

4. Claim: Given a graph $G=(V,E)$, with n vertices such that the degree of every vertex of G is at most k and in every connected component of G there is at least one vertex such that the degree of that vertex is less than k , the vertices of G can be colored with k distinctive colors such that the endpoints of every edge of G have distinct colors.

Proof: Prove by induction on n . Let $P(n)$ be "every graph with n vertices such that the degree of every vertex is at most k and in every connected component of G there is at least one vertex such that the degree of that vertex is less than k , can be colored using k distinct colors such that the endpoints of every edge of G have distinct colors." Let G be an arbitrary graph fitting the above definition.

Base Case: $P(1)$ holds because every connected component is made up of at most 1 vertex. We color those vertices with a single color. We can do so because $k \geq 0$.

Inductive Hypothesis: Assume $P(n-1)$ holds.

Inductive Step: Goal, prove $P(n)$ holds. Let V be a set of vertices chosen from G such that every vertex v that has degree k is $v \in V$ but not including vertices with degree k such that all of that vertex's neighbors are also in V . Let $G' = G - V$ and all edges incident to every $v \in V$. Note that removing V (and all edges incident to every $v \in V$) from G can only reduce the degree of every remaining vertex in G' . Specifically, every vertex of G' has at most degree $k-1$ and every $v \in V$ also has at most degree $k-1$. Since every connected component of G' has $n-1$ vertices, by the Inductive Hypothesis we color the vertices of G' with $k-1$ colors such that every edge has endpoints with distinct colors. Color every vertex of G , except the vertices of V , with the same vertex colors as those in G' . Now, to color $v \in V$ of G , note that every $v \in V$ of G has at most $k-1$ neighbors. Color every uncolored $v \in V$ of G with the one color not used by any of the neighbors of that v . \square

