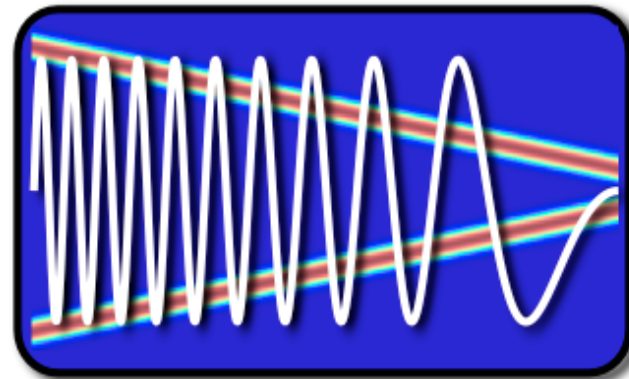


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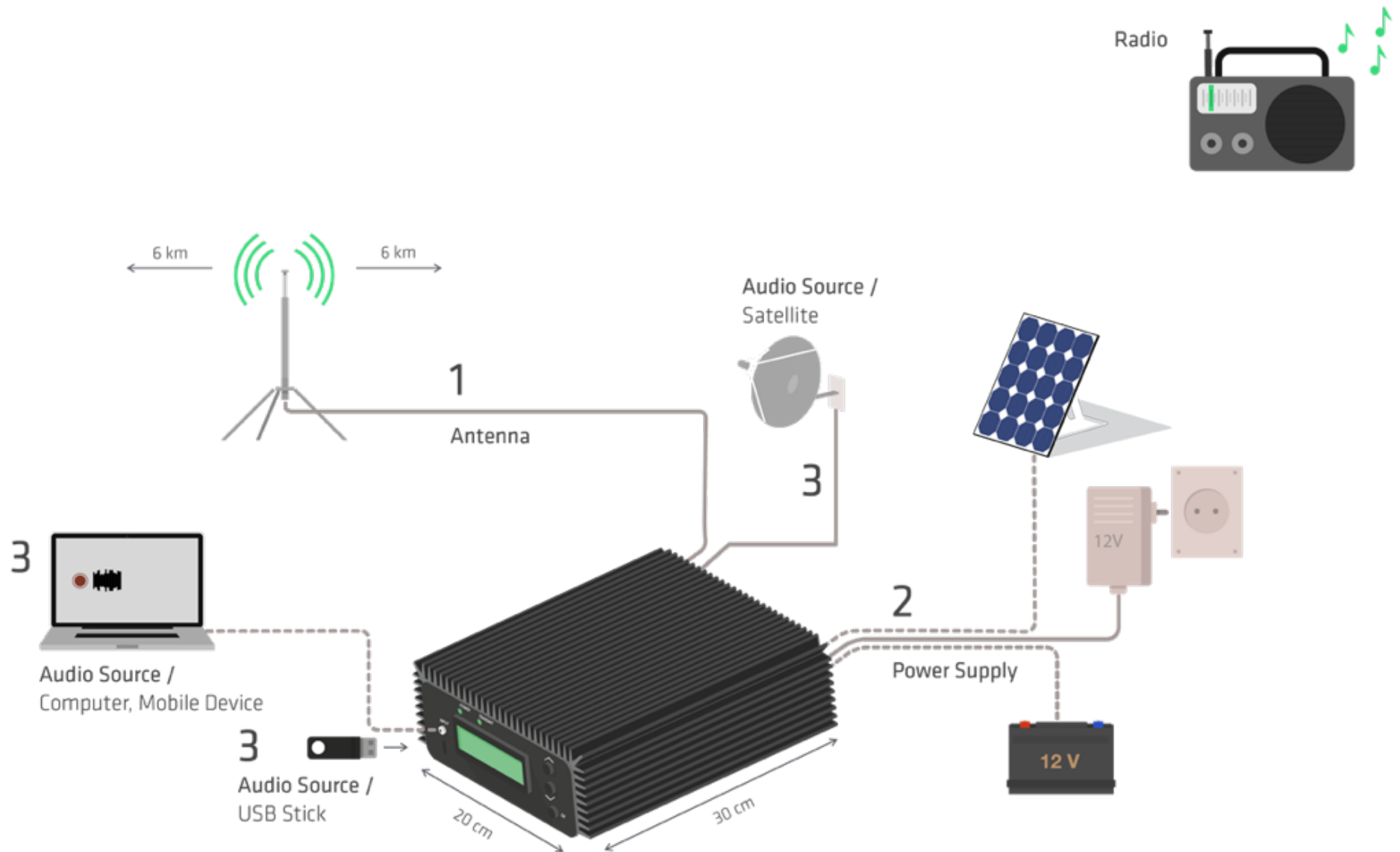


# Digital Signal Processing

## Lecture 5

# Interesting Project: Pocket FM

<http://www.pocket-fm.com>



## Info

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- Last time
  - Finished DTFT Ch. 2
  - 12min z-Transforms Ch. 3
- Today: DFT Ch. 8
- Reminders:
  - HW Due tonight
  - Lab Checkoff next week

## Motivation: Discrete Fourier Transform

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- Sampled Representation in time and frequency
  - Numerical Fourier Analysis requires discrete representation
  - But, sampling in one domain corresponds to periodicity in the other...
  - What about DFS (DFT)?
    - Periodic in “time” ✓
    - Periodic in “Frequency” ✓
  - What about non-periodic signals?
    - Still use DFS(T), but need special considerations

## Motivation: Discrete Fourier Transform

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- Efficient Implementations exist
  - Direct evaluation of DFT:  $O(N^2)$
  - Fast Fourier Transform (FFT):  $O(N \log N)$   
(ch. 9, next topic....)
  - Efficient libraries exist: FFTW
    - In Python:
      - > `X = np.fft.fft(x);`
      - > `x = np.fft.ifft(X);`
  - Convolution can be implemented efficiently using FFT
    - Direct convolution:  $O(N^2)$
    - FFT-based convolution:  $O(N \log N)$

## Discrete Fourier Series (DFS)

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- Definition:
  - Consider N-periodic signal:

$$\tilde{x}[n + N] = \tilde{x}[n] \quad \forall n$$

frequency-domain N-periodic representation:

$$\tilde{X}[k + N] = \tilde{X}[k] \quad \forall k$$

- “~” indicates periodic signal/spectrum

# Discrete Fourier Series (DFS)

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- Define:

$$W_N \triangleq e^{-j2\pi/N}$$

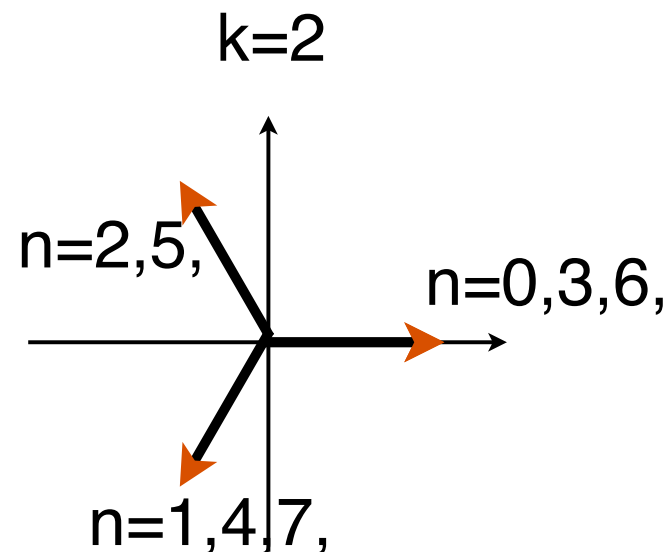
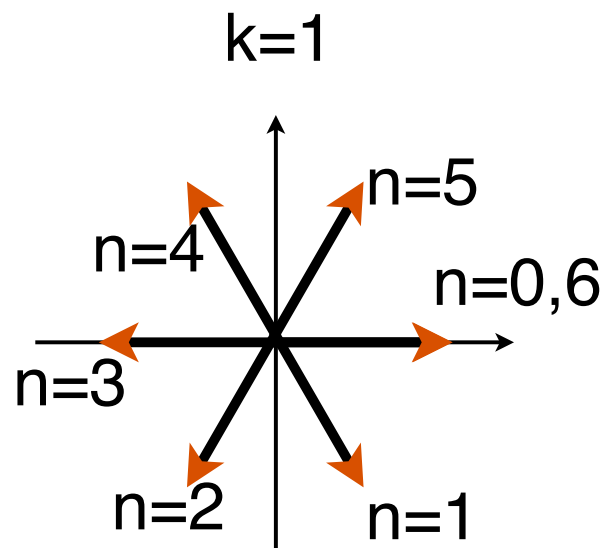
- DFS:

$$\begin{aligned}\tilde{x}[n] &= \frac{1}{N} \sum_{k=0}^{N-1} \tilde{X}[k] W_N^{-kn} \\ \tilde{X}[k] &= \sum_{n=0}^{N-1} \tilde{x}[n] W_N^{kn}\end{aligned}$$

Properties of  $W_N^{kn}$ ?

# Discrete Fourier Series (DFS)

- Properties of  $W_N$ :
  - $W_N^0 = W_N^N = W_N^{2N} = \dots = 1$
  - $W_N^{k+r} = W_N^k W_N^r$  or,  $W_N^{k+N} = W_N^k$
- Example:  $W_N^{kn}$  ( $N=6$ )





# Discrete Fourier Transform

---

- By Convention, work with **one** period:

$$x[n] \triangleq \begin{cases} \tilde{x}[n] & 0 \leq n \leq N - 1 \\ 0 & \text{otherwise} \end{cases}$$

$$X[k] \triangleq \begin{cases} \tilde{X}[k] & 0 \leq k \leq N - 1 \\ 0 & \text{otherwise} \end{cases}$$

Same same..... but different!

# Discrete Fourier Transform

---

- The DFT

$$x[n] = \frac{1}{N} \sum_{k=0}^{N-1} X[k] W_N^{-kn} \quad \text{Inverse DFT, synthesis}$$

$$X[k] = \sum_{n=0}^{N-1} x[n] W_N^{kn} \quad \text{DFT, analysis}$$

- It is understood that,

$$x[n] = 0 \quad \text{outside } 0 \leq n \leq N - 1$$

$$X[k] = 0 \quad \text{outside } 0 \leq k \leq N - 1$$

# Discrete Fourier Transform

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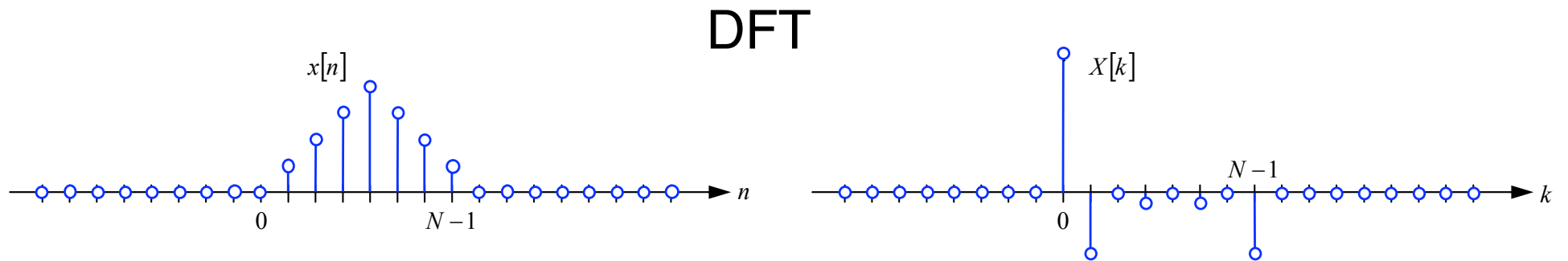
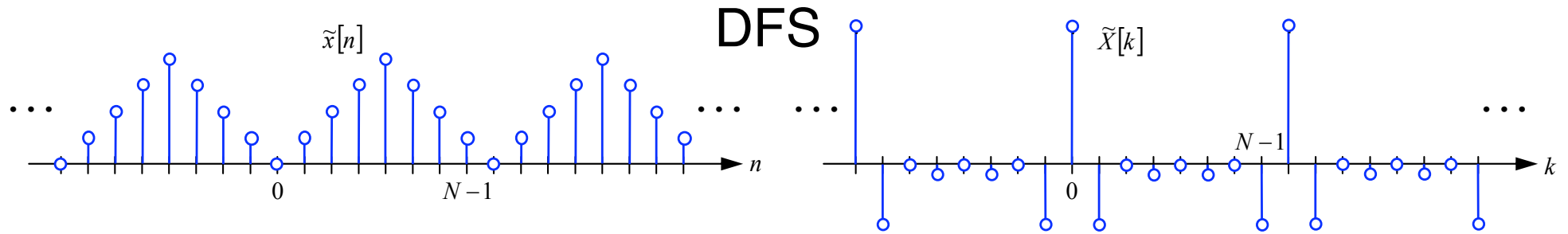
- Alternative formulation (not in book)  
Orthonormal DFT:

$$x[n] = \frac{1}{\sqrt{N}} \sum_{k=0}^{N-1} X[k] W_N^{-kn} \quad \text{Inverse DFT, synthesis}$$

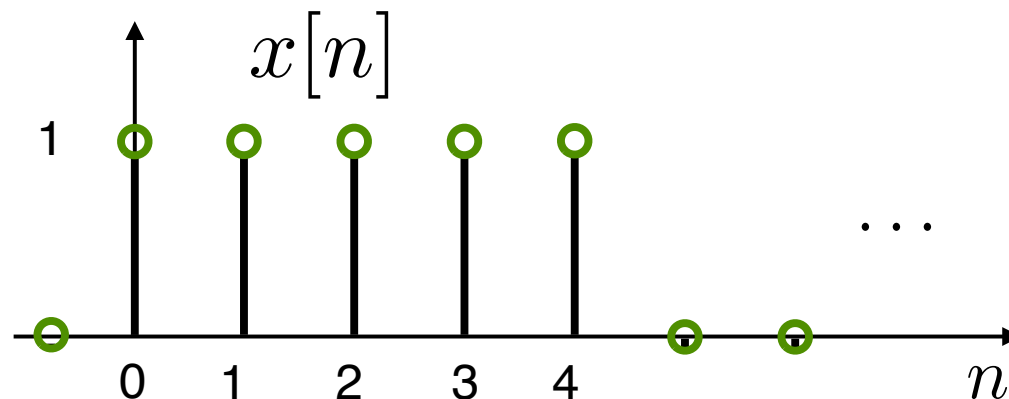
$$X[k] = \frac{1}{\sqrt{N}} \sum_{n=0}^{N-1} x[n] W_N^{kn} \quad \text{DFT, analysis}$$

Why use this or the other?

# Comparison between DFS/DFT



## Example



- Take  $N=5$

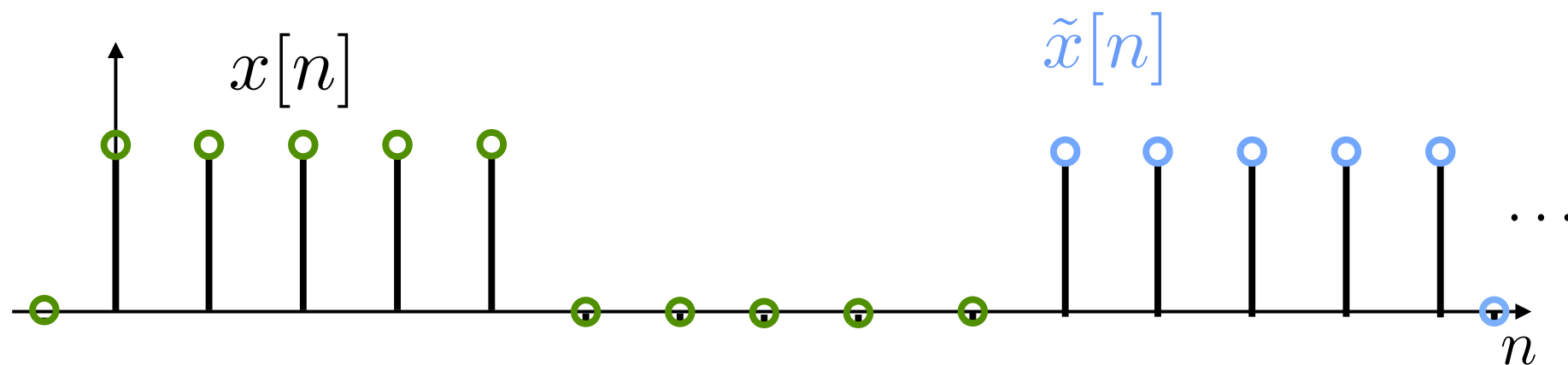
$$X[k] = \begin{cases} \sum_{n=0}^4 W_5^{nk} & k = 0, 1, 2, 3, 4 \\ 0 & \text{otherwise} \end{cases}$$
$$= 5\delta[k]$$

“5-point DFT”

## Example

- Q: What if we take  $N=10$ ?

A:  $X[k] = \tilde{X}[k]$  where  $\tilde{x}[n]$  is a period-10 seq.



$$X[k] = \begin{cases} \sum_{n=0}^4 W_{10}^{nk} & k = 0, 1, 2, \dots, 9 \\ 0 & \text{otherwise} \end{cases}$$

“10-point DFT”

## Example

---

- Show:

$$\begin{aligned} X[k] &= \sum_{n=0}^{\boxed{4}} W_{10}^{nk} \\ &= e^{-j \frac{4\pi}{10} k} \frac{\sin(\frac{\pi}{2} k)}{\sin(\frac{\pi}{10} k)} \end{aligned}$$

“10-point DFT”

## DFT vs DTFT

---

- For finite sequences of length  $N$ :
  - The  $N$ -point DFT of  $x[n]$  is:

$$X[k] = \sum_{n=0}^{N-1} x[n] W_N^{kn} = \sum_{n=0}^{N-1} x[n] e^{-j(2\pi/N)nk} \quad 0 \leq k \leq N-1$$

- The DTFT of  $x[n]$  is:

$$X(e^{j\omega}) = \sum_{n=0}^{N-1} x[n] e^{-j\omega n} \quad -\infty < \omega < \infty$$

What is similar?



## DFT vs DTFT

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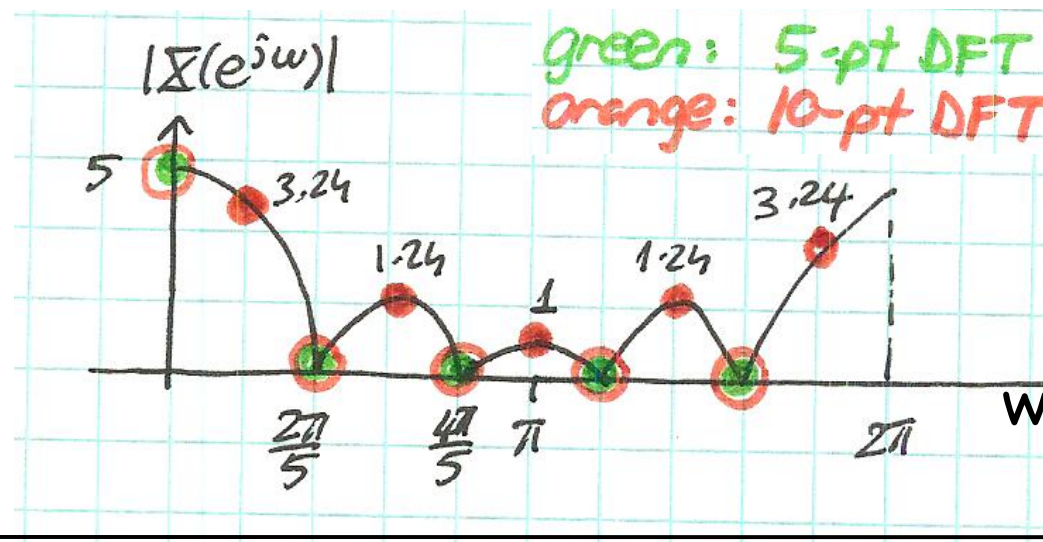
- The DFT are samples of the DTFT at  $N$  equally spaced frequencies

$$X[k] = X(e^{j\omega})|_{\omega=k\frac{2\pi}{N}} \quad 0 \leq k \leq N - 1$$

## DFT vs DTFT

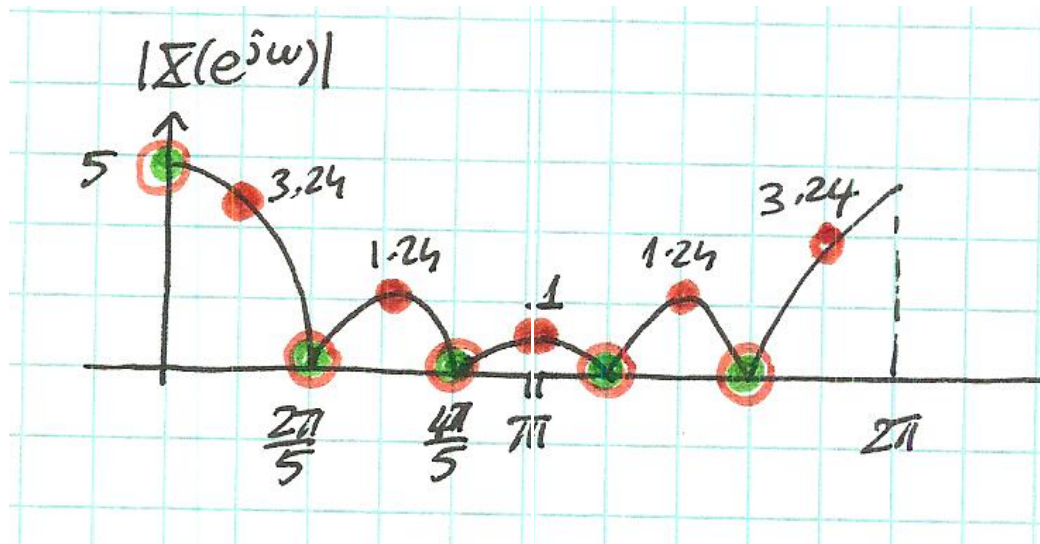
- Back to moving average example:

$$\begin{aligned} X(e^{j\omega}) &= \sum_{n=0}^4 e^{-j\omega n} \\ &= e^{-j2\omega} \frac{\sin(\frac{5}{2}\omega)}{\sin(\frac{\omega}{2})} \end{aligned}$$



# FFTSHIFT

- Note that  $k=0$  is  $\omega=0$  frequency
- Use `fftshift` to shift the spectrum so  $\omega=0$  in the middle.



## DFT and Inverse DFT

---

- Both computed similarly.....let's play:

$$\begin{aligned} N \cdot x^*[n] &= N \left( \frac{1}{N} \sum_{k=0}^{N-1} X[k] W_N^{-kn} \right)^* \\ &= \sum_{k=0}^{N-1} X^*[k] W_N^{kn} \\ &= \mathcal{DFT} \{ X^*[k] \} . \end{aligned}$$

- Also....

$$N \cdot x^*[n] = N \left( \mathcal{DFT}^{-1} \{ X[k] \} \right)^* .$$

## DFT and Inverse DFT

---

- So,

$$\mathcal{DFT} \{X^*[k]\} = N (\mathcal{DFT}^{-1} \{X[k]\})^*$$

or,

$$- \mathcal{DFT}^{-1} \{X[k]\} = \frac{1}{N} (\mathcal{DFT} \{X^*[k]\})^*$$

- Implement IDFT by:
  - Take complex conjugate
  - Take DFT
  - Multiply by 1/N
  - Take complex conjugate !

Why useful?

# DFT as Matrix Operator

**DFT:**

$$\begin{pmatrix} X[0] \\ \vdots \\ X[k] \\ \vdots \\ X[N-1] \end{pmatrix} = \begin{pmatrix} W_N^{00} & \cdots & W_N^{0n} & \cdots & W_N^{0(N-1)} \\ \vdots & \ddots & \vdots & \ddots & \vdots \\ W_N^{k0} & \cdots & W_N^{kn} & \cdots & W_N^{k(N-1)} \\ \vdots & \ddots & \vdots & \ddots & \vdots \\ W_N^{(N-1)0} & \cdots & W_N^{(N-1)n} & \cdots & W_N^{(N-1)(N-1)} \end{pmatrix} \begin{pmatrix} x[0] \\ \vdots \\ x[n] \\ \vdots \\ x[N-1] \end{pmatrix}$$

**IDFT:**

$$\begin{pmatrix} x[0] \\ \vdots \\ x[n] \\ \vdots \\ x[N-1] \end{pmatrix} = \frac{1}{N} \begin{pmatrix} W_N^{-00} & \cdots & W_N^{-0k} & \cdots & W_N^{-0(N-1)} \\ \vdots & \ddots & \vdots & \ddots & \vdots \\ W_N^{-n0} & \cdots & W_N^{-nk} & \cdots & W_N^{-n(N-1)} \\ \vdots & \ddots & \vdots & \ddots & \vdots \\ W_N^{-(N-1)0} & \cdots & W_N^{-(N-1)k} & \cdots & W_N^{-(N-1)(N-1)} \end{pmatrix} \begin{pmatrix} X[0] \\ \vdots \\ X[k] \\ \vdots \\ X[N-1] \end{pmatrix}$$

straightforward implementation requires  $N^2$  complex multiplies :-)

## DFT as Matrix Operator

---

- Can write compactly as:

$$\mathbf{X} = \mathbf{W}_N \mathbf{x}$$

$$\mathbf{x} = \frac{1}{N} \mathbf{W}_N^* \mathbf{X}$$

- So,

$$\mathbf{x} = \frac{1}{N} \mathbf{W}_N^* \mathbf{X} = \frac{1}{N} \mathbf{W}_N^* \mathbf{W}_N \mathbf{x} = \frac{1}{N} (N\mathcal{I}) \mathbf{x} = \mathbf{x}$$

WHY?

as expected.

## Properties of DFT

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- Inherited from DFS (EE120/20) so no need to be proved

- Linearity

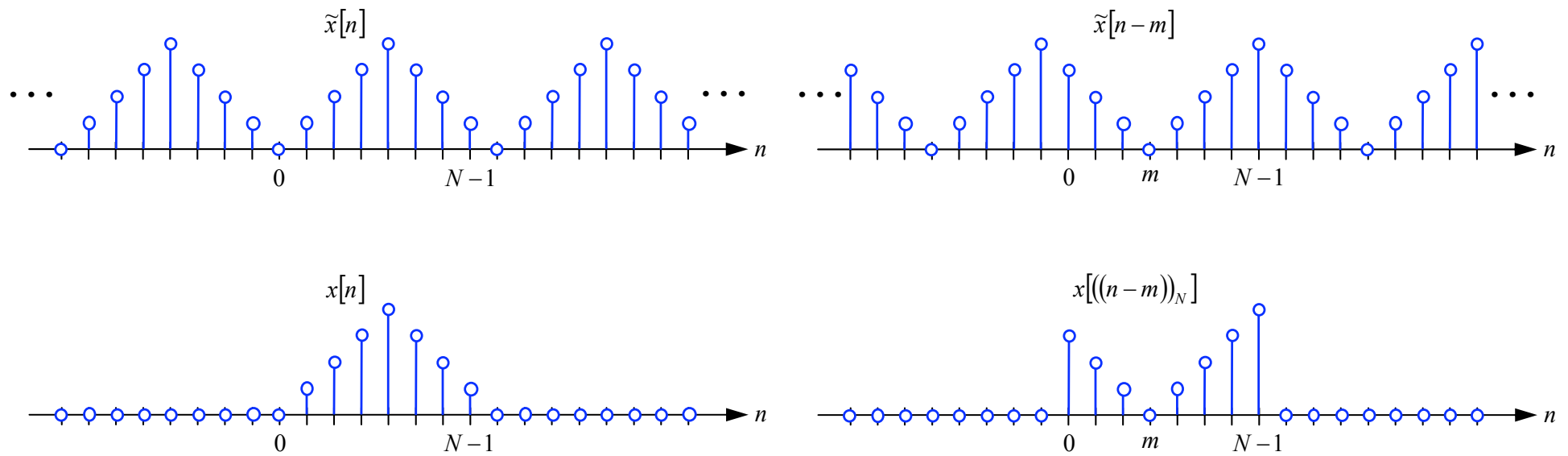
$$\alpha_1 x_1[n] + \alpha_2 x_2[n] \leftrightarrow \alpha_1 X_1[k] + \alpha_2 X_2[k]$$

- Circular Time Shift

$$x[((n - m))_N] \leftrightarrow X[k] e^{-j(2\pi/N)km} = X[k] W_N^{km}$$



# Circular shift



## Properties of DFT

---

- Circular frequency shift

$$x[n]e^{j(2\pi/N)nl} = x[n]W_N^{-nl} \leftrightarrow X[((k-l))_N]$$

- Complex Conjugation

$$x^*[n] \leftrightarrow X^*[((-k))_N]$$

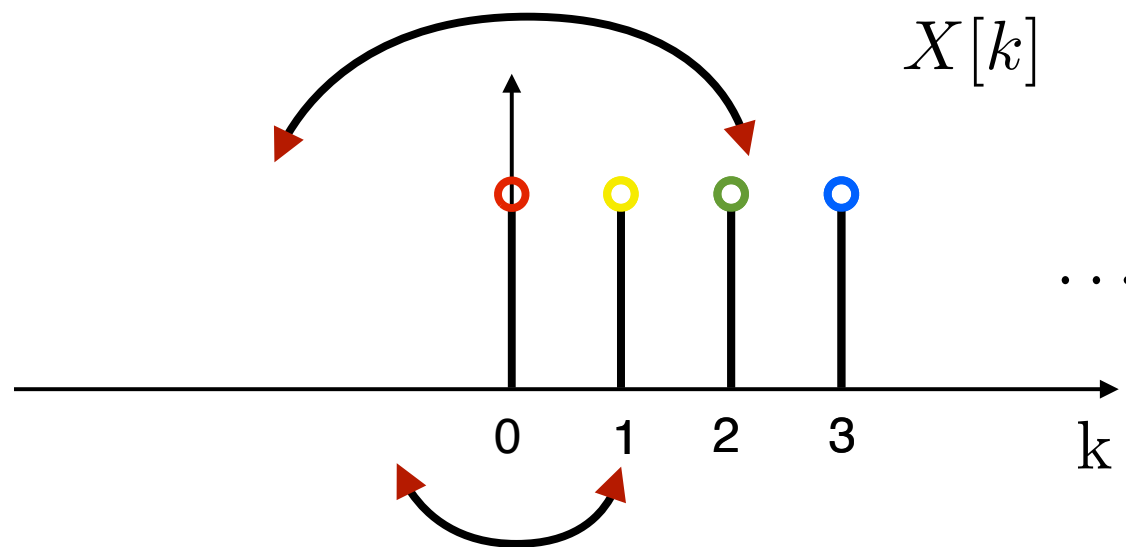
- Conjugate Symmetry for Real Signals

$$x[n] = x^*[n] \leftrightarrow X[k] = X^*[((-k))_N]$$

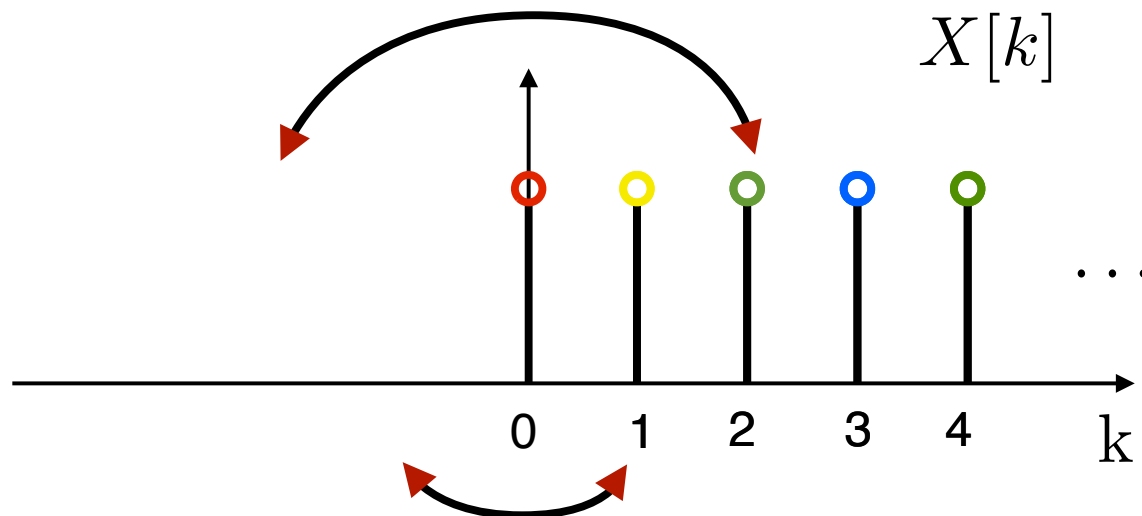
Show....

## Examples

- 4-point DFT  
–Symmetry

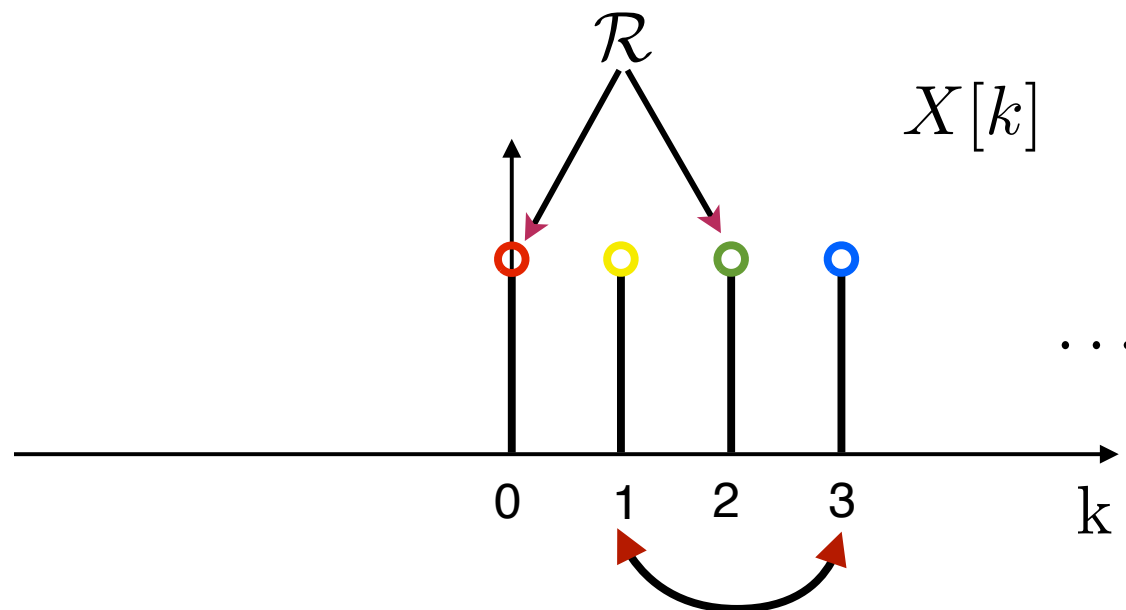


- 5-point DFT  
–Symmetry

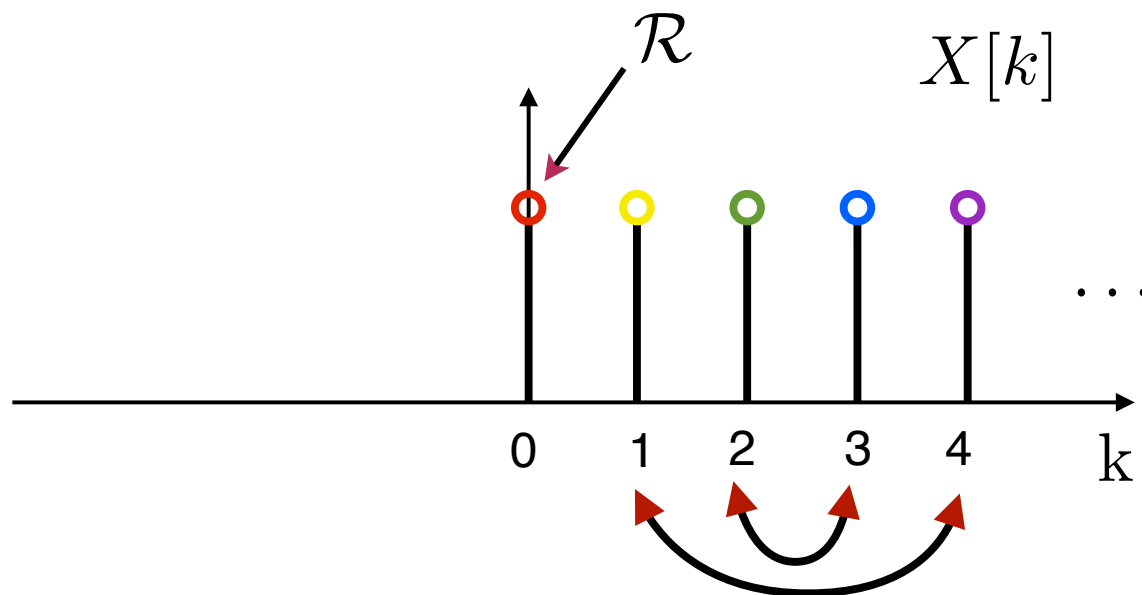


## Examples

- 4-point DFT
  - Symmetry



- 5-point DFT
  - Symmetry



## Properties of DFT

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- Parseval's Identity

$$\sum_{n=0}^{N-1} |x[n]|^2 = \frac{1}{N} \sum_{k=0}^{N-1} |X[k]|^2$$

- Proof (in matrix notation)

$$\mathbf{x}^* \mathbf{x} = \left( \frac{1}{N} \mathbf{W}_N^* \mathbf{X} \right)^* \left( \frac{1}{N} \mathbf{W}_N^* \mathbf{X} \right) = \frac{1}{N^2} \mathbf{X}^* \underbrace{\mathbf{W}_N \mathbf{W}_N^*}_{N \cdot \mathbf{I}} \mathbf{X} = \frac{1}{N} \mathbf{X}^* \mathbf{X}$$

## Circular Convolution Sum

---

- Circular Convolution:

$$x_1[n] \circledN x_2[n] \triangleq \sum_{m=0}^{N-1} x_1[m] x_2[((n - m))_N]$$

for two signals of length N

- Note: Circular convolution is commutative

$$x_2[n] \circledN x_1[n] = x_1[n] \circledN x_2[n]$$

## Properties of DFT

---

- **Circular Convolution:** Let  $x_1[n]$ ,  $x_2[n]$  be length  $N$

$$x_1[n] \circledast x_2[n] \leftrightarrow X_1[k] \cdot X_2[k]$$

Very useful!!! ( for linear convolutions with DFT)

- **Multiplication:** Let  $x_1[n]$ ,  $x_2[n]$  be length  $N$

$$x_1[n] \cdot x_2[n] \leftrightarrow \frac{1}{N} X_1[k] \circledast X_2[k]$$

# Linear Convolution

---

- Next....
  - Using DFT, circular convolution is easy
  - But, **linear** convolution is useful, not circular
  - So, show how to perform linear convolution with circular convolution
  - Used DFT to do linear convolution