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Falling Glass

- a) Given n floors, m glass sheets. At some floor, x , we have to find the worst case out of the floors above x ($n-x$) or the floors below x ($x-1$) which splits into two subproblems. In the subproblem for the floors above x , the glass pane doesn't break so we use the same number of sheets. In the subproblem for the floors below x , the glass pane breaks so we use one less sheet in that subproblem.

d)

e)

- f) To memoize this problem, I would use a table placed outside the recursion to store the minimum worst case for n floors and m sheets. While going through each subproblem, the algorithm will check if the minimum worst case has been stored and if not, it will recursively solve.

Rod Cutting

b)

length	1	2	3	4	5	6	7	8	9
price	1	3	9	9	10	15	21	25	26

Density: 1 1.5 3 2.25 2 2.5 3 3.13 2.8

Given a rod of length 10. Greedy solution would pick rod of length 8, then 2 which yields 28.

Optimal solution would be to pick rods of length 7 and 3 yielding 30.

