

## ECE3140

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Problem Set #3

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Q1. a) Average Response Time:  $\frac{(6+10+14+17+26)}{5} = \boxed{14.6}$

Maximum lateness:  $\boxed{3 \text{ (on } J_1)}$

Number of jobs that miss deadline =  $\boxed{5 \text{ (all of them)}}$

b) Average Response Time:  $\frac{(4+8+11+20+26)}{5} = \boxed{13.8}$

Maximum Lateness:  $\boxed{23 \text{ (on } J_1)}}$

Number of jobs that miss deadline:  $\boxed{1}$

Q2.  $\sigma = \langle J_1, J_2, J_3, \dots, J_n \rangle \quad C_i \leq C_k \quad 1 \leq i < k \leq n$

$\sigma$  Avg Response =  $\frac{1}{n} \sum_{i=1}^n (n+1-i) C_i$

$= \frac{1}{n} [nC_1 + (n-1)C_2 + \dots + (n+1-i)C_i + \dots + (n+1-k)C_k + \dots + C_n]$

swap  $i$  and  $k$

$\sigma_{\text{swap}}$  Avg Response =  $\frac{1}{n} [nC_1 + (n-1)C_2 + \dots + (n+1-i)C_k + \dots + (n+1-k)C_i + \dots + C_n]$

Compare two response times

$\left[ (n+1-i)C_i + \dots + (n+1-k)C_k \right] \stackrel{?}{\leq} \left[ (n+1-i)C_k + \dots + (n+1-k)C_i \right]$

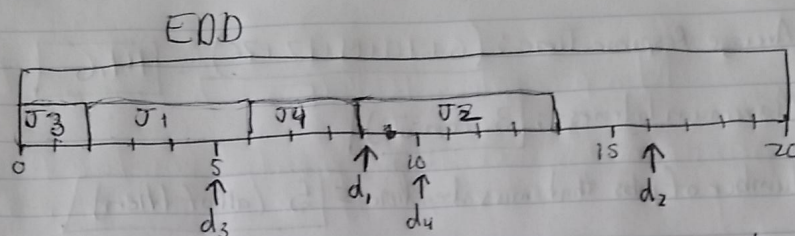
$-iC_i + \dots -kC_k \stackrel{?}{\leq} -iC_k + \dots -kC_i$

$C_k(k-i) \geq C_i(k-i)$

$\boxed{C_k \geq C_i} \rightarrow \text{means } k > i$

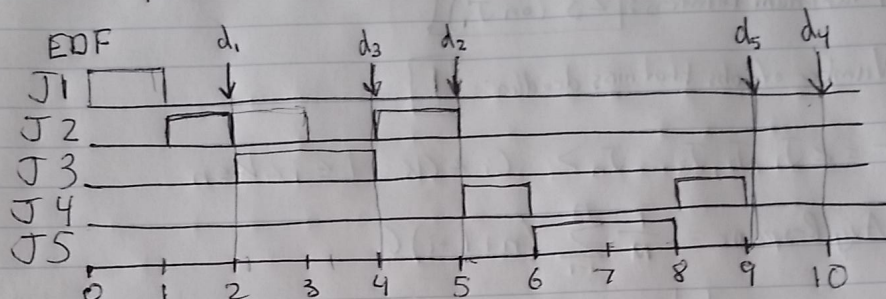
Given  $C_i < C_k$  for all  $1 \leq i < k \leq n$ , therefore, swapping any two jobs results in a new average response time that is not better than the original

Q3.

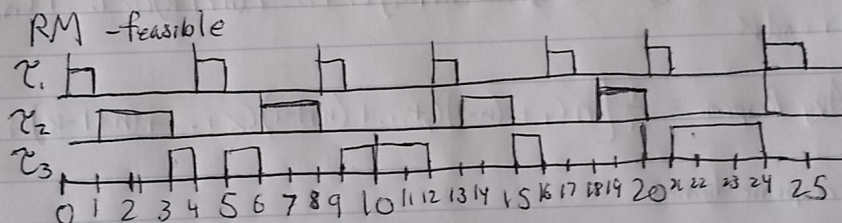


Each finish time is less than its corresponding deadline, so EDD produces a feasible schedule

Q4.

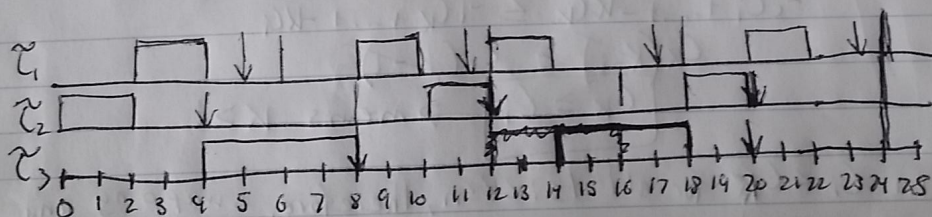


Q5.



Q6. EDF schedulability  $CPU \text{ utilization} = \left( \frac{1}{4} + \frac{2}{6} + \frac{3}{10} \right) = 0.88 < 1$   
therefore tasks are schedulable

Q7.





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Q8. Priority Inheritance protocol:  $P_{\tau_1} > P_{\tau_2} > P_{\tau_3}$

- task can be blocked at most for one critical section of each lower priority task

can be blocked by  $\tau_2$  for 3 seconds (resource A) and  $\tau_3$  for 6 seconds (resource C)

$$\tau_1 \text{ Maximum block time} = (3-1) + (6-1) = 7$$

(resources A and C)