

Define: $g^*(n)$ as minimum cost from root to $n \forall n \in \text{Nodes}$

Define: $g(n)$ as an easily computable approximation to g^*

Define: $h^*(n)$ as minimum cost from n to a goal $\forall n \in \text{Nodes}$

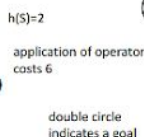
Define: $h(n)$ as an easily computable approximation to h^*

h is called a "heuristic function"

Define: $f(n) = g(n) + h(n)$

$QF(a, b) : \text{Sort}(\text{Append}(a, b), [])$

$[] =$ $h \rightarrow$ greedy (aka best first)
 $g \rightarrow$ uniform cost
 $f \rightarrow A/A^*$



Suppose $QF(\text{old}, \text{new}) : \text{Append}(\text{new}, \text{old})$; Depth First
 Suppose $QF(\text{old}, \text{new}) : \text{Append}(\text{old}, \text{new})$; Breadth First

Admissibility

A search algorithm (together with its heuristic function if needed) is *admissible* iff for all search problems:

1. If there exists a goal, the search will not fail.
2. If there are multiple goals the search will find an optimal one (best, least expensive to execute).

We have:

- A_1^* with some heuristic fcn h_1
- A_2^* with another heuristic fcn h_2
- A_1^* and A_2^* are admissible

Then we say

- A_1^* is *more informed* than A_2^*
- iff for all non-goal nodes n
- $h_1(n) > h_2(n)$

"More Informed" implies "guaranteed not to search more"

A set S is convex iff

$$\forall x, y \in S, \forall t \in [0, 1] \Rightarrow (1-t)x + ty \in S$$

Object constant	Individuals	\neg	negation	"not"
Variable	Properties	\wedge	conjunction	"and"
Function expression	Relations	\vee	disjunction	"or"
Predicate symbol		\Rightarrow	implication	"implies"
$\{x = y\}$	YES	\Leftrightarrow	equivalence	"if and only if"
$\{x = y, z = F(y)\}$	YES			
$\{x = y, z = F(y), x = A\}$	NO			
$\{x = y, z = F(y), y = A\}$	YES			
$\{x = y, y = F(z), z = G(x)\}$	NO			

There is some amount that I like CS440 and I like all other classes less

$\exists z \exists w [\text{Class}(w) \wedge \text{Name}(w, \text{"CS440"}) \wedge \text{Likes}(\text{Me}, w, z) \wedge \forall x \forall y [\text{Class}(x) \wedge \text{Likes}(\text{Me}, x, y) \wedge \text{Different}(x, w) \Rightarrow \text{Greater}(z, y)]]$

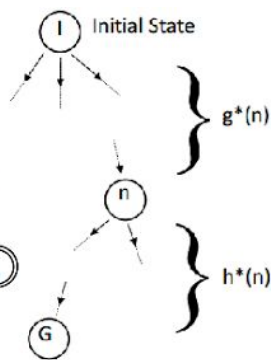
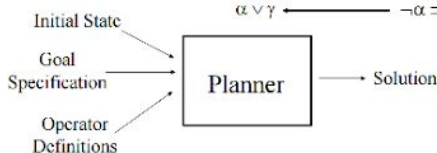
Likes(a,b,c) means "a likes b by amount c"
 Greater(a,b) means "a is larger than b"
 A unifier (also substitution, binding list*) is a set of pairings of variables with terms:

$$\{v_1 = e_1, v_2 = e_2, v_3 = e_3, \dots, v_n = e_n\}$$

such that

- each variable is paired at most once
- a variable's pairing term may not contain the variable directly or indirectly

$$\begin{array}{lcl} \alpha \vee \beta & \longrightarrow & \neg \alpha \Rightarrow \beta \\ \neg \beta \vee \gamma & \longrightarrow & \beta \Rightarrow \gamma \\ \alpha \vee \gamma & \longrightarrow & \neg \alpha \Rightarrow \gamma \end{array}$$



Admissibility of A*

Some authors use "A" if not met

- 1) $\forall n \forall n' \in \text{nodes}, \forall a \in \text{actions with } a(n) \rightarrow n'$

$$h(n) \leq \text{cost}(n, a, n') + h(n')$$

Triangle inequality, Monotonicity, or Consistency (for trees...)

Beam Search w/ beam width k

```
SEARCH(Problem P, Queuing Function QF):
  local: n /* current node */
        q /* nodes to explore */
  q ← singleton of Initial_State(P);
  Loop:
    if q = () return failure;
    n ← Pop(q);
    if n Solves P return n;
    q ← QF(q, Expand(n));
  end
  QF(a,b): Sort(Append(a,b), h);
           Delete all but the best k;
```

$$H_f = \begin{bmatrix} \frac{\partial^2 f}{\partial x_1 \partial x_1} & \frac{\partial^2 f}{\partial x_1 \partial x_2} \\ \frac{\partial^2 f}{\partial x_2 \partial x_1} & \frac{\partial^2 f}{\partial x_2 \partial x_2} \end{bmatrix} \text{ and } v^T H_f v > 0 \quad \forall v$$

A term denotes an individual in the universe of discourse
 variable
 object constant
 function expression

A function expression is an n-ary function symbol with n terms as arguments

An atom (also atomic sentence, atomic Well-Formed Formula) is an n-ary predicate symbol with n terms as arguments

A literal is an atom or a negated atom

Certain(Death) \wedge Certain(Taxes) \wedge
 $\forall x [(\text{Different}(x, \text{Death}) \wedge \text{Different}(x, \text{Taxes})) \Rightarrow \neg \text{Certain}(x)]$

Nothing is certain except death and taxes

p, q, r are WFFs

$p \wedge q = q \wedge p$
 $p \wedge (q \wedge r) = (p \wedge q) \wedge r$
 $p \Rightarrow q = \neg q \Rightarrow \neg p$
 $\neg(\neg p) = p$
 $p \Rightarrow q = \neg p \vee q$
 $p \vee p = p$
 $\neg(p \vee q) = \neg p \wedge \neg q$
 $\neg(p \wedge q) = \neg p \vee \neg q$
 $p \wedge (q \vee r) = (p \wedge q) \vee (p \wedge r)$
 $p \vee (q \wedge r) = (p \vee q) \wedge (p \vee r)$

MoveToBlock (x, z):

PC: Clr (x), Clr (z), On (x, y), Blk (x), Blk (z), Diff (x, z), Diff (y, z)

Effects: $\neg \text{On}(x, y), \neg \text{Clr}(z), \text{On}(x, z), \text{Clr}(y)$

MoveToTable (x, z):

PC: Clr (x), On (x, y), Blk (x), Tbl (z), Diff (y, z)

Effects: $\neg \text{On}(x, y), \text{On}(x, z), \text{Clr}(y)$

Generic Search Function

```
SEARCH(Problem P, Queuing Function QF):
  local: n /* current node */
        q /* nodes to explore */
  q ← singleton of Initial_State(P);
  Loop:
    if q = () return failure;
    n ← Pop(q);
    if n Solves P return n;
    q ← QF(q, Expand(n));
  end
```

Admissibility of A* (cont)

- 2) $\forall n \in \text{nodes}$

$$h(n) \leq h^*(n)$$

Informally: be optimistic

(or don't be pessimistic)

Why? Could you prove it?

Important General Principle:

Optimism Under Uncertainty

Admissibility does not depend on problem or tree!

Is "Uniform Cost" admissible?

Locally Optimizing Search: Hill Climbing

```
SEARCH(Problem P, Queuing Function QF):
  local: n /* current node */
        q /* nodes to explore */
  q ← singleton of Initial_State(P);
  Loop:
    if q = () return n;
    if First(q) worse_than n return n;
    n ← Pop(q);
    q ← QF(q, Expand(n));
  end
  QF(a,b): Sort(Append(a,b), U);
           Delete all but best;
```

Search Properties

- Tentative (don't throw away information)

Depth First Breadth First Uniform Cost
 Best First/Greedy A*

- Irrevocable (throw away information)

Hill Climbing Beam Optimizing Beam

- Exhaustive (will visit all nodes or find goal)

Breadth First Uniform Cost A*

- Admissible

Uniform Cost A* [if consistency & optimism]

Note quantifiers interact with negations!

$\neg \exists w P(w)$ is the same as $\forall x \neg P(w)$

Think DeMorgan...

$\neg (A \wedge B)$ is the same as $\neg A \vee \neg B$

Making explicit what is already true

Employing rules of inference

Modus Ponens A well-known rule of inference
 There are many others

$$\Theta \Rightarrow \Psi$$

$$\Theta$$

$$\Psi$$

Are there entailed WFFs that Modus Ponens cannot derive?

No, it is incomplete

One inference rule, resolution, is almost complete on its own

The MGU imposes the fewest constraints, specifying the weakest conditions for matching

MGU is unique assuming

order is not important

variable names are not important (alphabetic variants)

Applying the MGU to an expression yields a most general unification instance.

Variable substitutions are always interpreted with the unifier applied

FOPC – symbolic representations

- Quantification
- Inference
 - formal: possible worlds
 - algorithmic: inference rules like M.P.
- Unification

STRIPS operators

- Essentially FOPC but slightly reduced expressiveness
- Sometimes requires extra operators

Planning with STRIPS operators

- World states
- Abstract domain-independent procedures

Markov Decision Process (MDP)

- A finite set of states $S = \{s\}$
- A finite set of actions $A = \{a\}$
- Initial distribution over S (more general than R&N)
- Probabilistic transition model
- Probabilistic rewards (more general than R&N)

The world *may* be deterministic but we choose to model it as stochastic

P is a unary predicate; F is a unary function; x is a variable; Sam is a constant

$P(x)$	Atom, Literal, WFF
$P(Sam)$	Atom, Literal, WFF
$P(F(Sam))$	Atom, Literal, WFF
$P(x) \wedge P(Sam)$	WFF, conjunction of two atoms, conjunction of two literals
$\neg P(x) \wedge P(x)$	WFF, conjunction of a literal and an atom, conjunction of two literals
$\neg P(x)$	Literal, WFF, the negation of an atom
$\neg F(Sam)$	ill-formed, not a WFF at all
$P(\neg Sam)$	ill-formed, not a WFF at all
$P(x) \wedge F(Sam)$	ill-formed, not a WFF at all

We can learn the policy *directly*

Instead of T and U , learn Q

Q is the utility of performing an action in a state
 Q -learning is *model free* Reinforcement Learning
 Q -learning is *off policy*

- optimal greedy policy can be learned
- even if it is not followed during learning

General distinctions in learning

- Full or joint model vs. Conditional model
- Generative model vs. Discriminative model

Q function: $Q: A \times S \rightarrow \mathbb{R}$

$Q(a, s)$ - the expected utility of performing action a in state s

Q represents local and non-local information (much as U does)

But T (our model of the world's transition function) is not needed

The (converged) Q -table is independent of ϵ

Sometimes we may prefer a state-action choice that includes ϵ in some way

In ϵ -greedy, ϵ is reflected within SARSA's Q values

SARSA learns the best policy given our ϵ -greedy systematic departures from optimal

Convergence of both value iteration and policy iteration is based on *convexity* of the utility space

We will see gradient descent and stochastic gradient descent again with perceptrons & neural networks

But here in RL, both are neglecting information...

```
public void updatePolicy(double reward, int action, int oldState, int newState){
    qValue[oldState][action] = (1-learningRate)*qValue[oldState][action]+
    learningRate*(reward-discountFactor*qValue[newState][policy[newState]]);
    double max = 0, current = 0;
    for(int i=0; i<numOfActions; i++){
        current = qValue[oldState][i];
        if(current >= max){
            max = current;
            policy[oldState] = i;
        }
    }
}
```

Planning

- First-order representations
- Build in world dynamics
- Plan = sequence of actions
- Predict precise intermediate world states
- Sensing (and execution) play a small part
- Strong single agent assumption

Reinforcement learning

- Propositional (zeroth-order) representations
- Learn world dynamics
- Plan = action policy
- Tolerate arbitrarily bad behavior in the world
- Sensing during execution is crucial
- Weaker single agent assumption

Transition function

- $T: S \times A \times S \rightarrow [0,1]$
- Is T a probability distribution?
- $T(s, a, \cdot)$ denotes a probability distribution over next states
- $P(\cdot | s, a)$ as conditional probability (in R&N)

Reward function

- $Rw: S \times \mathcal{A} \rightarrow [0,1]$
- Each $Rw(s, \cdot)$ denotes a probability distribution over rewards

What do we care about with rewards?

$$R: S \rightarrow \mathbb{R} \quad U^\pi(s) \leftarrow U^\pi(s) + \alpha(R(s) + \gamma \cdot U^\pi(s') - U^\pi(s))$$

R maps states to expected rewards

Utility of a state s given a policy π with discount γ

$$U^\pi(s) = E \left[\sum_{t=0}^{\infty} \gamma^t R(s_t) \right]$$

U^π has a simple form: $U^\pi: S \rightarrow \mathbb{R}$
 (state utility is policy dependent)

Let π^* be the optimal policy

Given U^{π^*} and T we can write/implement the optimal policy concisely & efficiently

$T(i, j, k)$ can be estimated as the ratio:
 # times action j takes us from state i to state k
 divided by # times action j is tried in state i

$$U^\pi(s) \leftarrow (1-\alpha)U^\pi(s) + \alpha(R(s) + \gamma \cdot U^\pi(s'))$$

ϵ -Greedy Exploration

- Instead of greedy: $\arg \max_a \sum_s T(s, a, s') \cdot U^\pi(s')$

- Be greedy with probability $(1-\epsilon)$

- But with probability ϵ choose a random action

$$Q(a, s) \leftarrow Q(a, s) + \alpha \cdot (R(s) + \gamma \max_{a'} Q(a', s') - Q(a, s))$$

So:

- Choose an ϵ $0 < \epsilon \ll 1$

- At every decision get a random number $x: [0,1]$

- If $x < \epsilon$, choose randomly else be greedy with current T and U

Minimal regret grows as $\text{Sqrt}(N)$

- In fact

$$0.264 \cdot \text{Sqrt}(N) < \text{Regret} < 0.376 \cdot \text{Sqrt}(N)$$

Policy for TD Value Iteration?

– T, U

– $T: \mathcal{S} \times \mathcal{A} \times \mathcal{S}$

– $U: \mathcal{S}$

– $|\mathcal{S}|^2 \cdot |\mathcal{A}| + |\mathcal{S}|$ real numbers: $O(|\mathcal{S}|^2 \cdot |\mathcal{A}|)$

Policy for Q ?

– Q

– $Q: \mathcal{S} \times \mathcal{A}$

– $|\mathcal{S}| \cdot |\mathcal{A}|$ real numbers: $O(|\mathcal{S}| \cdot |\mathcal{A}|)$

Reinforcement Learning vs. Classical Planning

More robust (this is a big one...)

- Fewer & weaker *a priori* assumptions (esp. actions)
- Empirical model
- Fit (via parameter adjustment) to the observed world

A *LOT* of training is required

- Must see many and varied state transitions
 - Compared to no training for classical planning
- Scaling difficulties
- Propositional expressiveness (vs. first-order for classical planning)
 - Choosing / defining state distinctions can be challenging
 - Space & Time complexity is polynomial (but in what? The generally unknowable u -return mixing time & others...)
 - Generalizing to other similar problems can be difficult
 - Some domains require the very same problem solved repeatedly; others do not
 - Recognizing convergence? sufficiency?

$$\pi(s_{14}) = a_3$$

$$T(s_{14}, a_3) = s_{81}$$

We cannot anticipate the world's true transition function

Thus, our model of it must in part be learned

We model uncertainty as a distribution over next states

Our Transition Function

– $T: S \times A \times S \rightarrow [0,1]$

– where each $T(S, A, \cdot)$ is a distribution

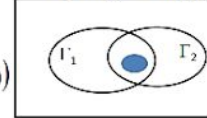
This approximates (in some sense [we hope]) the world's own transition function

$$E[X] = \sum_i x_i \cdot \Pr(X=x_i) \quad \begin{matrix} \Gamma_1: \exists x \text{ Bird}(x) \\ \Gamma_2: \exists x \text{ Flies}(x) \end{matrix}$$

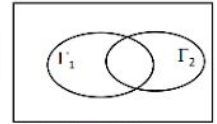
$$\pi^*(s) = \arg \max_a \sum_{s'} T(s, a, s') \cdot U^{\pi^*}(s')$$

$$U(s)_{\text{new}} \leftarrow U(s)_{\text{old}} + \alpha \cdot \text{error}$$

$\exists x [\text{Bird}(x) \wedge \text{Flies}(x)]$

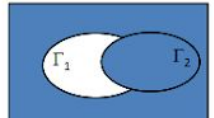


Possible Worlds



Possible Worlds

$\exists x [\text{Bird}(x) \Rightarrow \text{Flies}(x)]$



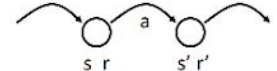
Possible Worlds

False - holds in no possible worlds (i.e., self contradictory)

True - holds in all possible worlds (i.e., tautology)

Satisfiable - holds in some possible worlds but not others

Entailed - holds in all remaining possible worlds after intersecting the axioms' possible worlds



$$Q(a, s) = R(s) + \gamma \sum_{s'} T(s, a, s') \max_{a'} Q(a', s')$$

Assume policy π chooses action a in s

Relate $U^\pi(s)$ and $U^\pi(s')$

$$U^\pi(s) = \gamma U^\pi(s') + R(s)$$

$$U^\pi(s) - \gamma U^\pi(s') = R(s)$$

If not equal then there is an error

$$\text{So error} = R(s) + \gamma U^\pi(s') - U^\pi(s)$$

RL, including Q -learning, employ a number of user-specified parameters

As in much of machine learning, understanding their influence and knowing how much to experiment is a key to success

γ - the discount rate

- domain parameter reflecting the relative importance of nearer rewards over more distant ones for the user
- it should not be used to influence the Q -learner

α - the learning rate

- reflects the relative confidence in the old and new information
- typically does not preclude convergence but can have a significant effect on the speed of convergence

ϵ - exploration probability

- employed to insure an acceptable amount of random exploration
- too low: slow convergence into a small good region around π^*
- too high: quickly converges into a large poor region around π^*

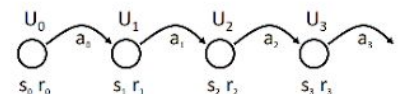
Q , an off-policy learner:

$$Q(a, s) \leftarrow Q(a, s) + \alpha \cdot (r_s + \gamma \max_{a'} Q(a', s') - Q(a, s))$$

SARSA, an on-policy learner:

$$Q(a, s) \leftarrow Q(a, s) + \alpha \cdot (r_s + \gamma \cdot Q(a', s') - Q(a, s))$$

Adaptive Dynamic Programming



Imagine successive value iteration...on s_0 ... on s_1 ... on s_2 :

Perform a_2 , update U_2 with r_2 and U_3

U_2 is now updated to a better value

We used the old U_2 to update U_1 , shouldn't it be changed as well?

What about U_0 ?

Fully appreciate each r