

Problem 1.

(35 pts) Consider a disk with block size $B = 512$ bytes. A block pointer is $P = 6$ bytes long, and a record pointer is $PR = 7$ bytes long. A file has $r = 100,000$ EMPLOYEE records of fixed length. Each record has the following fields:

NAME (25 bytes),

SSN (9 bytes),

DEPARTMENTCODE (9 bytes),

ADDRESS (40 bytes),

PHONE (9 bytes),

BIRTHDATE (8 bytes),

SEX (1 byte),

JOB CODE (5 bytes),

SALARY (4 bytes).

- a. Calculate the record size R in bytes.

$$\text{Record size } R = (25 + 9 + 9 + 40 + 9 + 8 + 1 + 5 + 4) = 110 \text{ bytes}$$

- b. Calculate the blocking factor bfr and the number of file blocks b , assuming an unspanned organization.

Blocking factor, “ bfr ” = $\text{floor}(B / R) = \text{floor}(512 / 110) = 4$ records/block

of file blocks, “ b ” = $\text{ceiling}(r / bfr) = (100,000 / 4) = 25000$ blocks needed for file

- c. Suppose that the file is ordered by the key field SSN and we want to construct a primary index on SSN. Calculate

- i. the index blocking factor bfr_i

Index record size, $R_i = (SSN + P) = (9 + 7) = 16$ bytes

Index blocking factor, $bfr_i = \text{floor}(B / R_i) = \text{floor}(512 / 16) = 32$

- ii. the number of first-level index entries and the number of first-level index blocks

of 1st-level index entries, $r_1 = \#$ of file blocks, $b = 25000$ entries

of 1st-level index blocks, $b_1 = \text{ceiling}(r_1 / bfr_i) = \text{ceiling}(25000 / 32) = 782$ blocks

- iii. the number of levels needed if we make it into a multilevel index

of 2nd-level index entries, $r_2 = \#$ of 1st-level blocks, $b_1 = 782$ entries

of 2nd-level index blocks, $b_2 = \text{ceiling}(r_2 / bfr_i) = \text{ceiling}(782 / 32) = 25$ blocks

of 3rd-level index entries, $r_3 = \#$ of 2nd-level blocks, $b_1 = 25$ entries

of 3rd-level index blocks, $b_3 = \text{ceiling}(r_3 / bfr_i) = \text{ceiling}(25 / 32) = 1$ blocks

CSE4701 - Databases**HW2 2020-03-04**

- iv. the total number of blocks required by the multilevel index, and

Total # of blocks for the index $b_i = b_1 + b_2 + b_3 = 782 + 25 + 1 = 808$ blocks

- v. the number of block accesses needed to search for and retrieve a record from the file—given its SSN value—using the primary index.

of block access to search a record = $x + 1 = 3 + 1 = 4$ block accesses

Problem 2. (15 pts) Given the same specifications of Problem 1, consider this time you are building a primary index on SSN using B-tree. Calculate ...

- i. the order p for the B-tree, ...

$$(block\ pointer \cdot p) + (key\ pointer + block\ pointer) \cdot (p - 1) \leq 512$$

$$(6 \cdot p) + (9 + 6) \cdot (p - 1) \leq 512$$

$$6p + 15p - 15 \leq 512$$

$$21p \leq 527$$

$$p = 25$$

- ii. the number of levels needed if blocks are approximately 69% full (round up for convenience), and...

$$p \cdot 0.69 = 25 \cdot 0.69 = 17.25 \approx 18$$

- need 25,000 blocks

	# of nodes	# of pointers	# of key values
Root:	1 node	18 ptrs	17 entries
Level 1:	18 nodes	18 x 18 ptrs	18 x 17 entries
Level 2:	18 x 18 nodes	18 x 18 x 18 ptrs	18 x 18 x 17 entries
Level 3:	18 x 18 x 18 nodes		18 x 18 x 18 x 17 = 99,144 entries

- $\log_{18}(25000) = 3.50357 \rightarrow 4$ levels

- iii. the worst-case number of blocks needed to search for and retrieve a record from the file—given its SSN value—using the B-tree you are estimating.

Worst-case block searches = # of levels + 1 = 4 levels + 1 = 5 blocks searches

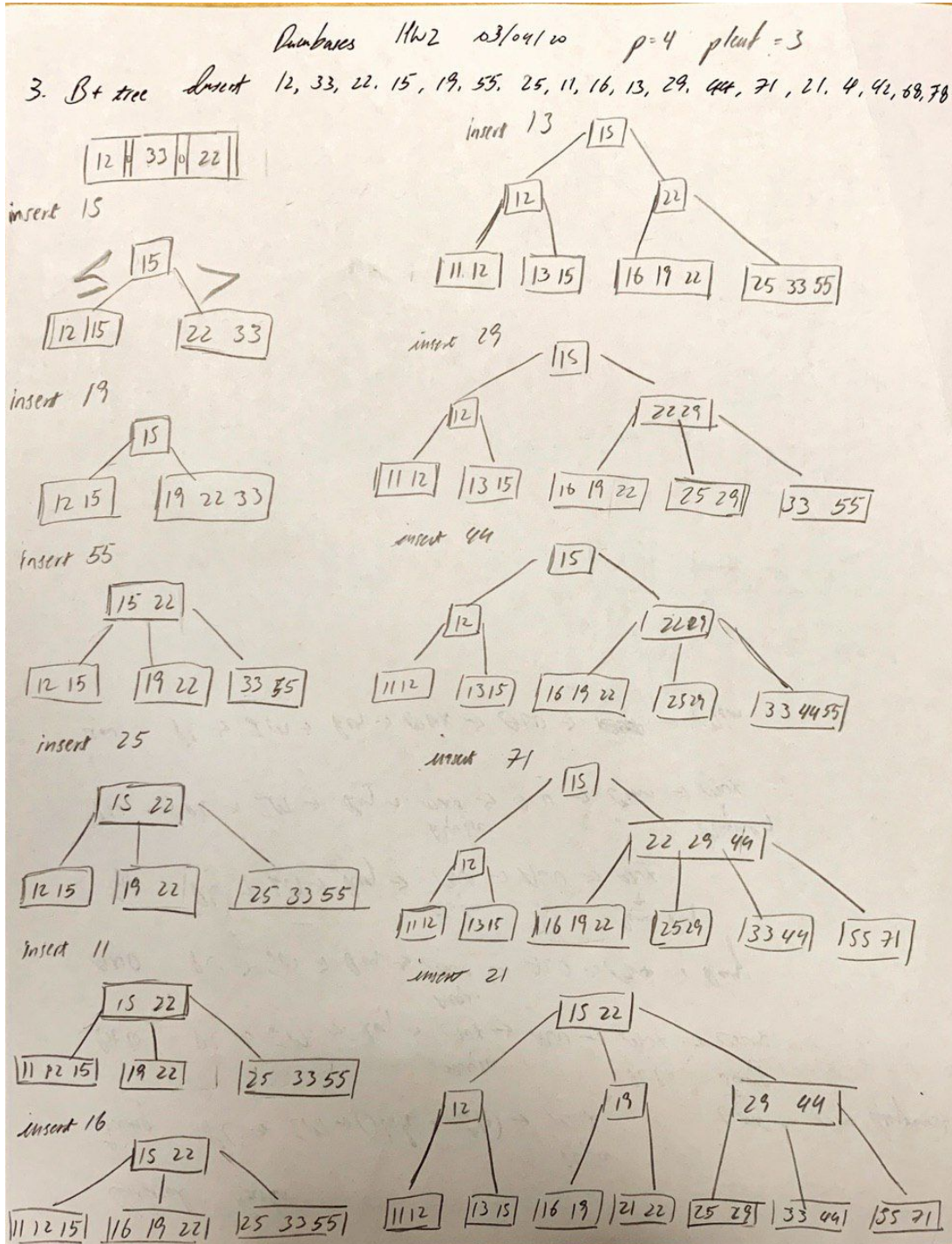
Problem 3. (25 pts) A PARTS file with Part# as key field includes records with the following Part# values: 12, 33, 22, 15, 19, 55, 25, 11, 16, 13, 29, 44, 71, 21, 4, 42, 68, 78.

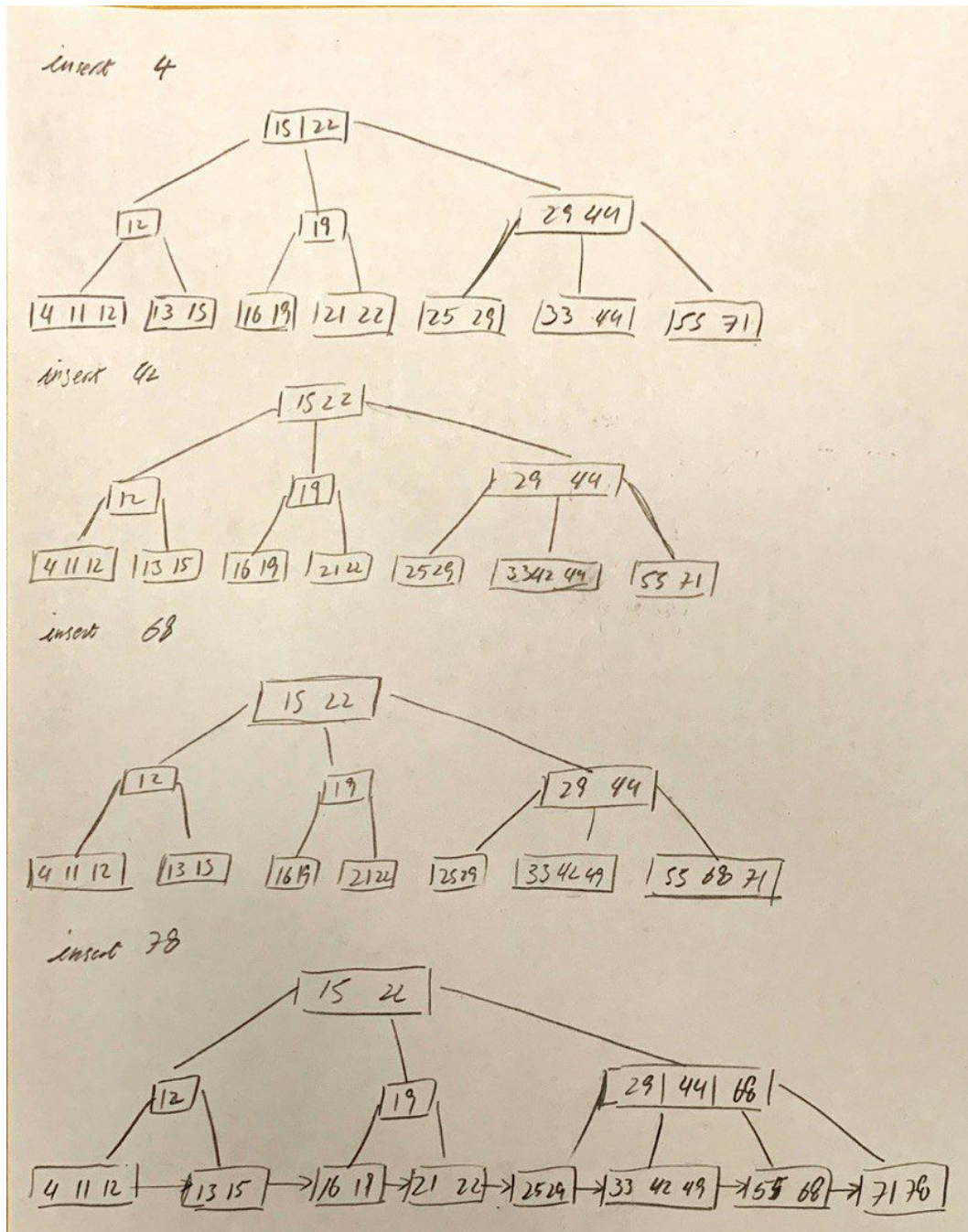
CSE4701 - Databases

HW2 2020-03-04

Suppose that the search field values are inserted in the given order in a B+-tree of order $p = 4$ and $pleaf = 3$

- iv. Show how the tree will expand (show all steps as in Fig 18.12 (6th ed)) and what the final tree will look like.





- v. What is the fill ratio of the final B+-tree you created? (Note: we learned 69% is the average fill ratio in class.) Here the fill ratio is defined as (# of filled key values)/(# of maximally allowed key values) regardless each key value resides in index node or sequence set node.

$$\text{Fill Ratio} = \frac{\# \text{ of nodes in tree}}{\# \text{ of possible nodes}} = \frac{25}{36} = 69.44\%$$