# UNIVERSITY OF EDINBURGH COLLEGE OF SCIENCE AND ENGINEERING SCHOOL OF INFORMATICS

### TYPES AND SEMANTICS FOR PROGRAMMING LANGUAGES

Saturday 1 st April 2017

00:00 to 00:00

### INSTRUCTIONS TO CANDIDATES

Answer QUESTION 1 and ONE other question.

Question 1 is COMPULSORY. If both QUESTION 2 and QUESTION 3 are answered, only QUESTION 2 will be marked.

All questions carry equal weight.

### CALCULATORS MAY NOT BE USED IN THIS EXAMINATION

Year 4 Courses

Convener: ITO-Will-Determine External Examiners: ITO-Will-Determine

THIS EXAMINATION WILL BE MARKED ANONYMOUSLY

# 1. THIS QUESTION IS COMPULSORY

Consider a type of trees defined as follows.

$$\begin{array}{c}
A \\
\hline
Tree A \\
\hline
Tree A
\end{array}$$
\_branch\_
\_
Tree A
\_
Tree A

Given a predicate *P* over *A*, we define predicates AllT and AnyT which hold when *P* holds for *every* leaf in the tree and when *P* holds for *some* leaf in the tree, respectively.

$$\begin{array}{c}
 & \text{AllT } P xt \\
 & \text{AllT } P \text{ (leaf } x)
\end{array}$$

$$\begin{array}{c}
 & \text{branch} \\
\hline
 & \text{AllT } P (xt \text{ branch } yt)
\end{array}$$

(a) Formalise the definitions above.

[12 marks]

(b) Prove AllT  $(\neg \ \circ P)$  xt implies  $\neg$  (AnyT P xt), for all trees xt.

[13 marks]

## 2. ANSWER EITHER THIS QUESTION OR QUESTION 3

You will be provided with a definition of intrinsically-typed lambda calculus in Agda. Consider constructs satisfying the following rules, written in extrinsically-typed style.

A computation of type  $Comp\ A$  returns either an error with a message msg which is a string, or an ok value of a term M of type A. Consider constructs satisfying the following rules:

Typing:

$$\underbrace{ \Gamma \vdash M \ \$ \ A }_{\text{Comp } A} \qquad \text{ok} \underbrace{ \begin{array}{c} \Gamma \vdash M \ \$ \ A \\ \hline \Gamma \vdash \text{ok} \ M \ \$ \ \text{Comp} \ A \end{array} }$$

$$\begin{array}{c} \Gamma \vdash M \ \text{\% Comp } A \\ \Gamma, \ x \ \text{\%} \ A \vdash N \ \text{\% Comp } B \\ \\ \text{letc} \hline \Gamma \vdash \text{letc} \ x \leftarrow M \ \text{in} \ N \ \text{\% Comp} \ B \end{array}$$

Values:

Reduction:

$$\xi \operatorname{-ok} \frac{M \longrightarrow M'}{\operatorname{ok} M \longrightarrow \operatorname{ok} M'} \qquad \xi \operatorname{-letc} \frac{M \longrightarrow M'}{\operatorname{letc} x \leftarrow M \operatorname{in} N \longrightarrow \operatorname{letc} x \leftarrow M' \operatorname{in} N}$$

$$\beta \operatorname{-error} \frac{\operatorname{letc} x \leftarrow (\operatorname{error} \operatorname{\mathit{msg}}) \operatorname{in} t \longrightarrow \operatorname{error} \operatorname{\mathit{msg}}}{\operatorname{letc} x \leftarrow (\operatorname{ok} V) \operatorname{in} N \longrightarrow N \ [x := V]}$$

- (a) Extend the given definition to formalise the evaluation and typing rules, including any other required definitions. [12 marks]
- (b) Prove progress. You will be provided with a proof of progress for the simply-typed lambda calculus that you may extend. [13 marks]

Please delimit any code you add as follows.

- -- begin
- -- end

# 3. ANSWER EITHER THIS QUESTION OR QUESTION 2

You will be provided with a definition of inference for extrinsically-typed lambda calculus in Agda. Consider constructs satisfying the following rules, written in extrinsically-typed style that support bidirectional inference.

Typing:

$$\begin{array}{c} \text{tt} \overline{ \Gamma \vdash \text{tt} \downarrow \top } \\ \\ \Gamma \vdash L \uparrow \top \\ \\ \Gamma \vdash M \downarrow A \\ \\ \hline \Gamma \vdash \text{caseT} \ L \ [\text{tt} \Rightarrow M \ ] \downarrow A \end{array}$$

(a) Extend the given definition to formalise the typing rules, and update the definition of equality on types.

[10 marks]

(b) Extend the code to support type inference for the new features.

[15 marks]

Please delimit any code you add as follows.

- -- begin
- -- end