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DOCTORAL THESIS

A Language of Polynomials

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in the*

Research Group Name
Department or School Name

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Declaration of Authorship

I, Eric UNG, declare that this thesis titled, “A Language of Polynomials” and the work presented in it are my own. I confirm that:

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“Thanks to my solid academic training, today I can write hundreds of words on virtually any topic without possessing a shred of information, which is how I got a good job in journalism.”

Dave Barry

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Abstract

Faculty Name
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Doctor of Philosophy

A Language of Polynomials

by Eric UNG

The Thesis Abstract is written here (and usually kept to just this page). The page is kept centered vertically so can expand into the blank space above the title too...

Acknowledgements

The acknowledgments and the people to thank go here, don't forget to include your project advisor...

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List of Abbreviations

LAH List Abbreviations **Here**
WSF What (it) Stands **For**

Physical Constants

Speed of Light $c_0 = 2.997\,924\,58 \times 10^8 \text{ m s}^{-1}$ (exact)

List of Symbols

a	distance	m
P	power	W (J s ⁻¹)
ω	angular frequency	rad

For/Dedicated to/To my...

k

Chapter 1

A Language of Polynomials

1.1 Introduction

This paper is on the re-framing of the one way function to a matrix multiplication problem - that of multiplying two 3×3 matrices to form a 6×6 matrix under a locally concatenative property.

1.2 Foundations

There exists a language such that it decides each monomial in the polynomial. In other words, there exists a set of deciders for each monomial in the polynomial where it decides if y is in the monomial. A decider in this term is not of the definition found originally in textbooks but one that is redefined in the below definition.

Given a polynomial

$$p(x) = ax^2 + bx + c$$

$$p(x) = 3x^2 + 4x + 5$$

$$p(2) = 3(2)^2 + 4(2) + 5$$

$$p(2) = 12 + 8 + 5$$

Let the decider be defined as the following:

Decider is a function $Decider < c \times x^{degree} > \equiv c \times x^{degree} = y$

such that $x_1 \times x_2 \times \dots \times x_n$ where n is equal to degree + constant is tested to be equivalent to y and x_1 is the start and x degree times is the finish then loop around x_1 to x degree times until it stops

For each state, x_i , i such that it is between 1 to n , x_i contains a subgroup of size n and for each subgroup, s_i , there exists another subgroup and so on and so forth such that there are n layers starting from x_i to 1. This is the same as saying that it is a rational expression.

A rational expression is a expression that satisfies the following.

$$A_n = A_{n-1} \cup \{E^* | E \in \mathcal{E}_{n-1}, (E, 1) = 0\}$$

It follows that each state $E \in s_{n-1}$ forms a subgroup.

A rational function is defined as the following:

$K[x]$ and $K[[x]]$. Let $K[[x]]$ describe a set of monomial deciders. S is an element of $K[[x]]$ meaning S is a monomial decider.

$$S = \sum_{n \geq 0} a_n x^n$$

Examples

Decider for ax^2 is $Decider(3, 2, 2) = 3(2)^2 \equiv 12$

Decider for bx is $Decider(4, 2, 1) = 4(2)^1 \equiv 8$

Decider for c is $Decider(1, 2, 1) = 5(2)^0 \equiv 5$

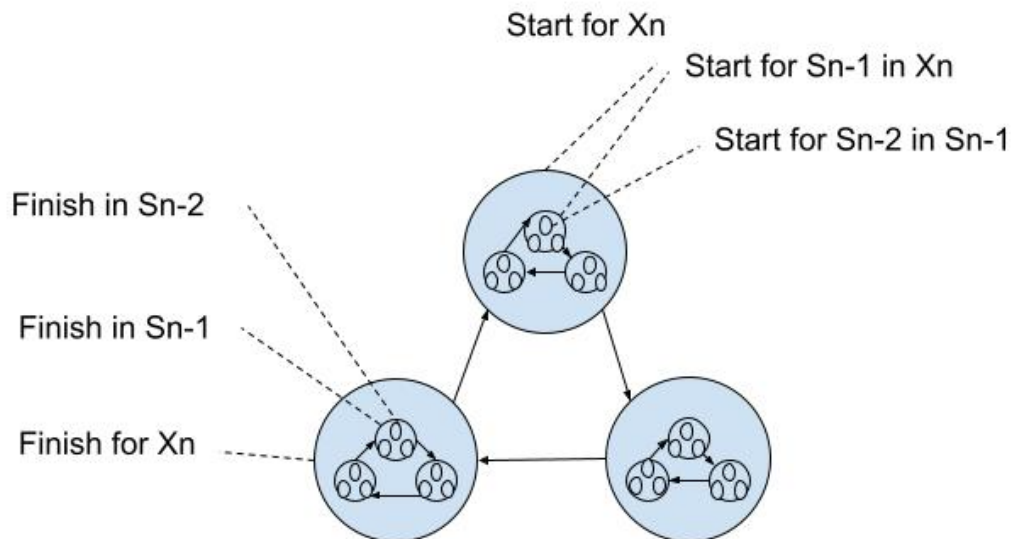


FIGURE 1.1: Decider X to the 3.

Figure 1.1 shows a monomial decider, $Decider < x^n >$ with that represents x^3 .

Theorem: A Decider is the equivalent to a cyclic automata

Proof: Remove the lowest level state, S_1 from the bottom then continue removing S_i from $i = 2$ to $n-1$ until you get only the states that are at X_n .

Remove the top state down from the S_{n-1} for each layer S_{n-1} to $S(n - n - 1)$. This preserves the start and finish state for the layer x_n . This is a cyclic automaton.

Figure 1.2 shows how the intuition for removing the state top down.

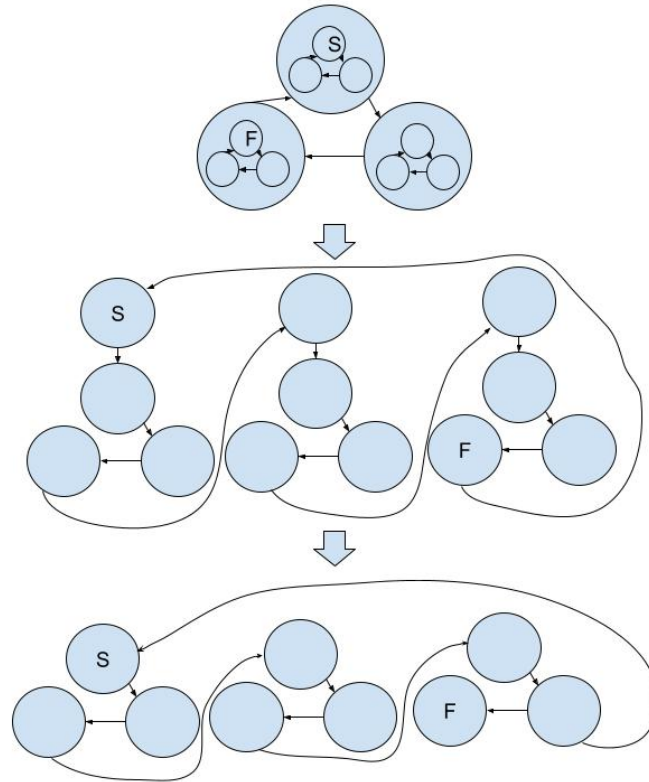


FIGURE 1.2: Top down removal for equivalence of decider and cyclic automata.

1.3 Monomial of One Degree

Given the definition of a decider:

A decider of at least one degree

$$Decider(3, x, 4) \equiv 3x^4 = y$$

$$\text{Contains } Decider(3, x, 3) \equiv 3x^3 = y$$

$$\text{Contains } Decider(3, x, 2) \equiv 3x^2 = y$$

$$\text{Contains } Decider(3, x, 1) \equiv 3x^1 = y$$

$$\text{Contains } Decider(3, x, 0) \equiv 3x^0 = y$$

Hence it can be generalized to:

$Decider(constant, degree, x)$ contains the sequence set

$Deciderconstant, degree, x, Deciderconstant, degree1, x, ..., Deciderconstant, 0, x$
 $\{start, ..., finish\}$

1.4 Addition

Given the first example:

$$p(x)=3x^2+4x+5$$

$$p(2)=3(2)^2+4(2)+5$$

$$p(2)=12+8+5$$

$$m_2=Decider(3,x,2)=3x^2$$

$$m_1=Decider(4,x,1)=4x$$

$$m_0=Decider(5,x,0)=5$$

Generalized to mx where x is the degree

Given polynomial functions, p_1 and p_2 , they are commutative

$$p_1(x) = m_a + \dots + m_0$$

$$p_2(x) = n_b + \dots + n_0$$

$$p_1(x) + p_2(x) = m_a + m[x+1] + n_b + n[y+1] + \dots + (m_x + n_y) + \dots + (m_0 + n_0) \text{ where } x=y$$

$$\text{Decider}(cx, x, dx)$$

$$+ \text{Decider}(cy, x, dy) = \text{Decider}(cx+cy, x, dx) = \text{Deciders}(cx+cy, x, dy) = dx=dy$$

1.5 Product

Given two monomials in the language, a and b , the product of a and b is also in the language.

Given $\text{Decider}(cx, x, dx)$ and $\text{Decider}(cy, x, dy)$ is in language L

Show that the product $\text{Decider}(cx+cy, x, dx \times dy)$ is in L

$$\text{Decider}(cx, x, dx) \times \text{Decider}(cy, x, dy)$$

$$= cx \times xdx \times cy \times xdy$$

$$= cx \times xdx + dy$$

$$= (cx+cy) \times xdx + dy \text{ is in } L$$

$$= \text{Decider}(cx+cy, x, dx+dy)$$

1.6 Problem with Matrices

Given a polynomial p of x , show that the monomial deciders represented in the language can't be contained in a finite matrix after a set number, n , such that xn .

$$axa \times bxb = nxn \text{ such that } a \neq b \text{ and } a, b < n$$

1.7 Multivariable Polynomials

A monomial with more than one variable can be treated the same way as handling single variables at different degrees.

$\text{Decider}(c, xyz, d) = c(xyz)d = \text{Decider}(c, x, d) \times \text{Decider}(c, y, d) \times \text{Decider}(c, z, d)$ where c is some constant $\text{Decider}(c, x, d_1) \times \text{Decider}(c, y, d_2) \times \text{Decider}(c, z, d_3) = cxd_1 \times yd_2 \times zd_3$ where c is some constant

$$\text{Given } f(x, y) = 3xy^2$$

$$\text{set } x=2, y=3$$

$$f(2, 3) = 3(2)(3)^2$$

$$f(2, 3) = 27$$

1.8 Generalized Monomial Deciders

A monomial decider can be represented in a more general graphic

$$6x6 = \text{Decider } 6, x, 6 = 6 \times x_{\text{start}}, x_2, x_3, x_4, x_5, x_{\text{finish}} \quad x6 = \text{Decider } 1, x, 6 = 1 \times x_{\text{start}}, x_2, x_3, x_4, x_5, x_{\text{finish}}$$

In both these examples, x_{start} is x_1 and x_{finish} is x_6 .

1.9 Concentric Monomial Deciders

Given a polynomial, p x , with a monomial decider represented as ax^n in p x and n in N and $a = 1$ such that p $x = x^n$, there is special property for these monomial deciders that can be illustrated below.

Decider $1, x, 4 = x = x_i$ such that i is in $start, 2, 3, finish$ and x_i contains x_i minus 1 where $i \neq 1$
 Decider $1, x, 2 = x_2 = x_i$ such that i is in $start, finish$ and x_i contains x_i minus 1 where $i \neq 1$
 In both these examples, x_i is a state in the monomial decider. Each state in a one constant monomial decider have a property of being concentric. This means that a state can be defined as, $x_2 = x_1, x_0$ so going from x_1 to x_0 then going to the next state x_3 or looping back to x_1 .

1.10 Constants

Given a constant, c , of p or better described in the example: $f(x) = 5$. Constants are seen as linear directed acyclic graphs.

$$f(x) = 5$$

$$Decider(5, x, 0) \equiv 5 \equiv start, 2, 3, 4, finish$$

There is no state in the decider where it loops back to the start. In other words, there is no x that represents a monomial in a constant.

1.11 Division

Division of monomial deciders

$$\begin{aligned} x^5/x &= x^4 \equiv Decider(1, x, 5)/Decider(1, x, 1) \quad Decider(1, x, 4) \quad x^5/x^2 = x^3 \equiv \\ &Decider(1, x, 5)/Decider(1, x, 2) \quad Decider(1, x, 3) \quad x^5/x^3 = x^2 \equiv Decider(1, x, 5)/Decider(1, x, 3) \quad Decider(1, x, 4) \\ x^5/x^4 &= x^1 \equiv Decider(1, x, 5)/Decider(1, x, 4) \quad Decider(1, x, 1) \quad x^5/x^5 = 1 \equiv Decider(1, x, 5)/Decider(1, x, 5) \end{aligned}$$

here is represented as a special kind of equivalence that we will get to later.

$$\begin{aligned} Decider(1, x, 5)/Decider(1, x, 1) &\equiv \text{sequence of permutations of } x_i, x_j \text{ such that the count of } i \text{ is 4 and } j \text{ is 1} \\ Decider(1, x, 5)/Decider(1, x, 2) &\equiv \text{sequence of permutations of } x_i, x_j \text{ such that the count of } i \text{ is 3 and } j \text{ is 2} \\ Decider(1, x, 5)/Decider(1, x, 3) &\equiv \text{sequence of permutations of } x_i, x_j \text{ such that the count of } i \text{ is 2 and } j \text{ is 3} \\ Decider(1, x, 5)/Decider(1, x, 4) &\equiv \text{sequence of permutations of } x_i, x_j \text{ such that the count of } i \text{ is 1 and } j \text{ is 4} \\ Decider(1, x, 5)/Decider(1, x, 5) &\equiv \text{sequence of permutations of } x_i, x_j \text{ such that the count of } i \text{ is 0 and } j \text{ is 5} \end{aligned}$$

Here, $x_i \neq x_j$, meaning x_i is a different representation than x_j

1.12 Multiple Divisions

Given multiple operations of division, this forms an interesting space. $x^5/x^2/x = x^5/x)/x^2$

$$Decider(1, x, 5)/Decider(1, x, 2)/Decider(1, x, 1) = Decider(1, x, 5)/Decider(1, x, 1)/Decider(1, x, 2) \text{ iff ignoring order of operations}$$

1.13 Equivalence

$Decider(1, x, 6)/Decider(1, x, 1)/Decider(1, x, 1) = Decider(1, x, 6)/Decider(1, x, 2)$ iff the order of operations is next to each other

Determining if y is in $f(x)$ is easy if we are given any monomial decider in the set of the language of polynomials and their representations has the possibility to give different representations if we consider them as representations of the function $f(x)$.

$Decider(1, x, 6)/Decider(1, x, 1)/Decider(1, x, 1) = \text{sequence of permutations of } x_i, x_j, x_k$ such that the count of i is 4 and j is 1 and k is 1

$Decider(1, x, 6)/Decider(1, x, 2) = \text{sequence of permutations of } x_i, x_j$ such that the count of i is 4 and j is 2

Theorem of Equivalence Decider $1,x,6$ / Decider $1,x,1$ / Decider $1,x,1$ = Decider $1,x,6$ / Decider $1,x,2$ such that there is some x such that the monomial represented by both deciders exists where $f(x) = y$

1.14 Theorem of Equivalence

$Decider(1, x, 6) / Decider(1, x, 1) / Decider(1, x, 1)$
 $\equiv Decider(1, x, 6) / Decider(1, x, 2)$

such that there is some x

such that the monomial represented by both deciders exists where $f(x) \equiv y$

1.15 Reversing

$Decider(1, x, 6) / Decider(1, x, 1) / Decider(1, x, 1) \equiv$ sequence of permutations of x_i, x_j, x_k such that the
 $Decider(1, x, 6) / Decider(1, x, 2) \equiv$ sequence of permutations of x_i, x_j such that the count of i is 4 and

Is shown that by the permutation of the order of operations that $Decider(1, x, 6) / Decider(1, x, 1) / Decider(1, x, 1)$ is not the same set as $Decider(1, x, 6) / Decider(1, x, 2)$

Theorem of Reversing

Given two deciders x, y in a decider of $m(x)$, $x \neq y$ iff sequence of all the states are not equivalent, representation-wise, from start to finish or it doesn't represent the same order of operations of the deciders being represented

1.16 Corollary of Reversing

A little more on the theorem of reversing Given a starting point, the paths a monomial decider takes to decide if y is in $f(x)$ is inherently unique to each representation. Start at the circle, S , and end at the circle, F . If the circle is white, it is 0. If it is blue, it is 1. As an example let's traverse some of the representations of x 4.

Corollary Given a decider, d , in Decider of $m(x)$ then there is path, p , that exists for d such that $p = \text{Path } d = s_1, s_2, \dots, s_i, \dots, s_n$ where i is count of the states in the decider of $m(x)$.

Example: Choose some x such that it is in Path Decider $1,x,6$ / Decider $1,x,2$ where $p = 001111$

1.17 Godel's Theorem

We see that there exists two statements from these theorems

1. $x = x$ from a theorem of equivalence
2. $x \neq x$ from a theorem of reversal

Example: Given some $d_1, d_2, d_3, \dots, \text{infinity}$ in deciders of $m(x)$ 1. $d_1 = d_2 = d_3 = \dots = \text{infinity}$ 2. $d_1 \neq d_2 \neq d_3 = \dots \neq \text{infinity}$

The different representations of a monomial through the language of monomial deciders will give us undecidability. This means we can come up with many formal definitions of the monomial decider and it will not be able to solve the problem of finding a specific representation of a monomial decider without having to guess or apply some sort of probability to it. Relating to the real line, given a real line a, b $a \leq b$, there is infinite choices between a and b because we can just make the number smaller. As long as $b \geq 0$, there requires some sort of probability of choosing some specific number that is between a and b .

1.18 Reframing The One Way Function

A probability exists to find a certain monomial decider in the set of its variations. $A/B = \text{Probability}$ where A is the monomial decider we want and B is the number of all the variations.

Example: d_1, d_2, \dots, d_6 in deciders $1, x, 6, D$, such that d_i are all distinct. Choose one of the deciders in D through probability. Probability of choosing d in D is $1/6$ so 0.16666667 . We'll call this picking a function and every time we call this function, the probability is multiplied such that it is n^k . As an example, if we call the picking function twice using the example above, we have $1/6 \cdot 1/6 = 1/36 = 0.02777778$.

This is formally known as the one way function.

1.19 Theorem of Infiniteness

Given that, if we can show for any language it abides a theorem of equivalence and a theorem of reversal and they both have infinite representations, we know that we need some sort of probability to choose a specific representation of a monomial decider. We'll call this a theorem of infiniteness because if we can show that some language is infinite, it will require some sort of guess to pick something unique out of all the things it represents, generates, or describes. In other words, you can't map infinity to infinity directly for you must map infinity to something discrete and something discrete to infinity.

Theorem A mapping must be from infinity to discrete representation to discrete representation to infinity.

Given any representation of infinity, Suppose a mapping infinity to infinity exists. Then this map is equivalent to infinity because infinity contains this map. Here, infinity is represented as $*$. $*$ contains $*$ — $> *$.

Intuition

Given Decider c, x, d and d is infinity, Decider c, x, d contains Decider c, x, d 1 d 1 is then infinity too.

1.19.1 References

The `bi-latex` package is used to format the bibliography and inserts references such as this one (**Reference1**). The options used in the `main.tex` file mean that the in-text citations of references are formatted with the author(s) listed with the date of the publication. Multiple references are separated by semicolons (e.g. (**Reference2**; **Reference1**)) and references with more than three authors only show the first author with *et al.* indicating there are more authors (e.g. (**Reference3**)). This is done automatically for you. To see how you use references, have a look at the `Chapter1.tex` source file. Many reference managers allow you to simply drag the reference into the document as you type.

Scientific references should come *before* the punctuation mark if there is one (such as a comma or period). The same goes for footnotes¹. You can change this but the most important thing is to keep the convention consistent throughout the thesis. Footnotes themselves should be full, descriptive sentences (beginning with a capital letter and ending with a full stop). The APA6 states: "Footnote numbers should be superscripted, [...], following any punctuation mark except a dash." The Chicago manual of style states: "A note number should be placed at the end of a sentence

¹Such as this footnote, here down at the bottom of the page.

or clause. The number follows any punctuation mark except the dash, which it precedes. It follows a closing parenthesis.”

The bibliography is typeset with references listed in alphabetical order by the first author’s last name. This is similar to the APA referencing style. To see how \LaTeX typesets the bibliography, have a look at the very end of this document (or just click on the reference number links in in-text citations).

A Note on bibtex

The bibtex backend used in the template by default does not correctly handle unicode character encoding (i.e. "international" characters). You may see a warning about this in the compilation log and, if your references contain unicode characters, they may not show up correctly or at all. The solution to this is to use the biber backend instead of the outdated bibtex backend. This is done by finding this in `main.tex`: `backend=bibtex` and changing it to `backend=biber`. You will then need to delete all auxiliary BibTeX files and navigate to the template directory in your terminal (command prompt). Once there, simply type `biber main` and biber will compile your bibliography. You can then compile `main.tex` as normal and your bibliography will be updated. An alternative is to set up your LaTeX editor to compile with biber instead of bibtex, see [here](#) for how to do this for various editors.

Chapter 2

Chapter Title Here

2.1 Main Section 1

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2.1.1 Subsection 1

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2.1.2 Subsection 2

Morbi rutrum odio eget arcu adipiscing sodales. Aenean et purus a est pulvinar pellentesque. Cras in elit neque, quis varius elit. Phasellus fringilla, nibh eu tempus venenatis, dolor elit posuere quam, quis adipiscing urna leo nec orci. Sed nec nulla auctor odio aliquet consequat. Ut nec nulla in ante ullamcorper aliquam at sed dolor. Phasellus fermentum magna in augue gravida cursus. Cras sed pretium lorem. Pellentesque eget ornare odio. Proin accumsan, massa viverra cursus pharetra, ipsum nisi lobortis velit, a malesuada dolor lorem eu neque.

2.2 Main Section 2

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Appendix A

Frequently Asked Questions

A.1 How do I change the colors of links?

The color of links can be changed to your liking using:

```
\hypersetup{urlcolor=red}, or  
\hypersetup{citecolor=green}, or  
\hypersetup{allcolor=blue}.
```

If you want to completely hide the links, you can use:

```
\hypersetup{allcolors=.}, or even better:  
\hypersetup{hidelinks}.
```

If you want to have obvious links in the PDF but not the printed text, use:

```
\hypersetup{colorlinks=false}.
```