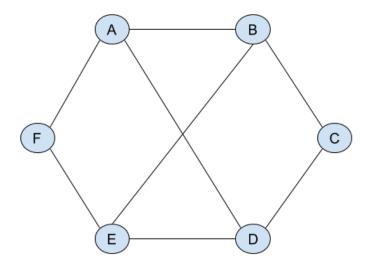
Computer Science Mentors 70

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1 Graph 101



- 1. Take a look at the following undirected graph.
 - (a) How many vertices are in this graph?
 - (b) What is the degree of vertex B?
 - (c) What is the total degree of this graph?
 - (d) Consider the traversal $A \to B \to C \to D \to A$. How would you categorize it (walk / cycle / simple path / tour)?
 - (e) Give an example of a simple path of length 4.
 - (f) Is it possible to construct a traversal that is a tour but not a simple cycle from this graph (can go through vertices twice, but not edges)? Why or why not?
 - (g) in terms of |E| and |V|, what is the sum of the degrees of the graph?
 - (h) What is the average degree of the graph?
- 2. Which of these graphs have Eulerian tours?
 - (a) The complete graph on 5 vertices (K_5) .
 - (b) The complete graph on 6 vertices (K_6) .
 - (c) The complete graph on 7 vertices (K_7) .
 - (d) The 3-dimensional hypercube.
 - (e) The 4-dimensional hypercube.

3.	path between every pair of its vertices.
	False Claim: If every vertex in an undirected graph has degree at least 1, then the graph is connected.
	<i>Proof.</i> We use induction on the number of vertices $n \ge 1$. let $P(n)$ be the proposition that if every vertex in an n -vertex graph has positive degree, then the graph is connected.
	Base case: A graph with 1 vertex doesn't have any positive-degree vertices so $P(1)$ is true vacuously.
	Inductive Hypothesis: Assume $P(n)$ holds. We want to show this implies $P(n+1)$.
	Inductive Step: Consider an n vertex graph that has positive degree. By the asumption $P(n)$, this graph is connected and there is a path from every vertex to every other vertex. Now add a new vertex to create an $n+1$ vertex graph. All that remains is to check that there is a path from v to every other vertex. Suppose we add this vertex v to an existing vertex u . Since the graph was previously connected, we already know there is a path from u to every other vertex in the graph. Therefore, when we connect v to u , we know there will be a path from v to every other vertex in the graph. This proves the claim for $P(n+1)$.
	(a) Give a counter-example to show the claim is false.
	(b) Since the claim is false, there must be an error in the proof. Explain the error.
	(c) How can we avoid this mistake?
	(d) What happens in the inductive step when you apply the fix?
4.	Consider all complete undirected graphs on an even number of vertices. Prove that such graphs can be partitioned into $\frac{n}{2}$ spanning trees that share no edge with another spanning tree.

2 Graph Coloring

1. Show that any tree is 2-colorable.

2. You are hosting a very exclusive party such that a guest is only allowed to come in if they are friends with you or someone else already at the party. After everyone has showed up, you notice that there are n people (including yourself); each person has at least one friend (of course), but no one is friends with everyone else. It is still quite a sad party, because among all the possible pairs of people, there are only a total of n-1 friendships. You want to play a game with two teams, and in order to kindle new friendships, you want to group the people (including yourself) such that within each team, no one is friends with each other. Is this possible? (Hint: How might the previous question be useful?)

- 3. Two knights are placed on diagonally opposite corners of a chessboard, one white and the other black. The knights take turn moving as in standard chess, with white moving first.
 - (a) Let every square on a chessboard represent a vertex of a graph, with edges between squares that are a knight's move away. Describe a 2-coloring of this graph.
 - (b) Show that the black knight can never be captured, even if it cooperates with the white knight.

3 Planarity

We say a graph is **planar** if it can be drawn on the plane without any edges crossing each other. If the graph is planar, we can use Euler's planar formula:

$$v + f = e + 2$$

Some other (less important) results follow from Euler's formula if a graph is planar:

Corollary 1:
$$e \le 3v - 6$$

Corollary 2:
$$f \le 2v - 4$$

1. Try to prove all three of the above results.

Hint 1: What is the minimum number of edges per face? What is the number of faces any edge can tough?

Hint 2: Take the graph with the most number of edges or faces per vertex. What does it look like? look at the first hint!

Hint 3: You should come up with a relation between e, f from the previous hints. Sub it into the expression!

the pe	ider a group of 6 friends sitting at a lerson sitting directly across from ther on u and v if and only if u and v are f	n. Consider the graph where	each individual is a vertex ar	
on or	n has $n \geq 2$ lightbulbs in a row, all tur turns a lightbulb off. Show that for a guration of lightbulbs being on or off	ny <i>n</i> , there is a a sequence o	of moves that Austin can mak	ce such that each possible
(a) E	Each of the \emph{n} lightbulbs is either on o	r off. How should we represe	nt the lightbulb states mathe	ematically?
(b) F	Frame the problem in terms of hyperc	ubes.		
(c) S	Solve the problem by showing a prope	erty of hypercubes.		
isn't t	vish to color the <i>edge</i> s of a <i>n-</i> dimens the problem in disc2b, where you colo s is necessary and sufficient. (You can	ored vertices so that vertices	s that share an edge are diffe	

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