

A Theoretical and Empirical Investigation into the Equivalence of Graph Neural Networks and the Weisfeiler-Leman Algorithm

From the faculty of Mathematics, Physics, and Computer Science approved for the purpose of obtaining
the academic degree of Bachelor of Sciences.

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1 Introduction

Yet to come!

2 Related Work

Yet to come!

3 Preliminaries

First, we introduce a couple of notions and definitions that will be used throughout this thesis. In particular, the definitions will play a crucial role in the theoretical part that follows. We start with some general notations, introduce a general graph definition, and familiarize the reader with the Weisfeiler-Leman algorithm. We will then introduce each framework independently, first the 1-WL+NN and then GNN. At the end, we will briefly introduce important properties of collections of functions computed by both methods.

3.1 General Notation

We first introduce a couple of notations and definitions that will be used throughout the thesis. With $[n]$, we denote the set $\{1, \dots, n\} \subset \mathbb{N}$ for any $n \in \mathbb{N}$ and with $\{\!\{ \dots \}\!\}$ we denote a multiset which is formally defined as a 2 tuple (X, m) with X being a set of all unique elements and $m : X \rightarrow \mathbb{N}_{\geq 1}$ a mapping that maps every element in X to its number of occurrences in the multiset.

3.2 Graphs

A graph G is a 3-tuple $G := (V, E, l)$ that consists of the set of all nodes V , the set of all edges $E \subseteq V \times V$ and a label function $l : M \rightarrow \Sigma$ with M being either $V, V \cup E$ or E and $\Sigma \subset \mathbb{N}$ a finite alphabet. Moreover, let \mathcal{G} be the set of all finite graphs. Note that our definition of the labeling function allows graphs with labels either on the nodes only, on the edges only, or on both the nodes and the edges. In some cases we refer to the values assigned by l as features instead of labels, however this is usually only the case when Σ is multidimensional. In addition, although we have defined it this way, the labeling function is optional, and in cases where no labeling function is given, we add the trivial labeling function $f_1 : V(G) \rightarrow \{1\}$.

Further, G can be either directed or undirected, depending on the definition of E , where $E \subseteq \{(v, u) \mid v, u \in V\}$ defines a directed and $E \subseteq \{(v, u), (u, v) \mid v, u \in V, v \neq u\}$ such that for every $(v, u) \in E$ also $(u, v) \in E$ defines an undirected graph. Additionally, we will use the notation $V(G)$ and $E(G)$ to denote the set of nodes of G and the set of edges of G respectively, as well as l_G to denote the label function of G . With $\mathcal{N}(v)$ for $v \in V(G)$ we denote the set of neighbors of v defined as $\mathcal{N}(v) := \{u \mid (u, v) \in E(G)\}$.

A coloring of a Graph G is a function $C : V(G) \rightarrow \mathbb{N}$ that assigns each node in the graph a color (here a positive integer). Further, a coloring C induces a partition P on the set of nodes, for which we define C^{-1} being the function that maps each color $c \in \mathbb{N}$ to its class of nodes with $C^{-1}(c) = \{v \in V(G) \mid C(v) = c\}$. In addition, we define $h_{G,C}$ as the histogram of graph G with coloring C , that maps every color in the image of C under $V(G)$ to the number of occurrences. In detail, $\forall c \in \mathbb{N} : h_{G,C}(c) := |\{v \in V(G) \mid C(v) = c\}| = |C^{-1}(c)|$

Permutation-invariance and -equivariance

We use S_n to denote the symmetric group over the elements $[n]$ for any $n \in \mathbb{N}_{\geq 1}$. S_n consists of all permutations over these elements. Let G be a graph with $V(G) = [n]$, applying a permutation $\pi \in S_n$ on G , is defined as $G_\pi := \pi \cdot G$ where $V(G_\pi) = \{\pi(1), \dots, \pi(n)\}$ and $E(G_\pi) = \{(\pi(v), \pi(u)) \mid (v, u) \in E(G)\}$. We will now introduce two key concepts for classifying functions on graphs.

Definition 1 (Permutation Invariant). Let $f : \mathcal{G} \rightarrow \mathcal{X}$ be an arbitrary function, then f is *permutation-invariant* if and only if for all $G \in \mathcal{G}$ where $n_G := |V(G)|$ and for every $\pi \in S_{n_G}$: $f(G) = f(\pi \cdot G)$.

Definition 2 (Permutation Equivariant). Let $f : \mathcal{G} \rightarrow \mathcal{X}$ be an arbitrary function, then f is *permutation-equivariant* if and only if for all $G \in \mathcal{G}$ where $n_G := |V(G)|$ and for every $\pi \in S_{n_G}$: $f(G) = \pi^{-1} \cdot f(\pi \cdot G)$.

3.3 Weisfeiler and Leman Algorithm

The Weisfeiler-Leman algorithm consists of two main parts, first the coloring algorithm and second the graph isomorphism test. We will introduce them in this section.

The Weisfeiler-Leman graph coloring algorithm

The 1-WL algorithm computes a node coloring of its input graph in each iteration. A color for a node is computed using only the coloring of its neighbors and the node itself. The algorithm will continue as long as it has not converged, and returns the final coloring of the graph.

Definition 3 (1-WL Algorithm). Let $G = (V, E, l)$ be a graph, then in each iteration i , the 1-WL computes a node coloring $C_i : V(G) \rightarrow \mathbb{N}$. In iteration $i = 0$, the initial coloring is $C_0 = l$ or if l is non-existing $\forall v \in V(G) : C_0(v) = c$ for an arbitrary constant $c \in \mathbb{N}$. For $i > 0$, the algorithm assigns a color to $v \in V(G)$ as follows:

$$C_i(v) = \text{RELABEL}(C_{i-1}(v), \{\{C_{i-1}(u) \mid u \in \mathcal{N}(v)\}\}),$$

where RELABEL injectively maps the above pair to a unique, previously not used, natural number. The algorithm terminates when the number of colors between two iterations does not change, meaning the algorithm terminates after iteration i if the following condition is satisfied:

$$\forall v, w \in V(G) : C_i(v) = C_i(w) \iff C_{i+1}(v) = C_{i+1}(w).$$

Upon terminating we define $C_\infty := C_i$ as the stable coloring, such that $1\text{-WL}(G) := C_\infty$.

The colorings computed in each iteration always converge to the final one, such that the algorithm always terminates. In more detail, Grohe [2017] showed that it always holds after at most $|V(G)|$ iterations. For an illustration of this coloring algorithm, see Figure 2. Moreover, based on the work of Paige and Tarjan [1987] about efficient refinement strategies, Cardon and Crochemore [1982] proved that the stable coloring C_∞ can be computed in time $\mathcal{O}(|V(G)| + |E(G)| \cdot \log |V(G)|)$.

The Weisfeiler-Leman Graph Isomorphism Test

Definition 4 (1-WL Isomorphism Test). To determine if two graphs $G, H \in \mathcal{G}$ are non-isomorphic ($G \not\cong H$), one applies the 1-WL coloring algorithm on both graphs “in parallel” and checks after each iteration if the occurrences of each color are equal, else the algorithm would terminate and conclude non-isomorphic. Formally, the algorithm concludes non-isomorphic in iteration i if there exists a color c such that:

$$|\{v \in V(G) \mid c = C_i(v)\}| \neq |\{v \in V(H) \mid c = C_i(v)\}|.$$

Note that this test is only sound and not complete for the *graph isomorphism problem*. Counterexamples where the algorithm fails to distinguish non-isomorphic graphs can be easily constructed, see Figure 1 which was discovered and proven by Cai et al. [1992].

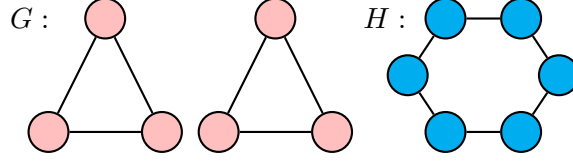


Figure 1: An example of two graphs G and H that are non-isomorphic but cannot be distinguished by the 1-WL isomorphism test.

3.4 1-WL+NN

As seen in the previous section, the 1-WL algorithm is quite powerful in identifying substructures. With the 1-WL+NN framework, we define functions that utilize this structural information to derive further application-specific insights. We do this by combining well-known machine learning techniques and the algorithm.

Definition 5 (1-WL+NN). We say a function \mathcal{B} working over a set of graphs $\mathcal{X} \subseteq \mathcal{G}$ is computable by 1-WL+NN, if it can be decomposed into a multilayer perceptron MLP, an encoding function f_{enc} and the 1-WL algorithm. In particular, when dealing with graph task \mathcal{B} can be decomposed as:

$$\mathcal{B} : \mathcal{X} \rightarrow \mathbb{R}, G \mapsto \text{MLP} \circ f_{\text{enc}}(\{\{1\text{-WL}(G)(v) \mid v \in V(G)\}\}),$$

where $1\text{-WL}(G)$ is the coloring computed by the 1-WL algorithm when applied on G . For node or edge tasks we define it as below:

$$\mathcal{B} : \mathcal{X} \rightarrow \mathcal{Y}, G \mapsto G' := (V(G), E(G), \text{MLP} \circ 1\text{-WL}(G)),$$

where a graph G is mapped to a structurally identical graph G' with the key difference being that the label function is replaced with the coloring computed by the 1-WL algorithm and further processed by a MLP.

As a concrete example of a collection of functions computable by 1-WL+NN we will introduce the collection of functions working from \mathcal{X} to \mathbb{R} that use the so called *counting-encoding* function as its encoding function. We will define this particular encoding function in the following:

Definition 6 (Counting Encoding Functions). The *counting-encoding* function $f_{\text{count},k}$ is parametrized by a $k \in \mathbb{N}$ and works as follows:

$$f_{\text{count},k}(\{\{C(v) \mid v \in V(G)\}\}) := \begin{bmatrix} |\{v \in V(G) \mid C(v) = 1\}| \\ \vdots \\ |\{v \in V(G) \mid C(v) = k\}| \end{bmatrix} \in \mathbb{N}^k.$$

It maps the occurrences of the colors $c \in [k]$ of its input to a vector over \mathbb{N}^k , such that each i .th component of this vector encodes the number of occurrences of the color i in the input.

To illustrate how this encoding function works and why we coined it *counting-encoding*, we will quickly introduce an example graph G . In Figure 2, we give a visual representation of G and its stable coloring after applying the 1-WL algorithm to it. The *counting-encoding* function $f_{\text{count},k}$ counts through all colors $i \in [k]$ and sets each i .th component of the output vector to the number of occurrences in the final coloring. Therefore, the respective color histogram $h_{G,C_\infty} = \{\{2, 2, 3, 4\}\}$ of G is being mapped to $(0, 2, 1, 1)^T \in \mathbb{N}^4$ by $f_{\text{count},4}$ and to $(0, 2)^T \in \mathbb{N}^2$ by $f_{\text{count},2}$ for example, since color 2 appears two times, while color 3 and 4 occur only once.

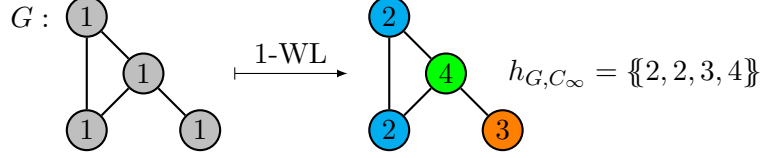


Figure 2: An example of the final coloring computed by applying the 1-WL algorithm on the graph G . The graph G consists of 4 nodes with all their labels being initially set to 1. Note that each label corresponds to a color, which we have also plotted for illustration purposes.

3.5 Graph Neural Networks (Message Passing)

A Graph Neural Network (GNN) is a composition of multiple layers, where each layer computes a new feature for each node and edge. The GNN layer thus technically obtains a new graph that is structurally identical to the previous one, but contains new feature information. After an input graph has been passed through all layers, there can be an additional final function, aggregating the computed information into a fixed size output. With this, it is possible to apply a GNN to every graph, regardless of its size, as the “computation” will only take place on the nodes and edges of the graph.

Note that in the following we will restrict the definition to only consider node features, however, one can easily extend it to also include edge features.

Definition 7 (Graph Neural Network). Let $G = (V, E, l)$ be an arbitrary graph. A Graph Neural Network (GNN) is a composition of multiple layers where each layer t is represented by a function $f^{(t)}$ that works over the set of nodes $V(G)$. To begin with, we need a function $f^{(0)} : V(G) \rightarrow \mathbb{R}^{1 \times d}$ that is consistent with l , that translates all labels into a vector representation. Further, for every $t > 0$, $f^{(t)}$ is of the format:

$$f^{(t)}(v) = f_{\text{merge}}^{W_{1,t}}(f^{(t-1)}(v), f_{\text{agg}}^{W_{2,t}}(\{f^{(t-1)}(w) \mid w \in \mathcal{N}(v)\})),$$

where $f_{\text{merge}}^{W_{1,t}}$ and $f_{\text{agg}}^{W_{2,t}}$ are arbitrary differentiable functions with $W_{1,t}$ and $W_{2,t}$ their respective parameters. Additionally, $f_{\text{agg}}^{W_{2,t}}$ has to be permutation-invariant.

Depending on the objective, whether the GNN is tasked with a graph or a node task, the last layer differs. In the case of graph tasks, we add a permutation-invariant aggregation function to the end, here called READOUT, working as follows:

$$\text{READOUT}(\{f^{(t)}(v) \mid v \in V(G)\}),$$

that aggregates over every node and computes a fixed-size output vector for the entire graph, e.g. a label for graph classification. In order to ensure that we can train the GNN in an

end-to-end fashion, we require READOUT to be also differentiable. Let \mathcal{A} be an instance of the described GNN framework. Further, let $K \in \mathbb{N}$ be the number of layers of the GNN, \mathcal{G} the set of all graphs, \mathcal{Y} the task-specific output set (e.g. labels of a classification task), then the overall function computed by \mathcal{A} is:

$$\mathcal{A} : \mathcal{G} \rightarrow \mathcal{Y}, G \mapsto \text{READOUT}(\{\{f^{(K)}(v) \mid v \in V(G)\}\}),$$

if \mathcal{A} is configured for a graph task, otherwise:

$$\mathcal{A} : \mathcal{G} \rightarrow \mathcal{Y}, G \mapsto G' := (V(G), E(G), f^{(K)}),$$

where a graph G is mapped to a structurally identical graph G' with the key difference being the label function is $f^{(K)}$.

Note that, as we require all aggregation functions to be permutation-invariant, the total composition \mathcal{A} is permutation-invariant, and with similar reasoning, it is also differentiable. This enables us to train \mathcal{A} like any other machine learning method in an end-to-end fashion, regardless of the underlying encoding used for graphs. This definition and use of notation are inspired by Morris et al. [2019] and Xu et al. [2019].

To demonstrate what kind of functions are typically used, we provide functions used by Hamilton et al. [2017] for a node classification:

$$\begin{aligned} f_{\text{merge}}^{W_{1,t}}(v) &= \text{ReLU}(W_{\text{merge}} \cdot \text{concat}(f^{(t-1)}(v), f_{\text{agg}}^{W_{2,t}}(v))) \\ f_{\text{agg}}^{W_{2,t}}(v) &= \text{MAX}(\{\text{ReLU}(W_{\text{pool}} \cdot f^{(t-1)}(u) + b) \mid u \in \mathcal{N}(v)\}) \end{aligned}$$

where ReLU is a rectified linear unit element wise activation function, MAX the element-wise max operator; W_{merge} , W_{pool} are trainable matrices, b a trainable vector and concat the concatenation function.

3.6 Important for later

In this section, we introduce a formal definition of multilayer perceptron as it is required in a later proof, as well as the $\simeq_{1\text{WL}}$ relation. Additionally, two very important properties for collections of functions.

Definition 8 (Multilayer Perceptron). Multilayer perceptrons are a class of functions from \mathbb{R}^n to \mathbb{R}^m , with $n, m \in \mathbb{N}$. In this thesis, we define a multilayer perceptron as a finite sequence, such that a multilayer perceptron MLP is defined as $\text{MLP} := (\text{MLP})_{i \in [k]}$ where k is the number of layers. For every $i \in [k]$, the i .th layer of the MLP is the i .th item in the finite sequence $(\text{MLP})_i$. Further, all layers are recursively defined as:

$$\begin{aligned} (\text{MLP})_0(v) &:= v \\ (\text{MLP})_{i+1}(v) &:= \sigma_i(W_i \cdot (\text{MLP})_i(v) + b_i), \quad \forall i \in [k-1] \end{aligned}$$

where σ_i is an element wise activation function, W_i is the weight matrix and b_i the bias vector of layer i . Note, that for each W_i , the succeeding W_{i+1} must have the same number of columns as W_i has rows, in order to be well-defined. Similarly, for every layer i , W_i and b_i have to have the same number of rows. Following this definition, when applying a MLP on input $v \in \mathbb{R}^n$ it is $\text{MLP}(v) := (\text{MLP})_k(v)$.

Definition 9 (1-WL Relation). For any graphs $G, H \in \mathcal{X}$ we will denote $G \simeq_{1\text{WL}} H$ if the 1-WL isomorphism test can not distinguish both graphs. Note that due to the soundness of this algorithm, if $G \not\simeq_{1\text{WL}} H$, we always can conclude that $G \neq H$.

The $\simeq_{1\text{WL}}$ relation can further be classified as an equivalence relation, as it is reflexive, symmetric and transitive. With this, we introduce a notation of its equivalence classes. Let $\mathcal{X} \subseteq \mathcal{G}$ and $G \in \mathcal{X}$, then we denote with $\mathcal{X}/\simeq_{1\text{WL}}(G) := \{G' \in \mathcal{X} \mid G \simeq_{1\text{WL}} G'\}$ its equivalence class.

Definition 10 (1-WL-Discriminating). Let \mathcal{C} be a collection of permutation invariant functions from \mathcal{X} to \mathbb{R} . We say \mathcal{C} is **1-WL-Discriminating** if for all graphs $G_1, G_2 \in \mathcal{X}$ for which the 1-WL isomorphism test concludes non-isomorphic ($G_1 \not\simeq_{1\text{WL}} G_2$), there exists a function $h \in \mathcal{C}$ such that $h(G_1) \neq h(G_2)$.

Definition 11 (GNN-Approximating). Let \mathcal{C} be a collection of permutation invariant functions from \mathcal{X} to \mathbb{R} . We say \mathcal{C} is **GNN-Approximating** if for all permutation-invariant functions \mathcal{A} computed by a GNN, and for all $\epsilon \in \mathbb{R}$ with $\epsilon > 0$, there exists $h_{\mathcal{A}, \epsilon} \in \mathcal{C}$ such that $\|\mathcal{A} - h_{\mathcal{A}, \epsilon}\|_{\infty} := \sup_{G \in \mathcal{X}} |\mathcal{A}(G) - h_{\mathcal{A}, \epsilon}(G)| < \epsilon$

4 Theoretical Connection

This section is the main part of our theoretical investigation of the two frameworks. We will present 4 intriguing theorems, which will be proven separately afterwards. These results will form the basis for the empirical part that follows. The first two theorems will establish an equivalence between the two frameworks when the input set of graphs is finite. The last two theorems will go one step further and establish a connection for continuous functions computed by 1-WL+NN and GNNs and prove a somewhat weaker connection between them.

Throughout the first two theorems we will concentrate on a finite collection of graphs which we will denote with $\mathcal{X} \subset \mathcal{G}$.

Theorem 12 (Finite Case: “GNN \subseteq 1-WL+NN”). Let \mathcal{C} be a collection of functions from \mathcal{X} to \mathbb{R} computable by GNNs, then \mathcal{C} is also computable by 1-WL+NN.

Theorem 13 (Finite Case: “1-WL+NN \subseteq GNN”). Let \mathcal{C} be a collection of functions from \mathcal{X} to \mathbb{R} computable by 1-WL+NN, then \mathcal{C} is also computable by GNNs.

With these theorems, we showed the equivalence between both frameworks. Specifically, every function computed by 1-WL+NN is also computable by a GNN, and vice versa. Notice that, we did not leverage any constraints on the encoding of graphs throughout the first two theorems and their corresponding proves, but rather kept it general.

Having established a connection between the two frameworks on a finite subset of graphs, we wanted to further demonstrate the expressive power of 1-WL-Discriminating by investigating a connection between the two frameworks for continuous feature spaces and continuous functions. However, since the graph isomorphism problem is an inherently discrete problem, the 1-WL algorithm is only defined as a discrete function operating on discrete colors, such that extending the definition of the 1-WL algorithm to continuous functions is not very trivial and, to our knowledge, has not yet been widely investigated. We therefore assume in the proofs and the following theorems that there exists a continuous version of the 1-WL algorithm.

We define the set of graphs with continuous features using the following definition:

add that
graph
wise
classification
implies
node
classification

Definition 14. Let X be a compact subset of \mathbb{R} including 0. We decode graphs with n nodes as a matrix $G \in X^{n \times n}$, where $G_{i,i}$ decodes the label of node i for $i \in [n]$, and $G_{i,j}$ with $i \neq j \in [n]$ decodes an edge from node i to j and a corresponding edge features. We say that there is an edge between node i and j if and only if $G_{i,j} \neq 0$. Further, if G encodes an undirected graph, G is symmetric. For simplicity, we denote $\mathcal{X} := X^{n \times n}$ throughout the next two theorems and their respective proofs.

Theorem 15 (Continuous Case: “GNN \subseteq 1-WL+NN”). Let \mathcal{C} be a collection of continuous functions from \mathcal{X} to \mathbb{R} computable by 1-WL+NN. If \mathcal{C} is 1-WL-Discriminating, then there exists a collection of functions \mathcal{C}' computable by 1-WL+NN that is GNN-Approximating.

Theorem 16 (Continuous Case: “1-WL+NN \subseteq GNN”). Let \mathcal{C} be a collection of continuous functions from \mathcal{X} to \mathbb{R} that is GNN-Approximating, then \mathcal{C} is also 1-WL-Discriminating.

Since we only wanted to show the expressive power of 1-WL-Discriminating and made the major assumption of the existence of a continuous 1-WL algorithm, we have included the proofs of each theorem in the Appendix in subsubsection 1.2.

Immediately from the last theorem follows the corollary:

Corollary 17. There exists a collection \mathcal{C} of function from \mathcal{X} to \mathbb{R} computable by GNNs that is 1-WL-Discriminating.

Proof. The collection of all functions computable by GNNs is trivially GNN-Approximating, such that we can use Theorem 16 which concludes the proof. \square

By putting both theorems into perspective, we can now argue that even for continuous functions 1-WL+NN and GNNs can compute almost the same functions. Each framework can approximate the other framework arbitrarily well. From this we can conclude that the ability 1-WL-Discriminating is very powerful, so we can assume that 1-WL+NN is sufficiently expressive for the upcoming empirical part.

4.1 Proof of Theorem 12

We will prove Theorem 12 by introducing a couple of small lemmas, which combined prove the theorem. In detail, in Lemma 18 we show the existence of a collection computed by 1-WL+NN that is 1-WL-Discriminating. In Lemmas 19 to 21 we derive properties of 1-WL+NN functions we will use throughout Lemmas 22, 23 and 28 with which we prove the theorem. We took great inspiration for Lemmas 22, 23 and 28 from the proof presented in section 3.1 in the work of Chen et al. [2019].

Lemma 18. There exists a collection \mathcal{C} of functions from \mathcal{X} to \mathbb{R} computable by 1-WL+NN that is 1-WL-Discriminating.

Proof. We define f_c for $c \in \mathbb{N}$ as the encoding functions that returns the number of nodes colored as c . With this we can construct the collection of functions \mathcal{C} as follows:

$$\mathcal{C} := \{\mathcal{B}_c : \mathcal{X} \rightarrow \mathbb{R}, G \mapsto \text{MLP}_{\text{id}} \circ f_c(\{1\text{-WL}(G)(v) \mid v \in V(G)\}) \mid c \in \mathbb{N}\},$$

where MLP_{id} is a dummy multilayer perceptron that just returns its input. Since every function $\mathcal{B}_c \in \mathcal{C}$ is composed of the 1-WL algorithm, an encoding function and a multilayer perceptron, each function is computable by 1-WL+NN, and consequently also the whole collection.

Let $G_1, G_2 \in \mathcal{X}$ with $G_1 \not\approx_{1\text{WL}} G_2$. Further, let C_1, C_2 be the final colorings computed by the 1-WL algorithm when applied on G_1, G_2 respectively. Due to $G_1 \not\approx_{1\text{WL}} G_2$, there exists a color $c \in \mathbb{N}$ such that $h_{G_1, C_1}(c) \neq h_{G_2, C_2}(c)$. Such that $\mathcal{B}_c \in \mathcal{C}$ exists with $\mathcal{B}_c(G_1) \neq \mathcal{B}_c(G_2)$. \square

Lemma 19 (1-WL+NN Equivalence). Let \mathcal{B} be a function over \mathcal{X} computable by 1-WL+NN, then for every pair of graphs $G_1, G_2 \in \mathcal{X}$: if $G_1 \simeq_{1\text{WL}} G_2$ then $\mathcal{B}(G_1) = \mathcal{B}(G_2)$.

Proof. Let \mathcal{B} be an arbitrary function over \mathcal{X} computable by 1-WL+NN, then \mathcal{B} is composed as follows: $\mathcal{B}(\cdot) = \text{MLP} \circ f_{\text{enc}}\{\{1\text{-WL}(\cdot)(v) \mid v \in V(G)\}\}$. Further, let $G_1, G_2 \in \mathcal{X}$ be arbitrary graphs with $G_1 \simeq_{1\text{WL}} G_2$, then by definition of the relation $\simeq_{1\text{WL}}$ we know that $1\text{-WL}(G_1) = 1\text{-WL}(G_2)$. With this the equivalence follows immediately. \square

Lemma 20 (1-WL+NN Permutation Invariance). Let \mathcal{B} be a function over \mathcal{X} computable by 1-WL+NN, then \mathcal{B} is permutation invariant.

Proof. Let $G_1, G_2 \in \mathcal{X}$ be arbitrary graphs with $G_1 \simeq G_2$ and \mathcal{B} an arbitrary function computable by 1-WL+NN. Since the 1-WL algorithm is sound, we know that $G_1 \simeq G_2$ implies $G_1 \simeq_{1\text{WL}} G_2$. Using Lemma 19, we can therefore conclude that: $\mathcal{B}(G_1) = \mathcal{B}(G_2)$. \square

Lemma 21 (1-WL+NN Composition). Let \mathcal{C} be a collection of functions computable by 1-WL+NN. Further, let $h_1, \dots, h_n \in \mathcal{C}$ and MLP^\bullet an multilayer perceptron, then the function \mathcal{B} composed of $\mathcal{B}(\cdot) := \text{MLP}^\bullet(h_1(\cdot), \dots, h_n(\cdot))$ is also computable by 1-WL+NN.

Proof. Assume the above and let f_1, \dots, f_n be the encoding functions, as well as $\text{MLP}_1, \dots, \text{MLP}_n$ be the multilayer perceptrons used by h_1, \dots, h_n respectively. The idea of this proof is, we construct an encoding function f^* that maps a coloring C_∞ to a concatenation of the vectors obtained when applying each encoding function f_i individually. Additionally, we construct a multilayer perceptron MLP^* that takes in this concatenation of vectors and simulates all $\text{MLP}_1, \dots, \text{MLP}_n$ simultaneously on their respective section of the encoding vector of f^* , and applies afterwards the given MLP^\bullet on the concatenation of the output of all MLP_i 's. See Figure 3 for a sketch of the proof idea. A complete proof can be found in the Appendix, as this proof is very technical and not that interesting.

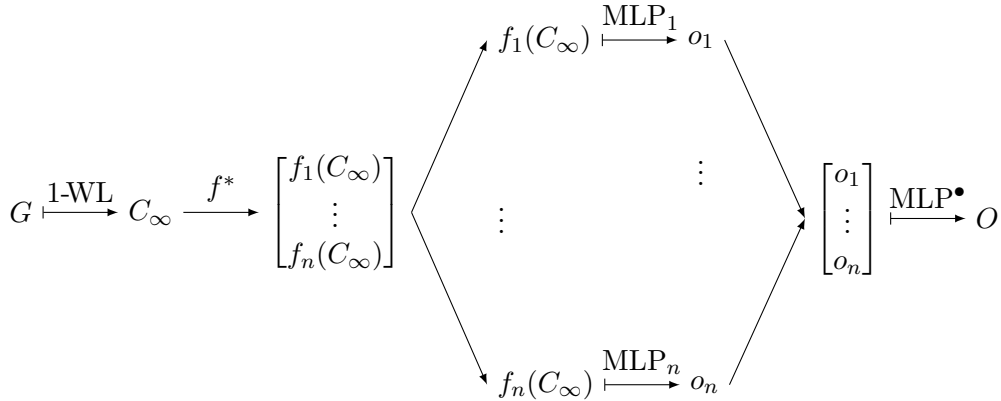


Figure 3: Proof idea for Lemma 21, how the constructed functions f^* and MLP^* will work on input G . Here we denote with C_∞ the final coloring computed by the 1-WL algorithm when applied on G . Further, we let o_i be the output computed by MLP_i on input $f_i(C_\infty)$.

□

Lemma 22. Let \mathcal{C} be a collection of functions from \mathcal{X} to \mathbb{R} computable by 1-WL+NN that is 1-WL-Discriminating. Then for all $G^* \in \mathcal{X}$, there exists a function h_{G^*} from \mathcal{X} to \mathbb{R} computable by 1-WL+NN, such that for all $G \in \mathcal{X}$: $h_{G^*}(G) = 0$ if and only if $G \simeq_{1\text{WL}} G^*$.

Proof. Assume the above. For any $G_1, G_2 \in \mathcal{X}$ with $G_1 \not\simeq_{1\text{WL}} G_2$ let $h_{G_1, G_2} \in \mathcal{C}$ be the function distinguishing them, with $h_{G_1, G_2}(G_1) \neq h_{G_1, G_2}(G_2)$. We define the function \bar{h}_{G_1, G_2} working over \mathcal{X} as follows:

$$\begin{aligned} \bar{h}_{G_1, G_2}(\cdot) &= |h_{G_1, G_2}(\cdot) - h_{G_1, G_2}(G_1)| \\ &= \max(h_{G_1, G_2}(\cdot) - h_{G_1, G_2}(G_1)) + \max(h_{G_1, G_2}(G_1) - h_{G_1, G_2}(\cdot)) \end{aligned} \quad (0.1)$$

Note, that in the formula above “ $h_{G_1, G_2}(G_1)$ ” is a fixed constant and the resulting function \bar{h}_{G_1, G_2} is non-negative. Let $G_1 \in \mathcal{X}$ now be fixed, we will construct the function h_{G_1} with the desired properties as follows:

$$h_{G_1}(x) = \sum_{G_2 \in \mathcal{X}, G_1 \not\simeq_{1\text{WL}} G_2} \bar{h}_{G_1, G_2}(x).$$

Since \mathcal{X} is finite, the sum is finite and therefore well-defined. Next, we will prove that for a fixed graph $G_1 \in \mathcal{X}$, the function h_{G_1} is correct on input $G \in \mathcal{X}$:

1. If $G_1 \simeq_{1\text{WL}} G$, then for every function \bar{h}_{G_1, G_2} of the sum with $G_1 \not\simeq_{1\text{WL}} G_2$, we know, using Lemma 19, that $\bar{h}_{G_1, G_2}(G)$ is equal to $\bar{h}_{G_1, G_2}(G_1)$ which is by definition 0, such that $h_{G_1}(G) = 0$.
2. If $G_1 \not\simeq_{1\text{WL}} G$, then $\bar{h}_{G_1, G}(G)$ is a summand of the overall sum, and since $\bar{h}_{G_1, G}(G) > 0$, we can conclude $h_{G_1}(G) > 0$ due to the non-negativity of each function \bar{h}_{G_1, G_2} .

This function can be encoded via a multilayer perceptron by replacing the max operator in Equation 0.1 by the activation function ReLU. Therefore, we can conclude with Lemma 21 that for every graph G , h_G is also 1-WL+NN computable. □

Lemma 23. Let \mathcal{C} be a collection of functions from \mathcal{X} to \mathbb{R} computable by 1-WL+NN so that for all $G^* \in \mathcal{X}$, there exists $h_{G^*} \in \mathcal{C}$ satisfying $h_{G^*}(G) = 0$ if and only if $G \simeq_{1\text{WL}} G^*$ for all $G \in \mathcal{X}$. Then for every $G^* \in \mathcal{X}$, there exists a function φ_{G^*} computable by 1-WL+NN such that for all $G \in \mathcal{X}$: $\varphi_{G^*}(G) = \mathbb{1}_{G \simeq_{1\text{WL}} G^*}$.

Proof. Assuming the above. Due to \mathcal{X} being finite, we can define for every graph G^* the constant:

$$\delta_{G^*} := \frac{1}{2} \min_{G \in \mathcal{X}, G \not\simeq_{1\text{WL}} G^*} |h_{G^*}(G)| > 0.$$

With this constant, we can use a so-called “bump” function working from \mathbb{R} to \mathbb{R} that will be similar to the indicator function. We define this function for parameter $a \in \mathbb{R}$ with $a > 0$ as:

$$\psi_a(x) := \max\left(\frac{x}{a} - 1, 0\right) + \max\left(\frac{x}{a} + 1, 0\right) - 2 \cdot \max\left(\frac{x}{a}, 0\right).$$

The interesting property of ψ_a is that it maps every value x to 0, except when x is being drawn from the interval $(-a, a)$. In particular, it maps x to 1 if and only if x is equal to 0. See Figure 4 in the Appendix for a plot of the relevant part of this function with exemplary values for a .

We use these properties to define for every graph $G^* \in \mathcal{X}$ the function $\varphi_{G^*}(\cdot) := \psi_{\delta_{G^*}}(h_{G^*}(\cdot))$. We will quickly demonstrate that this function is equal to the indicator function, for this let G^* be fixed and G , an arbitrary graph from \mathcal{X} , the input:

1. If $G \simeq_{1\text{WL}} G^*$, then $h_{G^*}(G) = 0$ resulting in $\varphi_{G^*}(G) = \psi_{\delta_{G^*}}(0) = 1$.
2. If $G \not\simeq_{1\text{WL}} G^*$ then $h_{G^*}(G) \neq 0$, such that $|h_{G^*}(G)| > \delta_{G^*}$ so that $h_{G^*}(G) \notin (-\delta_{G^*}, \delta_{G^*})$ resulting in $\varphi_{G^*}(G) = 0$.

Note that we can encode φ_{G^*} via a multilayer perceptron with a single layer, where δ_{G^*} is a constant and the max operator is replaced by the non-linear activation function ReLU. With Lemma 21 we can therefore conclude that φ_{G^*} is computable by 1-WL+NN for every graph $G^* \in \mathcal{X}$. \square

Lemma 24. Let \mathcal{C} be a collection of functions from \mathcal{X} to \mathbb{R} computable by 1-WL+NN so that for all $G^* \in \mathcal{X}$, there exists $\varphi_{G^*} \in \mathcal{C}$ satisfying $\forall G \in \mathcal{X} : \varphi_{G^*}(G) = \mathbb{1}_{G \simeq_{1\text{WL}} G^*}$, then every function computable by a GNN is also computable by 1-WL+NN.

Proof. Assume the above. For any function \mathcal{A} computed by an GNN that works over \mathcal{X} to \mathbb{R} , we show that it can be decomposed as follows for any $G \in \mathcal{X}$ as input:

$$\begin{aligned}
\mathcal{A}(G) &= \left(\frac{1}{|\mathcal{X}/\simeq_{1\text{WL}}(G)|} \sum_{G^* \in \mathcal{X}} \mathbb{1}_{G \simeq_{1\text{WL}} G^*} \right) \cdot \mathcal{A}(G) \\
&= \frac{1}{|\mathcal{X}/\simeq_{1\text{WL}}(G)|} \sum_{G^* \in \mathcal{X}} \mathcal{A}(G^*) \cdot \mathbb{1}_{G \simeq_{1\text{WL}} G^*} \\
&= \sum_{G^* \in \mathcal{X}} \frac{\mathcal{A}(G^*)}{|\mathcal{X}/\simeq_{1\text{WL}}(G^*)|} \cdot \varphi_{G^*}(G)
\end{aligned} \tag{0.2}$$

with $\mathcal{X}/\simeq_{1\text{WL}}(G^*)$ we denote the set of all graphs G over \mathcal{X} that are equivalent to G^* according to the $\simeq_{1\text{WL}}$ relation.

Since \mathcal{A} is permutation-invariant, and GNNs are at most as good as the 1-WL algorithm in distinguishing non-isomorphic graphs, we can use the fact that for every graph $G, H \in \mathcal{X}$ with $G \simeq_{1\text{WL}} H$: $\mathcal{A}(G) = \mathcal{A}(H)$. Therefore, we can decompose \mathcal{A} as stated in Equation 0.2 via a multilayer perceptron with $\frac{\mathcal{A}(G^*)}{|\mathcal{X}/\simeq_{1\text{WL}}(G^*)|}$ being constants and $\varphi_{G^*} \in \mathcal{C}$ encoding the indicator function. Combined with the Lemma 21, we can conclude that \mathcal{A} is computable by 1-WL+NN. Important to note, we can only do this since \mathcal{X} is finite, making the overall sum finite and the cardinality of $\mathcal{X}/\simeq_{1\text{WL}}(G^*)$ well-defined for all graphs. \square

4.2 Proof of Theorem 13

In this section we will prove the converse direction. We start with Lemma 25, where we introduce an upper bound that we will use in Lemma 26 to show that there exists a collection of GNN-computable functions that is 1-WL-Discriminating. After that, we will prove a composition lemma with Lemma 27 which is similar to the one we introduced in the previous section. From this point on, the proof continues as in the previous section and concludes the property to be proved in Lemma 28.

Lemma 25. Let G be an arbitrary graph with $n := |V(G)|$ the number of nodes and $C : V(G) \rightarrow \mathbb{N}$ an arbitrary coloring of the graph G . Then the total number of possible tuples of the form:

$$(C(v), \{C(u) \mid u \in \mathcal{N}(v)\}),$$

for all $v \in V(G)$ can be upper bounded by:

$$n \cdot \sum_{i=0}^{n-1} \binom{n+i-1}{i}.$$

Proof. Assume the above. For the first entry of the tuple, there exist at most n different colors, since there are n nodes. For the second entry, each node $v \in V(G)$ can have between 0 and $n - 1$ neighbors, such that the total number of possibilities is the sum over each cardinality of a multiset with n different colors. In the end, we soundly combine both results by multiplying both together. \square

Lemma 26 (GNN 1-WL-Discriminating). There exists a collection \mathcal{C} of functions from \mathcal{X} to \mathbb{R} computable by GNNs that is 1-WL-Discriminating. Meaning for every $G_1, G_2 \in \mathcal{X}$ with $G_1 \not\sim_{1\text{WL}} G_2$ there exists $\mathcal{A} \in \mathcal{C}$ such that $\mathcal{A}(G_1) \neq \mathcal{A}(G_2)$.

Proof. Since \mathcal{X} is finite, we define $n := \max\{|V(G)| \mid G \in \mathcal{X}\}$ to be the maximum number of nodes of a graph in \mathcal{X} , and $k := \max\{l_G(v) \mid v \in V(G), G \in \mathcal{X}\}$ to be the largest label of any node of a graph in \mathcal{X} . Using Lemma 25, we can compute an upper bound m using n for the number of distinct tuples. Note that, this bound holds true for all graphs in \mathcal{X} . We construct a GNN with n layers working as follows for any $v \in V(G)$:

$$\begin{aligned} f^{(0)}(v) &:= l_G(v), \\ f^{(t)}(v) &:= h_{m,t}(f^{(t-1)}(v), \{f^{(t-1)}(u) \mid u \in \mathcal{N}(v)\}), \quad 0 < t < n. \end{aligned}$$

Here $h_{m,t}$ is a function that maps the tuples injectively to an integer of the set:

$$\{i \in \mathbb{N} \mid k + (t - 1) \cdot m + 1 \leq i \leq k + t \cdot m\}$$

such that each layer uses its own coloring domain. This function exists as by the soundness of the upper bound of Lemma 25 the cardinality of its co-domain is greater or equal than the one of its domain. With this, we ensure that each layer, maps a tuple to a new, previously not used, color. Therefore, every layer of this GNN, computes a single iteration of the 1-WL algorithm. Further, since the 1-WL algorithm converges after at most $|V(G)| =: n$ iterations, we set the number of layers to n . Thereby we ensure that the coloring computed after n layers when applied on any graph $G \in \mathcal{X}$ is similarly expressive as the coloring computed by the 1-WL algorithm when applied on G .

We define the collection \mathcal{C} of functions computable by GNNs that is 1-WL-Discriminating as:

$$\mathcal{C} := \{\mathcal{A} : \mathcal{X} \rightarrow \mathbb{R}, G \mapsto \text{Readout}_c(\{f^{(n)}(v) \mid v \in V(G)\}) \mid c \in \mathbb{N}\},$$

where Readout_c is the READOUT function that returns the number of nodes colored as c in the coloring of $f^{(n)}$. \square

Similar to the proof in the previous section, we will use Lemma 27 to introduce the ability to construct GNNs that take in as input multiple GNNs and then apply a multilayer perceptron to the combined output.

Lemma 27 (Composition GNN). Let \mathcal{C} be a collection of function computable by GNNs. Further, let $\mathcal{A}_1, \dots, \mathcal{A}_n \in \mathcal{C}$ and MLP a suitable multilayer perceptron, then $\hat{\mathcal{A}}(\cdot) := \text{MLP}(\mathcal{A}_1(\cdot), \dots, \mathcal{A}_n(\cdot))$ is also computable by a GNN.

Proof. Assume the above. Further, for the ease of readability we make the assumptions that the node features computed in each layer by each \mathcal{A}_i is one dimensional. Note that, one can easily extend this proof to work over arbitrary feature dimensions.

Further, we will for any $x \in \mathbb{R}^d$ use the notation $x[i]$ to indicate the i .th element of the vector x . Additionally, we use the following notation to indicate the merge and aggregation function used in layer t by \mathcal{A}_i as $f_{\text{merge},i}^{(t)}$ and $f_{\text{agg},i}^{(t)}$. Similarly, does Readout_i indicate the READOUT function of \mathcal{A}_i .

We will prove the lemma by giving a construction of $\hat{\mathcal{A}}$. Let K be the maximum number of layers of all $\mathcal{A}_1, \dots, \mathcal{A}_n$. We construct the GNN $\hat{\mathcal{A}}$ with K layers, with the input layer working as follows on $v \in G$:

$$\hat{f}^{(0)}(v) := \begin{bmatrix} f_1^{(0)}(v) \\ \vdots \\ f_n^{(0)}(v) \end{bmatrix},$$

and each other layer $0 < t \leq K$ uses the merge $f_{\text{merge}}^{(t)}$ and aggregation $f_{\text{agg}}^{(t)}$ functions as defined below:

$$\begin{aligned} \hat{f}_{\text{merge}}^{(t)}(\hat{f}^{(t-1)}(v), \text{Agg}) &= \begin{bmatrix} f_{\text{merge},i}^{(t)}(\hat{f}^{(t-1)}(v)[1], \text{Agg}[1]) \\ \vdots \\ f_{\text{merge},n}^{(t)}(\hat{f}^{(t-1)}(v)[n], \text{Agg}[n]) \end{bmatrix}, \text{ and} \\ \hat{f}_{\text{agg}}^{(t)}(\{\{\hat{f}^{(t-1)}(w) \mid w \in \mathcal{N}(v)\}\}) &= \begin{bmatrix} f_{\text{agg},1}^{(t)}(\{\{f^{(t-1)}(w)[1] \mid w \in \mathcal{N}(v)\}\}) \\ \vdots \\ f_{\text{agg},n}^{(t)}(\{\{f^{(t-1)}(w)[n] \mid w \in \mathcal{N}(v)\}\}) \end{bmatrix}. \end{aligned}$$

Note that, not all \mathcal{A}_i will have K layers. For these cases we define the missing functions as follows:

$$\begin{aligned} f_{\text{merge},i}^{(t)}(\hat{f}^{(t-1)}(v), \text{Agg}) &:= \hat{f}^{(t-1)}(v), \text{ and} \\ f_{\text{agg},i}^{(t)}(\{\{f^{(t-1)}(w) \mid w \in \mathcal{N}(v)\}\}) &:= 0. \end{aligned}$$

These functions do not change anything, and only forward the result of the actual computation of \mathcal{A}_i to the last layer. Finally, we construct the READOUT function of $\hat{\mathcal{A}}$ as follows:

$$\text{Readout}(\{\{\hat{f}^{(t)}(v) \mid v \in V(G)\}\}) := \text{MLP} \circ \begin{bmatrix} \text{Readout}_i(\{\{f_{\text{merge},1}^{(t)}(v) \mid v \in V(G)\}\}) \\ \vdots \\ \text{Readout}_n(\{\{f_{\text{merge},n}^{(t)}(v) \mid v \in V(G)\}\}) \end{bmatrix}.$$

□

As a result, we are in the same starting position as at the beginning of the proof in subsection 4.1. That is, we have shown with Lemma 26 that there is a collection of functions that can distinguish any pair of graphs distinguishable by the 1-WL algorithm, and with Lemma 27 we have shown that the composition of multiple GNNs and a multilayer perceptron is still computable by a GNN. Such that we can simply apply the results of Lemmas 22 and 23 to GNNs, such that we know that for any fixed $G^* \in \mathcal{X}$ the indicator function φ_{G^*} working over \mathcal{X} with:

$$\varphi_{G^*}(G) := \begin{cases} 1, & \text{if } G \simeq_{1\text{WL}} G^* \\ 0, & \text{else} \end{cases},$$

is computable by a GNN.

Lemma 28. Let \mathcal{C} be a collection of functions from \mathcal{X} to \mathbb{R} computable by GNNs so that for all $G^* \in \mathcal{X}$, there exists $\varphi_{G^*} \in \mathcal{C}$ satisfying $\forall G \in \mathcal{X} : \varphi_{G^*}(G) = \mathbb{1}_{G \simeq_{1\text{WL}} G^*}$, then every function computable by a 1-WL+NN is also computable by a GNN.

Proof. Assume the above. For any function \mathcal{B} computed by 1-WL+NN that works over \mathcal{X} to \mathbb{R} , we show that it can be decomposed as follows for any $G \in \mathcal{X}$ as input:

$$\begin{aligned} \mathcal{B}(G) &= \left(\frac{1}{|\mathcal{X}/\simeq_{1\text{WL}}(G)|} \sum_{G^* \in \mathcal{X}} \mathbb{1}_{G \simeq_{1\text{WL}} G^*} \right) \cdot \mathcal{B}(G) \\ &= \frac{1}{|\mathcal{X}/\simeq_{1\text{WL}}(G)|} \sum_{G^* \in \mathcal{X}} \mathcal{B}(G^*) \cdot \mathbb{1}_{G \simeq_{1\text{WL}} G^*} \\ &= \sum_{G^* \in \mathcal{X}} \frac{\mathcal{B}(G^*)}{|\mathcal{X}/\simeq_{1\text{WL}}(G^*)|} \cdot \varphi_{G^*}(G) \end{aligned} \tag{0.3}$$

with $\mathcal{X}/\simeq_{1\text{WL}}(G^*)$ we denote the set of all graphs G over \mathcal{X} that are equivalent to G^* according to the $\simeq_{1\text{WL}}$ relation. With Lemma 19, we know that for any $G_1, G_2 \in \mathcal{X}$ with $G_1 \simeq_{1\text{WL}} G_2$: $\mathcal{B}(G_1) = \mathcal{B}(G_2)$, such this decomposition is sound.

We can encode \mathcal{B} as stated in Equation 0.3 via a multilayer perceptron with $\frac{\mathcal{B}(G^*)}{|\mathcal{X}/\simeq_{1\text{WL}}(G^*)|}$ being constants and $\varphi_{G^*} \in \mathcal{C}$ encoding the indicator function. Combined with the Lemma 27, we can conclude that \mathcal{B} is computable by 1-WL+NN. Important to note, we can only do this since \mathcal{X} is finite, making the overall sum finite and the cardinality of $\mathcal{X}/\simeq_{1\text{WL}}(G^*)$ well-defined for all graphs. \square

5 Empirical Investigation

Yet to come!

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1 Appendix

1.1 Figures and graphs

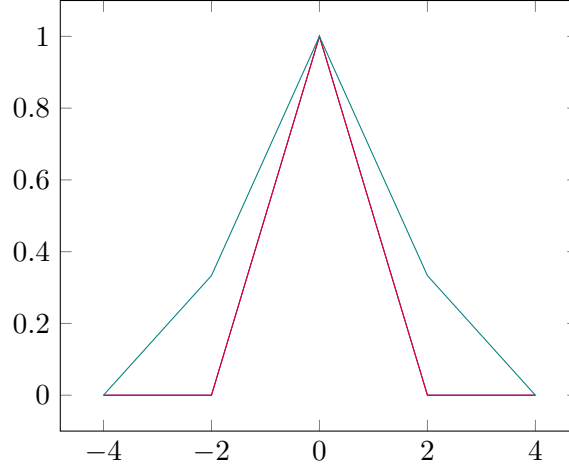


Figure 4: Illustration of the so-called “bump” function $\psi_a(x)$ used in the proof of Lemma 23. Here the colors of the displayed functions correspond to the parameter a set to $a := 1$ in blue, $a := 2$ in red and $a := 3$ in green.

1.2 Proofs

Lemma 21: Composition Lemma for 1-WL+NN

Proof of Lemma 21. Let \mathcal{C} be a collection of functions computed by 1-WL+NN, $h_1, \dots, h_n \in \mathcal{C}$, and MLP^\bullet a multilayer perceptron. Further, let f_1, \dots, f_n be the encoding functions, as well as $\text{MLP}_1, \dots, \text{MLP}_n$ be the multilayer perceptrons used by h_1, \dots, h_n respectively. As outlined above, we will now construct f^* and MLP^* , such that for all graphs $G \in \mathcal{X}$:

$$\text{MLP}^\bullet(h_1(G), \dots, h_n(G)) = \text{MLP}^* \circ f^*(\{\{1\text{-WL}(G)(v) \mid v \in V(G)\}\})$$

with which we can conclude that the composition of multiple functions computable by 1-WL+NN, is in fact also 1-WL+NN computable.

We define the new encoding function f^* to work as follows on an arbitrary input multiset M :

$$f^*(M) := \text{concat}\left(\begin{bmatrix} f_1(M) \\ \vdots \\ f_n(M) \end{bmatrix}\right),$$

where concat is the concatenation function, concatenating all encoding vectors to one single vector.

Using the decomposition introduced in Definition 8, we can decompose each MLP_i for $i \in [n]$ at layer $j > 0$ as follows: $(\text{MLP}_i)_j(v) := \sigma_{i,t}(W_j^i \cdot (\text{MLP}_i)_{j-1}(v) + b_j^i)$. Using this notation we construct MLP^* as follows:

the activation function might not be the same everywhere

$$(\text{MLP}^*)_0(v) := v$$

$$(\text{MLP}^*)_{j+1}(v) := \sigma_j^*(W_j^* \cdot (\text{MLP}^*)_j(v) + \text{concat}(\begin{bmatrix} b_j^1 \\ \vdots \\ b_j^n \end{bmatrix})) \quad , \forall j \in [k]$$

$$(\text{MLP}^*)_{j+k+1}(v) := (\text{MLP}^\bullet)_{j+1}(v) \quad , \forall j \in [k^\bullet - 1]$$

where k is the maximum number of layers of the set of MLP_i 's, k^\bullet is the number of layers of the given MLP^\bullet and σ an element wise activation function. Thereby, we define in the first equation line, that the start of the sequence is the input, with the second line, we construct the “simultaneous” execution of the MLP_i 's, and in the last equation line, we add the layers of the given MLP^\bullet to the end. Further, we define the weight matrix W_j^* as follows:

$$W_j^* := \begin{bmatrix} W_j^1 & 0 & \dots & 0 \\ 0 & W_j^2 & \ddots & \vdots \\ \vdots & \ddots & \ddots & 0 \\ 0 & \dots & 0 & W_j^n \end{bmatrix},$$

such that we build a new matrix where each individual weight matrix is placed along the diagonal. Here we denote with 0, zero matrices with the correct dimensions, such that W_j^* is well-defined. Important to note, should for an MLP_i , W_j^i not exist, because it has less than j layers, we use for W_j^i the identity matrix I_m where m is the dimension of the output computed by MLP_i . And we define the overall activation function σ_j^* as following:

$$\sigma_j^*(v) := \begin{bmatrix} \sigma_{1,j}(v[1]) \\ \vdots \\ \sigma_{1,j}(v[d_1]) \\ \vdots \\ \sigma_{n,j}(v[d_1 + \dots + d_{n-1} + 1]) \\ \vdots \\ \sigma_{n,j}(v[d_1 + \dots + d_n]) \end{bmatrix},$$

where d_i is the dimension of the output of MLP_i at layer j , and for the ease of readability we denote the i .th component of vector v here with $v[i]$. Thereby, we construct an activation function that applies each respective activation function of the MLP_i 's individually to their respective computation. \square

Theorem 15: Continuous Case: “GNN \subseteq 1-WL+NN”

We will prove Theorem 15 by first introducing a definition and then 2 lemmas using that definition which collectively prove the entire theorem. In particular, we define the property that a collection of functions is able to find any $\simeq_{1\text{WL}}$ equivalence class. Then, in Lemma 30, we prove that 1-WL-Discriminating somehow implies being able to locate any $\simeq_{1\text{WL}}$ -equivalence class. And finally, in Lemma 31, we prove that being able to locate any $\simeq_{1\text{WL}}$ equivalence class further implies being somehow GNN-approximating. The overall proof is inspired by the work of [3].

Definition 29 (Locating $\simeq_{1\text{WL}}$ equivalence classes). Let \mathcal{C} be a collection of continuous functions from \mathcal{X} to \mathbb{R} . If for all $\epsilon \in \mathbb{R}$ with $\epsilon > 0$ and for every graph $G^* \in \mathcal{X}$ the function h_{G^*} exists in \mathcal{C} , with the following properties:

1. for all $G \in \mathcal{X} : h_{G^*}(G) \geq 0$,
2. for all $G \in \mathcal{X}$ with $G \simeq_{1\text{WL}} G^* : h_{G^*}(G) = 0$, and
3. there exists a constant $\delta_{G^*} > 0$, such that for all $G \in \mathcal{X}$ with $h_{G^*}(G) < \delta_{G^*}$ there exists a graph $G' \in \mathcal{X} / \simeq_{1\text{WL}}(G)$ in the equivalence class of G such that $\|G' - G^*\|_2 < \epsilon$

we say \mathcal{C} is able to locate every $\simeq_{1\text{WL}}$ equivalence class.

One can interpret this function h_{G^*} as a kind of loss function that measures the similarity between its input graph to G^* . It yields no loss for the input G , if G is indistinguishable from G^* by the 1-WL algorithm ($G \simeq_{1\text{WL}} G^*$), only a small loss if G is close to a graph in the $\simeq_{1\text{WL}}$ equivalence class of G^* (the loss is upper bounded by δ_{G^*}), and an arbitrary loss otherwise.

Lemma 30. Let \mathcal{C} be a collection of continuous functions from \mathcal{X} to \mathbb{R} computable by 1-WL+NN. If \mathcal{C} is 1-WL-Discriminating, then there exists a collection of functions \mathcal{C}' computable by 1-WL+NN that is able to locate every $\simeq_{1\text{WL}}$ equivalence class on \mathcal{X} .

Proof. Let $G^* \in \mathcal{X}$ be fixed and $\epsilon > 0$ be given. Since \mathcal{C} is 1-WL-Discriminating, we know that for every $G \in \mathcal{X}$ with $G \not\simeq_{1\text{WL}} G^*$, there exists a function $h_{G,G^*} \in \mathcal{C}$ such that $h_{G,G^*}(G) \neq h_{G,G^*}(G^*)$. We use this property to construct for each $G \in \mathcal{X}$ with $G \not\simeq_{1\text{WL}} G^*$ a set A_{G,G^*} as follows:

$$A_{G,G^*} := \{G' \in \mathcal{X} \mid h_{G,G^*}(G') \in (h_{G,G^*}(G) \pm \frac{|h_{G,G^*}(G) - h_{G,G^*}(G^*)|}{2})\},$$

where $(a \pm b)$ is the open set $(a - b, a + b)$ over \mathbb{R} . One can use Figure 5 below for a better understanding, when a graph G' is contained in A_{G,G^*} and when not. With the illustration one can easily see, that $G \in A_{G,G^*}$ and $G^* \notin A_{G,G^*}$.

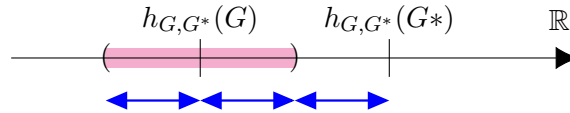


Figure 5: An illustration to better understand the proof. The set A_{G,G^*} consists of all graphs G that are mapped by h_{G,G^*} into the pink area, which size depends on the distance between the image of G and G^* under h_{G,G^*} . Moreover, the blue distances all have the same length.

Furthermore, by assumption h_{G,G^*} is continuous, such that A_{G,G^*} is an open set. For every $G \in \mathcal{X}$ with $G \simeq_{1\text{WL}} G^*$ we define A_{G,G^*} as follows:

$$A_{G,G^*} := \{G' \in \mathcal{X} \mid \|G' - G\|_2 < \epsilon\},$$

where $\|\cdot\|_2$ is the l_2 .

Thus, $\{A_{G,G^*}\}_{G \in \mathcal{X}}$ is an open cover of \mathcal{X} . Since \mathcal{X} is compact, there exists a finite subset \mathcal{X}_0 such that $\{A_{G,G^*}\}_{G \in \mathcal{X}_0}$ also covers \mathcal{X} . Hence, $\forall G \in \mathcal{X} \exists G_0 \in \mathcal{X}_0 : G \in A_{G_0,G^*}$.

We define the desired function h_{G^*} as follows:

$$h_{G^*}(\cdot) = \sum_{\substack{G_0 \in \mathcal{X}_0 \\ G_0 \not\sim_{1WL} G^*}} \bar{h}_{G_0, G^*}(\cdot), \quad (.4)$$

where we define \bar{h}_{G_0, G^*} almost exactly the same as in a previous proof:

$$\bar{h}_{G_0, G^*}(\cdot) = |h_{G_0, G^*}(\cdot) - h_{G_0, G^*}(G^*)| \quad (.5)$$

$$= \max(h_{G_0, G^*}(\cdot) - h_{G_0, G^*}(G^*)) + \max(h_{G_0, G^*}(G^*) - h_{G_0, G^*}(\cdot)) \quad (.6)$$

Note that, “ $h_{G_0, G^*}(G^*)$ ” is a constant in the definition of $\bar{h}_{G_0, G^*}(\cdot)$. We will shortly proof, that $h_{G^*}(\cdot)$ fulfills the desired properties on input $G \in \mathcal{X}$:

1. By construction, any \bar{h}_{G_0, G^*} is non-negative, such that the sum over these functions is also non-negative.
2. If $G \simeq_{1WL} G^*$, using Lemma 19 we know that $h_{G^*}(G) = h_{G^*}(G^*)$, and by definition $h_{G^*}(G^*) = 0$, such that we can conclude $h_{G^*}(G) = 0$.
3. Let δ_{G^*} be:

$$\delta_{G^*} := \frac{1}{2} \min_{\substack{G_0 \in \mathcal{X} \\ G_0 \not\sim_{1WL} G^*}} |h_{G_0, G^*}(G_0) - h_{G_0, G^*}(G^*)|.$$

Prove by contraposition: Assume that for every graph $G' \in \mathcal{X} / \simeq_{1WL}(G)$: $\|G' - G^*\|_2 \geq \epsilon$. Hence, $G \notin \bigcup_{G' \in \mathcal{X} / \simeq_{1WL}(G^*)} A_{G', G^*}$ (not in the union of l_2 balls of size ϵ around all graphs of the equivalence class of G^*). However, since $\{A_{G, G^*}\}_{G \in \mathcal{X}_0}$ is a cover of \mathcal{X} , there must exist a $G_0 \in \mathcal{X}_0$ with $G_0 \not\sim_{1WL} G^*$ such that $G \in A_{G_0, G^*}$. Thus, by definition of A_{G_0, G^*} we know that $h_{G_0, G^*}(G) \in (h_{G_0, G^*}(G_0) \pm \frac{|h_{G_0, G^*}(G_0) - h_{G_0, G^*}(G^*)|}{2})$, which when reformulated states:

$$|h_{G_0, G^*}(G) - h_{G_0, G^*}(G_0)| < \frac{1}{2} |h_{G_0, G^*}(G_0) - h_{G_0, G^*}(G^*)|. \quad (.7)$$

Using this, we can prove $\bar{h}_{G_0, G^*}(G) \geq \delta_{G^*}$, which implies $h_{G^*}(G) \geq \delta_{G^*}$ and concludes the proof:

$$\begin{aligned} \bar{h}_{G_0, G^*}(G) &= |h_{G_0, G^*}(G) - h_{G_0, G^*}(G^*)| \\ &\geq |h_{G_0, G^*}(G_0) - h_{G_0, G^*}(G^*)| - |h_{G_0, G^*}(G) - h_{G_0, G^*}(G_0)| \end{aligned} \quad (.8)$$

$$\geq \frac{1}{2} |h_{G_0, G^*}(G_0) - h_{G_0, G^*}(G^*)| \quad (.9)$$

$$\geq \frac{1}{2} \min_{\substack{G_0 \in \mathcal{X} \\ G_0 \not\sim_{1WL} G^*}} |h_{G_0, G^*}(G_0) - h_{G_0, G^*}(G^*)| =: \delta_{G^*}$$

To understand these inequalities, it helps to visualize them, see therefore Figure 6. We will try to give an explanation now, using the colored distances depicted in Figure 6. In Equation .8, we use the fact that the red distance is always greater than the green minus the blue distance. Further, using Equation .7, we know that the blue distance is always smaller than half of than the green distance. Using this fact, it is easy to see that in Equation .9 the green minus the blue distance is always greater than or equal to the half of the green distance.

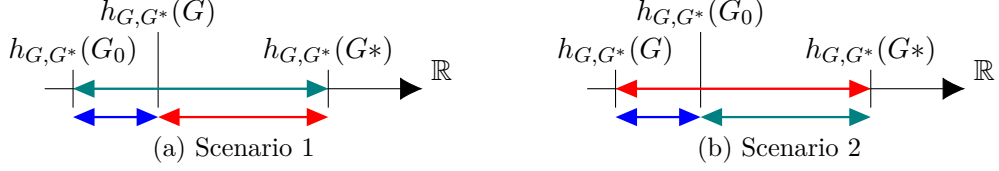


Figure 6: An illustration of two general scenarios to better understand the above used inequalities. Note that, there exists two more general scenarios, with $h_{G_0,G^*}(G^*) < h_{G_0,G^*}(G_0)$, but since we use the absolute function as our measure of distance, these scenarios are equivalent to the ones depicted due to the symmetry of the absolute function. In Table 1 below, we listed all colors used to decode the distances in the figures with their corresponding term in the inequalities.

Color	Term
red	$ h_{G_0,G^*}(G) - h_{G_0,G^*}(G^*) $
green	$ h_{G_0,G^*}(G_0) - h_{G_0,G^*}(G^*) $
blue	$ h_{G_0,G^*}(G) - h_{G_0,G^*}(G_0) $

Table 1: Color legend for the proof of Lemma 30

Consequently, we know that if $h_{G^*}(G) < \delta_G$ it means that $G \in \bigcup_{G' \in \mathcal{X}/\simeq_{1\text{WL}}(G^*)} A_{G',G^*}$, which implies that there exists $G' \in \mathcal{X}/\simeq_{1\text{WL}}(G)$ with $\|G' - G^*\|_2 < \epsilon$.

As a final note, we can construct a multilayer perceptron $\text{MLP}_{h_{G^*}}$ computing the function $h_{G^*}(\cdot)$ as in Equation .4. The $\text{MLP}_{h_{G^*}}$ gets as input a vector of the output of the finite set of functions $\{\bar{h}_{G_0,G^*}\}_{G_0 \in \mathcal{X}_0, G_0 \not\simeq_{1\text{WL}} G^*}$ applied on the input graph. Further, Equation .4 can be encoded by replacing the max operator by the non-linear activation function ReLU. Using Lemma 21, we can conclude that $h_{G^*}(\cdot)$ is computable by 1-WL+NN. \square

Lemma 31. Let \mathcal{C} be a collection of continuous functions from \mathcal{X} to \mathbb{R} . If \mathcal{C} is able to locate every $\simeq_{1\text{WL}}$ equivalence class on \mathcal{X} , then there exists a collection of functions \mathcal{C}' computable by 1-WL+NN that is able to locate every $\simeq_{1\text{WL}}$ equivalence class on \mathcal{X} .

Proof. Let \mathcal{A} be a continuous function from \mathcal{X} to \mathbb{R} computable by a GNN. Since \mathcal{X} is compact, this implies that \mathcal{A} is uniformly continuous on \mathcal{X} , which further implies that $\forall \epsilon > 0 \exists r > 0$ such that $\forall G_1, G_2 \in \mathcal{X}$, if $\|G_1 - G_2\|_2 < r$, then $|\mathcal{A}(G_1) - \mathcal{A}(G_2)| < \epsilon$. Further, since \mathcal{A} is GNN computable, we know that $\forall G_1, G_2 \in \mathcal{X}$, if $G_1 \simeq_{1\text{WL}} G_2 : \mathcal{A}(G_1) = \mathcal{A}(G_2)$, hence if $G' \in \mathcal{X}/\simeq_{1\text{WL}}(G_1)$ with $\|G' - G_2\|_2 < r$ exists, then $|\mathcal{A}(G_1) - \mathcal{A}(G_2)| < \epsilon$.

Throughout the rest of the proof, let $\epsilon > 0$ be fixed. Using the property described above, we know that for this ϵ there exists a constant r , we fix this as well. Further, by the assumptions of the lemma we try to prove, we know that for any $\epsilon' > 0$ there exists $h_G \in \mathcal{C}$ for any $G \in \mathcal{X}$ with the above described properties. We choose $\epsilon' := r$ for all $h_G \in \mathcal{C}$ throughout the proof.

For any $G \in \mathcal{X}$, we define the set $h_G^{-1}(a) := \{G' \in \mathcal{X} \mid h_G(G') \in [0, a)\}$, as the set of graphs that are mapped into the open interval $[0, a)$ by h_G . We illustrated this set in Figure 7 for better understanding.

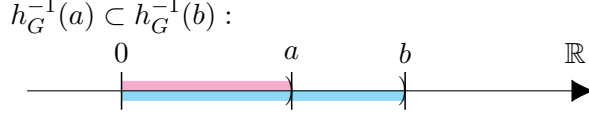


Figure 7: Illustration of the set $h_G^{-1}(\cdot)$ for two arbitrary values a, b with $a < b$. The pink area visualizes the open interval $[0, a)$, and similarly in blue for $[0, b)$.

Then, by definition of h_G , there exists a constant δ_G such that:

$$h_G^{-1}(\delta_G) \subseteq \bigcup_{G' \in \mathcal{X} / \simeq_{1\text{WL}}(G)} \{G'' \in \mathcal{X} \mid \|G'' - G'\|_2 < r\}.$$

Since h_G is continuous by definition, $h_G^{-1}(\delta_G)$ is an open set. Hence, $\{h_G^{-1}(\delta_G)\}_{G \in \mathcal{X}}$ is an open cover of \mathcal{X} , and further as \mathcal{X} is compact, there exists a finite subset $\mathcal{X}_0 \subseteq \mathcal{X}$ such that $\{h_{G_0}^{-1}(\delta_{G_0})\}_{G_0 \in \mathcal{X}_0}$ is also a cover of \mathcal{X} . For each $G_0 \in \mathcal{X}_0$, we construct the function φ_{G_0} from \mathcal{X} to \mathbb{R} as follows:

$$\varphi_{G_0}(\cdot) := \max(\delta_{G_0} - h_{G_0}(\cdot), 0).$$

The function has two important properties, for once it is non-negative and for any $G \in \mathcal{X} : \varphi_{G_0}(G) > 0$, if and only if $G \in h_{G_0}^{-1}(\delta_{G_0})$, thereby acting as a sort of weak indicator function. Building up on this property, we construct for each $G_0 \in \mathcal{X}_0$ the function ψ_{G_0} from \mathcal{X} to \mathbb{R} as follows:

$$\psi_{G_0}(\cdot) := \frac{\varphi_{G_0}(\cdot)}{\sum_{G' \in \mathcal{X}_0} \varphi_{G'}(\cdot)},$$

which is well-defined, because $\{h_{G_0}^{-1}(\delta_{G_0})\}_{G_0 \in \mathcal{X}_0}$ is a cover of \mathcal{X} , such that we can conclude that for any input graph G on $\psi_{G_0}(\cdot)$ there exists a set $h_{G_0}^{-1}(\delta_{G_0})$ with $G \in h_{G_0}^{-1}(\delta_{G_0})$ implying $\varphi_{G_0}(G) > 0$, thus making the denominator not 0. The function ψ_{G_0} has two important properties, for once it is non-negative, because φ_G for all $G \in \mathcal{X}$ is non-negative, and for any $G \in \mathcal{X} : \psi_{G_0}(G) > 0$, if and only if $G \in h_{G_0}^{-1}(\delta_{G_0})$.

Further, we can observe that the set of functions $\{\psi_{G_0}\}_{G_0 \in \mathcal{X}_0}$ is a partition of unity on \mathcal{X} with respect to the open cover $\{h_{G_0}^{-1}(\delta_{G_0})\}_{G_0 \in \mathcal{X}_0}$, because:

1. For any $G \in \mathcal{X}$ the set of functions mapping G not to 0 is finite, as the set of all functions $\{\psi_{G_0}\}_{G_0 \in \mathcal{X}_0}$ is finite, since \mathcal{X}_0 is finite.
2. For any $G \in \mathcal{X} : \sum_{G_0 \in \mathcal{X}_0} \psi_{G_0}(G) = 1$, since:

$$\sum_{G_0 \in \mathcal{X}_0} \psi_{G_0}(G) = \sum_{G_0 \in \mathcal{X}_0} \frac{\varphi_{G_0}(G)}{\sum_{G' \in \mathcal{X}_0} \varphi_{G'}(G)} = \frac{\sum_{G_0 \in \mathcal{X}_0} \varphi_{G_0}(G)}{\sum_{G' \in \mathcal{X}_0} \varphi_{G'}(G)} = 1.$$

We can use this property to decompose the given function \mathcal{A} as follows on any input $G \in \mathcal{X}$:

$$\mathcal{A}(G) = \mathcal{A}(G) \cdot \left(\sum_{G_0 \in \mathcal{X}_0} \psi_{G_0}(G) \right) = \sum_{G_0 \in \mathcal{X}_0} \mathcal{A}(G) \cdot \psi_{G_0}(G).$$

Recall the property from the beginning of the proof that for any $G \in \mathcal{X}$ if $G \in h_{G_0}^{-1}(\delta_{G_0})$, then there exists a $G' \in \mathcal{X} / \simeq_{1\text{WL}}(G)$ with $\|G' - G_0\|_2 < r$, which implies that $|\mathcal{A}(G) - \mathcal{A}(G_0)| < \epsilon$. With this, we can construct a function $\hat{\mathcal{A}}$ on \mathcal{X} that approximates \mathcal{A} within accuracy ϵ :

$$\hat{\mathcal{A}}(\cdot) = \sum_{G_0 \in \mathcal{X}_0} \mathcal{A}(G_0) \cdot \psi_{G_0}(\cdot).$$

Note that, “ $\mathcal{A}(G_0)$ ” is a constant here. To prove that $\hat{\mathcal{A}}$ approximates \mathcal{A} within accuracy ϵ we need to show that $\sup_{G \in \mathcal{X}} |\mathcal{A}(G) - \hat{\mathcal{A}}(G)| < \epsilon$. Let $G \in \mathcal{X}$ be arbitrary, then:

$$\begin{aligned}
|\mathcal{A}(G) - \hat{\mathcal{A}}(G)| &= \left| \sum_{G_0 \in \mathcal{X}_0} \mathcal{A}(G) \cdot \psi_{G_0}(G) - \sum_{G_0 \in \mathcal{X}_0} \mathcal{A}(G_0) \cdot \psi_{G_0}(G) \right| \\
&= \sum_{G_0 \in \mathcal{X}_0} |\mathcal{A}(G) - \mathcal{A}(G_0)| \cdot \psi_{G_0}(G) \\
&< \sum_{G_0 \in \mathcal{X}_0} \epsilon \cdot \psi_{G_0}(G) \\
&= \epsilon \cdot \sum_{G_0 \in \mathcal{X}_0} \psi_{G_0}(G) \\
&= \epsilon \cdot 1
\end{aligned} \tag{.10}$$

In Equation .10, we use the fact that applies to all $G_0 \in \mathcal{X}_0$ on input G :

- If $\psi_{G_0}(G) > 0$, then we know that $G \in h_{G_0}^{-1}(\delta_{G_0})$, such that we can upper bound $|\mathcal{A}(G) - \mathcal{A}(G_0)| < \epsilon$.
- If $\psi_{G_0}(G) = 0$, we know that no matter what $|\mathcal{A}(G) - \mathcal{A}(G_0)|$ is, the summand is 0, such that we can just assume $|\mathcal{A}(G) - \mathcal{A}(G_0)| < \epsilon$ without loss of generality.

In the end, we give a short explanation, that $\hat{\mathcal{A}}$ is computable by 1-WL+NN. For this we construct a multilayer perceptron with three layers and then conclude with Lemma 21 the computability. Let us therefore break down the whole construction of $\hat{\mathcal{A}}$:

$$\hat{\mathcal{A}}(\cdot) = \sum_{G_0 \in \mathcal{X}_0} \mathcal{A}(G_0) \cdot \frac{1}{\sum_{G' \in \mathcal{X}_0} \max(\delta_{G'} - h_{G'}(\cdot), 0)} \cdot \max(\delta_{G_0} - h_{G_0}(\cdot), 0).$$

We construct a multilayer perceptron $\text{MLP}_{\hat{\mathcal{A}}}$ that takes in as input a vector of the output of the finite set of functions $\{h_{G_0}\}_{G_0 \in \mathcal{X}_0}$ applied on the input graph. In the first layer we compute each $\max(\delta_{G_0} - h_{G_0}(\cdot))$ term, where δ_{G_0} is a constant, in particular the bias, and the max operator is replaced by the activation function ReLU. In the second layer, we compute the sum of the denominator of the fraction, to which we apply the activation function $f(x) := \frac{1}{x}$. In the last layer we compute the overall sum where $\mathcal{A}(G_0)$ is a constant. \square

Theorem 16: Continuous Case: “1-WL+NN \subseteq GNN”

In this section we will shortly prove that GNN-Approximating implies 1-WL-Discriminating. Similar to the proof of Theorem 12, we will take advantage of the fact that our Definition 7 of GNNs does not impose any constraints on the readout function, apart from the fact that it must be permutation invariant and computable. We have drawn inspiration for this proof from the work of [3].

Proof. \square

Proof of Theorem 16. Let \mathcal{C} be a collection of continuous functions from \mathcal{X} to \mathbb{R} that is GNN approximating. Let $G_1, G_2 \in \mathcal{X}$ with $G_1 \not\sim_{1\text{WL}} G_2$, then we define the function h_{G_1, G_2} on input $G \in \mathcal{X}$:

$$h_{G_1, G_2}(G) = \min_{\pi \in S_n} \|\pi^T G \pi - G_1\|_2,$$

Not yet adapted to the new definition of the readout functions

New proof will be very similar to the finite case

where $\|\cdot\|_2$ is the l_2 norm. Note that, h_{G_1, G_2} is computable as the number of permutations π in S_n are finite. Since h_{G_1, G_2} is permutation invariant, we can construct a GNN with 0 layers and h_{G_1, G_2} as its READOUT function, thereby constructing a GNN computing h_{G_1, G_2} and consequently showing that the function is GNN computable. We choose $\epsilon := \frac{1}{2} \cdot h_{G_1, G_2}(G_2) > 0$, then there exists a function $h_\epsilon \in \mathcal{C}$ approximating h_{G_1, G_2} within ϵ accuracy. With this, we have a function $h_\epsilon \in \mathcal{C}$ that can distinguish G_1 and G_2 , since:

$$\begin{aligned}
|h_\epsilon(G_1) - h_\epsilon(G_2)| &> |(h_{G_1, G_2}(G_1) - \epsilon) - (h_{G_1, G_2}(G_2) + \epsilon)| \\
&= |h_{G_1, G_2}(G_2) - 2 \cdot \epsilon| && \text{by definition } h_{G_1, G_2}(G_1) = 0 \\
&= |h_{G_1, G_2}(G_2) - 2 \cdot \frac{1}{2} \cdot h_{G_1, G_2}(G_2)| = 0.
\end{aligned}$$

□