

WHY LOG PROBABILITIES ?

CODING THEORY

RELATIVE ENTROPY

WANT TO SEND SYMBOL  $x$  ON CHANNEL

$x$	$P(X=x_i)$	$-\log P(x_i)$	
$x_1$	$\frac{1}{2}$	1	} BITS
$x_2$	$\frac{1}{4}$	2	
$x_3$	$\frac{1}{8}$	3	
$x_4$	$\frac{1}{8}$	3	

MEASURE OF SURPRISE

$$-\log 1 = 0 \quad \text{NO SURPRISE}$$

$$-\log 0 = \infty \quad \text{INFINITE "}$$

$$-\log 2^i = i \quad \text{BITS}$$

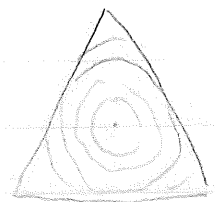
ENTROPY EQUALS EXPECTED SURPRISE

$$H(X) := \sum_i p(x_i) \log_2 \frac{1}{p(x_i)}$$

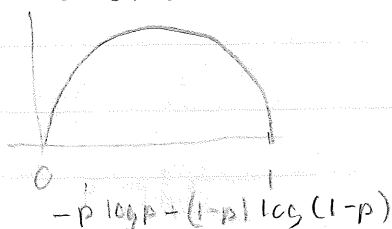
$$= \frac{1}{2} \cdot 1 + \frac{1}{4} \cdot 2 + \frac{1}{8} \cdot 3 + \frac{1}{8} \cdot 3$$

$$= 1 \frac{3}{4}$$

3 DIM

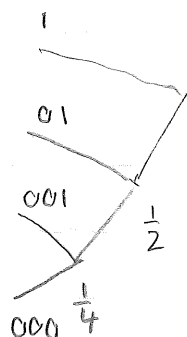


2 DIM



HUFFMAN CODE

$x_1$	$\frac{1}{2}$
$x_2$	$\frac{1}{4}$
$x_3$	$\frac{1}{8}$
$x_4$	$\frac{1}{8}$



LOOP



PICK SMALLEST TWO  
COMBINE BOTH INTO ONE  
SUM THEIR PROBS

ENTROPY = EXPECTED CODELENGTH

X

$x_1$	$\frac{1}{3}$
$x_2$	$\frac{1}{3}$
$x_3$	$\frac{1}{3}$



$$H(X) = \log_2 3 = 1.58 \text{ BITS}$$

$$\begin{aligned} \text{EXPECTED CODELENGTH} &= \frac{1}{3} (1 + 2 + 2) \\ &= \frac{5}{3} \\ &= 1.66 \text{ BITS} \end{aligned}$$

CODE: ASSIGNS SYMBOLS A BITSTRING (CODEWORD)

- ANY SEQUENCE OF CODEWORDS MUST  
BE UNIQUELY DECODABLE

$$L(C) = \sum_i p(x_i) l_C(x_i) \quad \leftarrow \text{EXPECTED CODELENGTH}$$

$\uparrow$  CODE                       $\uparrow$  LENGTH OF CODEWORD FOR  $x_i$

OPTIMAL CODE  $C^*$

- MINIMUM  $L(C)$

THM:  $H(X) \leq L(C^*) \leq H(X) + 1$

THM: HUFFMAN CODES ARE OPTIMAL

MORE INFO:

FIRST FIVE CHAPTERS OF  
COVER & THOMAS

# RELATIVE ENTROPY PROBABILITY VECTORS

$$\Delta(\bar{p}, \bar{q}) = \sum_i p_i \ln \frac{p_i}{q_i}$$

 $H(\bar{p})$ 

SYMBOL USED FOR  
DIVERGENCES

$$= \sum_i p_i \ln \frac{1}{q_i} - \sum_i p_i \ln \frac{1}{p_i}$$

EXPECTED CODELENGTH  
OF BEST CODEBOOK  
FOR  $\bar{q}$

EXPECTED CODELENGTH  
OF BEST CODEBOOK  
FOR  $\bar{p}$

BOTH EXPECTATIONS ARE WRT  $\bar{p}$

$$\Delta(\bar{p}, (\frac{1}{n})) = \sum_i p_i \ln \frac{p_i}{\frac{1}{n}}$$

$$= \sum_i p_i \ln p_i + \sum_i p_i \ln n$$

$$= \ln n - H(\bar{p})$$

$$\geq 0$$

= 0 AT CORNERS OF SIMPLEX

## MAPLE PLOTS

$$\Delta(\bar{p}, \bar{q})$$

↑  
VAR

$$\Delta(\bar{p}, \bar{q})$$

↑  
VAR

NOT TOO STEEP AT  
BOUNDARY

VERY STEEP

BARRIERS

FOR SIMPLEX