 1 is a red-black tree, while the ones in Figure 2 and 3 are not.



For each given binary search tree, you are supposed to tell if it is a legal red-black tree.

Input Specification:

Each input file contains several test cases. The first line gives a positive integer K (≤ 30) which is the total number of cases. For each case, the first line gives a positive integer N (≤ 30), the total number of nodes in the binary tree. The second line gives the preorder traversal sequence of the tree. While all the keys in a tree are positive integers, we use negative signs to represent red nodes. All the numbers in a line are separated by a space. The sample input cases correspond to the trees shown in Figure 1, 2 and 3.

JENNY'S LECTURES

Insertion in Red-Black Tree

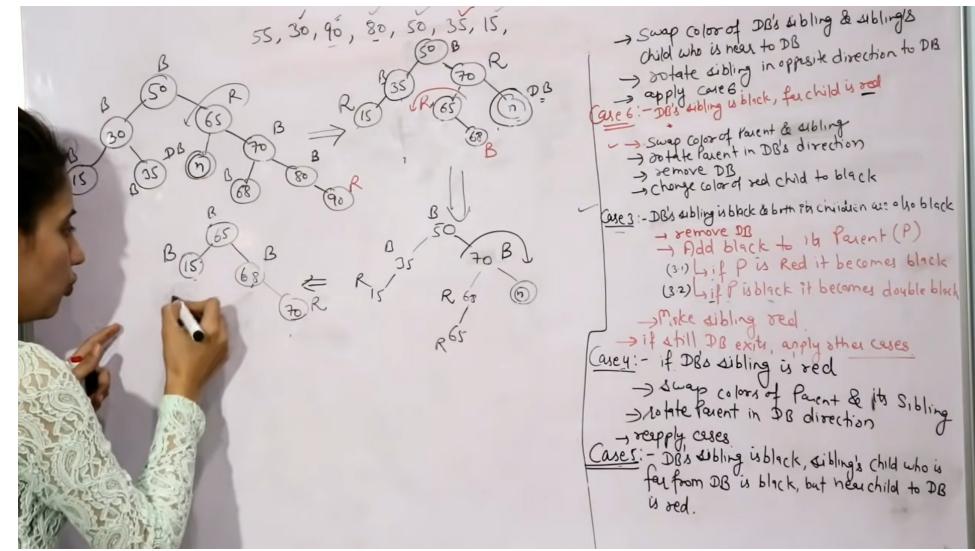
10, 18, 7, 15, 16, 30, 25, 40, 60, 2, 1, 70

① If tree is empty, create newnode as root node with color black
 ② If tree is not empty, create newnode as leaf node with color Red.
 ③ If parent of newnode is black then exit.
 ④ If parent of newnode is Red, then check the color of parent's sibling of newnode:
 (a) If color is black or null then do suitable rotation & recolor
 (b) if color is Red then recolor & also check if parent's parent of newnode is not root node then recolor it & recheck
 → root = Black
 → no two adjacent red nodes
 → Count no of black nodes in each path.

JENNY'S LECTURES

Deletion in Red-Black Tree

Step 1:- Perform BST deletion
 Step 2:- if node to be deleted is red,
 Case 1:- just delete it.
 Case 2:- if root is DB, just remove DB
 Case 3:- if DB's sibling is black, & both its children are black
 → remove DB
 → Add black to 1b Parent (P)
 (1) If P is Red it becomes black
 (2) If P is black it becomes double black
 → Make sibling red
 → if still DB exists, apply other cases.
 Case 4:- if DB's sibling is red
 → Swap colors of Parent & its Sibling
 → Rotate Parent in DB direction
 → resupply cases
 Cases:- DB's sibling is black, sibling's child who is far from DB is black, but near child to DB is red.



Expression Tree

Expression trees

- Binary trees could be used to represent arithmetic expressions
- Each operator (+ - * /) may have one or two operands
- The left operand is the left subtree of the operator, right operand is the right subtree

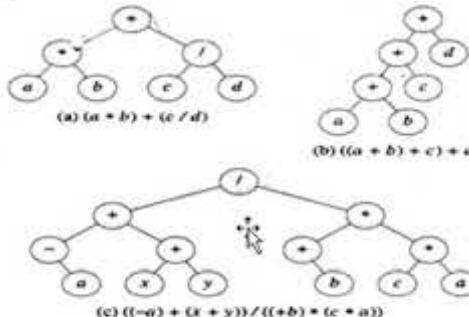


Figure 11.5 Expression trees

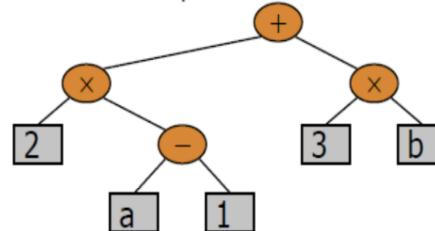
Arithmetic Expression Tree

- Binary tree associated with an arithmetic expression is a proper binary tree where:
 - internal nodes: operators
 - external nodes: operands

Example:

arithmetic expression tree for the expression

$(2 \times (a - 1) + (3 \times b))$



Prefix, Postfix, Infix

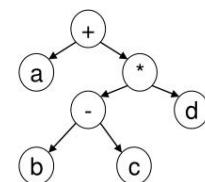
Prefix, Infix, and Postfix Forms

- A preorder traversal of an expression tree yields the prefix (or polish) form of the expression.
 - In this form, every operator appears before its operand(s).
- An inorder traversal of an expression tree yields the infix form of the expression.
 - In this form, every operator appears between its operand(s).
- A postorder traversal of an expression tree yields the postfix (or reverse polish) form of the expression.
 - In this form, every operator appears after its operand(s).

Prefix form: $+ a * - b c d$

Infix form: $a + b - c * d$

Postfix form: $a b c - d * +$



Infix Traversal



```
Algorithm infix(val tree <tree pointer>)
    if (tree not empty)
        if (tree→token is an operand)
            print (tree → token)
        else
            print (open parenthesis)
            infix(tree →left)
            print (tree →token)
            infix(tree →right)
            print (close parenthesis)
        end if
    end if
    return
end infix
```

Postfix Traversal



- Menggunakan postorder traversal seperti pada tree.
- Tidak membutuhkan kurung

```
Algorithm postfix(val tree <tree pointer>)
    if (tree not empty)
        postfix(tree →left)
        postfix(tree →right)
        print (tree →token)
    end if
    return
end postfix
```

Heap

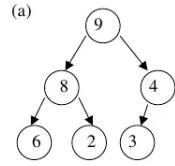
What is a Heap Data Structure?

- A Heap data structure is a binary tree with the following properties :

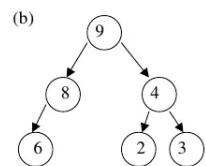
- It is a **complete binary tree**; that is, each level of the tree is completely filled, except possibly the bottom level. At this level, it is filled from left to right.
- It satisfies the heap-order property: The data item stored in each node is greater than or equal to the data items stored in its children.

- Examples:

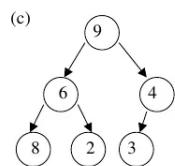
(a)



(b)



(c)



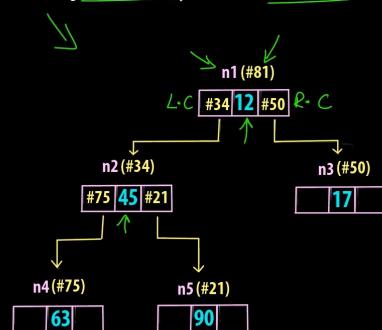
- In the above examples **only (a) is a heap**, (b) is not a heap as it is not complete and (c) is complete but does not satisfy the second property defined for heaps.

Application of Heap

- Heapsort**: One of the best sorting methods being in-place and with no quadratic worst-case scenarios.
- Finding the min, max, both the min and max, median, or even the k-th largest element can be done in linear time using heaps.
- Priority Queue:: Priority queues can be efficiently implemented using Binary Heap because it supports insert(), delete() and extractmax(), decreaseKey() operations in $O(\log n)$ time.
- Graph algorithms like [Prim's Algorithm](#) and [Dijkstra's algorithm](#).

HEAP Data Structure - Definition, Types, Applications, Implementation

1) Creating Nodes in Memory & link them in Tree Structure (min - heap)



2) Arranging the Heap DS elements in an Array structure (min - heap)

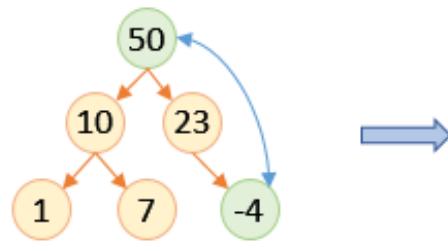
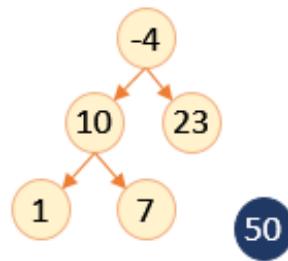
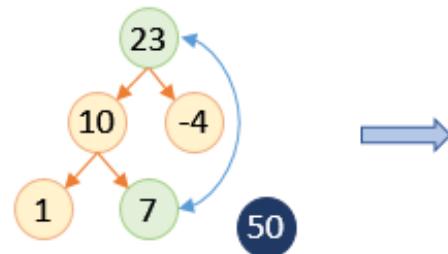
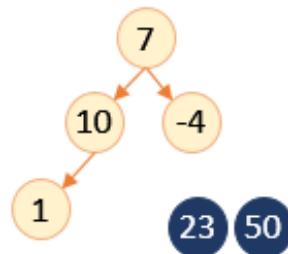
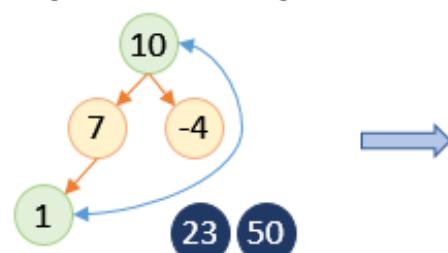
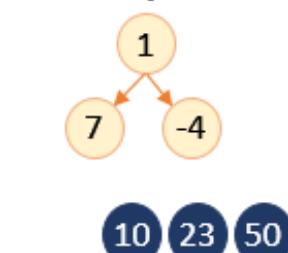
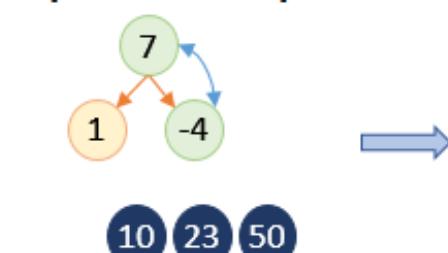
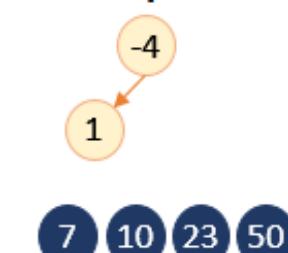
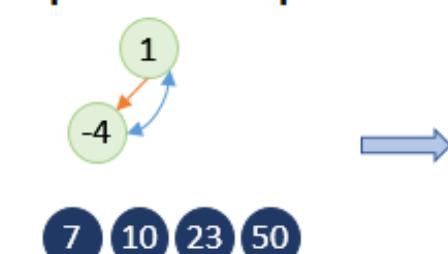
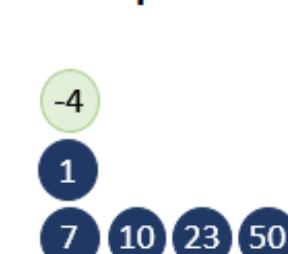
- The representation is done as -
- The root element will be at A[0]
 - Below table shows indexes of other nodes for the ith node i.e., A[i]:

- A [(i-1)/2] Returns the PARENT node
- A [(2*i)+1] Returns the LEFT child node
- A [(2*i)+2] Returns the RIGHT child node

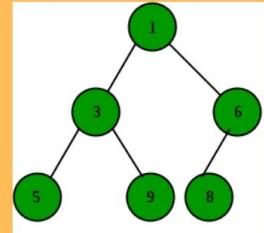
A = 16 | 54 | 19 | 67 | 87

The image you are requesting does not exist or is no longer available.

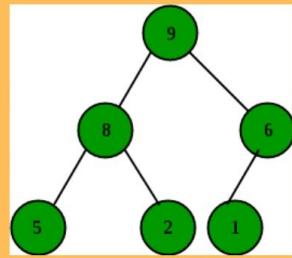
imgur.com

Step 1: Initial Max Heap**Step 2****Step 3: Max Heap****Step 4****Step 5: Max Heap****Step 6****Step 7: Max Heap****Step 8****Step 9: Max Heap****Step 10****Priority Queue**

Heap Data Structure



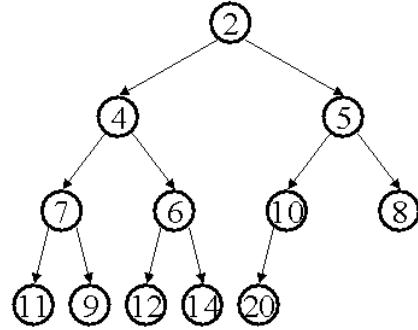
**Min Heap
Max Heap**



Priority Queue

Binary Heap Priority Q Data Structure

- Heap-order property
 - parent's key is less than children's keys
 - result: minimum is always at the top
- Structure property
 - complete tree with fringe nodes packed to the left
 - result: depth is always $O(\log n)$; next open location always known



How do we find the minimum?

```
In [ ]: #include <iostream>
#include <math.h>
using namespace std;

/*
swap(int &x, int &y){
    int temp = x;
    x = y;
    y = temp;
}
*/

class MinHeap{
public:
    int* heapArray;
    int capacity;
    int heapSize;

    MinHeap(int cap);
```

```

        int linearSearch(int val);
        int parent(int i);
        int leftchild(int i);
        int rightchild(int i);
        void insertion(int val);
        void minHeapify(int i);
        void deletion(int val);
        void decreaseKey(int i, int newVal);
        int getMin();
        int extractMin();
        void heapSort();
        int height();
        void printHeapArray();
    };

MinHeap::MinHeap(int cap){
    capacity = cap;
    heapSize = 0;
    heapArray = new int[cap];
}

int MinHeap::linearSearch(int val){
    for(int i=0; i<heapSize; i++){
        if(heapArray[i]==val){
            return i;
        }
    }
    return -1;
}

int MinHeap::parent(int i) {
    return (i-1)/2;
}
int MinHeap::leftchild(int i){
    return (2*i+1);
}
int MinHeap::rightchild(int i){
    return (2*i+2);
}

void MinHeap::insertion(int val){
    if(heapSize == capacity){
        cout << "\nMin Heap Overflow" << endl;
        return;
    }
    for(int i=0; i<heapSize; i++){
        if(heapArray[i]==val){
            cout << "Value Already exist. Insert Another Value." << endl;
            return;
        }
    }
    heapSize++;
    int i = heapSize - 1;
    heapArray[i] = val;

    while(i!=0){
        if(heapArray[parent(i)] > heapArray[i]){
            swap(heapArray[parent(i)], heapArray[i]);
            i = parent(i); // continue to compare until i==0
        }
    }
}

```

```

        return;
    }

void MinHeap::minHeapify(int i){
    int left = leftchild(i);
    int right = rightchild(i);
    int smallest = i;
    if ((left < heapSize) && (heapArray[left] < heapArray[smallest])){
        smallest = left;
    }
    if((right < heapSize) && (heapArray[right] < heapArray[smallest])){
        smallest = right;
    }
    if(smallest!=i){
        swap(heapArray[i], heapArray[smallest]);
        minHeapify(smallest);
    }
}

void MinHeap::decreaseKey(int i, int newVal){
    heapArray[i] = newVal;
    while(i>=0 && (heapArray[parent(i)] > heapArray[i])){
        swap(heapArray[parent(i)], heapArray[i]);
        i = parent(i); // continue to compare until i==0
    }
}

void MinHeap::deletion(int val){
    int index = linearSearch(val);
    if(index == -1){
        cout << "Value " << val << " Not Found";
    }
    else {
        decreaseKey(index, INT_MIN);
        extractMin();
        cout << "Delete Value " << val << endl;
    }
}

int MinHeap::getMin(){
    return heapArray[0];
}

int MinHeap::extractMin(){
    if(heapSize <=0)
        return INT_MAX;
    if(heapSize==1){
        heapSize--;
        return heapArray[0];
    }
    int min = heapArray[0];
    heapArray[0] = heapArray[heapSize-1];
    heapSize--;
    minHeapify(0);
    return min;
}

void MinHeap::heapSort(){
    for(int i=0; i<capacity; i++){
}

```

```
        cout << extractMin() << " ";
    }

int MinHeap::height(){
    return ceil(log2(heapSize+1)) - 1;
}

void MinHeap::printHeapArray(){
    for(int i=0; i<heapSize; i++){
        cout << heapArray[i] << " ";
    }
    cout << endl << endl;
}

int main(){

    int cap;
    cout << "Enter the Size of Min Heap: ";
    cin >> cap;
    MinHeap obj(cap);
    cout << "Min Heap Created" << endl << endl;

    int option, val;

    do{
        cout << "What operation do you want to perform?" << endl << endl;
        cout << "1. Insert" << endl;
        cout << "2. Search" << endl;
        cout << "3. Delete" << endl;
        cout << "4. Get Min" << endl;
        cout << "5. Extract Min" << endl;
        cout << "6. Height of Heap" << endl;
        cout << "7. Min Heap Sort" << endl;
        cout << "8. Print the Heap" << endl;
        cout << "9. Clear Screen" << endl;
        cout << "0. Exit Program" << endl;

        cout << "\nSelect Option Number or 0 to Exit." << endl;
        cin >> option;

        switch(option){
            case 0:
                break;
            case 1:
                cout << "Insert Operation" << endl;
                cout << "Enter Value to Insert in Heap: ";
                cin >> val;
                obj.insertion(val);
                cout << endl << endl;
                break;
            case 2:
                cout << "Search Operation" << endl;
                cout << "Enter Value to Search in Heap: ";
                cin >> val;
                if(obj.linearSearch(val)==-1){
                    cout << "Not Found";
                }
                else {

```

```
        cout << val << " Found at index " << obj.linearSearch(val);
    }
    cout << endl << endl;
    break;
case 3:
    cout << "Delete Operation" << endl;
    cout << "Enter Value to Delete in Heap: ";
    cin >> val;
    obj.deletion(val);
    cout << endl << endl;
    break;
case 4:
    cout << "Get Min Operation" << endl;
    cout << "Min Value: " << endl;
    cout << obj.getMin() << endl;
    cout << endl << endl;
    break;
case 5:
    cout << "Extract Min Value: ";
    cout << obj.extractMin() << endl;
    cout << endl << endl;
    break;
case 6:
    cout << "The Height of Heap: ";
    cout << obj.height() << endl;
    cout << endl << endl;
    break;
case 7:
    cout << "\nUnsorted Array: " << endl;
    obj.printHeapArray();
    cout << "Sorted Array: " << endl;
    obj.heapSort();
    cout << endl << endl;
    break;
case 8:
    cout << "Print Heap Array: " << endl;
    obj.printHeapArray();
    cout << endl << endl;
    break;
case 9:
    system("cls");
    break;
default:
    cout << "Please Enter the Proper Option.";
    break;
}
} while(option!=0);
return 0;
}
```

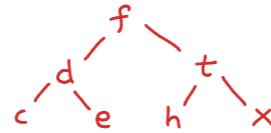
Treeep

data structure: treap (tree + heap)

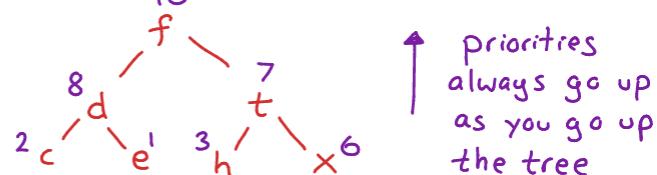
JULIA EVANS
@b0rk

Every node has a value (abc) and a priority (123)

A treap is a binary search tree:



where every node has a (randomly assigned) priority



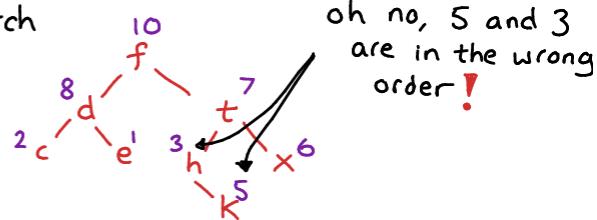
To search: for a value (like "h"), do a binary search.

Inserting: a new value is where it gets clever!

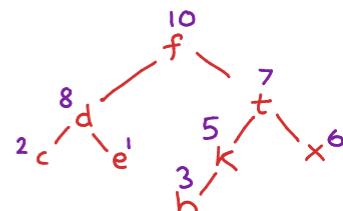
Let's insert \textcircled{K}

① Choose a random priority: 5

② Do a binary search to insert $\textcircled{K} 5$ into the tree



③ Do tree rotations until the priorities are in the right order!



Splay Tree

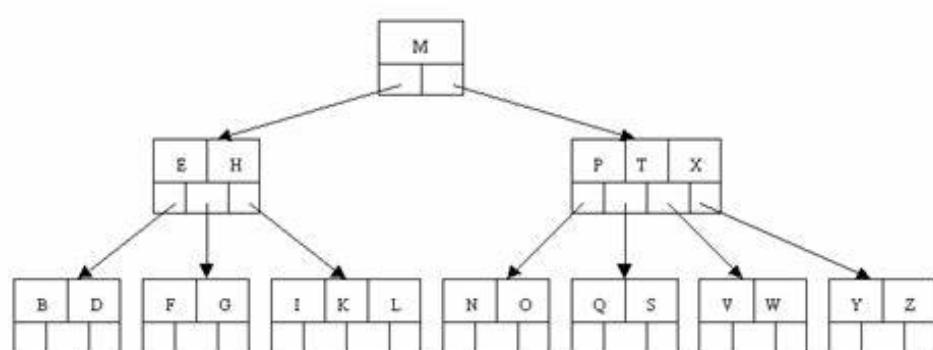
What is Splay Tree?

- Splay trees are self-adjusting binary search trees.
- A node in a Binary_Search_Tree is "splayed" when it is moved to the root of the tree by one or more "rotations".
- Whenever a node is accessed (Find, Insert, Remove, etc.), it is splayed, thereby making it the root.
- In addition to moving the accessed node to the root, the height of the tree may be shortened.

SPLAY TREE TIME COMPLEXITY

Sorted List	Search	Insertion	Deletion
with arrays	$O(\log n)$	$O(n)$	$O(n)$
with linked list	$O(n)$	$O(n)$	$O(n)$
With Splay trees	$O(\log n)$	$O(\log n)$	$O(\log n)$

B Tree

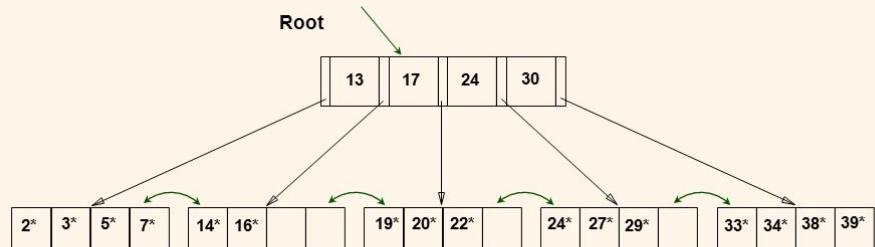


B⁺ Tree

Example B+ Tree



- ❖ Search begins at root, and key comparisons direct it to a leaf (as in ISAM).
- ❖ Search for 5*, 15*, all data entries $\geq 24^*$...



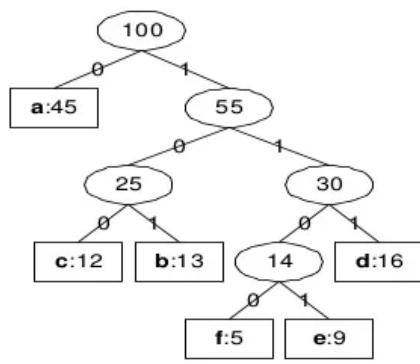
* Based on the search for 15*, we know it is not in the tree!

Database Management Systems 3ed, R. Ramakrishnan and J. Gehrke

10

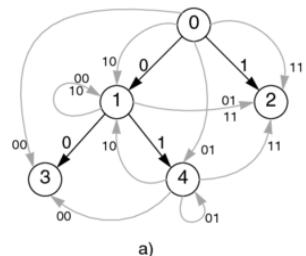
Huffman Code

Huffman Codes



- To encode:
 - Search the tree for the character to encode
 - As you progress, add "0" or "1" to right of code
 - Code is complete when you find character
- To decode a code:
 - Proceed through bit string left to right
 - For each bit, proceed left or right as indicated
 - When you reach a leaf, that is the decoded character

A Huffman Code Tree



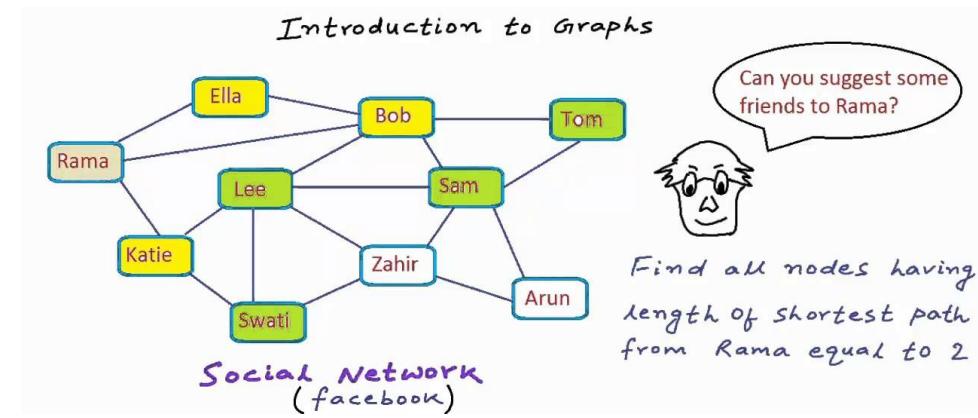
start-nodes	words			
	00	01	10	11
0	3	4	1	2
1	1	2	1	2
2	3	4	1	2
3	3	4	1	2
4	3	4	1	2

a) b)

from node	word	symbols
0	00	A
0	01	B
0	10	C
0	11	C, C
1	00	A
1	01	A, C
1	10	B
1	11	B, C

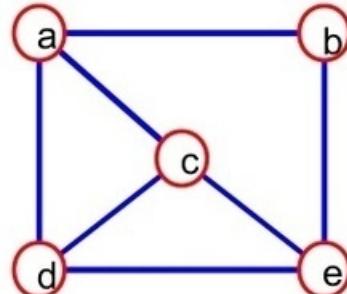
c)

Graph



What is a Graph?

- A graph $G = (V, E)$ is composed of:
 - V : set of vertices
 - E : set of edges connecting the vertices in V
- An edge $e = (u, v)$ is a pair of vertices
- Example:

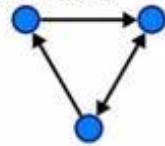


$$V = \{a, b, c, d, e\}$$

$$E = \{(a, b), (a, c), (a, d), (b, e), (c, d), (c, e), (d, e)\}$$

Directed and undirected graphs

- A graph is said as **directed graph** whose definition makes reference to edges which are directed. i.e, edges which are ordered pair of vertices.



- A graph is said as **undirected graph** whose definition makes reference to unordered pairs of vertices as edges is known as an undirected graph.

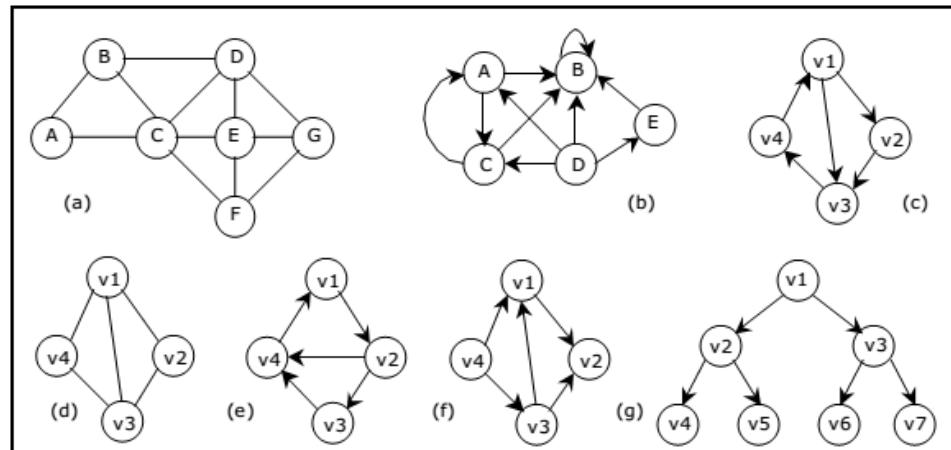
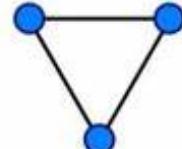
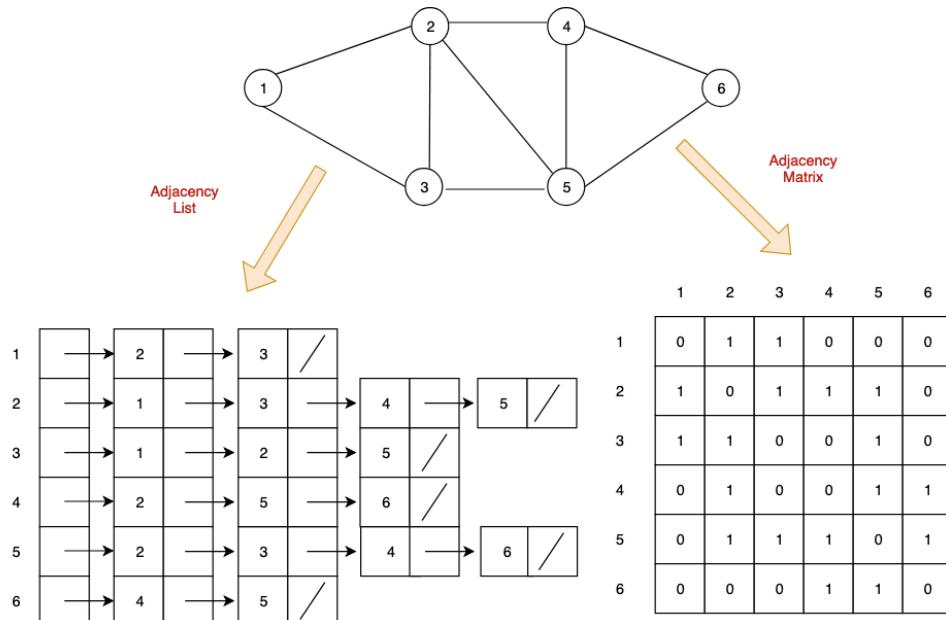
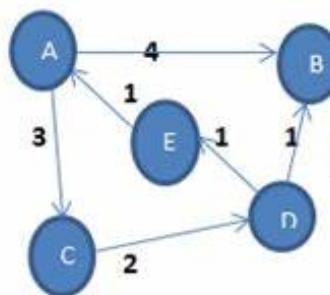


Figure 6.5.1 Various Graphs

Graph Presentation





Weighted Graph

	A	B	C	D	E
A	0	4	3	0	0
B	0	0	0	0	0
C	0	0	0	2	0
D	0	1	0	0	1
E	1	0	0	0	0

Adjacency matrix

```

In [ ]: #include <iostream>
using namespace std;

const int MAX_NODE_CNT = 10;

void distFromSource(
    const bool adjacent[][MAX_NODE_CNT],
    int dist[],
    int nodeCnt,
    int source
);

int main(){
    int nodeCnt = 5;
    bool adjacent[MAX_NODE_CNT][MAX_NODE_CNT] = {
        {1, 1, 0, 0, 1},
        {1, 1, 1, 0, 0},
        {0, 1, 1, 1, 0},
        {0, 0, 1, 1, 1},
        {1, 0, 0, 1, 1}
    };
    int dist[MAX_NODE_CNT] = {0};
    int source = 0;

    distFromSource(adjacent, dist, nodeCnt, source);

    cout << "\nThe complete result: \n";
    for (int i = 0; i < nodeCnt; i++){
        cout << dist[i] << " ";
    }

    return 0;
}

void distFromSource(const bool adjacent[][MAX_NODE_CNT],
                    int dist[], int nodeCnt, int source)
{
    for (int i = 0; i < nodeCnt; i++){
        dist[i] = nodeCnt;
    }

    dist[source] = 0;
    int curDist = 1;
    int complete = 1;

    while (complete < nodeCnt){
        for (int i = 0; i < nodeCnt; i++){
            if (dist[i] == curDist){
                for (int j = 0; j < nodeCnt; j++){
                    if (adjacent[i][j] == 1 && dist[j] > curDist + 1){
                        dist[j] = curDist + 1;
                    }
                }
            }
        }
        curDist++;
        if (curDist == nodeCnt) {
            complete = 1;
        } else {
            complete = 0;
        }
    }
}
  
```

```

        for (int i = 0; i < nodeCnt; i++){
            if(dist[i] == curDist - 1){
                for (int j = 0; j < nodeCnt; j++){
                    if((adjacent[i][j] == true)
                        && (dist[j] == nodeCnt)){
                        dist[j] = curDist;
                        complete++;
                    }
                }
            }
            curDist++;
        }
    }
}

```

Graph Implementation - Structure & Operations

Graph Implementation -

- Our graph would represent network of connected cities.
- Undirected & Weighted.

Operations -

1. Add Vertex (state)
2. Update Vertex(State name)
3. Delete Vertex
4. Add Edge
5. Update Edge(weight)
6. Delete Edge
7. Get All Neighbors of a Vertex
8. Check if 2 Vertices are connected/neighbors
9. Search a vertex
10. Print Full Graph

Graph (visual representation)

STL - Standard Template Library

The Standard Template Library (STL) is a set of C++ template classes to provide common programming data structures and functions such as lists, stacks, arrays, etc. It is a library of container classes, algorithms, and iterators.

Vector -

Vectors are same as dynamic arrays with the ability to resize itself automatically when an element is inserted or deleted, with their storage being handled automatically by the container. Vector elements are placed in contiguous storage so that they can be accessed and traversed using iterators. In vectors, data is inserted at the end.

Iterators -

Iterators are used to point at the memory addresses of STL containers. They are primarily used in sequence of numbers, characters etc.

List -> Doubly Linked List

Lists are sequence containers that allow non-contiguous memory allocation. They are similar to Doubly Linked List

SIMPLE SNIPPETS

Graph Implementation - Structure & Operations

Graph Implementation -

- Our graph would represent network of connected cities.
- Undirected & Weighted.

Operations -

1. Add Vertex (state)
2. Update Vertex(State name)
3. Delete Vertex
4. Add Edge
5. Update Edge(weight)
6. Delete Edge
7. Get All Neighbors of a Vertex
8. Check if 2 Vertices are connected/neighbors
9. Search a vertex
10. Print Full Graph

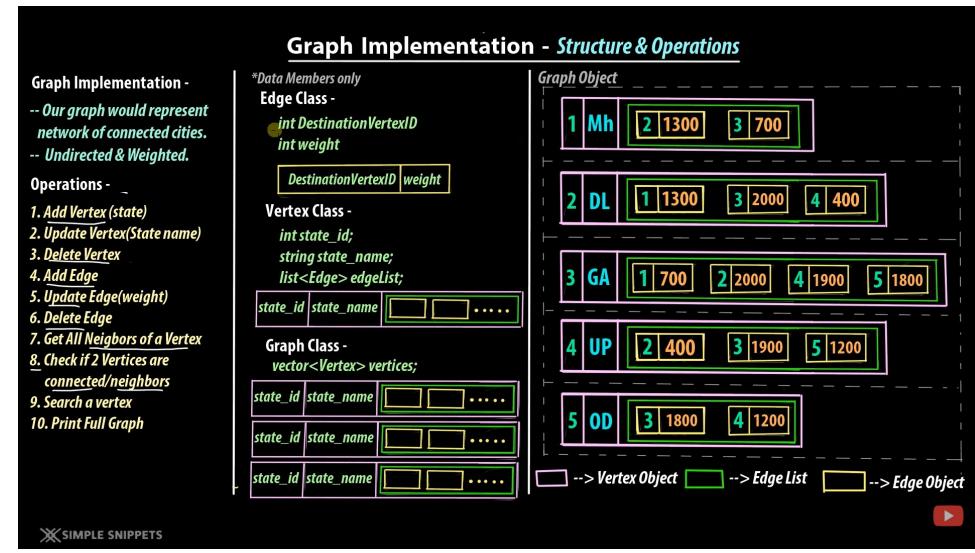
Graph (visual representation)

Graph Object

Graph Object	1 Mh	2 DL	3 GA	4 UP	5 OD
	[2 1300]	[1 1300]	[1 700]	[2 400]	[3 1800]
	[3 700]	[3 2000]	[2 2000]	[3 1900]	[4 1200]
	[4 400]	[4 1900]	[4 1900]	[5 1200]	[4 1200]
	[5 1200]	[5 1800]	[5 1800]		

--> Vertex Object --> Edge List --> Edge Object

SIMPLE SNIPPETS



```
In [ ]: #include <iostream>
#include <vector>
#include <list>
#include <iterator>
using namespace std;

class Edge {
private:
    int DestinationVertexID;
    int weight;
public:
    Edge(); // constructor
    Edge(int destVID, int w);
    void setEdgeValues(int destVID, int w);
    void setWeight(int w);
    int getDestinationVertexID();
    int getWeight();
};

Edge::Edge() : DestinationVertexID(0), weight(0) {}

Edge::Edge(int destVID, int w){
    DestinationVertexID = destVID;
    weight = w;
}

void Edge::setEdgeValues(int destVID, int w){
    DestinationVertexID = destVID;
    weight = w;
}

void Edge::setWeight(int w){
    weight = w;
}

int Edge::getDestinationVertexID(){
    return DestinationVertexID;
}

int Edge::getWeight(){
    return weight;
}
```

```

class Vertex {
private:
    int state_id;
    string state_name;
public:
    Vertex();
    Vertex(int id, string name);
    int getStateID();
    string getStateName();
    void setStateID(int id);
    void setStateName(string sname);
    list<Edge> edgeList;
    list<Edge> getEdgeList();
    void printEdgeList();
};

Vertex::Vertex() : state_id(0), state_name("") {}

Vertex::Vertex(int id, string sname){
    state_id = id;
    state_name = sname;
}

int Vertex::getStateID(){
    return state_id;
}

string Vertex::getStateName(){
    return state_name;
}

void Vertex::setStateID(int id){
    state_id = id;
}

void Vertex::setStateName(string sname){
    state_name = sname;
}

list<Edge> Vertex::getEdgeList(){
    return edgeList;
}

void Vertex::printEdgeList(){
    cout << "[";
    for(auto it=edgeList.begin(); it!=edgeList.end(); it++){
        cout << it->getDestinationVertexID()<< "(" << it->getWeight() << ") -->";
    }
    cout << "]" << endl;
}

class Graph {
public:
    vector<Vertex> vertices;
    bool checkIfVertexExistByID(int vid);
    bool checkIfEdgeExistByID(int fromVertex, int toVertex);
}

```

```

        Vertex getVertexByID(int vid);
        void addVertex(Vertex newVertex);
        void addEdgeByID(int fromVertex, int toVertex, int weight);
        void updateEdgeByID(int fromVertex, int toVertex, int newWeight);
        void deleteEdgeByID(int fromVertex, int toVertex);
        void deleteVertexByID(int vid);
        void updateVertexByID(int vid, string sname);
        void printGraph();
    };

    Vertex Graph::getVertexByID(int vid){
        Vertex temp;
        for(int i=0; i<vertices.size(); i++){
            temp = vertices.at(i);
            if(temp.getStateID() == vid){
                return temp;
            }
        }
        return temp;
    }

    bool Graph::checkIfVertexExistByID(int vid){
        for(int i=0; i<vertices.size(); i++){
            if(vertices.at(i).getStateID() == vid)
                return true;
        }
        return false;
    }

    void Graph::addVertex(Vertex newVertex){
        bool check = checkIfVertexExistByID(newVertex.getStateID());
        if(check){
            cout << "Vertex with this ID already exist" << endl << endl << endl;
        }
        else {
            vertices.push_back(newVertex);
            cout << "New Vertex Added Successfully" << endl << endl << endl;
        }
    }

    bool Graph::checkIfEdgeExistByID(int fromVertex, int toVertex){
        Vertex v = getVertexByID(fromVertex);
        list<Edge> e;
        e = v.getEdgeList();
        for(auto it = e.begin(); it != e.end(); it++){
            if(it->getDestinationVertexID() == toVertex){
                return true;
            }
        }
        return false;
    }

    void Graph::addEdgeByID(int fromVertex, int toVertex, int weight){
        bool check1 = checkIfVertexExistByID(fromVertex);
        bool check2 = checkIfVertexExistByID(toVertex);
        if (check1 && check2){
            bool check3 = checkIfEdgeExistByID(fromVertex, toVertex);
            if(check3){
                cout << "Edge between" << getVertexByID(fromVertex).getStateName()
                    << "(" << fromVertex << ")" and " << getVertexByID(fromVertex).getStateName()
                    << " has been added successfully" << endl;
            }
            else {
                cout << "Edge between" << getVertexByID(fromVertex).getStateName()
                    << "(" << fromVertex << ")" and " << getVertexByID(toVertex).getStateName()
                    << " has been added successfully" << endl;
            }
        }
    }
}

```

```

        << "(" << toVertex << ") already existed";
    }
    else {
        for(int i=0; i<vertices.size(); i++){
            if(vertices.at(i).getStateID() == fromVertex){
                Edge e(toVertex, weight);
                vertices.at(i).edgeList.push_back(e);
            }
            else if(vertices.at(i).getStateID() == toVertex){
                Edge e(fromVertex, weight);
                vertices.at(i).edgeList.push_back(e);
            }
        }
        cout << "Edge between " << fromVertex << " and "
        << toVertex << "added edge successfully" << endl << endl;
    }
    cout << "Invalid Vertex ID entered." << endl << endl;
}

void Graph::updateEdgeByID(int fromVertex, int toVertex, int newWeight){
    bool check = checkIfEdgeExistByID(fromVertex, toVertex);
    if(check){
        for(int i=0; i<vertices.size(); i++){
            if(vertices.at(i).getStateID() == fromVertex){
                for(auto it=vertices.at(i).edgeList.begin(); it!=vertices.at(i).edgeList.end();
                    if(it->getDestinationVertexID() == toVertex){
                        it->setWeight(newWeight);
                        break;
                    }
                }
            }
            else if(vertices.at(i).getStateID() == toVertex){
                for(auto it=vertices.at(i).edgeList.begin(); it!=vertices.at(i).edgeList.end();
                    if(it->getDestinationVertexID() == fromVertex){
                        it->setWeight(newWeight);
                        break;
                    }
                )
            }
        }
        cout << "Edge Weight Updated Successfully" << endl << endl;
    }
    else {
        cout << "Edge Between " << getVertexByID(fromVertex).getStateID()
        << "and " << getVertexByID(toVertex).getStateID() << " not existed";
    }
}

void Graph::deleteEdgeByID(int fromVertex, int toVertex){
    bool check = checkIfEdgeExistByID(fromVertex, toVertex);
    if(check){
        for(int i=0; i<vertices.size(); i++){
            if(vertices.at(i).getStateID() == fromVertex){
                for(auto it=vertices.at(i).edgeList.begin(); it!=vertices.at(i).edgeList.end();
                    if(it->getDestinationVertexID() == toVertex){
                        vertices.at(i).edgeList.erase(it);
                        break;
                    }
                )
            }
        }
    }
}

```

```

        else if(vertices.at(i).getStateID() == toVertex){
            for(auto it=vertices.at(i).edgeList.begin(); it!=vertices.at(i).edgeList.end();
                if(it->getDestinationVertexID() == fromVertex){
                    vertices.at(i).edgeList.erase(it);
                    break;
                }
            }
        }
        cout << "Edge Deleted Successfully" << endl << endl;
    }
    else {
        cout << "Edge does not exist" << endl << endl;
    }
}

void Graph::deleteVertexByID(int vid){
    bool check = checkIfVertexExistByID(vid);
    if(check){
        int vertexIndex = 0;
        for(int i=0; i<vertices.size(); i++){
            if(vertices.at(i).getStateID() == vid){
                vertexIndex = i;
            }
        }
        for(auto it=vertices.at(vertexIndex).edgeList.begin(); it!=vertices.at(vertexIndex).edgeList.end();
            deleteEdgeByID(it->getDestinationVertexID(), vid);
        }
        vertices.erase(vertices.begin() + vertexIndex);
        cout << "Vertex Deleted Successfully" << endl;
    }
    else{
        cout << "Invalid Vertex ID entered." << endl;
    }
}

void Graph::updateVertexByID(int vid, string sname){
    bool check = checkIfVertexExistByID(vid);
    if(check){
        for(int i=0; i<vertices.size(); i++){
            if(vertices.at(i).getStateID() == vid){
                vertices.at(i).setStateName(sname);
                break;
            }
        }
        cout << "Vertex(State) Updated Successfully" << endl << endl;
    }
    else {
        cout << "Invalid Vertex ID entered." << endl;
    }
}

void Graph::printGraph(){
    for(int i=0; i<vertices.size(); i++){
        Vertex temp;
        temp = vertices.at(i);
        cout << temp.getStateName() << "(" << temp.getStateID() << ") --> ";
        temp.printEdgeList();
        cout << endl << endl;
    }
}

```

```
        }

    }

}

int main(){

    int option;
    int id, id1, id2, weight;
    Graph g;
    string sname;
    Vertex v;

    do {
        cout << "\nWelcome to Learning Data Structure and Algorithm" << endl;
        cout << "Topic: Graph" << endl << endl;

        cout << "What operation do you want to perform? "
        << "Select Option Number or Enter 0 to exit." << endl << endl;

        cout << "1. Add Vertex" << endl;
        cout << "2. Update Vertex" << endl;
        cout << "3. Delete Vertex" << endl;
        cout << "4. Add Edge" << endl;
        cout << "5. Update Edge" << endl;
        cout << "6. Delete Edge" << endl;
        cout << "7. Check if 2 Vertices are Neighbors" << endl;
        cout << "8. Print All Neighbors of a Vertex" << endl;
        cout << "9. Print Graph" << endl;
        cout << "10. Clear Screen" << endl;
        cout << "0. Exit Program" << endl;
        cout << endl;
        cin >> option;

        cout << endl;

        switch (option) {
            case 0:
                break;
            case 1:
                cout << "Add Vertex" << endl;
                cout << "Enter State ID: ";
                cin >> id;
                cout << "Enter State Name: ";
                cin >> sname;
                v.setStateID(id);
                v.setStateName(sname);
                g.addVertex(v);
                break;
            case 2:
                cout << "Update Vertex" << endl;
                cout << "Enter ID of Vertex(State) to Update: ";
                cin >> id;
                cout << "Enter the new State Name: ";
                cin >> sname;
                g.updateVertexByID(id, sname);
                break;
            case 3:
                cout << "Delete Vertex" << endl;
                cout << "Enter ID of Vertex(State) to Delete: ";
        }
    }
}
```

```

        cin >> id;
        g.deleteVertexByID(id);
        break;
    case 4:
        cout << "Add Edge" << endl;
        cout << "Enter ID of Source Vertex(State): ";
        cin >> id1;
        cout << "Enter ID of Destination Vertex(State): ";
        cin >> id2;
        cout << "Enter Weight of Edge: ";
        cin >> weight;
        g.addEdgeByID(id1, id2, weight);
        break;
    case 5:
        cout << "Update Edge" << endl;
        cout << "Enter ID of Source Vertex(State): ";
        cin >> id1;
        cout << "Enter ID of Destination Vertex(State): ";
        cin >> id2;
        cout << "Enter Updated Weight of Edge: ";
        cin >> weight;
        g.updateEdgeByID(id1, id2, weight);
        break;
    case 6:
        cout << "Delete Edge" << endl;
        cout << "Enter ID of Source Vertex(State): ";
        cin >> id1;
        cout << "Enter ID of Destination Vertex(State): ";
        cin >> id2;
        g.deleteEdgeByID(id1, id2);
        break;
    case 7:
        cout << "Check if 2 Vertices are Neighbors" << endl;
        break;
    case 8:
        cout << "Print All Neighbors of a Vertex " << endl;
        break;
    case 9:
        cout << "Print Graph" << endl;
        g.printGraph();
    case 10:
        break;
    default:
        cout << "Please enter the proper option" << endl;
        break;
    }
} while ( option != 0);

return 0;
}

```

Graph Algorithm

Complexity
Terminology of graphs
The “big O” notation
Graph algorithms

Graphs/networks

- In graph theory, we talk about **graphs/networks**.
- A graph has **nodes (vertices)** and **edges (arcs/links)**.
 - A typical interpretation: Nodes are locations and arcs are roads.
- This graph has 9 nodes and 13 edges.
- Two nodes are adjacent** if there is an edge between them.
 - We say they are **neighbors**.
 - A node's **degree** is its number of neighbors.

Complexity
Terminology of graphs
The “big O” notation
Graph algorithms

Paths

- A **path (route)** from node s to node t is a set of directed edges $(s, v_1), (v_1, v_2), \dots$, and (v_{k-1}, v_k) , and (v_k, t) such that s and t are connected.
 - s is called the **source** and t is called the **destination** of the path.
 - Sometimes we write a path as $(s, v_1, v_2, \dots, v_k, t)$.
 - Direction matters!
- There are at least two paths from node 8 to node 9: $(8, 1, 5, 9)$ and $(8, 7, 1, 2, 3, 9)$.

Complexity
Terminology of graphs
The “big O” notation
Graph algorithms

Cycles

- A **cycle** (equivalent to circuit in some textbooks) is a path whose destination node is the source node.
 - A path is a **simple path** if it is not a cycle.
 - A graph is an **acyclic graph** if it contains no cycle.
- There is a cycle $(1, 2, 3, 9, 6)$.

Complexity
Terminology of graphs

The “big O” notation
Graph algorithms

Weights

- An edge may have a weight.
 - A weight may be a distance, a cost per unit item shipped, etc.
 - A **weighted graph** is a graph whose edges are weighted.
- In this network, we may use edge weights to represent distances.
 - The distance of the path $(8, 1, 5, 9)$ is 36. That of $(8, 7, 1, 2, 3, 9)$ is 56.
- A node may also have a weight.

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Complexity
Terminology of graphs

The “big O” notation
Graph algorithms

Graph algorithms

- As graphs can represent many things (logistic networks, power networks, social networks, etc.), there are many interesting issues.
 - How to find a shortest path from a node to another node?
 - How to link all nodes while minimizing the weights of selected edges?
 - How to check whether there is a cycle?
 - How to find the node with the maximum degree (number of neighbors)?
 - How to select the minimum number of nodes such that all nodes are either selected or adjacent to a selected node?
- Algorithms that solve these issues on graphs are **graph algorithms**.
- Below we give some examples demonstrating how to use an adjacency matrix.

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DFS

Introduction
Calculator

Applications of stacks
Graph search

Implementation of a stack
Recursion

Implementing depth-first search

- How to implement DFS?
- It is very natural to implement DFS with a **stack**.
 - When we want to visit a new node, find an out-neighbor from the **top** node of the stack. Note that the out-neighbor should be unvisited (why?).
 - Once we find one node to visit, **push** it into the stack.
 - When we reach a dead end, **pop** a node from the stack. This allows us to **backtrack** to the previous node.
- BFS may be implemented with a queue (to be introduced in the future).

Data Structures – Stacks 37 / 78 Ling-Chieh Kung (NTU IM)

Introduction Applications of stacks Implementation of a stack
Calculator Graph search Recursion

Implementing DFS with a stack

- Let's implement DFS with a stack containing nodes.
- First, push s .
- In each iteration, we **peek** (not pop!) the top node u and **push** a node v into the stack if (u, v) exists (i.e., v is an **out-neighbor** of u) and v is **unvisited**.
 - If no such v exists, **pop** the top node and continue.
 - If there are multiple candidates, there should be a way to pick one out-neighbor among them.
- Keep going until **t is pushed** into the stack or the stack becomes **empty**.
 - In the former case, the nodes in the stack gives a path from s to t .
 - In the latter case, there is no path from s to t .

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Introduction Applications of stacks Implementation of a stack
Calculator Graph search Recursion

Remarks

- DFS may also be used to identify **cycles** in a graph:
 - Whenever we try to visit a node that **has been visited**, we know there is a cycle.
- However, how to **know whether a node has been visited**?
 - We may need the support of some other data structures (such as an array) to record the history of visit.

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FOR MORE [HTTP://SCANTREE.COM](http://scantree.com)

Depth-First Search: The Code

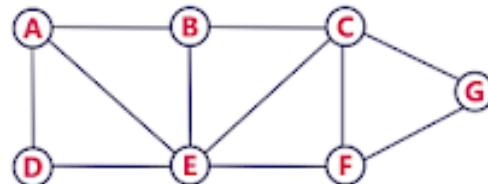
```
DFS(G)
{
    for each vertex u ∈ G->V
    {
        u->color = WHITE;
    }
    time = 0;
    for each vertex u ∈ G->V
    {
        if (u->color == WHITE)
            DFS_Visit(u);
    }
}
```

```
DFS_Visit(u)
{
    u->color = GREY;
    time = time+1;
    u->d = time;
    for each v ∈ u->Adj[]
    {
        if (v->color == WHITE)
            DFS_Visit(v);
    }
    u->color = BLACK;
    time = time+1;
    u->f = time;
}
```

Running time: $O(n^2)$ because call `DFS_Visit` on each vertex, and the loop over `Adj[]` can run as many as $|V|$ times

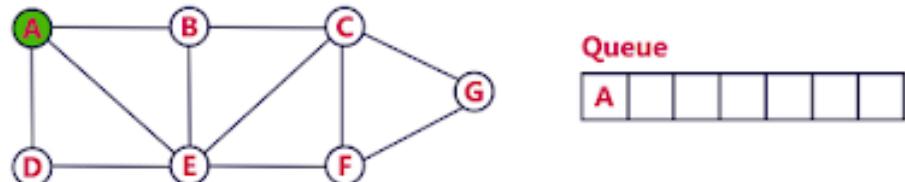
BFS

Consider the following example graph to perform BFS traversal



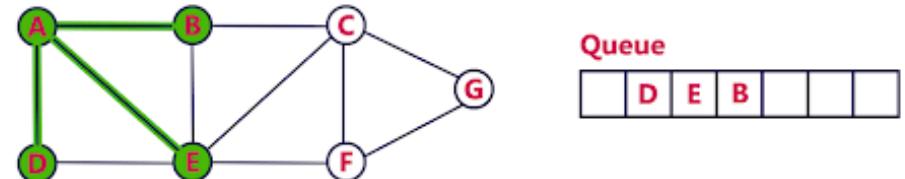
Step 1:

- Select the vertex **A** as starting point (visit **A**).
- Insert **A** into the Queue.



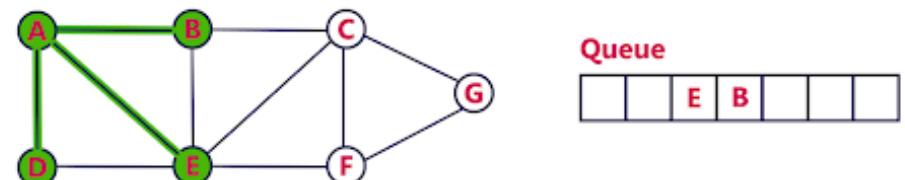
Step 2:

- Visit all adjacent vertices of **A** which are not visited (**D, E, B**).
- Insert newly visited vertices into the Queue and delete A from the Queue..



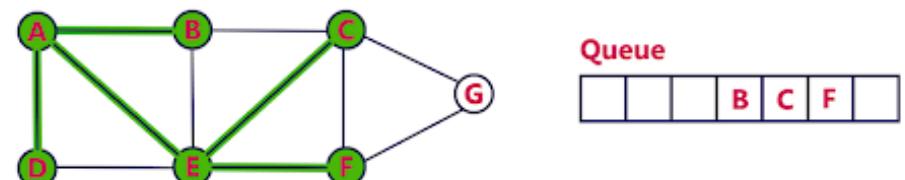
Step 3:

- Visit all adjacent vertices of **D** which are not visited (there is no vertex).
- Delete D from the Queue.



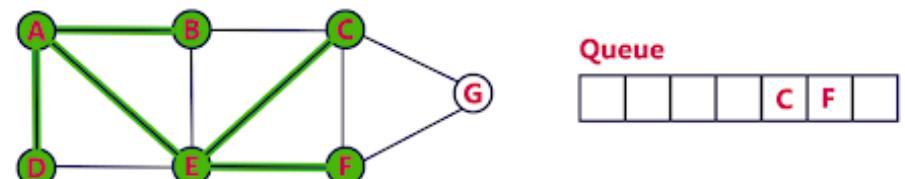
Step 4:

- Visit all adjacent vertices of **E** which are not visited (**C, F**).
- Insert newly visited vertices into the Queue and delete E from the Queue..



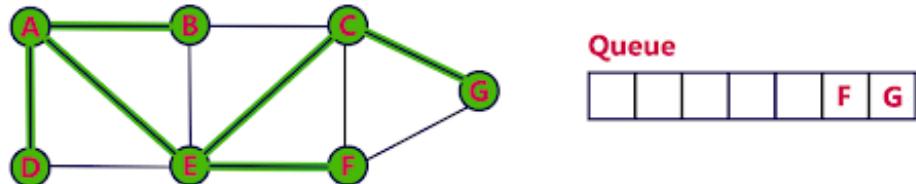
Step 5:

- Visit all adjacent vertices of **B** which are not visited (there is no vertex).
- Delete **B** from the Queue.

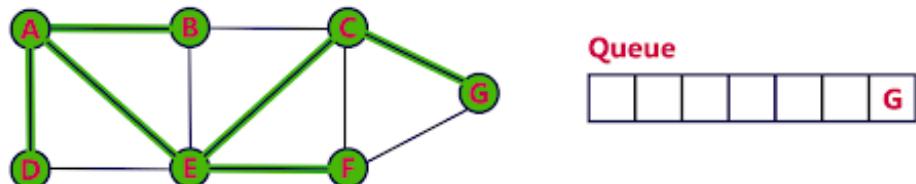


Step 6:

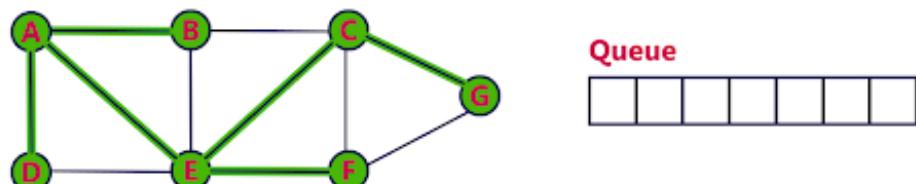
- Visit all adjacent vertices of **C** which are not visited (**G**).
- Insert newly visited vertex into the Queue and delete **C** from the Queue.

**Step 7:**

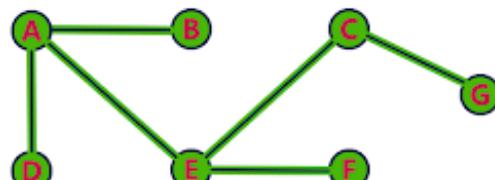
- Visit all adjacent vertices of **F** which are not visited (**there is no vertex**).
- Delete **F** from the Queue.

**Step 8:**

- Visit all adjacent vertices of **G** which are not visited (**there is no vertex**).
- Delete **G** from the Queue.

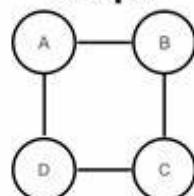


- Queue became Empty. So, stop the BFS process.
- Final result of BFS is a Spanning Tree as shown below...

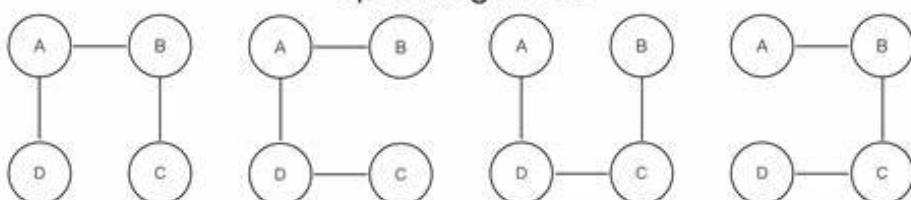


Graph - Spanning Tree

Graph



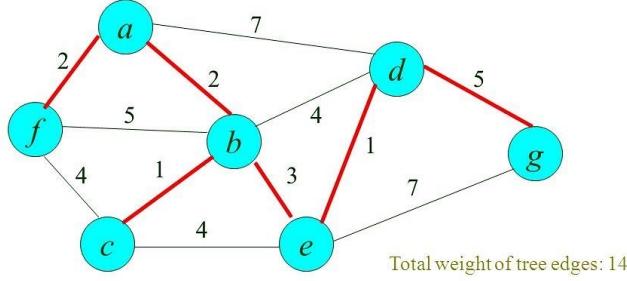
Spanning Trees



Minimum Spanning Tree (MST)

Problem Select edges in a connected and undirected graph to

- ★ form a tree that connects all the vertices (*spanning tree*).
- ★ minimize the total edge weight of the spanning tree.



Complexity

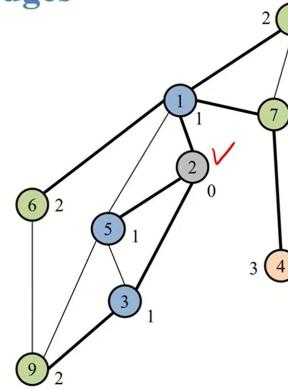
Terminology of graphs

The “big O” notation

Graph algorithms

Minimum number of edges

- Those “shortest paths” (thick lines in the graph) together form a spanning tree.

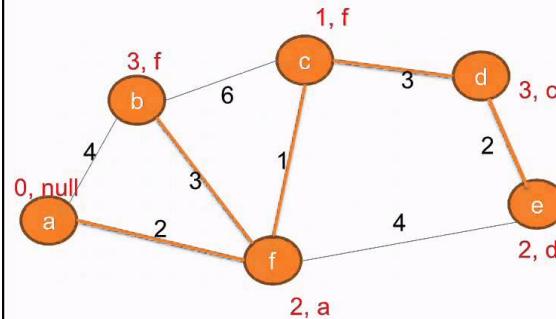


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- Removing one edge from the ST will make it disconnected
- Adding one edge to the ST will create a loop
- If each edge has distinct weight then there will be only one ?
⑤ 14, 13, 11, 9, 9
- Unique MST
- A Complete undirected graph can have $\frac{n(n-1)}{2}$ no. of ST
- Every connected & undirected graph has atleast one ST
- Disconnected graph does not have any ST
- From a complete graph by removing max $(n-1)$ edges we can construct a ST.

Prim's Algorithm for Minimum Spanning Tree (MST)

Prim Algorithm



$A = \{ (a, f), (c, f), (b, f), (c, d), (d, e) \}$

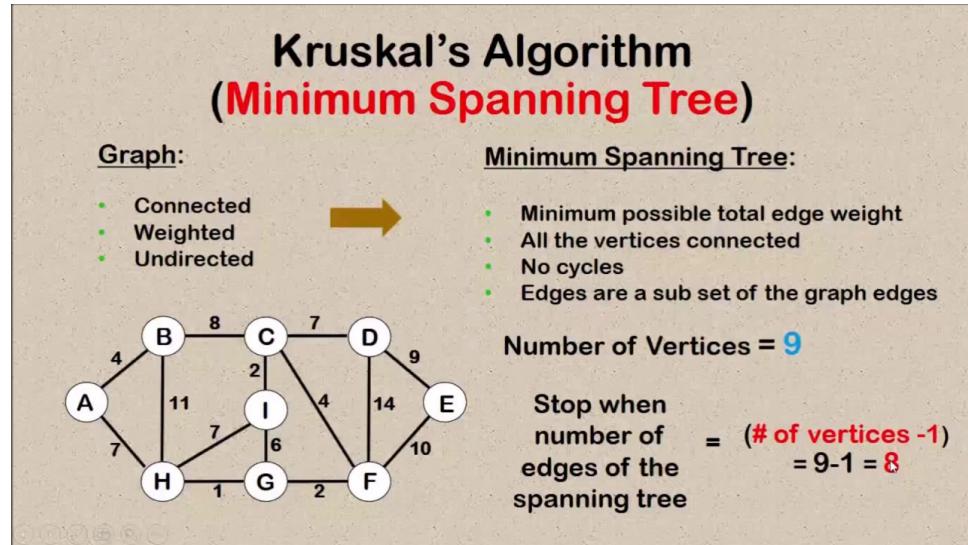
$Q = \{ \}$

```

 $A = \emptyset$ 
foreach  $v \in V:$ 
     $KEY[v] = \infty$ 
     $PARENT[v] = \text{null}$ 
 $KEY[r] = 0$ 
 $Q = V$ 
while  $Q \neq \emptyset:$ 
     $u = \min(Q)$  by  $KEY$  value
     $Q = Q - u$ 
    if  $PARENT(u) \neq \text{null}:$ 
         $A = A \cup (u, PARENT(u))$ 
    foreach  $v \in \text{Adj}(u):$ 
        if  $v \in Q$  and  $w(u,v) < KEY[v]:$ 
             $PARENT[v] = u$ 
             $KEY[v] = w$ 
return  $A$ 

```

Kruskals Algorithm for Minimum Spanning Tree- Greedy method

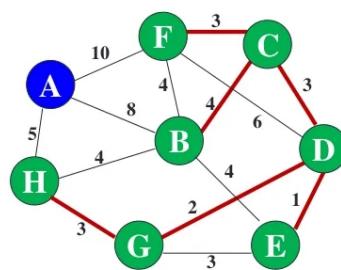


Lecture Slides By Adil Aslam

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Walk-Through

Select first $|V|-1$ edges which do not generate a cycle



edge	d_v	
(D,E)	1	✓
(D,G)	2	✓
(E,G)	3	✗
(C,D)	3	✓
(G,H)	3	✓
(C,F)	3	✓
(B,C)	4	✓

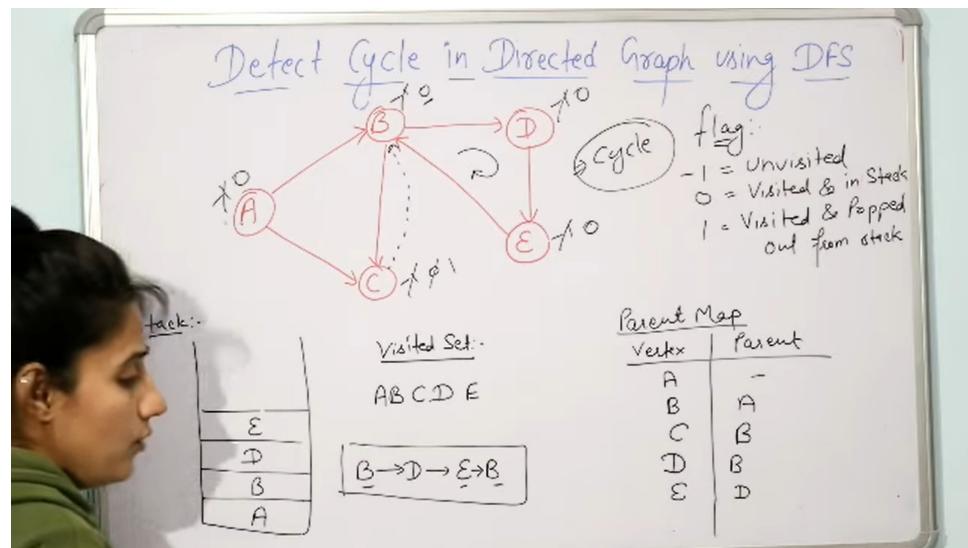
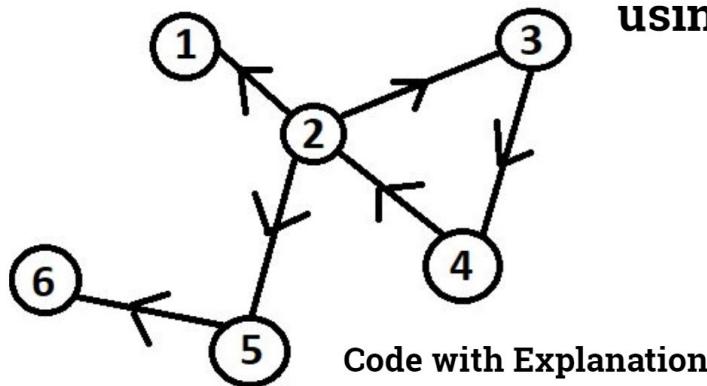
edge	d_v	
(B,E)	4	
(B,F)	4	
(B,H)	4	
(A,H)	5	
(D,F)	6	
(A,B)	8	
(A,F)	10	

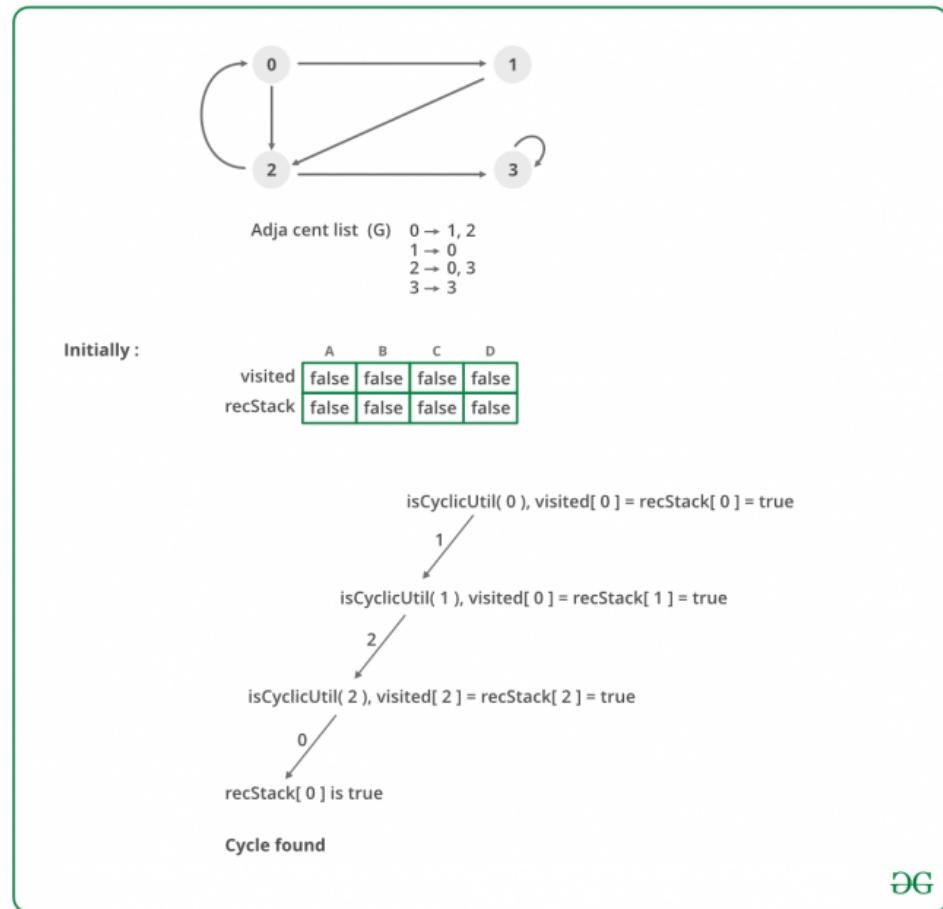
Detect Cycle in Directed Graph

Detect Cycle in a Directed Graph

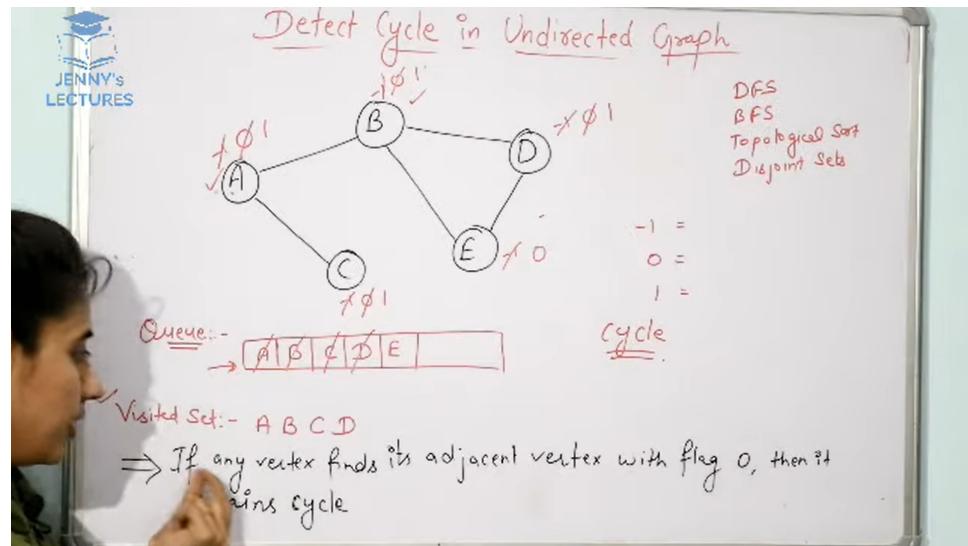
using

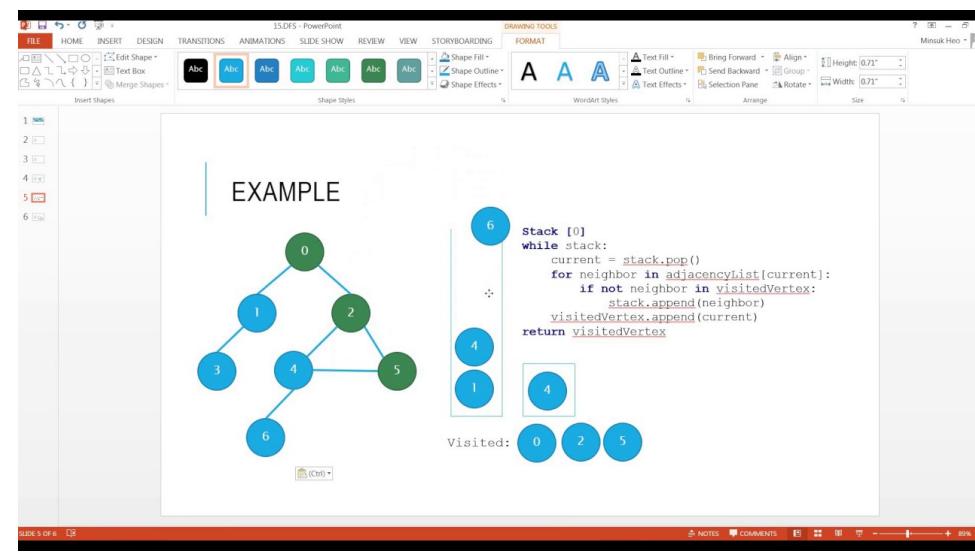
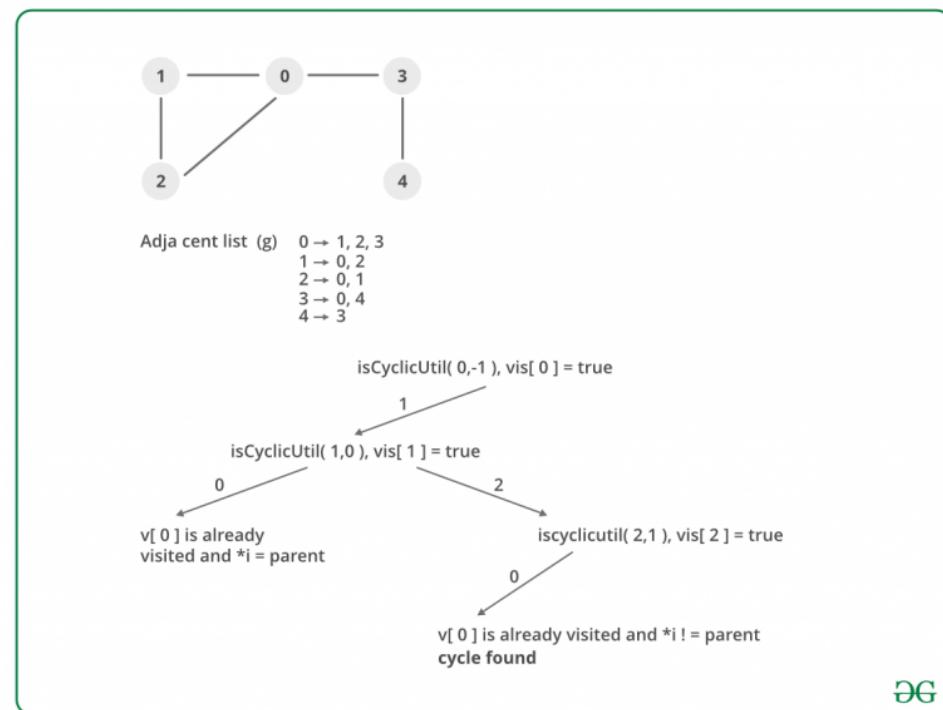
DFS





Detect Cycle in Undirected Graph

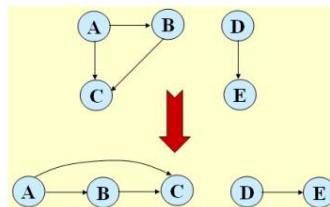




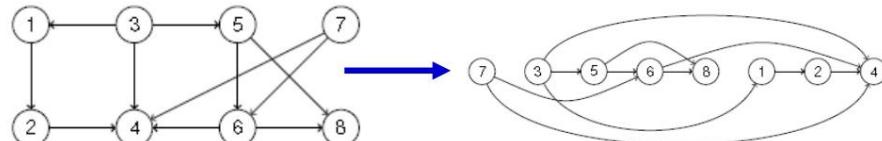
Topological Sorting for Graph

Definition of Topological Sort

- Given a directed graph $G = (V, E)$ a topological sort of G is an ordering of V such that for any edge (u, v) , u comes before v in the ordering.
- Example1:



- Example2:

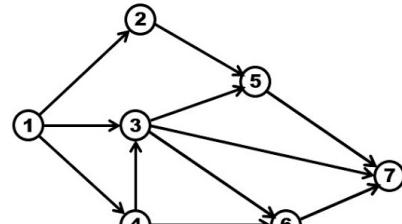


4

TOPOLOGICAL SORTING

Consider the following Graph:

Indegree [1] = 0	Check [1] = 0
Indegree [2] = 1	Check [2] = 0
Indegree [3] = 2	Check [3] = 0
Indegree [4] = 1	Check [4] = 0
Indegree [5] = 2	Check [5] = 0
Indegree [6] = 2	Check [6] = 0
Indegree [7] = 3	Check [7] = 0



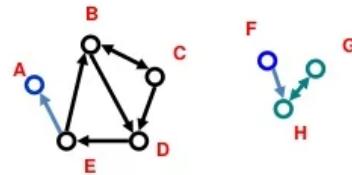
TS

Connected Components

Connected Components

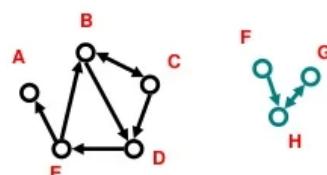
- Strongly connected components
 - Each node within the component can be reached from every other node in the component by following directed links

Strongly connected components
 B C D E
 A
 G H
 F



- Weakly connected components:
 every node can be reached from every other node by following links in either direction

Weakly connected components
 A B C D E
 G H F



- In undirected networks one talks simply about "connected components"

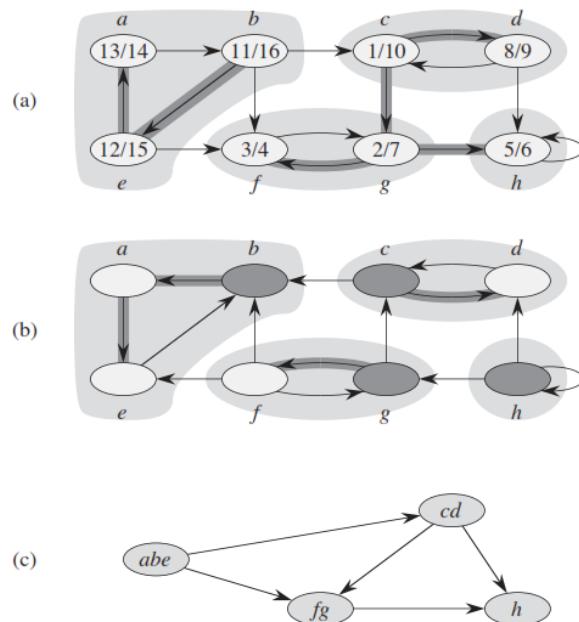


Figure 22.9 (a) A directed graph G . Each shaded region is a strongly connected component of G . Each vertex is labeled with its discovery and finishing times in a depth-first search, and tree edges are shaded. (b) The graph G^T , the transpose of G , with the depth-first forest computed in line 3 of STRONGLY-CONNECTED-COMPONENTS shown and tree edges shaded. Each strongly connected component corresponds to one depth-first tree. Vertices b , c , g , and h , which are heavily shaded, are the roots of the depth-first trees produced by the depth-first search of G^T . (c) The acyclic component graph G^{SCC} obtained by contracting all edges within each strongly connected component of G so that only a single vertex remains in each component.

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Connected Components

```

Connected_Components(G)
for each vertex v ∈ N
    flag[v] = -1
    count = 0
    for( int v = 0; v < N, v++)
        if( flag[v] == -1)
            { DFS(v, flag)
                count++;
            }
        cout << count;
    }

DFS(int v, int flag)
{
    flag[v] = 1; cout << v;
    for each adjacent node u of v
        if( flag[u] == -1)
            DFS(u, flag)
}

```

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Connected Components

```

Connected_Components(G)
for each vertex v ∈ N
    flag[v] = -1
    count = 0
    for( int v = 0, v < N, v++)
        if( flag[v] == -1)
            { DFS(v, flag)
                count++;
            }
        cout << count;
    }

DFS(int v, int flag)
{
    flag[v] = 1; cout << v;
    for each adjacent node u of v
        if( flag[u] == -1)
            DFS(u, flag)
}

```

O/P:
0 1 2 3
4 5

$u = 1$

$N = 9$
 $G(V, E)$
 $\text{count} = 1/2$

$v = 4$

$\text{count} = 1/2$

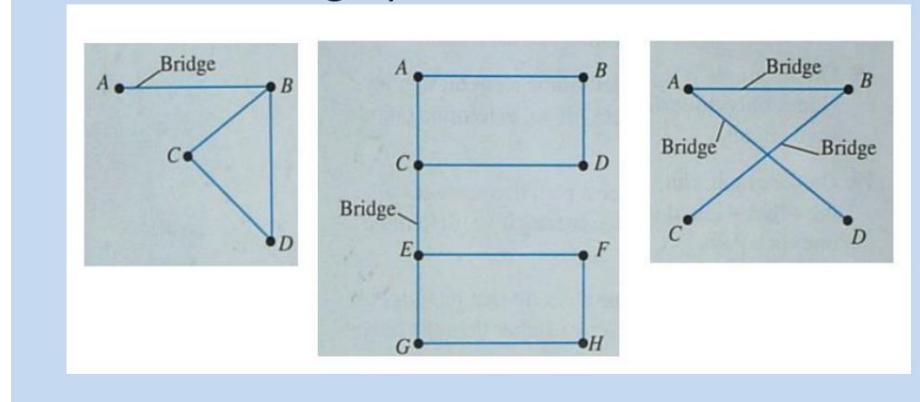
Find Bridges(Cut Edge) in a Graph

Mathematical Decision Making

Graph Theory

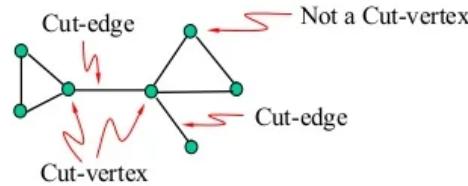
T.Serino

A **bridge** is an edge that if removed from a connected graph would create a disconnected graph.



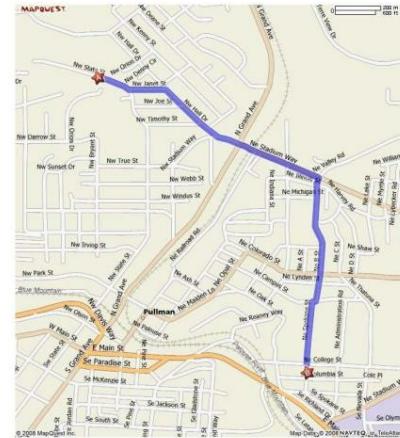
Cut-edge, Cut-vertex 1.2.12

- A ***cut-edge*** or ***cut-vertex*** of a graph is an edge or vertex whose deletion increases the number of components



SSSP(Shortest Path Algorithms)

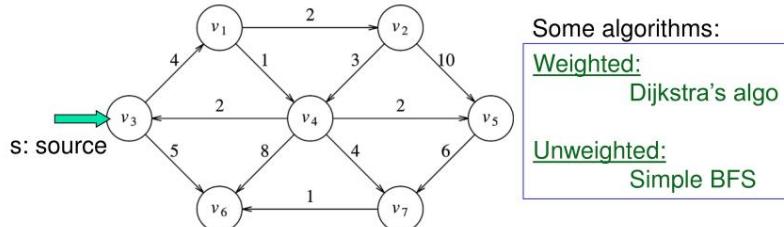
- 
- ### Shortest-Path Algorithms
- Find the “shortest” path from point A to point B
 - “Shortest” in time, distance, cost, ...
 - Numerous applications
 - Map navigation
 - Flight itineraries
 - Circuit wiring
 - Network routing



Shortest Path Problems

Single-source shortest path problem

- Given a weighted graph $G=(V,E)$, and a **source vertex s** , find the minimum weighted path from s to every other vertex in G



4

Complexity	The "big O" notation
Terminology of graphs	Graph algorithms

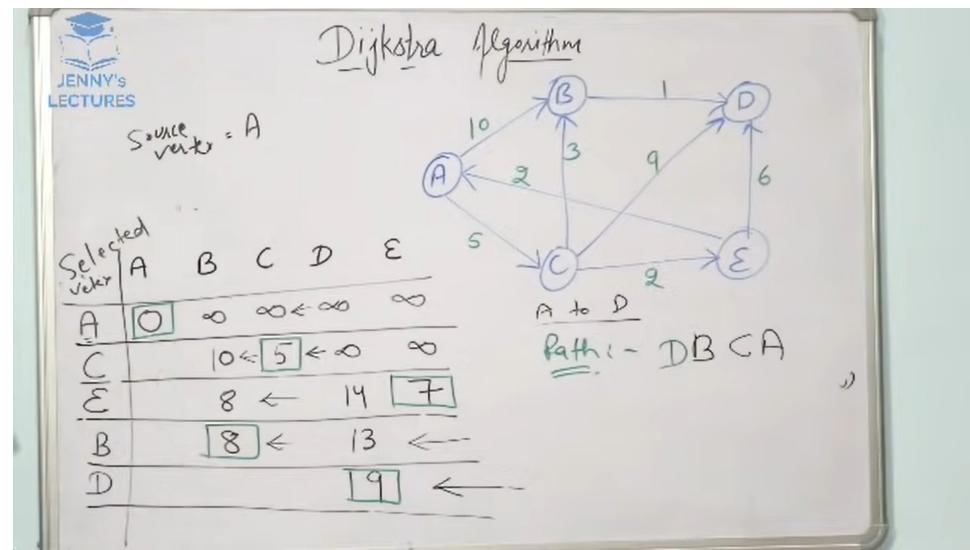
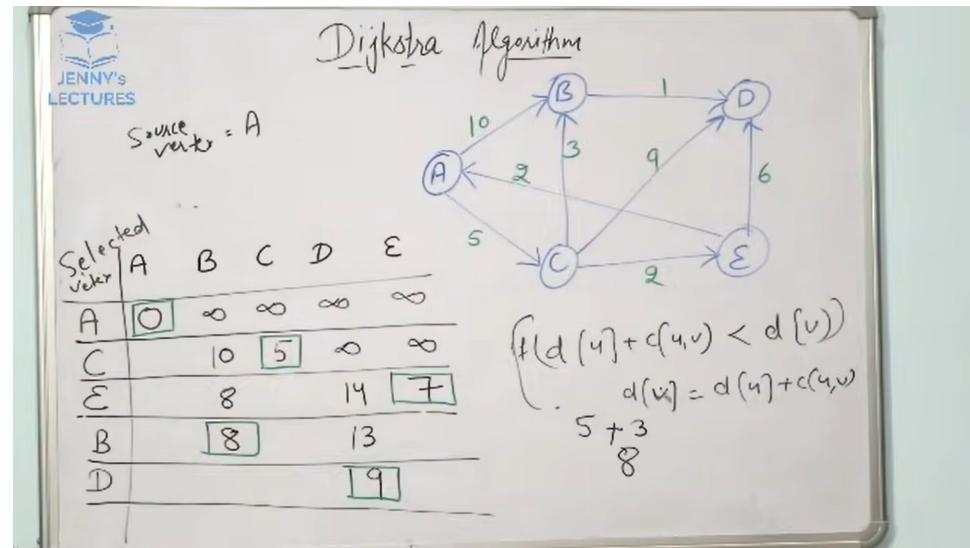
Minimum number of edges

- To find the distances from s to all nodes, we use **breadth-first search (BFS)**.
- Let all nodes have weights representing their distances from s .
 - First, we label s as 0 and all other nodes as ∞ .
 - We then find the neighbors of s . Label them as 1.
 - For each node whose label is 1, find its neighbors that are currently labeled as ∞ . Label them as 2.
 - Continue until all nodes are labeled.
- The graph should be **connected** (i.e., there is a path from s to any other node).

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Dijkstra Algorithm

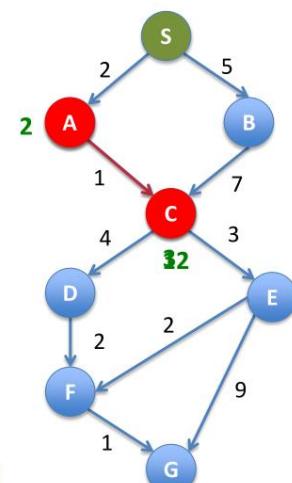
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Example: Single-Source Shortest-Path

- **Problem Formulation**
 - Compute shortest distance from source node S to every other node
- **Many algorithms**
 - Bellman-Ford (1957)
 - Dijkstra (1959)
 - Chaotic relaxation (Miranker 1969)
 - Delta-stepping (Meyer et al. 1998)
- **Common structure**
 - Each node has label *dist* with known shortest distance from S
- **Key operation**
 - relax-edge(u,v)

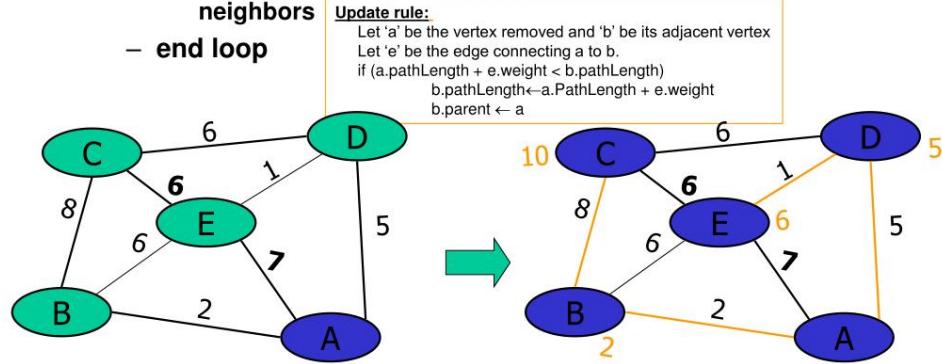
$\text{if } \text{dist}(A) + W_{Ac} < \text{dist}(C)$
 $\text{dist}(C) = \text{dist}(A) + W_{Ac}$



Single Source Shortest Path

Dijkstra's Algorithm

- Label Path length to each vertex as ∞ (or 0 for the start vertex)
- Loop while there is a vertex
 - Remove up the vertex with minimum path length
 - Check and if required update the path length of its adjacent neighbors
- end loop



Dijkstra's Rules

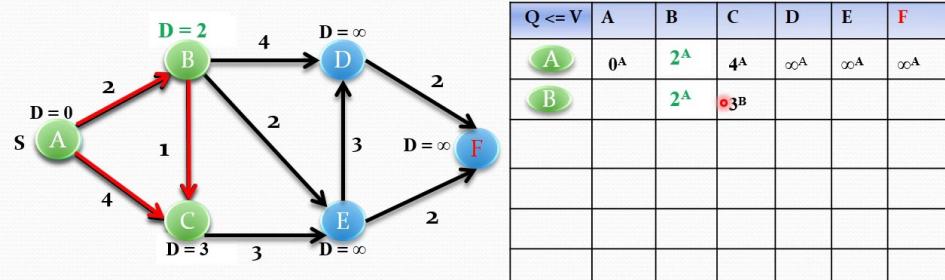
Rule 1: Make sure there is no negative edges. Set distance to source vertex as zero and set all other distances to infinity.

Rule 2: Relax all vertices adjacent to the current vertex.

Rule 3: Choose the closest vertex as next current vertex.

Rule 4: Repeat Rule2 and Rule 3 until the queue or reach the destination.

If $(D[Cj] + D[AdjEdgej]) < D[Adj]$ {Update Adj's D with new shortest path}



Bellman Ford Algorithm



Bellman-Ford algorithm

```

 $d[s] \leftarrow 0$ 
for each  $v \in V - \{s\}$  do  $d[v] \leftarrow \infty$  } initialization

for  $i \leftarrow 1$  to  $|V| - 1$ 
  do for each edge  $(u, v) \in E$ 
    do if  $d[v] > d[u] + w(u, v)$ 
      then  $d[v] \leftarrow d[u] + w(u, v)$  } relaxation step

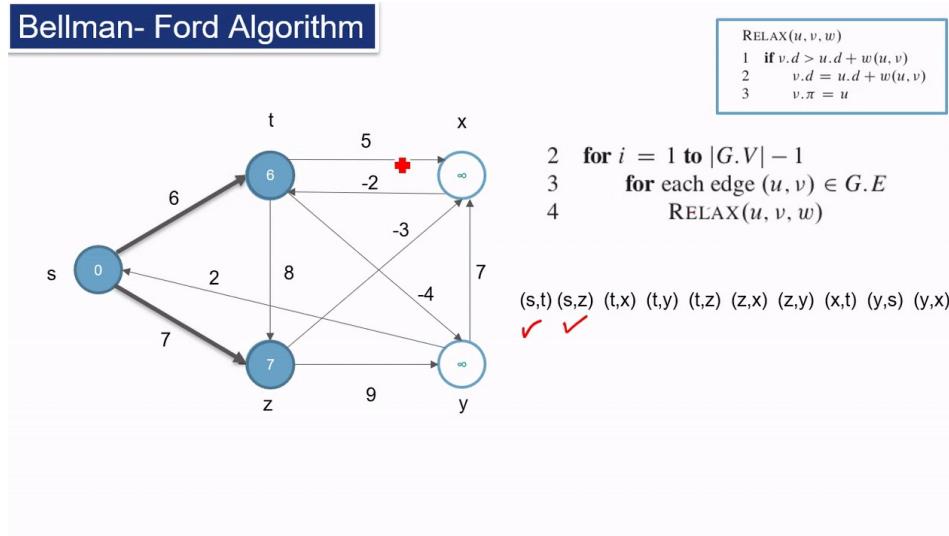
for each edge  $(u, v) \in E$ 
  do if  $d[v] > d[u] + w(u, v)$ 
    then report that a negative-weight cycle exists

At the end,  $d[v] = \delta(s, v)$ , if no negative-weight cycles.  

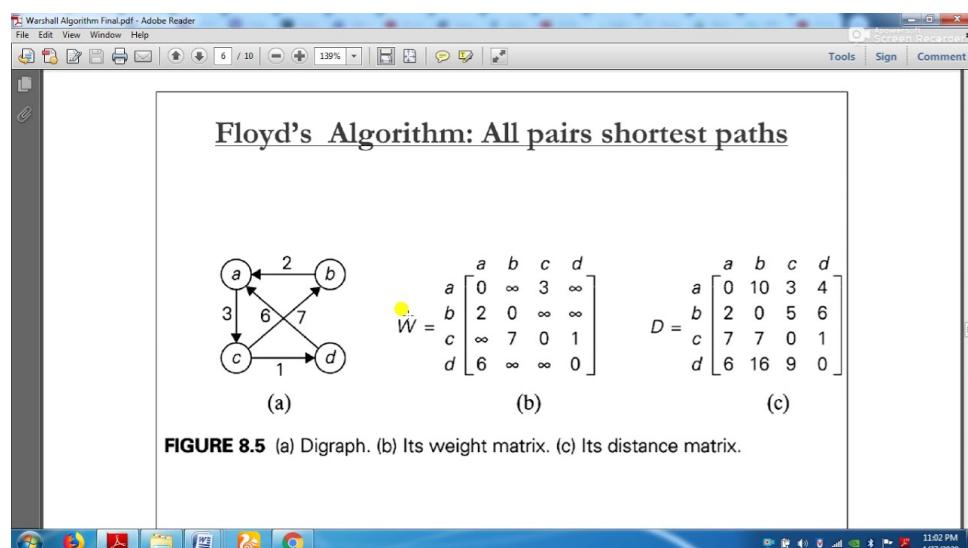
Time =  $O(VE)$ .
  
```

Algorithms

L18.4



Floyd Warshall Algorithm (All Pair Shortest Path algorithm)



FLOYD-WARSHALL(W)

1. $n \leftarrow \text{rows}[W]$
2. $D^{(0)} \leftarrow W$
3. **for** $k \leftarrow 1$ **to** n
4. **do for** $i \leftarrow 1$ **to** n
5. **do for** $j \leftarrow 1$ **to** n
6. $d_{ij}^{(k)} \leftarrow \min(d_{ij}^{(k-1)}, d_{ik}^{(k-1)} + d_{kj}^{(k-1)})$
7. **return** $D^{(n)}$

Search Algorithm

Algorithms and complexity Recursion **Searching and sorting**

Linear search vs. binary search

- In binary search, the number of instructions to be executed is roughly proportional to $\log_2 n$.
- So binary search is **much more efficient** than linear search!
 - The difference is huge if the array is large.
 - However, binary search is possible only if the array is sorted.
 - Is it worthwhile to sort an array before we search it?
- It is natural to implement binary search with **recursion**.
 - A subproblem is to search for the element in one half of the array.
- Binary search can also be implemented with repetition.
 - Is it natural to do so?

Programming Design – Algorithms and Recursion 35 / 43 Ling-Chieh Kung (NTU IM)

Linear Search

```
In [ ]: #include <iostream>
using namespace std;

void linearSearch(int arr[], int n);

int main(){
    int arr[5] = {1, 22, 34, 5, 7};
    int searchNum;

    cout << "Please enter a number to search: ";
    cin >> searchNum;

    linearSearch(arr, searchNum);

    return 0;
}

void linearSearch(int arr[], int searchNum){
    int found = -1;
```

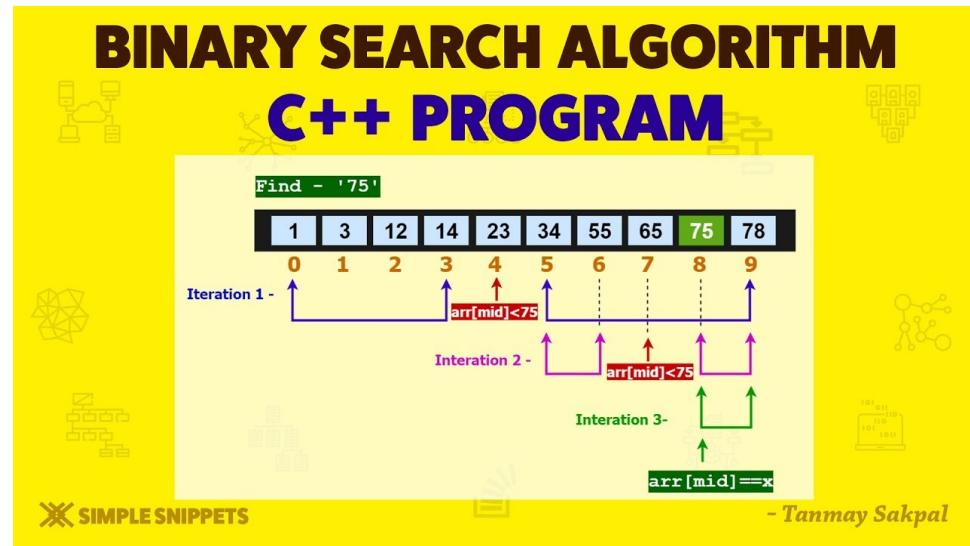
```

for(int i = 0; i < 5; i++){
    if(arr[i] == searchNum){
        cout << "Found at index: " << i << endl;
        found = 1;
    }
}

if(found == -1){
    cout << "Not Found" << endl;
}
}

```

Binary Search in sorted Array



```

In [ ]: #include <iostream>
using namespace std;

int binarySearch(int arr[], int left, int right, int n);

int main(){

    int myArr[10];
    int searchNum;
    int index;

    cout << "Please enter 10 numbers in ASCENDING order" << endl;
    for(int i=0; i<10; i++){
        cin >> myArr[i];
    }
    cout << endl;

    cout << "Please enter a number to search: ";
    cin >> searchNum;
    cout << endl;

    index = binarySearch(myArr, 0, 9, searchNum);

    if(index == -1){
        cout << "No Match Found" << endl;
    }
    else {

```

```

        cout << "Found at index: " << index << endl;
    }

    return 0;
}

int binarySearch(int arr[], int left, int right, int x){

    while(left <= right){
        int mid = left + (right - left)/2;
        if(arr[mid] == x){
            return mid;
        }
        else if(arr[mid] < x){
            left = mid + 1;
        }
        else {
            right = mid - 1;
        }
    }

    return -1;
}

```

Algorithms and complexity Recursion Searching and sorting

Binary search: pseudocode

```

binarySearch(a sorted array A, search in between from and to, search for p)
if n = 1
    return true if Afrom = p; return false otherwise
else
    let median be floor((from + to) / 2)
    if p = Amedian
        return true
    else if p < Amedian
        return binarySearch(A, from, median, p)
    else
        return binarySearch(A, median + 1, to, p)

```

Programming Design – Algorithms and Recursion 34 / 43

Ling-Chieh Kung (NTU IM)

In []: #include <iostream>
using namespace std;

```

int binarySearch(int arr[], int left, int right, int n);

int main(){

    int myArr[10];
    int searchNum;
    int index;

    cout << "Please enter 10 numbers in ASCENDING order" << endl;
    for(int i=0; i<10; i++){
        cin >> myArr[i];
    }
    cout << endl;
}
```

```
cout << "Please enter a number to search: ";
cin >> searchNum;
cout << endl;

index = binarySearch(myArr, 0, 9, searchNum);

if(index == -1){
    cout << "No Match Found" << endl;
}
else {
    cout << "Found at index: " << index << endl;
}

return 0;
}

int binarySearch(int arr[], int left, int right, int x){

while(left <= right){
    int mid = left + (right - left)/2;
    if(arr[mid] == x){
        return mid;
    }
    else if(arr[mid] < x){
        return binarySearch(arr, mid + 1, right, x);
    }
    else {
        return binarySearch(arr, left, mid - 1, x);
    }
}

return -1;
}
```

Sort Algorithm

Sorting Algorithm	Time Complexity			Space Complexity
	Best Case	Average Case	Worst Case	
Insertion Sort	$\Omega(n)$	$\Theta(n^2)$	$O(n^2)$	$O(1)$
Selection Sort	$\Omega(n^2)$	$\Theta(n^2)$	$O(n^2)$	$O(1)$
Bubble Sort	$\Omega(n)$	$\Theta(n^2)$	$O(n^2)$	$O(1)$
Merge Sort	$\Omega(n \log(n))$	$\Theta(n \log(n))$	$O(n \log(n))$	$O(n)$
Quicksort	$\Omega(n \log(n))$	$\Theta(n \log(n))$	$O(n^2)$	$O(\log(n))$
Heapsort	$\Omega(n \log(n))$	$\Theta(n \log(n))$	$O(n \log(n))$	$O(1)$
Counting Sort	$\Omega(n+k)$	$\Theta(n+k)$	$O(n+k)$	$O(k)$
Radix Sort	$\Omega(nk)$	$\Theta(nk)$	$O(nk)$	$O(n+k)$

|>

C++'s Sort Methods

- Designed to optimize the average performance.
- `std::sort()`
 - Modified Quick sort.
 - Heap Sort
 - when the number of subdivision exceed `clogn`
 - Insertion Sort
 - when the segment size becomes small
- `std::stable_sort()`
 - Merge Sort.
 - Insertion Sort
 - when the segment size becomes small
- `std::partial_sort()`
 - Heap Sort.

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Sorting in Array

Here are the two C++ implementations to use for your tracing exercise. First is a Selection Sort

```

1 template<class T> void sort(T* a, int n)
2 { for (int i=0; i < n-1; i++)
3     { int min=i;
4      for (int j=i+1; j < n; j++)
5          if (a[j] < a[min]) min = j;
6      swap(a[min], a[i]);
7    }
8 }
```

The second is an Insertion Sort:

```

1 template<class T> void sort(T* a, int n)
2 { for (int i=1, j; i < n; i++)
3     { T temp = a[i];
4      for (j=i; j > 0 && a[j-1] > temp; j--)
5          a[j] = a[j-1];
6      a[j] = temp;
7    }
8 }
```

Selection Sort

SELECTION SORT ALGORITHM (C++ PROGRAM)

```

void selectionSort(int arr[])
{
    for(int i=0;i<4;i++)
    {
        int min = i;

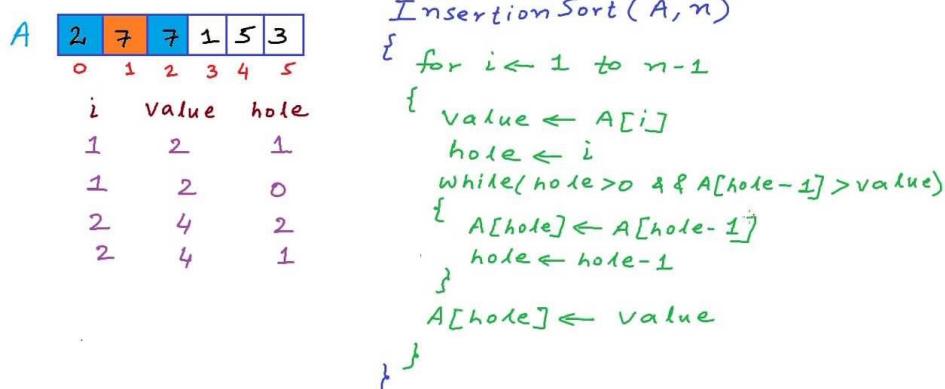
        for(int j=i+1;j<5;j++)
        {
            if(arr[j]<arr[min])
            {
                min=j;
            }
        }

        if(min!=i)
        {
            int temp=arr[min];
            arr[min] = arr[i];
            arr[i] = temp;
        }
    }
}

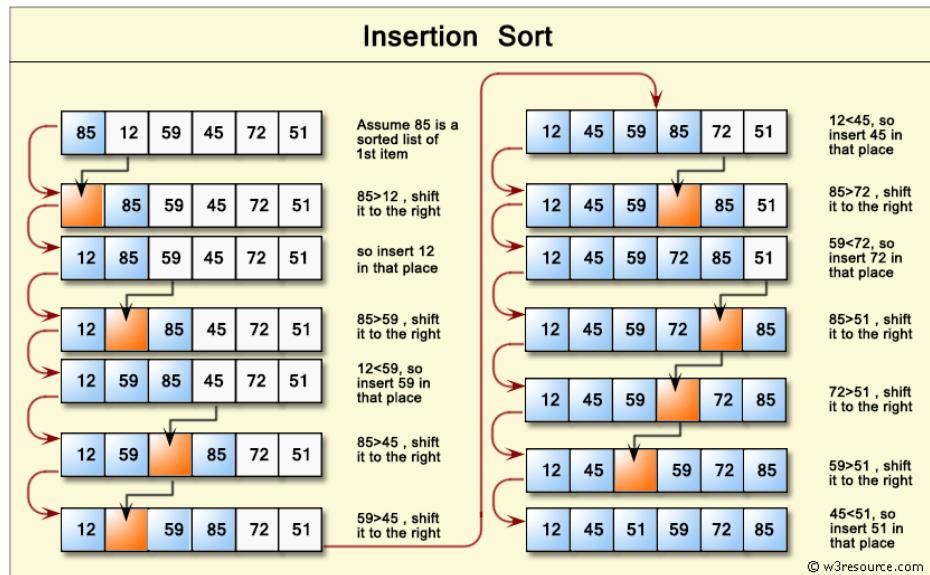
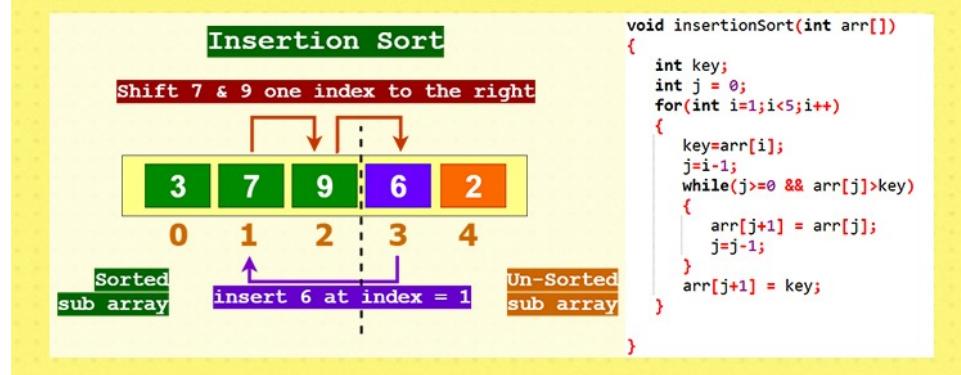
```

- Tanmay Sakpal

Insertion Sort



INSERTION SORT ALGORITHM



(Non-repetitive) insertion sort

- The pseudocode:

```

insertionSort(a non-repetitive array A, the array length n, an index cutoff < n)
// at any time, A[1..cutoff] is sorted and A[cutoff+1..n] is unsorted
if A[cutoff+1] < A[1..cutoff]
    let p be 1
else
    find p such that A[p-1] < A[cutoff+1] < A[p]
    insert A[cutoff+1] to A[p] and shift A[p..cutoff] to A[(p+1)..(cutoff+1)]
if cutoff+1 < n
    insertionSort(A, n, cutoff+1)
  
```



- What if A is repetitive?

Bubble Sort

BUBBLE SORT ALGORITHM

Bubble Sort
Larger elements shift to the end

```

void bubbleSort(int a[])
{
    for(int i=0; i<5; i++)
    {
        for(int j=0; j<(5-i-1); j++)
        {
            if(a[j]>a[j+1])
            {
                int temp = a[j];
                a[j] = a[j+1];
                a[j+1] = temp;
            }
        }
    }
}
  
```

OPTIMIZED BUBBLE SORT ALGORITHM

Bubble Sort
Larger elements shift to the end

```

void optimizedBubbleSort(int a[])
{
    int rounds=0;
    for(int i=0; i<5; i++)
    {
        rounds++;
        int flag = false;
        for(int j=0; j<(5-i-1); j++)
        {
            if(a[j]>a[j+1])
            {
                flag=true;
                int temp = a[j];
                a[j] = a[j+1];
                a[j+1] = temp;
            }
        }
        if(flag==false)
        {
            break;
        }
    }
    cout<<"No of rounds : "<<rounds<<endl;
}
  
```

In []: `#include <iostream>`
`using namespace std;`

```

void display(int* arr, int n);

void selectionSort(int* arr, int n);

void insertionSort(int* arr, int n);
  
```

```
void bubbleSort(int* arr, int n);

int main(){

    int option;
    int arraySize;

    do {
        cout << "Welcome to Learning C++ Data Structure and Algorithm" << endl;
        cout << "Topic: Sorting Algorithm" << endl << endl;

        cout << "1. Selection Sort" << endl;
        cout << "2. Insertion Sort" << endl;
        cout << "3. Bubble Sort" << endl;
        cout << "0. Exit" << endl;
        cout << endl;

        cout << "Select Your Option Number: ";
        cin >> option;
        cout << endl;

        if(option==0){
            break;
        }

        cout << "Enter your array size: ";
        cin >> arraySize;
        int arr[arraySize];
        cout << "\nEnter " << arraySize << " numbers in random order: " << endl;

        for(int i = 0; i < arraySize; i++){
            cin >> arr[i];
        }
        cout << endl;

        switch(option){
            //case 0://Quit
            //    break;
            case 1://Selection Sort
                selectionSort(arr, arraySize);
                display(arr, arraySize);
                break;
            case 2://Insertion Sort
                insertionSort(arr, arraySize);
                display(arr, arraySize);
                break;
            case 3://Bubble Sort
                bubbleSort(arr, arraySize);
                display(arr, arraySize);
                break;
            default:
                cout << "Enter the proper option number" << endl;
                cout << endl;
                break;
        }
    } while (option != 0);

    return 0;
}
```

```

void display(int* arr, int n){
    cout << "The Sorted Array: " << endl;
    for(int i=0; i<n; i++){
        cout << arr[i] << " ";
    }
    cout << endl << endl << endl;
}

void selectionSort(int* arr, int n){
    for(int i=0; i < n-1; i++){
        int min = i;
        for(int j=i+1; j<n; j++){
            if(arr[min] > arr[j]){
                min = j;
            }
        }
        if (min != i){
            int temp = arr[min];
            arr[min] = arr[i];
            arr[i] = temp;
        }
    }
}

void insertionSort(int* arr, int n){
    for(int i=1; i<=n-1; i++){
        int key = arr[i];
        int hole = i;
        while(hole>=0 && arr[hole-1]>key){
            arr[hole] = arr[hole-1];
            hole = hole-1;
        }
        arr[hole] = key;
    }
}

/*
void bubbleSort(int* arr, int n){
    for(int i=0; i<n-1; i++){
        for(int j=0; j<n-i-1; j++){
            if(arr[j] > arr[j+1]){
                int temp = arr[j];
                arr[j] = arr[j+1];
                arr[j+1] = temp;
            }
        }
    }
}
*/
//Optimized Bubble Sort Algorithm
void bubbleSort(int* arr, int n){
    for(int i=0; i<n-1; i++){
        bool swapFlag = false;
        for(int j=0; j<n-i-1; j++){
            if(arr[j] > arr[j+1]){
                swapFlag = true;
            }
        }
        if(swapFlag == false)
            break;
    }
}

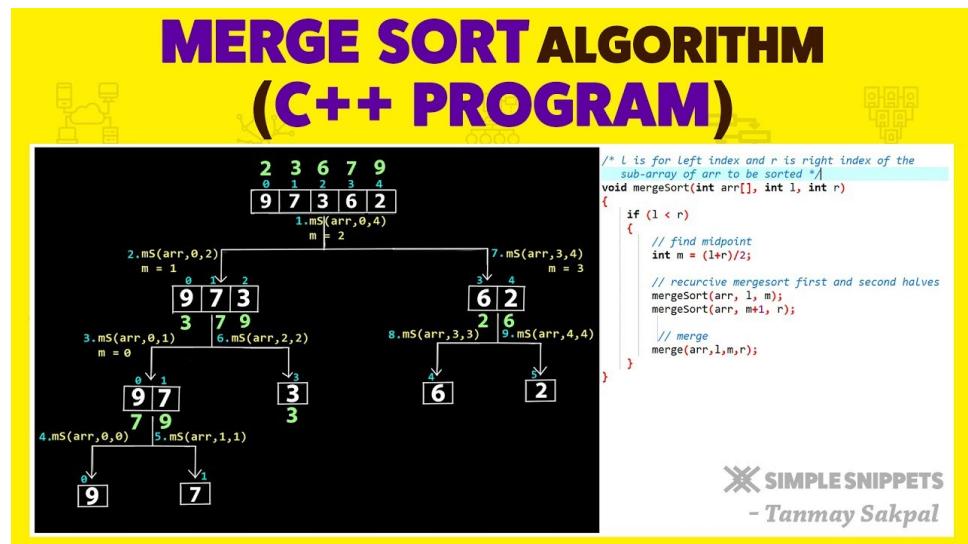
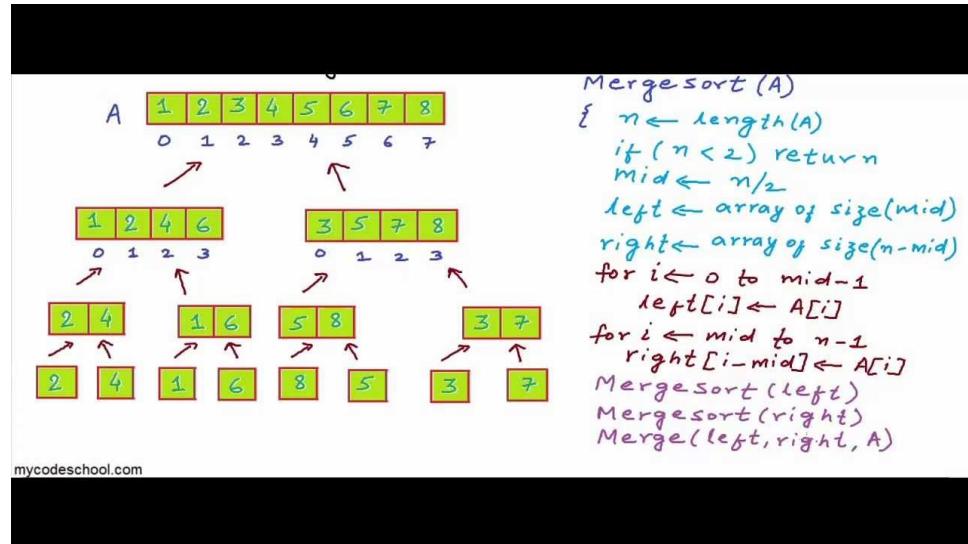
```

```

        int temp = arr[j];
        arr[j] = arr[j+1];
        arr[j+1] = temp;
    }
}
if(swapFlag==false){
    break;
}
}
}

```

Merge Sort

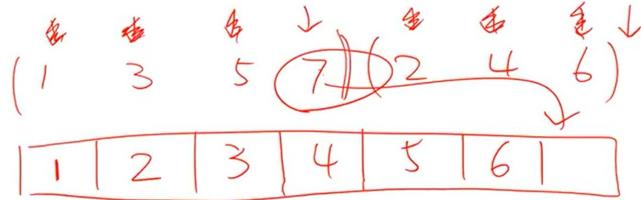


Mergesort (Merge sort): pseudocode

```

mergeSort(an array A, the array length n)
    let median be floor((1 + n) / 2)
    ✓ mergeSort(A[1..median], median) // now A[1..median] is sorted
    ✓ mergeSort(A[(median + 1)..n], n - median + 1) // now A[(median + 1)..n] is sorted
    ✓ merge A[1..median] and A[(median + 1)..n] // how?

```



Mergesort (Merge sort)

- Interestingly, insertion sort is a special way of running mergesort.
 - Not splitting the array into two halves.
 - Instead, splitting it into $A[1..n-1]$ and $A[n]$.
- Once we use the “smart split”, the efficiency is improved a lot!
 - Insertion sort: Roughly proportional to n^2 .
 - Merge sort: Roughly proportional to $n \log n$.
- A simple observation can make a huge difference!

```

In [ ]: #include <iostream>
using namespace std;

void display(int* arr, int n);

void mergeSort(int* arr, int l, int r, int n);

int main(){
    int arraySize;

    cout << "Welcome to Learning C++ Data Structure and Algorithm" << endl;
    cout << "Topic: MergeSorting Algorithm" << endl << endl;

    cout << "Enter your array size: ";
    cin >> arraySize;
    int arr[arraySize];
    cout << "\nEnter " << arraySize << " numbers in random order: " << endl;

    for(int i = 0; i < arraySize; i++){
        cin >> arr[i];
    }
    cout << endl;

    mergeSort(arr, 0, (arraySize-1), arraySize);
}

```

```
        display(arr, arraySize);

    return 0;
}

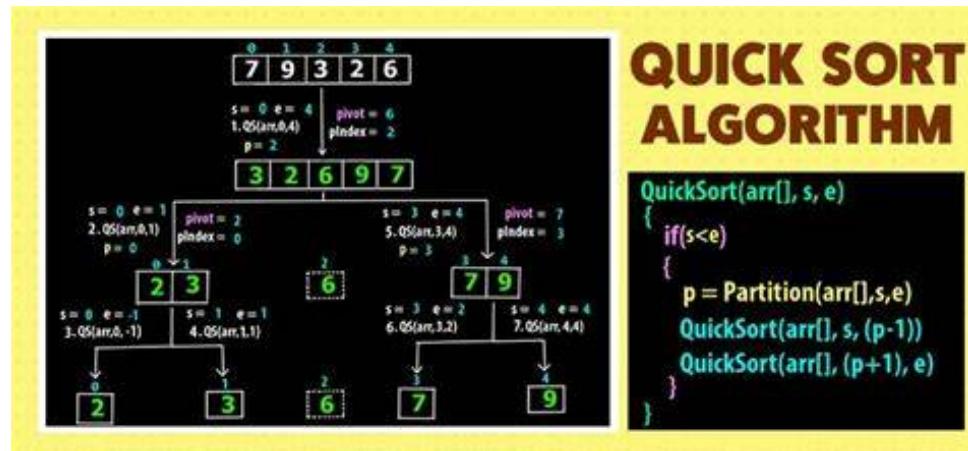
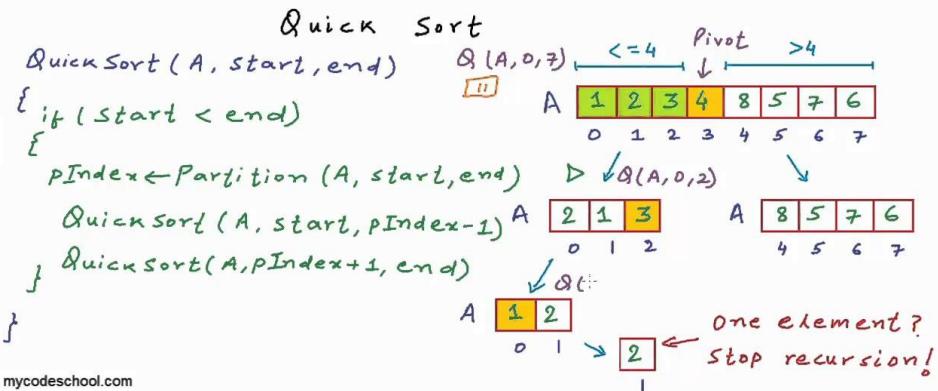
void display(int* arr, int n){
    cout << "The Sorted Array: " << endl;
    for(int i=0; i<n; i++){
        cout << arr[i] << " ";
    }
    cout << endl << endl << endl;
}

void merge(int* arr, int l, int m, int r, int n){
    int i = l;
    int j = m+1;
    int k = l;
    int temp[n];

    while(i<=m && j<=r){
        if(arr[i] <= arr[j]){
            temp[k] = arr[i];
            i++;
        }
        else {
            temp[k] = arr[j];
            j++;
        }
        k++;
    }
    if(i>m){
        while(j<=r){
            temp[k] = arr[j];
            j++;
            k++;
        }
    }
    else {
        while(i<=m){
            temp[k] = arr[i];
            i++;
            k++;
        }
    }
    for(k=l; k<=r; k++){
        arr[k] = temp[k];
    }
}

void mergeSort(int* arr, int l, int r, int n){
    if(l<r){
        int m = l + (r - l) / 2;
        mergeSort(arr, l, m, n);
        mergeSort(arr, m+1, r, n);
        merge(arr, l, m, r, n);
    }
}
```

Quick Sort



```

In [ ]: #include <iostream>
using namespace std;

void display(int* arr, int n);

void quickSort(int* arr, int l, int r);

int main(){

    int arraySize;

    cout << "Welcome to Learning C++ Data Structure and Algorithm" << endl;
    cout << "Topic: QuickSorting Algorithm" << endl << endl;

    cout << "Enter your array size: ";
    cin >> arraySize;
    int arr[arraySize];
    cout << "\nEnter " << arraySize << " numbers in random order: " << endl;

    for(int i = 0; i < arraySize; i++){
        cin >> arr[i];
    }
    cout << endl;

    quickSort(arr, 0, (arraySize-1));

    display(arr, arraySize);
}

```

```

    return 0;
}

void display(int* arr, int n){
    cout << "The Sorted Array: " << endl;
    for(int i=0; i<n; i++){
        cout << arr[i] << " ";
    }
    cout << endl << endl << endl;
}

int partition(int* arr, int l, int r){
    int pivot = arr[r];
    int pIndex = l;
    for(int i=l; i<r; i++){
        if(arr[i]<=pivot){
            int temp = arr[i];
            arr[i] = arr[pIndex];
            arr[pIndex] = temp;
            pIndex++;
        }
    }
    int temp = arr[r];
    arr[r] = arr[pIndex];
    arr[pIndex] = temp;

    return pIndex;
}

void quickSort(int* arr, int l, int r){
    if(l<r){
        int pIndex = partition(arr, l, r);
        quickSort(arr, l, pIndex-1);
        quickSort(arr, pIndex+1, r);
    }
}

```

Counting Sort

COUNTING SORT ALGORITHM WITH C++ CODE

Pseudocode - CountingSort()

```

Pseudocode-
CountingSort()
{
    1. take input_array[size]
    2. create output_array[size]
    3. take range (or no of unique elements)
    4. for(int i=0 to i<range)
        4.1 count_array[i] = 0
    5. for(int i=0 to i<size)
        5.1 ++count_array[input_array[i]]
    6. for(int i=1 to i<range)
        6.1 count_array[i] = count_array[i]+count_array[i-1]
    7. for(int i=0 to i<size)
        7.1 output_array[i] = count_array[input_array[i]] = input_array[i]
    8. for(i=0 to i<size)
        8.1 input_array[i] = output_array[i]
}

```



Key-indexed counting

Task: sort an array $a[]$ of N integers between 0 and $R-1$

Plan: produce sorted result in array $\text{temp}[]$

1. Count frequencies of each letter using key as index
2. Compute frequency cumulates
3. Access cumulates using key as index to find record positions.
4. Copy back into original array

```

int N = a.length;
int[] count = new int[R];

count frequencies → for (int i = 0; i < N; i++)
    count[a[i]+1]++;
compute cumulates → for (int k = 1; k < 256; k++)
    count[k] += count[k-1];
move records → for (int i = 0; i < N; i++)
    temp[count[a[i]+1]] = a[i];
copy back → for (int i = 0; i < N; i++)
    a[i] = temp[i];

```

	$a[]$	$\text{temp}[]$
0	a	a
1	a	a
2	b	b
3	b	b
4	b	b
5	c	c
6	d	d
7	d	d
8	e	e
9	f	f
10	f	f
11	f	f
	$\text{count}[]$	
a	2	
b	5	
c	6	
d	8	
e	9	
f	12	

6

```

In [ ]: #include <iostream>
#include <string>
using namespace std;

void display(int* arr, int n);

void countingSort(int* arr, int n);

int main(){

    cout << "Welcome to Learning C++ Data Structure and Algorithm" << endl;
    cout << "Topic: CountingSorting Algorithm" << endl << endl;

    int arraySize = 0;

    cout << "Enter your array size: ";
    cin >> arraySize;
    int arr[arraySize];
    cout << "\nEnter " << arraySize << " numbers in random order in range 0-9: " << endl;

    for(int i = 0; i < arraySize; i++){
        cin >> arr[i];
    }
    cout << endl;

    countingSort(arr, arraySize);

    display(arr, arraySize);

    return 0;
}

void display(int* arr, int n){
    cout << "The Sorted Array: " << endl;
    for(int i=0; i<n; i++){
        cout << arr[i] << " ";
    }
}

```

```

    }
    cout << endl << endl << endl;
}

void countingSort(int* arr, int n){
    int countArray[10] = {0};
    int outputArray[n];

    // to take a count of all elements in the input array
    for(int i=0; i<n; i++){
        ++countArray[arr[i]];
    }

    // update countArray to get the position in the outputArray
    for(int i=1; i<10; i++){
        countArray[i] = countArray[i] + countArray[i-1];
    }

    // placing inputArray elements into outputAarray
    for(int i=n-1; i>=0; i--){
        outputArray[--countArray[arr[i]]] = arr[i];
    }

    for(int i=0; i<n; i++){
        arr[i] = outputArray[i];
    }
}

```

Radix Sort

The image shows a yellow-themed graphic for the Radix Sort algorithm. On the left, there is a black-bordered box containing the C++ code for Radix Sort. On the right, the word "RADIX" is written in large, bold, purple letters. Below it, "SORT" is in a smaller, bold, purple font. Underneath that, "ALGORITHM" is in a bold, black font. To the right of "ALGORITHM", there is a small icon of a computer circuit board. Below the main title, "C++" is written in large, bold, purple letters. Underneath "C++", "PROGRAM" is written in a bold, black font. At the bottom right, there is a logo for "SIMPLE SNIPPETS" with a stylized 'X' icon.

```

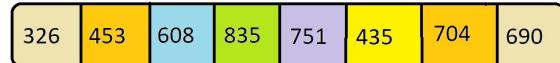
void CountingSort(int arr[], int size, int div)
{
    int output[size];
    int count[10] = {0};

    for (int i = 0; i < size; i++)
        count[(arr[i]/div)%10]++;
    for (int i = 1; i < 10; i++)
        count[i] += count[i - 1];
    for (int i = size - 1; i >= 0; i--)
    {
        output[count[(arr[i]/div)%10] - 1] = arr[i];
        count[(arr[i]/div)%10]--;
    }
    for (int i = 0; i < size; i++)
        arr[i] = output[i];
}

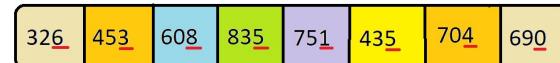
void RadixSort(int arr[], int size)
{
    int m = getMax(arr, size);
    for (int div = 1; m/div > 0; div *= 10)
        CountingSort(arr, size, div);
}

```

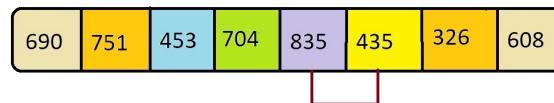
Let's Say The Given Array Is This :-



First, Consider The One's Place :-

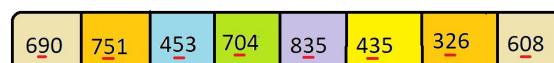


Now Sort the above array on the basis of digits on one's place

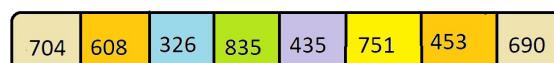


Observe That 835 has before 90 this is because it appeared before in the original array.

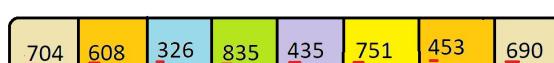
Now Consider the 10's Place :-



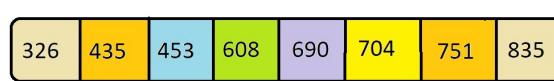
Now Sort the above array on the basis of digits on 10 's place



Now Consider the 100's Place :-



Now Sort the above array on the basis of digits on 100's place



Array Is Now Sorted

```
In [ ]: #include <iostream>
#include <string>
using namespace std;

void display(int* arr, int n);

void radixSort(int* arr, int n);

int main(){

    cout << "Welcome to Learning C++ Data Structure and Algorithm" << endl;
    cout << "Topic: RadixSorting Algorithm" << endl << endl;

    int arraySize = 0;
    int countSize = 10;

    cout << "Enter your array size: ";
    cin >> arraySize;
    int arr[arraySize];
    cout << "\nEnter " << arraySize << " numbers in random order: " << endl;

    for(int i = 0; i < arraySize; i++){
        cin >> arr[i];
    }

    cout << "Your sorted array is: ";
    display(arr, arraySize);
}
```

```

        }

        cout << endl;

        radixSort(arr, arraySize);

        display(arr, arraySize);

        return 0;
    }

void display(int* arr, int n){
    cout << "The Sorted Array: " << endl;
    for(int i=0; i<n; i++){
        cout << arr[i] << " ";
    }
    cout << endl << endl << endl;
}

void countSort(int* arr, int n, int pos){
    int countArray[10] = {0}; // initialize with "0"
    int outputArray[n];

    // to take a count of all elements in the input array
    for(int i=0; i<n; i++){
        ++countArray[(arr[i]/pos) % 10];
    }

    // update countArray to get the position in the outputArray
    for(int i=1; i<10; i++){
        countArray[i] = countArray[i] + countArray[i-1];
    }

    // placing inputArray elements into outputAarray
    for(int i=n-1; i>=0; i--){
        outputArray[--countArray[(arr[i]/pos) % 10]] = arr[i];
    }

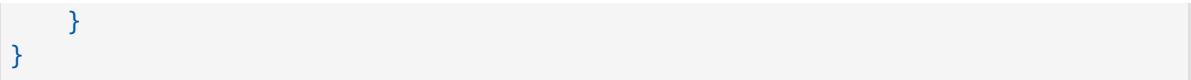
    for(int i=0; i<n; i++){
        arr[i] = outputArray[i];
    }
}

int getMax(int* arr, int n){
    int maxNumber = arr[0];
    for(int i=1; i<n; i++){
        if(maxNumber<arr[i]){
            maxNumber = arr[i];
        }
    }
    return maxNumber;
}

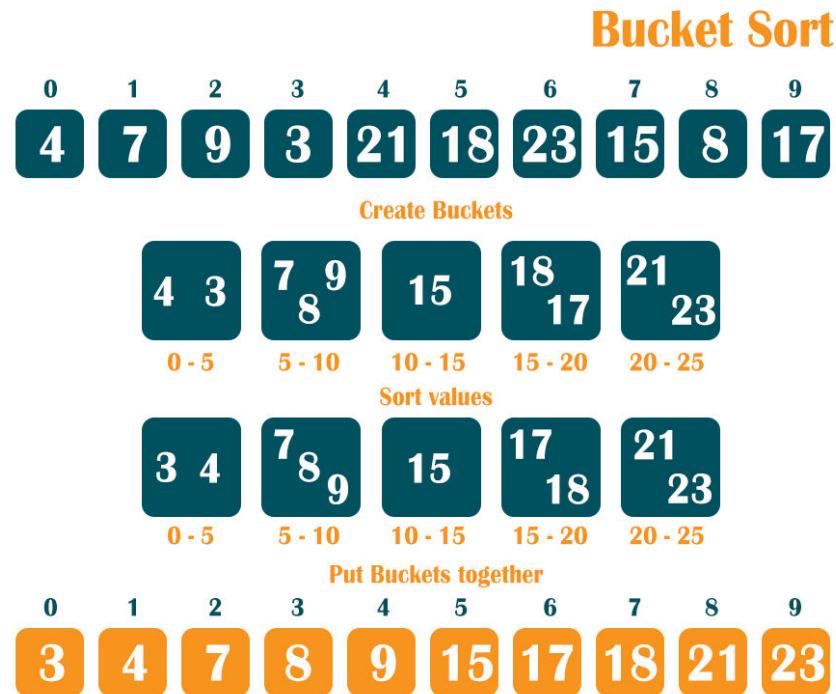
void radixSort(int* arr, int n){
    int maxNumber = getMax(arr, n);

    for(int pos=1; maxNumber/pos>0; pos *= 10){
        countSort(arr, n, pos);
    }
}

```



Bucket Sort



Analysis of Bucket Sort

- Bucket-Sort(A)
1. Let B[0....n-1] be a new array
 2. n = length[A]
 3. for i = 0 to n-1
 4. make B[i] an empty list
 5. for i = 1 to n
 6. do insert A[i] into list B[floor of n A[i]]
 7. for i = 0 to n-1
 8. do sort list B[i] with Insertion-Sort
 9. Concatenate lists B[0], B[1],...,B[n-1] together in order
- Step 5 and 6 takes O(n) time

Step 7 and 8 takes O(n log(n/k)) time

Step 9 takes O(k) time

In total Bucket sort takes : $O(n)$ (if $k=\Theta(n)$)

Shell Sort

SHELL SORT ALGORITHM (C++ PROGRAM)

Shell Sort

Swap elements which are at interval positions from each other

$n(\text{size of array}) = 9$

$\text{gap/interval} = n/2 = 4$

170	75	802	90	66	2	24	45	7
0	1	2	3	4	5	6	7	8

Compare these 2 elements that are at gap/interval distance form each other & perform swapping/shifting based on bubble or insertion sort

```
void ShellSort(int arr[], int n)
{
    for(int gap=n/2;gap>0;gap /= 2)
    {
        for(int j = gap;j<n;j+=1)
        {
            int temp = arr[j];
            int i=0;

            for(i=j;(i>=gap)&&(arr[i-gap]>temp);i-=gap)
            {
                arr[i] = arr[i-gap];
            }
            arr[i] = temp;
        }
    }
}
```

 SIMPLE SNIPPETS

```
In [ ]: #include <iostream>
#include <string>
using namespace std;

void display(int* arr, int n);

void shellSort(int* arr, int n);

int main(){

    cout << "Welcome to Learning C++ Data Structure and Algorithm" << endl;
    cout << "Topic: shellSorting Algorithm" << endl << endl;

    int arraySize = 0;
    int countSize = 10;

    cout << "Enter your array size: ";
    cin >> arraySize;
    int arr[arraySize];
    cout << "\nEnter " << arraySize << " numbers in random order: " << endl;

    for(int i = 0; i < arraySize; i++){
        cin >> arr[i];
    }
    cout << endl;

    shellSort(arr, arraySize);

    display(arr, arraySize);

    return 0;
}

void display(int* arr, int n){
    cout << "The Sorted Array: " << endl;
    for(int i=0; i<n; i++){
        cout << arr[i] << " ";
    }
    cout << endl << endl << endl;
}
```

```
void countSort(int* arr, int n, int pos){
    int countArray[10] = {0}; // initialize with "0"
    int outputArray[n];

    // to take a count of all elements in the input array
    for(int i=0; i<n; i++){
        ++countArray[(arr[i]/pos) % 10];
    }

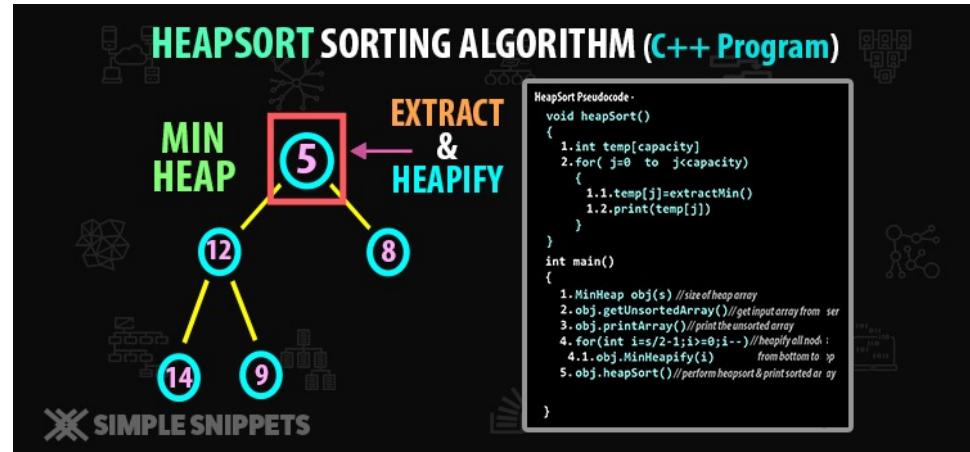
    // update countArray to get the position in the outputArray
    for(int i=1; i<10; i++){
        countArray[i] = countArray[i] + countArray[i-1];
    }

    // placing inputArray elements into outputAarray
    for(int i=n-1; i>=0; i--){
        outputArray[--countArray[(arr[i]/pos) % 10]] = arr[i];
    }

    for(int i=0; i<n; i++){
        arr[i] = outputArray[i];
    }
}

void shellSort(int* arr, int n){
    for(int gap=n/2; gap>0; gap/=2){
        for(int j=gap; j<n; j++){
            for(int i=j-gap; i>=0; i-=gap){
                if(arr[i+gap] > arr[i]){
                    break;
                }
                else{
                    int temp = arr[i];
                    arr[i] = arr[i+gap];
                    arr[i+gap] = temp;
                }
            }
        }
    }
}
```

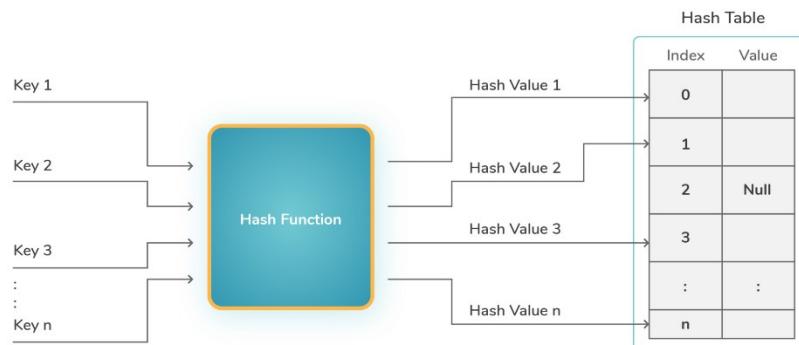
Heap Sort



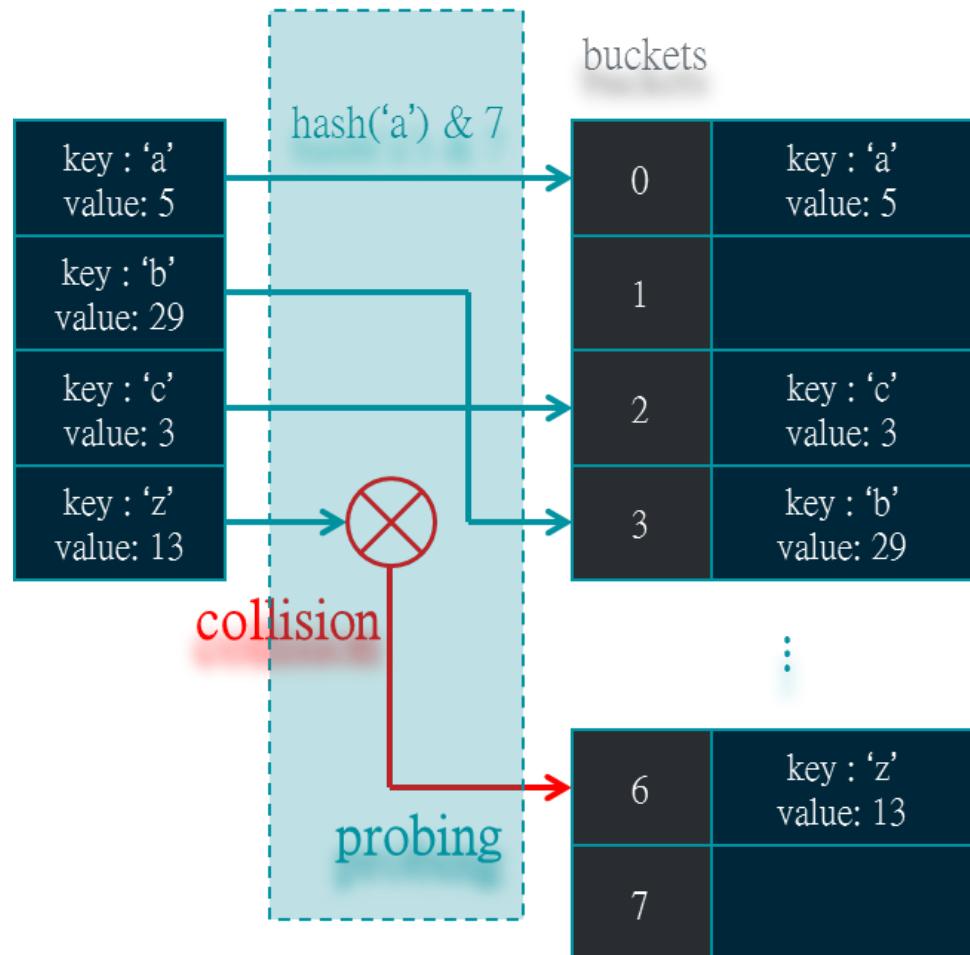
Hashing

Hash Table

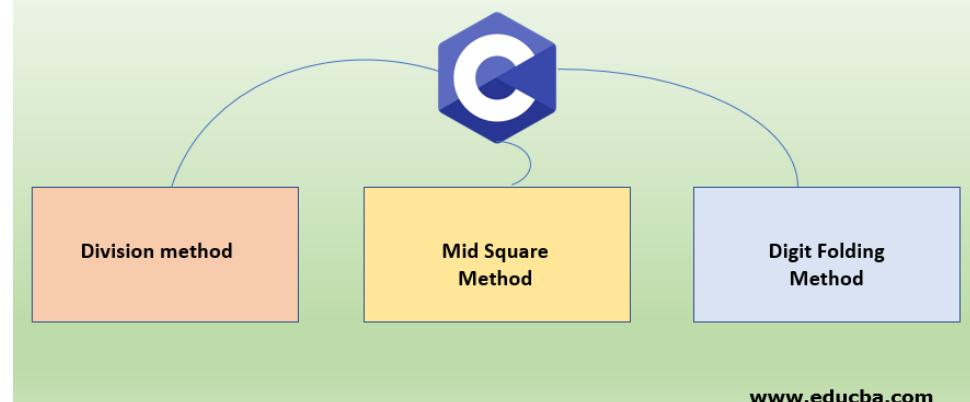
InterviewBit



Hash Function



Hashing function in C

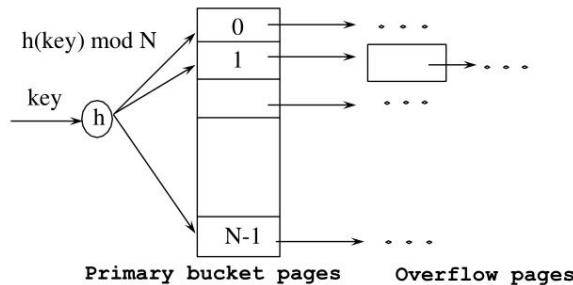


Static Hash



Static Hashing

- # primary pages fixed, allocated sequentially, never de-allocated; overflow pages if needed.
- $h(k) \text{ MOD } N =$ bucket to which data entry with key k belongs. ($N = \# \text{ of buckets}$)



Static Hashing

- ❖ A **bucket** is a unit of storage containing one or more records (a bucket is typically a disk block).
- ❖ In a **hash file organization** we obtain the bucket of a record directly from its search-key value using a **hash function**.
- ❖ Hash function h is a function from the set of all search-key values K to the set of all bucket addresses B .
- ❖ Hash function is used to locate records for access, insertion as well as deletion.
- ❖ Records with different search-key values may be mapped to the same bucket; thus entire bucket has to be searched sequentially to locate a record.

2

-- Memo End --