Adaptive Hashing for Model Counting



Jonathan Kuck, Tri Dao, Shengjia Zhao, Burak Bartan, Ashish Sabharwal, and Stefano Ermon

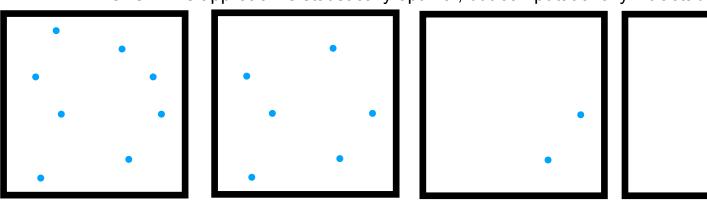
kuck@stanford.edu, trid@stanford.edu, sjzhao@stanford.edu, bbartan@stanford.edu, ashishs@allenai.org, ermon@stanford.edu



TLDR: Hashing-based counting algorithms^{1,2} currently operate identically for any shape of input instance. We adapt hash functions to specific instances using previously found solutions and use a novel aggregation strategy across independent repetitions to improve statistical efficiency. Our methods improve lower bounds by a median factor of 213. over 1198 input instances.

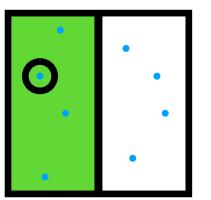
Background: Approximate Counting

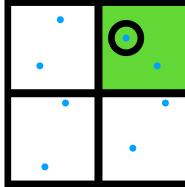
The number of elements in a set can be estimated by randomly removing each element with probability 1/2. Repeating this operation k times, until only 1 element remains, gives 2k as an estimate of the set's size. This approach is statistically optimal, but computationally intractable.

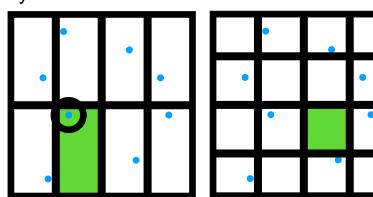


Background: Fixed Hashing

Hash functions provide a computationally efficient alternative. All elements in the set are hashed and a randomly selected bin is checked for elements. Repeating this process with increasingly more bins, until the random bin is empty, gives an estimate of the set size. Computational efficiency comes at the cost of low statistical efficiency as the number of elements may vary across bins

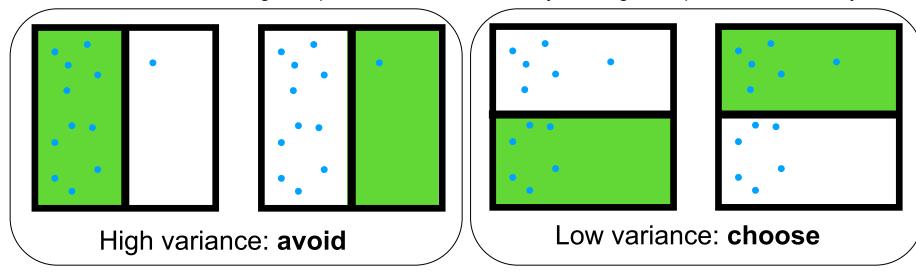






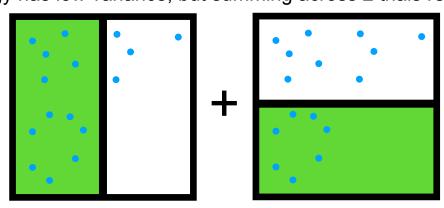
Adaptive Hashing

We choose low variance, cheap binning strategies for specific instances based on previously found solutions, resulting in improved statistical efficiency with high computational efficiency.



Aggregation Across Trials

We sample independent binning strategies and aggregate the solutions. This improves statistical efficiency, but requires additional parallel computation. Note that in this example neither binning strategy has low variance, but summing across 2 trials reduces the variance.



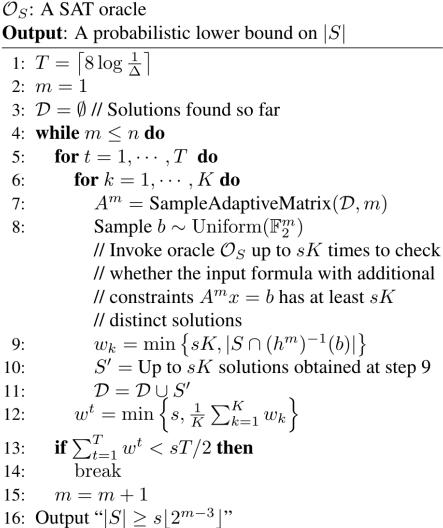
Algorithm

Algorithm 2 New, Adaptive Lower Bound **Inputs**: s: Solution cutoff

 Δ : Failure probability

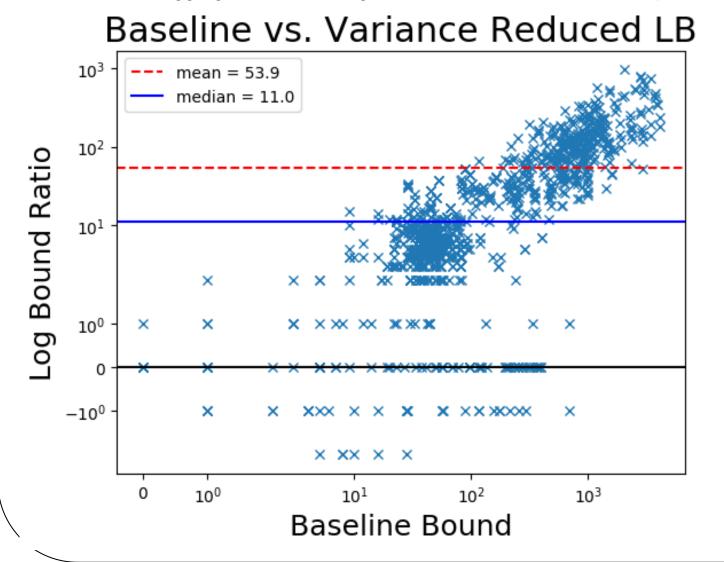
K: Number of repetitions per trial

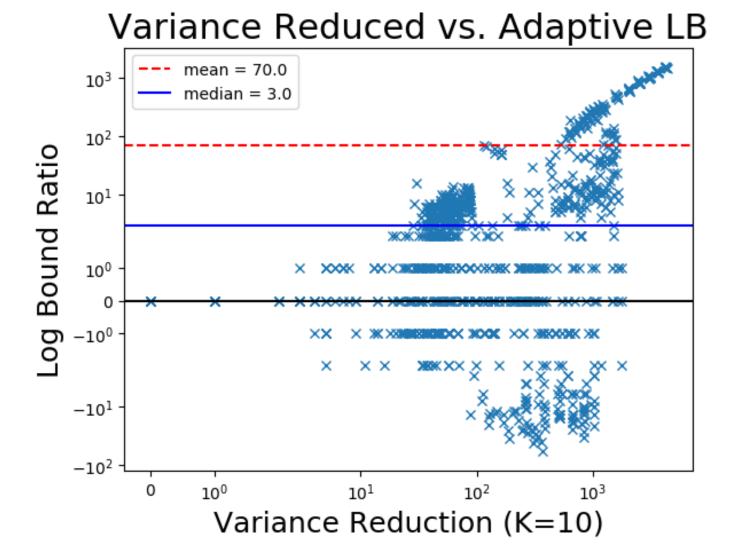
Output: A probabilistic lower bound on |S|

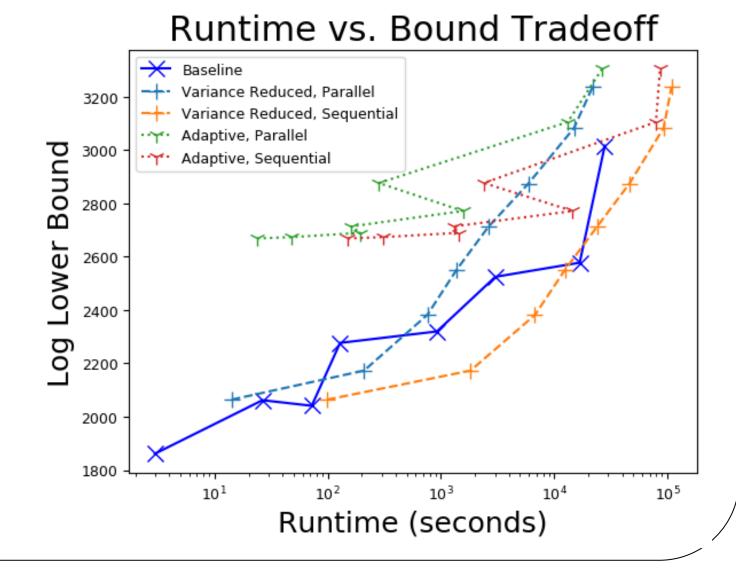


Empirical Study

Our baseline family of fixed hashed functions was inspired by LDPC codes, the current state of the art in ultra-efficiency. Left: aggregation improves the lower bound. Middle: adaptation with aggregation improves the lower bound compared with aggregation alone. Right: tradeoff in statistical and computational efficiency for our method and the baseline on a single challenging instance. We can improve both the lower bound and runtime by orders of magnitude.







- 1. Achlioptas, Dimitris, Zayd Hammoudeh, and Panos Theodoropoulos. "Fast and Flexible Probabilistic Model Counting." International Conference on Theory and Applications of Satisfiability Testing 2018.
- 2. Chakraborty, Supratik, Kuldeep S. Meel, and Moshe Y. Vardi. "Algorithmic Improvements in Approximate Counting for Probabilistic Inference: From Linear to Logarithmic SAT Calls." IJCAI 2016.