

# Random graphs Network Science Lecture 3

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#### Network models



#### Empirical network features:

- Power-law (heavy-tailed) degree distribution
- Small average distance (graph diameter)
- Large clustering coefficient (transitivity)

#### Generative models:

- Random graph model (Erdos & Renyi, 1959)
- "Small world" model (Watts & Strogatz, 1998)
- Preferential attachement model (Barabasi & Albert, 1999)

## Random graph model



Graph  $G\{E, V\}$ , nodes n = |V|, edges m = |E| Erdos and Renyi, 1959.

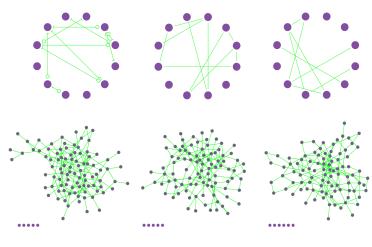
#### Random graph models

- G(n, m) a randomly selected graph from the set of  $C_N^m$  graphs,  $N = \frac{n(n-1)}{2}$ , with n nodes and m edges
- G(n,p) each pair out of  $N=\frac{n(n-1)}{2}$  pairs of nodes is connected with probability p,m random number

$$\langle m \rangle = p \frac{n(n-1)}{2}$$
  $\langle k \rangle = \frac{1}{n} \sum_{i} k_{i} = \frac{2\langle m \rangle}{n} = p (n-1) \approx pn$  
$$\rho = \frac{\langle m \rangle}{n(n-1)/2} = p$$

# Random graph model G(n, p)





$$n = 12, p = 1/6$$

$$n = 100, p = 0.03$$

## Random graph model G(n, p)



• In G(n, p) model, probability for a network to have m links is given by binomial distribution:

$$P(m) = C_N^m p^m (1-p)^{N-m}$$

where 
$$N = \frac{n(n-1)}{2}$$

- $p^m$  probability that m links are present  $(1-p)^{N-m}$  probability that other links are not  $C_N^m$  number of ways to select m links out of all N,  $C_N^m = \frac{N!}{m!(N-m)!}$
- expected number of links

$$\langle m \rangle = \sum_{m=0}^{N} mP(m) = pN = p\frac{n(n-1)}{2}$$

## Degree distribution



• Probability that *i*-th node has a degree  $k_i = k$  is given by Binomial distribution:

$$P(k_i = k) = P(k) = C_{n-1}^k p^k (1-p)^{n-1-k}$$

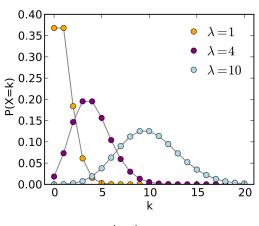
 $p^k$  - probability that connects to k nodes (has k-edges)  $(1-p)^{n-k-1}$  - probability that does not connect to any other node  $C_{n-1}^k$  - number of ways to select k nodes out of all to connect to,  $C_{n-1}^k = \frac{(n-1)!}{k!(n-k-1)!}$ 

• Binomial distribution, when  $\langle \mathbf{k} \rangle << N$  or  $n \to \infty$  and  $p \to 0$  at fixed  $\langle \mathbf{k} \rangle$ , is well approximated by Poisson distribution:

$$P(k) = \frac{\langle k \rangle^k e^{-\langle k \rangle}}{k!} = \frac{\lambda^k e^{-\lambda}}{k!}, \ \langle k \rangle = pn = \lambda$$

#### Poisson Distribution

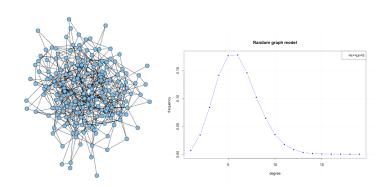




$$P(k_i = k) = \frac{\lambda^k e^{-\lambda}}{k!}, \ \lambda = \langle k \rangle = pn$$

## Random graph





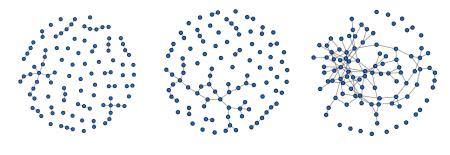
$$\langle \mathbf{k} \rangle = \mathbf{pn} = 5$$

## Random graph model



#### Consider $G_{n,p}$ as a function of p

- p = 0, empty graph  $\langle k \rangle = 0$
- p=1, complete (full) graph  $\langle k \rangle = n-1$
- $n_G$  -largest connected component,  $s = \frac{n_G}{n}$



#### Phase transition



Let u – fraction of nodes that do not belong to GCC. The probability that a node does not belong to GCC

$$u = \frac{n - n_G}{n} = P(k = 0) + P(k = 1) \cdot u + P(k = 2) \cdot u^2 + P(k = 3) \cdot u^3 \dots =$$

$$= \sum_{k=0}^{\infty} P(k)u^k = \sum_{k=0}^{\infty} \frac{\lambda^k e^{-\lambda}}{k!} u^k = e^{-\lambda} e^{\lambda u} = e^{\lambda(u-1)}$$

Let s -fraction of nodes belonging to GCC (size of GCC)

$$u = e^{\lambda(u-1)}$$

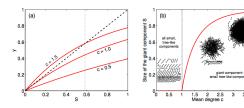
$$s = 1 - u, 1 - s = e^{-\lambda s}$$

$$\lambda=pn=\langle k \rangle$$
 when  $\lambda \to \infty, \; s \to 1$  when  $\lambda \to 0, \; s \to 0$ 

#### Phase transition



$$s = 1 - e^{-\lambda s}$$



non-zero solution exists when (at s = 0):

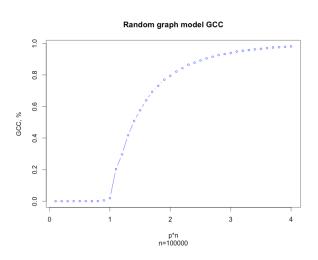
$$\lambda e^{-\lambda s} > 1$$

critical value:

$$\lambda_{c}=1$$
 
$$\lambda_{c}=\langle \mathbf{k}\rangle=p_{c}n=1,\ p_{c}=\frac{1}{n}$$

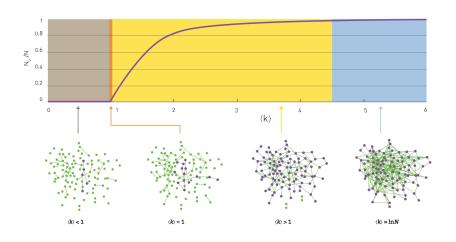
## **Numerical simulations**





## **Evolution of random network**





from A-L. Barabasi, 2016

#### Phase transition



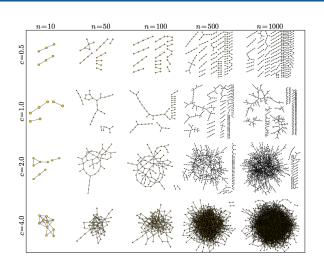
Graph G(n, p), for  $n \to \infty$ , critical value  $p_c = 1/n$ 

- Subcritical regime:  $p < p_c$ ,  $\langle k \rangle < 1$  there is no components with more than  $O(\ln n)$  nodes, largest component is a tree
- Critical point:  $p=p_{cr}\langle k\rangle=1$  the largest component has  $O(n^{2/3})$  nodes
- Supercritical regime:  $p>p_c$ ,  $\langle k\rangle>1$  gigantic component has all  $O((p-p_c)n)$  nodes
- Connected regime:  $p >> \ln n/n$ ,  $\langle k \rangle > \ln n$  gigantic component has all O(n) nodes

Critical value:  $\langle k \rangle = p_c n = 1$ - on average one neighbor for a node

## Numerical simulation





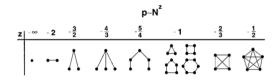
Clauset, 2014

## Threshold probabilities



#### Graph G(n, p)

Threshold probabilities when different subgraphs of k-nodes and l-edges appear in a random graph  $p_s \sim n^{-k/l}$ 



#### When $p > p_s$ :

- $p_s \sim n^{-k/(k-1)}$ , having a tree with k nodes
- $p_s \sim n^{-1}$ , having a cycle with k nodes
- $p_s \sim n^{-2/(k-1)}$ , complete subgraph with k nodes

Barabasi, 2002

# Clustering coefficient



 Clustering coefficient (probability that two neighbors link to each other):

$$C_i(k) = \frac{\text{\#of links between NN}}{\text{\#max number of links NN}} = \frac{pk(k-1)/2}{k(k-1)/2} = p$$

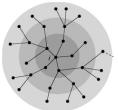
$$C = p = \frac{\langle k \rangle}{n}$$

• when  $n \to \infty$ ,  $C \to 0$ 

## Graph diameter



 G(n, p) is locally tree-like (GCC) (no loops; low clustering coefficient)



on average, the number of nodes d steps away from a node

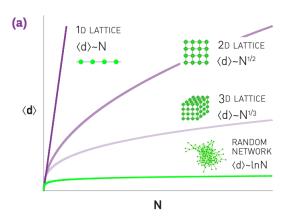
$$n = 1 + \langle \mathbf{k} \rangle + \langle \mathbf{k} \rangle^2 + ... \langle \mathbf{k} \rangle^D = \frac{\langle \mathbf{k} \rangle^{D+1} - 1}{\langle \mathbf{k} \rangle - 1} \approx \langle \mathbf{k} \rangle^D$$

• around  $p_c$ ,  $\langle k \rangle^D \sim n$ ,

$$D \sim rac{\ln n}{\ln \langle k 
angle}$$

## Graph diameter





## Random graph model



Node degree distribution function - Binomial/Poisson:

$$P(k) = \frac{\lambda^k e^{-\lambda}}{k!}, \ \lambda = pn = \langle k \rangle$$

Average path length:

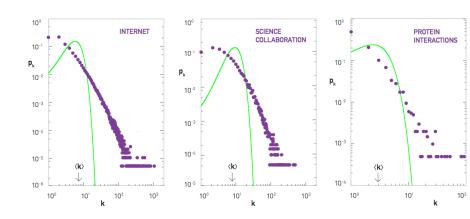
$$\langle L \rangle = \frac{\ln n}{\ln \langle k \rangle}$$

Clustering coefficient:

$$C = \frac{\langle k \rangle}{n}$$



#### Degree distribution in real networks



## Configuration model



- Random network with a predefined degree sequence:  $D = \{k_1, k_2, k_3...k_n\}$ , n nodes and  $m = 1/2 \sum_i k_i$  edges.
- Construct by randomly matching two stubs and connecting them by an edge.



- Can contain self loops and multiple edges
- Probability that two nodes i and j are connected

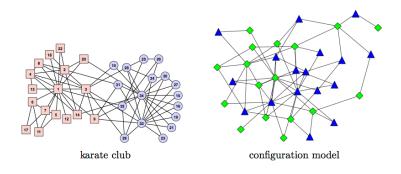
$$p_{ij} = \frac{k_i k_j}{2m - 1}$$

Will be a simple graph for special "graphical degree sequence"

## Configuration model



#### Can be used as a "null model" for comparative network analysis



Clauset, 2014

## References



- On random graphs I, P. Erdos and A. Renyi, Publicationes Mathematicae 6, 290–297 (1959).
- On the evolution of random graphs, P. Erdos and A. Renyi, Publication of the Mathematical Institute of the Hungarian Academy of Sciences, 17-61 (1960)