

Parallel Reduction of Stoughton's paper:

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data  $\Rightarrow$  :  $\Lambda \rightarrow \Lambda \rightarrow \text{Set}$  where
 $\Rightarrow \mathbf{v}$  :  $(x : \mathbf{V}) \rightarrow \mathbf{v} x \Rightarrow \mathbf{v} x$ 
 $\Rightarrow \lambda$  :  $(x : \mathbf{V})\{M M' : \Lambda\} \rightarrow M \Rightarrow M' \rightarrow \lambda x M \Rightarrow \lambda x M'$ 
 $\Rightarrow \cdot$  :  $\{M M' N N' : \Lambda\} \rightarrow M \Rightarrow M' \rightarrow N \Rightarrow N' \rightarrow M \cdot N \Rightarrow M' \cdot N'$ 
 $\Rightarrow \beta$  :  $(x : \mathbf{V})\{M M' N N' : \Lambda\} \rightarrow M \Rightarrow M' \rightarrow N \Rightarrow N' \rightarrow \lambda x M \cdot N \Rightarrow M' \bullet \mathbf{v} \prec + (x, N')$ 
 $\Rightarrow \alpha$  :  $\{M N N' : \Lambda\} \rightarrow M \Rightarrow N \rightarrow N \sim \alpha N' \rightarrow M \Rightarrow N'$ 

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We already have the following result in the paper:

The alpha rule of this parallel relation definition brings high non determinism in the construction evidence. In fact, infinite of them exists for any two related lambda terms. So we define a predicate that holds only if a parallel reduction has no alpha rules in it.

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no $\alpha$  :  $\{M N : \Lambda\} \rightarrow M \Rightarrow N \rightarrow \text{Set}$ 
no $\alpha$  ( $\Rightarrow \mathbf{v}$  _) =  $\top$ 
no $\alpha$  ( $\Rightarrow \lambda$  _  $M \Rightarrow N$ ) = no $\alpha$   $M \Rightarrow N$ 
no $\alpha$  ( $\Rightarrow \cdot$   $M \Rightarrow N$   $M' \Rightarrow N'$ ) = no $\alpha$   $M \Rightarrow N$   $\times$  no $\alpha$   $M' \Rightarrow N'$ 
no $\alpha$  ( $\Rightarrow \beta$   $x$   $M \Rightarrow N$   $M' \Rightarrow N'$ ) = no $\alpha$   $M \Rightarrow N$   $\times$  no $\alpha$   $M' \Rightarrow N'$ 
no $\alpha$  ( $\Rightarrow \alpha$  _ _) =  $\perp$ 

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Next, we prove that given a parallel reduction $M \Rightarrow N$ then it exists a term P , alpha equivalent to N , such that $M \Rightarrow P$ with no alpha rules in the derivations.

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lemma $\Rightarrow$ no $\alpha$  :  $\{M N : \Lambda\} \rightarrow M \Rightarrow N \rightarrow \exists (\lambda P \rightarrow P \sim \alpha N \times \Sigma (M \Rightarrow P) (\lambda M \Rightarrow P \rightarrow \text{no}\alpha M \Rightarrow P))$ 

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Proof. By induction on parallel relation.

var. rule: The proof is direct.

abs. rule: Hypotheses: $\lambda \frac{M \Rightarrow N}{\lambda x M \Rightarrow \lambda x N}$

Thesis: $\exists P, \lambda x N \sim P, \lambda x M \Rightarrow P$ with no α rules.

Proof:

$$\text{ind. hyp} \frac{\frac{M \Rightarrow N}{\exists P', P' \sim N, M \Rightarrow P' \text{ with no } \alpha \text{ rules}}}{\lambda x M \Rightarrow \lambda x P' \text{ with no } \alpha \text{ rules}} \quad \frac{P' \sim N}{\lambda x P' \sim \lambda x N}$$

$$\Rightarrow \lambda \frac{\lambda x M \Rightarrow \lambda x P' \text{ with no } \alpha \text{ rules} \quad \lambda x P' \sim \lambda x N}{\exists P = \lambda x P', \lambda x N \sim P, \lambda x M \Rightarrow P \text{ with no } \alpha \text{ rules}}$$

app. rule: Hypotheses: $\Rightarrow \text{app} \frac{M \Rightarrow M' \quad N \Rightarrow N'}{MN \Rightarrow M'N'}$

Thesis: $\exists P, MN \sim P, MN \Rightarrow P$ with no α rules.

Proof: ind. hyp $\frac{M \Rightarrow M'}{\exists Q, Q \sim M', M \Rightarrow Q \text{ with no } \alpha \text{ rules}}$

$$\text{ind. hyp} \frac{N \Rightarrow N'}{\exists R, R \sim N', M \Rightarrow R \text{ with no } \alpha \text{ rules}}$$

Then exists $P = QR$ such that $M'N' \sim P$ by alpha application rule, and $MN \Rightarrow P$ by parallel application rules with no α rules in the derivation.

$$\beta \text{ rule: Hypotheses: } \Rightarrow \beta \frac{M \Rightarrow M' \quad N \Rightarrow N'}{\lambda x MN \Rightarrow M'[x := N']}$$

Thesis: $\exists P, M'[x := N'] \sim P, \lambda x MN \Rightarrow P$ with no α rules.

$$\text{Proof: ind. hyp} \frac{M \Rightarrow M'}{\exists Q, Q \sim M', M \Rightarrow Q \text{ with no } \alpha \text{ rules}}$$

$$\text{ind. hyp} \frac{N \Rightarrow N'}{\exists R, R \sim N', N \Rightarrow R \text{ with no } \alpha \text{ rules}}$$

$$(1) \quad \text{subst. lemma} \frac{M' \sim Q \quad N' \sim R}{M'[x := N'] \sim Q[x := R]}$$

$$(2) \quad \Rightarrow \beta \frac{M \Rightarrow Q \text{ no } \alpha \text{ rules} \quad N \Rightarrow R \text{ no } \alpha \text{ rules}}{\lambda x MN \sim Q[x := R] \text{ no } \alpha \text{ rules}}$$

Then exists $P = Q[x := R]$ such that $M'[x := N'] \sim P$ by (1), $\lambda x MN \Rightarrow P$ with no α rules by (2).

$$\alpha \text{ rule: Hypotheses: } \Rightarrow \alpha \frac{M \Rightarrow N \quad N \sim Q}{M \Rightarrow Q}$$

Thesis: $\exists P, Q \sim P, M \Rightarrow P$ with no α rules.

$$\text{Proof: } \frac{\text{ind. hyp} \frac{M \Rightarrow N}{\exists R, R \sim N, M \Rightarrow R \text{ with no } \alpha \text{ rules}}}{\exists P = R / \text{ by transitive of } \sim Q \sim P, M \Rightarrow P \text{ with no } \alpha \text{ rules}}$$

□

With previous result we are able to prove that the parallel relation is also alpha compatible in the left term of the relation.

$$\text{lemma} \Rightarrow \alpha \text{left} : \{M N P : \Lambda\} \rightarrow M \sim \alpha N \rightarrow N \Rightarrow P \rightarrow M \Rightarrow P$$

To prove this result we will use the next similar one, it adds the extra premise that the parallel relation hypothesis has no alpha rules in its derivation.

$$\text{lemma} \Rightarrow \text{no} \alpha \text{left} : \{M N P : \Lambda\} \rightarrow M \sim \alpha N \rightarrow (N \Rightarrow P : N \Rightarrow P) \rightarrow \text{no} \alpha N \Rightarrow P \rightarrow M \Rightarrow P$$

With this weakened lemma and the previously proved lemma $\Rightarrow \text{no} \alpha$ we directly prove the desired result.

$$\text{Proof. prev. lemma} \frac{M \sim N \quad \text{lemma} \Rightarrow \text{no} \alpha \frac{N \Rightarrow P}{\exists Q, Q \sim P, N \Rightarrow Q \text{ with no } \alpha \text{ rules}}}{\Rightarrow \alpha \frac{M \Rightarrow Q}{M \Rightarrow P} \quad Q \sim P} \quad \square$$

Next, we prove our auxiliary lemma.

Proof. By structural induction on the left term M of the alpha-equivalence relation hypothesis

var. case: $M = x$ this case is direct.

app. case: $M = MM'$ it has two subcases depending on the last rule on the parallel relation derivation, which could be an application or beta rule (alpha rule case is discarded by the extra hypothesis).

$$\begin{array}{c} \text{app. rule: Hypotheses:} \quad \sim \text{ app. } \frac{M \sim N \quad M' \sim N'}{MM' \sim NN'} \\ \Rightarrow \text{ app. } \frac{N \Rightarrow P \text{ with no } \alpha \text{ rules} \quad N' \Rightarrow P' \text{ with no } \alpha \text{ rules}}{NN' \Rightarrow PP'} \end{array}$$

Tesis: $MM' \Rightarrow PP'$

Proof:

$$\begin{array}{c} \text{ind. hyp. } \frac{M \sim N \quad N \Rightarrow P \text{ no } \alpha \text{ rules}}{M \Rightarrow P} \quad \text{ind. hyp. } \frac{M' \sim N' \quad N' \Rightarrow P' \text{ no } \alpha \text{ rules}}{M' \Rightarrow P'} \\ \Rightarrow \text{ app. } \frac{M \Rightarrow P \quad M' \Rightarrow P'}{MM' \Rightarrow PP'} \end{array}$$

β rule: Hypotheses: $M = \lambda x MM'$

$$\begin{array}{c} \sim \text{ app. } \frac{\lambda x M \sim \lambda y N \quad M' \sim N'}{\lambda x MM' \sim \lambda y NN'} \\ \Rightarrow \text{ app. } \frac{N \Rightarrow P \text{ with no } \alpha \text{ rules} \quad N' \Rightarrow P' \text{ with no } \alpha \text{ rules}}{\lambda y NN' \Rightarrow P[y := P']} \end{array}$$

Tesis: $\lambda x MM' \Rightarrow P[y := P']$

Proof:

$$\begin{array}{c} \text{ind. hyp. } \frac{\lambda x M \sim \lambda y N \quad \text{subst. lemma } \Rightarrow \frac{N \Rightarrow P}{N[y := x] \Rightarrow P[y := x]}}{M \sim N[y := x]} \quad \text{lemma} \Rightarrow \text{no } \alpha \quad \frac{\exists Q, Q \sim P[y := x], N[y := x] \Rightarrow Q \text{ no } \alpha \text{ rules}}{M \Rightarrow Q} \\ \Rightarrow \alpha \quad \frac{M \Rightarrow Q \quad Q \sim P[y := x]}{M \Rightarrow P[y := x]} \quad \text{ind. hyp. } \frac{M' \sim N' \quad N' \Rightarrow P'}{M' \Rightarrow P'} \\ \Rightarrow \beta \quad \frac{M \Rightarrow P[y := x] \quad \lambda x MM' \Rightarrow P[y := x][x := P']}{\lambda x MM' \Rightarrow P[y := x][x := P']} \\ \frac{x \# \lambda x M \quad \lambda x M \sim \lambda y N}{x \# \lambda y N} \quad \frac{\lambda y N \Rightarrow \lambda y P}{x \# \lambda y P} \\ \Rightarrow \alpha \quad \frac{\lambda x MM' \Rightarrow P[y := x][x := P'] \quad \frac{P[y := x][x := P'] \sim P[y := P']}{x \# \lambda y P}}{\lambda x MM' \Rightarrow P[y := P']} \end{array}$$

abs. case: $M = \lambda x M$

$$\text{Hypotheses: } \lambda x M \sim \lambda y N, \Rightarrow \text{ abs. } \frac{N \Rightarrow P}{\lambda y N \Rightarrow \lambda y P}$$

Tesis: $\lambda x M \Rightarrow \lambda y P$

Proof:

$$\begin{array}{c} \text{ind. hyp. } \frac{\lambda x M \sim \lambda y N \quad \text{subst. lemma } \Rightarrow \frac{N \Rightarrow P}{N[y := x] \Rightarrow P[y := x]}}{M \sim N[y := x]} \quad \text{lemma} \Rightarrow \text{no } \alpha \quad \frac{\exists Q, Q \sim P[y := x], N[y := x] \Rightarrow Q \text{ no } \alpha \text{ rules}}{M \Rightarrow Q} \end{array}$$

$$\begin{array}{c}
\Rightarrow_{\alpha} \frac{M \Rightarrow Q \quad Q \sim P[y := x]}{M \Rightarrow P[y := x]} \quad \frac{x \# \lambda x M \quad \lambda x M \sim \lambda y N}{x \# \lambda y N} \quad \frac{\lambda y N \Rightarrow \lambda y P}{\lambda x P[y := x] \sim \lambda y P} \\
\Rightarrow_{\text{abs.}} \frac{\lambda x M \Rightarrow \lambda x P[y := x]}{\lambda x M \Rightarrow \lambda y P} \\
\Rightarrow_{\alpha} \frac{\lambda x M \Rightarrow \lambda x P[y := x] \quad \lambda x P[y := x] \sim \lambda y P}{\lambda x M \Rightarrow \lambda y P}
\end{array}$$

□

Finally, we present the diamond property of the parallel reduction.

diam \Rightarrow : diamond $_ \Rightarrow _$

To prove this result we use a similar approach to previous lemma, proving first a weakened version of it, which adds that the parallel reduction premises have no α rules in its derivations.

$$\begin{array}{l}
\text{diam}\Rightarrow\text{no}\alpha : \{M \ N \ P : \bigwedge\} \rightarrow (M \Rightarrow N : M \Rightarrow N) \rightarrow (M \Rightarrow P : M \Rightarrow P) \\
\rightarrow \text{no}\alpha \ M \Rightarrow N \rightarrow \text{no}\alpha \ M \Rightarrow P \\
\rightarrow \exists (\lambda \ Q \rightarrow N \Rightarrow Q \times P \Rightarrow Q)
\end{array}$$

With previous lemma and the lemma $\Rightarrow\text{no}\alpha$ result the proof is direct.

$$\begin{array}{c}
\text{Proof. lemma}\Rightarrow\text{no}\alpha \frac{M \Rightarrow N}{\exists N', N' \sim N, M \Rightarrow N' \text{ with no } \alpha \text{ rules}} \\
\text{lemma}\Rightarrow\text{no}\alpha \frac{M \Rightarrow N'}{M \Rightarrow P} \\
\text{diam}\Rightarrow\text{no}\alpha \frac{\exists P', P' \sim P, M \Rightarrow P' \text{ with no } \alpha \text{ rules} \quad M \Rightarrow N' \text{ with no } \alpha \text{ rules} \quad M \Rightarrow P' \text{ with no } \alpha \text{ rules}}{M \Rightarrow P} \\
\text{lemma}\Rightarrow\alpha\text{left} \frac{N \sim N' \quad N' \Rightarrow Q \quad \exists Q, N' \Rightarrow Q \Rightarrow Q \quad \text{lemma}\Rightarrow\alpha\text{left} \frac{P \sim P' \quad P' \Rightarrow Q}{P \Rightarrow Q}}{N \Rightarrow Q \quad \exists Q, N \Rightarrow Q, P \Rightarrow Q}
\end{array}$$

□

Next, we prove the auxiliary lemma.

Proof. By induction on the parallel reduction relations. Since α rules can not occur by hypotheses, the possible cases are reduced to the following ones:

var-var rules: The proof is direct.

$$\begin{array}{c}
\text{abs-abs rules: Hypotheses: } \Rightarrow_{\text{abs}} \frac{M \Rightarrow N \text{ with no } \alpha \text{ rules}}{\lambda x M \Rightarrow \lambda x N} \\
\Rightarrow_{\text{abs}} \frac{M \Rightarrow P \text{ with no } \alpha \text{ rules}}{\lambda x M \Rightarrow \lambda x P}
\end{array}$$

Thesis: $\exists Q, \lambda x N \Rightarrow Q, \lambda x P \Rightarrow Q$

Proof:

$$\begin{array}{c}
\text{ind. hyp. } \frac{M \Rightarrow N \text{ with no } \alpha \text{ rules} \quad M \Rightarrow P \text{ with no } \alpha \text{ rules}}{\exists Q', N \Rightarrow Q', P \Rightarrow Q'} \\
\Rightarrow_{\text{abs}} \frac{N \Rightarrow Q' \quad \lambda x N \Rightarrow \lambda x Q'}{\lambda x N \Rightarrow \lambda x Q'} \quad \Rightarrow_{\text{abs}} \frac{P \Rightarrow Q' \quad \lambda x P \Rightarrow \lambda x Q'}{\lambda x P \Rightarrow \lambda x Q'} \\
\frac{\lambda x N \Rightarrow \lambda x Q' \quad \lambda x P \Rightarrow \lambda x Q'}{\exists Q = \lambda x Q', \lambda x N \Rightarrow Q, \lambda x P \Rightarrow Q}
\end{array}$$

$$\begin{array}{l} \text{app-app rules: Hypotheses: } \Rightarrow_{\text{app}} \frac{M \Rightarrow N \text{ with no } \alpha \text{ rules} \quad M' \Rightarrow N' \text{ with no } \alpha \text{ rules}}{MM' \Rightarrow NN'} \\ \Rightarrow_{\text{app}} \frac{M \Rightarrow P \text{ with no } \alpha \text{ rules} \quad M' \Rightarrow P' \text{ with no } \alpha \text{ rules}}{MM' \Rightarrow PP'} \end{array}$$

Thesis: $\exists Q, NN' \Rightarrow Q, PP' \Rightarrow Q$

Proof:

$$\begin{array}{l} \text{ind. hyp. } \frac{M \Rightarrow N \text{ with no } \alpha \text{ rules} \quad M \Rightarrow P \text{ with no } \alpha \text{ rules}}{\exists Q', N \Rightarrow Q', P \Rightarrow Q'} \\ \text{ind. hyp. } \frac{M' \Rightarrow N' \text{ with no } \alpha \text{ rules} \quad M' \Rightarrow P' \text{ with no } \alpha \text{ rules}}{\exists Q'', N' \Rightarrow Q'', P' \Rightarrow Q''} \\ \Rightarrow_{\text{app}} \frac{N \Rightarrow Q' \quad N' \Rightarrow Q''}{NN' \Rightarrow Q'Q''} \quad \Rightarrow_{\text{app}} \frac{P \Rightarrow Q' \quad P' \Rightarrow Q''}{PP' \Rightarrow Q'Q''} \\ \hline \exists Q = Q'Q'', NN' \Rightarrow Q, PP' \Rightarrow Q \end{array}$$

$$\begin{array}{l} \text{beta-beta rules: Hypotheses: } \Rightarrow_{\beta} \frac{M \Rightarrow N \text{ with no } \alpha \text{ rules} \quad M' \Rightarrow N' \text{ with no } \alpha \text{ rules}}{\lambda x MM' \Rightarrow N[x := N']} \\ \Rightarrow_{\beta} \frac{M \Rightarrow P \text{ with no } \alpha \text{ rules} \quad M' \Rightarrow P' \text{ with no } \alpha \text{ rules}}{\lambda x MM' \Rightarrow P[x := P']} \end{array}$$

Thesis: $\exists Q, N[x := N'] \Rightarrow Q, P[x := P'] \Rightarrow Q$

Proof:

$$\begin{array}{l} \text{ind. hyp. } \frac{M \Rightarrow N \text{ with no } \alpha \text{ rules} \quad M \Rightarrow P \text{ with no } \alpha \text{ rules}}{\exists Q', N \Rightarrow Q', P \Rightarrow Q'} \\ \text{ind. hyp. } \frac{M' \Rightarrow N' \text{ with no } \alpha \text{ rules} \quad M' \Rightarrow P' \text{ with no } \alpha \text{ rules}}{\exists Q'', N' \Rightarrow Q'', P' \Rightarrow Q''} \end{array}$$

Applying the substitution lemma of parallel reduction (lemma \Rightarrow).

$$\text{lemma}\Rightarrow \frac{N \Rightarrow Q' \quad \frac{N' \Rightarrow Q''}{[x := N'] \Rightarrow [x := Q'']}}{N[x := N'] \Rightarrow Q'[x := Q'']}$$

In a similar way we have that $P[x := P'] \Rightarrow Q'[x := Q'']$, then exists $Q = Q'[x := Q'']$ that satisfies the thesis.

beta-app rules: Similar to the two previous cases combined.

app-beta rules: Simetric to previous case.

□

Beta contraction

$$\begin{array}{l} \text{data } _ \triangleright _ : \Lambda \rightarrow \Lambda \rightarrow \text{Set where} \\ \triangleright \beta : \{M N : \Lambda\} \{x : V\} \\ \rightarrow \lambda x M \cdot N \triangleright M \bullet \iota \prec + (x, N) \end{array}$$

Contextual closure of beta contraction.

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data ctx (r : Rel) : Rel where
  ctxinj : ∀ {t t'} → r t t' → ctx r t t'
  ctx·l : ∀ {t₁ t₁' t₂} → ctx r t₁ t₁' → ctx r (t₁ · t₂) (t₁' · t₂)
  ctx·r : ∀ {t₁ t₂ t₂'} → ctx r t₂ t₂' → ctx r (t₁ · t₂) (t₁ · t₂')
  ctxλ : ∀ {x t₁ t₁'} → ctx r t₁ t₁' → ctx r (λ x t₁) (λ x t₁')

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infix 4 _→β_
_→β_ : Rel
_→β_ = ctx _▷_

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Alpha beta union relation.

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_→α_ : Rel
_→α_ = _→β_ ∪ _~α_

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Reflexive clausure of previous relation

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_→α*_ : Rel
_→α*_ = star _→α_

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To end the proof of \rightarrow_α relation confluence we the following two results:

- $\rightarrow_\alpha \subseteq \Rightarrow$
- $\Rightarrow \subseteq \rightarrow_\alpha^*$

Both proofs are a direct induction on the left relations on both inclusions. In the beta case of the second proof we use the following two substitutions lemmas for the \rightarrow_α^* relation:

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lemma→α*Λ : {M M' : Λ}{σ : Σ}
  → M →α* M'
  → M • σ →α* M' • σ

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Proof. This lemma is proved by induction on the reflexive and transitive closure of \rightarrow_α relation. The reflexive and transitive cases of the closure are direct. The only remaining case is the one step case, which can be an α or β step.

- α case: As the substitution normalises α -convertible terms the result is a direct reflexive step.
- β case: We do an induction on the closure cases, all are direct, except from the abstraction and β -contraction cases that we discuss next.
 - abstraction case:

$$\begin{array}{lll}
 (\lambda x M) \sigma & \rightarrow_\alpha^* & \{\text{by substitution definition, with } y = \chi \sigma (\lambda x M) \} \\
 \lambda y (M \sigma \prec+(x, y)) & \rightarrow_\alpha^* & \{\text{by inductive hypotheses and that the} \} \\
 & & \{\rightarrow_\alpha^* \text{ relation is preserved by the application of an abstraction} \} \\
 \lambda y (M' \sigma \prec+(x, y)) & \rightarrow_\alpha^* & \{\text{by being } \alpha\text{-convertible, with } y' = \chi \sigma (\lambda x M') \} \\
 \lambda y' (M' \sigma \prec+(x, y')) & \rightarrow_\alpha^* & \{\text{by substitution definition} \} \\
 ((\lambda x M') \sigma & &
 \end{array}$$

– β case:

$$\begin{array}{rcl}
((\lambda x M)N)\sigma & \rightarrow_{\alpha}^* & \{\text{by substitution definition, with } y = \chi \sigma (\lambda x M) \} \\
(\lambda y (M\sigma \prec+(x, y)))(N\sigma) & \rightarrow_{\alpha}^* & \{\beta\text{-contraction} \} \\
(M\sigma \prec+(x, y))\iota \prec+(y, N\sigma) & \rightarrow_{\alpha}^* & \{\text{substitution lemma corollary 1} \} \\
(M\sigma \prec+(x, N\sigma)) & \rightarrow_{\alpha}^* & \{\text{corollary 1 prop 7 Stoughton} \} \\
(M\iota \prec+(x, N))\sigma & & \}
\end{array}$$

□

Finally we prove the remaining lemma to complete the relations inclusions proof.

$$\begin{array}{l}
\text{lemma} \rightarrow \alpha^* \Sigma : \{M : \Lambda\} \{ \sigma \sigma' : \Sigma \} \\
\rightarrow \sigma \rightarrow \alpha^* \sigma' \downarrow M \\
\rightarrow M \bullet \sigma \rightarrow \alpha^* M \bullet \sigma'
\end{array}$$

Proof. By induction on terms. The only non direct case is the following abstraction case:

$$\begin{array}{rcl}
(\lambda x M)\sigma & \rightarrow_{\alpha}^* & \{\text{by substitution definition, with } y = \chi \sigma (\lambda x M) \} \\
\lambda y (M\sigma \prec+(x, y)) & \rightarrow_{\alpha}^* & \{\text{by inductive hypotheses and that the} \} \\
& & \{\rightarrow_{\alpha}^* \text{ relation is preserved by the application of an abstraction} \} \\
\lambda y (M\sigma' \prec+(x, y)) & \rightarrow_{\alpha}^* & \{\text{by being } \alpha\text{-convertible, with } y' = \chi \sigma (\lambda x M') \} \\
\lambda y' (M\sigma' \prec+(x, y')) & \rightarrow_{\alpha}^* & \{\text{by substitution definition} \} \\
(\lambda x M)\sigma' & & \}
\end{array}$$

□