

**Problem Set 1 Solution:**

1. Define the predicate:  $P(n)$ :  $4^n + 15n - 1$  is divisible by 9.

Base case: for  $n = 1$ ,  $4^1 + 15 \times 1 - 1 = 18$  which is divisible by 9.

Induction step: Assume  $P(k)$ , i.e.,  $4^k + 15k - 1$  is divisible by 9, then we need to show  $P(k+1)$ , i.e.,  $4^{k+1} + 15(k+1) - 1$  is divisible by 9.

$$\begin{aligned} & 4^{k+1} + 15(k+1) - 1 \\ &= 4 \cdot 4^k + 15k + 14 \\ &= 4(4^k + 15k - 1) - 45k + 18 \\ &= 4(4^k + 15k - 1) - 9(5k - 2) \end{aligned}$$

The first term in the above result is divisible by 9 by the induction hypothesis; the second term is divisible by 9 because of the factor 9. Therefore  $4^{k+1} + 15(k+1) - 1$  is divisible by 9, i.e.,  $P(k+1)$  is true, which completes the induction step.

2. Attached separately.

3. Consider the Fibonacci-esque function  $g$ :

$$g(n) = \begin{cases} 1, & \text{if } n = 0 \\ 3, & \text{if } n = 1 \\ g(n-2) + g(n-1) & \text{if } n > 1 \end{cases}$$

Use complete induction to prove that if  $n$  is a natural number greater than 1, then  $2^{n/2} \leq g(n) \leq 2^n$ .

**Sample solution:** Proof, using complete induction.

**inductive step:** Let  $n$  be a typical natural number greater than 1 and assume  $H(n)$ : Every natural number  $i \in \{2, \dots, n-1\}$  satisfies  $2^{i/2} \leq g(i) \leq 2^i$ .

**show that inductive conclusion follows:** We'll derive  $C(n)$ :  $2^{n/2} \leq g(n) \leq 2^n$ .

**Base cases:**  $1 < n < 4$ :  $g(2) = 4$  and  $g(3) = 7$  # by the definition of  $g(2)$  and  $g(3)$ .

$$2^{2/2} = 2 \leq 4 = g(2) \leq 4 = 2^2 \quad \text{and} \quad 2^{3/2} = 2\sqrt{2} \leq 7 = g(3) \leq 8 = 2^3$$

$C(2)$  and  $C(3)$  follow from our assumptions in this case.

**Case  $n \geq 4$ :** By assumption  $H(n)$  we have #  $n \geq 4$  implies  $2 \leq n-2, n-1 < n$ :

$$2^{(n-2)/2} \leq g(n-2) \leq 2^{n-2} \quad \text{and} \quad 2^{(n-1)/2} \leq g(n-1) \leq 2^{n-1}.$$

Substituting these inequalities into the definition of  $g(n)$  # by definition of  $g(n)$ ,  $n \geq 4 > 0$ :

$$\begin{aligned} g(n) &= g(n-2) + g(n-1) \geq 2^{(n-2)/2} + 2^{(n-1)/2} = (1 + \sqrt{2})2^{(n-2)/2} \geq 2 \times 2^{(n-2)/2} = 2^{n/2} \\ g(n) &= g(n-2) + g(n-1) \leq 2^{n-2} + 2^{n-1} = (1 + 2)2^{n-2} \leq 2^2 \times 2^{n-2} = 2^n \end{aligned}$$

$C(n)$  follows from our assumptions in this case.

In all cases  $H(n)$  implies  $C(n)$ .

4. If  $n = 7^k$  then we will get that  $L(n) = L(7^k) = 1 + L(7^{k-1}) = 1 + 1 + L(7^{k-2}) = \dots = k + L(1) = k = \log_7 n$ . Rather than proving it now, we will prove for the general case that  $L(n) = \lfloor \log_7 n \rfloor$  for  $n > 0$  and  $L(0) = 0$ . We show this by complete induction. Since  $L(n) = 0 = \lfloor \log_7 n \rfloor$  for  $0 < n < 7$  and since  $L(0) = 0$  we need only to check  $L(n)$  for  $n > 6$ .

$$L(n) = 1 + L(\lfloor n/7 \rfloor) + 1 = \lfloor \log_7 \lfloor n/7 \rfloor \rfloor + 1 = (\lfloor \log_7 n \rfloor - 1) + 1 = \lfloor \log_7 n \rfloor.$$

For  $T$ , we have the recursion,

$T(n) = 3$  if  $n < 7$  and otherwise  $T(n) = 4 + T(\lfloor n/7 \rfloor)$ . Unwinding will give  $T(n) = 4 * \log_7 n + T(1)$  for  $n$  which is a power of 7. We guess  $T(n) = 4 * \lfloor \log_7 n \rfloor + 3$  for a general  $n$ . Proving this by induction is essentially the same as with  $L$ .

5. We show that each element  $n$  of  $F$  can be written as  $7k$  for some integer  $k$ .

Base case:  $7 = 7 \times 1$ , so  $7 \bmod 7$  is 0 as required.

Inductive case: suppose that  $m \in F$  and that  $m = 7p$  for some integer  $p$ . Also suppose that  $n \in F$  and that  $n = 7q$  for some integer  $q$ . We now verify that the second rule generates integers divisible by 7:

$$\begin{aligned} m + n &= 7p + 7q \\ &= 7(p + q) \end{aligned}$$

which is divisible by 7.