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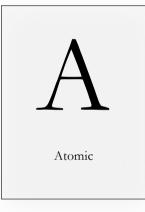




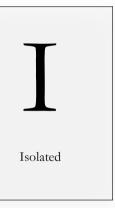


In this session

- What is ACID compliance
- Isolation levels in SQL Server
- Durability in SQL Server

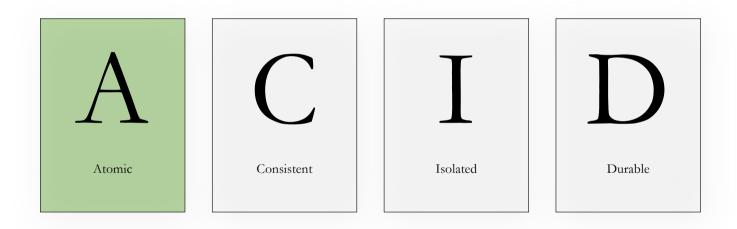




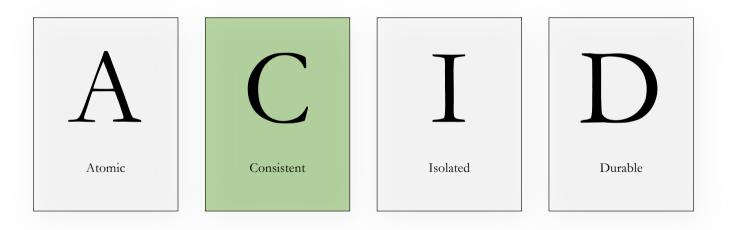




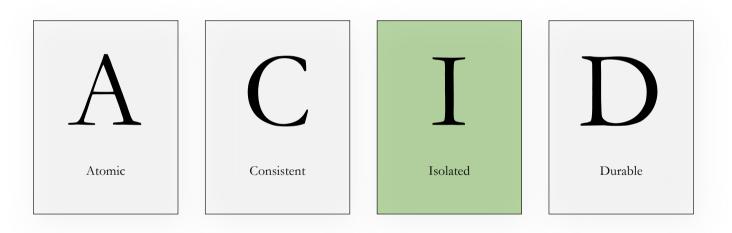
Durable



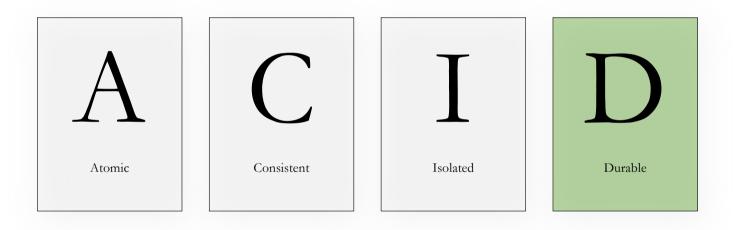
All or nothing. The *business* logic.



A single truth. The *database* logic.



No interference when changing data.



Keep all promises.

Who am I?

- SQL Server developer since 1997
- Consultant
- Organizer of Data Saturday Stockholm
- Data Platform MVP
- Dog person

@dhmacher on (almost) all the socials.



Isolation: Quirks and features of

multi-user databases

- no Nobel literature prize for you.

Transaction UPDATE

Read Uncommitted (a.k.a. NOLOCK)

SELECT

- no Nobel literature prize for you.

Transaction UPDATE

Read Committed Wait for lock to release SELECT

- no Nobel literature prize for you.

 DML does not respect READ UNCOMMITTED, and cannot make changes to uncommitted data.

- no Nobel literature prize for you.

Transaction UPDATE

Read Uncommitted (a.k.a. NOLOCK)

Wait for lock to release

UPDATE



Non-repeatable reads



Non-repeatable reads





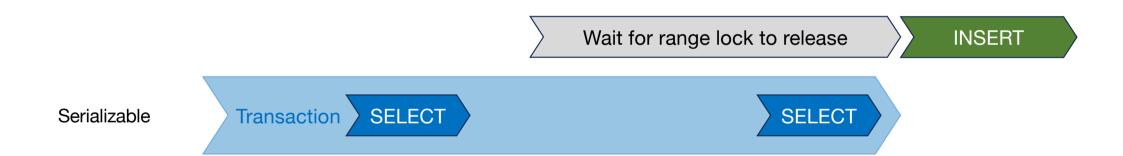
Phantom reads

Your database is gaslighting you.

Phantom reads



Phantom reads





Read uncommitted

Read committed

Repeatable read

Serializable

Read uncommitted

Read committed

Repeatable read

Serializable

Dirty reads

Non-repeatable

Phantom reads

Read uncommitted

Read committed

Repeatable read

Serializable

Dirty reads

Committed reads

Non-repeatable reads

Phantom reads

Read Read Repeatable uncommitted committed read Serializable

Dirty reads Committed reads

Non-repeatable reads Repeatable reads

Phantom reads

Read uncommitted committed read Serializable

Dirty reads

Committed reads

Non-repeatable reads

Phantom reads

So... Serialize all the things, then?

Less locking



More locking

- Less isolation
- Less predictable

- Better isolation
- More predictable

So... Serialize all the things, then?

Less locking



More locking

- Less isolation
- Less predictable
- Better concurrency

- Better isolation
- More predictable
- Lower concurrency

So... Serialize all the things, then?

Less locking



More locking

- Less isolation
- Less predictable
- Better concurrency
- Fewer conflicts

- Better isolation
- More predictable
- Lower concurrency
- More conflicts

Repeatable Read Read Serializable read uncommitted committed Dirty reads Committed reads Non-repeatable reads Repeatable reads Phantom reads When two transactions conflict: Deadlock less likely to block more likely to block .

Transaction 1

UPDATE a

Transaction 2

Transaction 1

UPDATE a

Transaction 2

UPDATE b

Transaction 1

UPDATE a

UPDATE b

Transaction 2

UPDATE b

Transaction 1

UPDATE a

UPDATE b: blocked by 2

Transaction 2

UPDATE b

Transaction 1

UPDATE a

UPDATE b: blocked by 2

Transaction 2

UPDATE b

UPDATE a

Transaction 1

UPDATE a

UPDATE b: blocked by 2

Transaction 2

UPDATE b

UPDATE a: blocked by 1

Transaction 1

UPDATE a

UPDATE b: blocked by 2

Deadlock: Transaction 1 and 2 are waiting on each other.

Transaction 2

UPDATE b

UPDATE a: blocked by 1

SQL Server will fail the transaction with the least amount of work to roll back:

Transaction 1

UPDATE a

UPDATE b: blocked by 2

Rolled back

Transaction 2

UPDATE b

UPDATE a: blocked by 1

Transaction 1

Rolled back

Transaction 1 is no longer blocking a.

Transaction 2

UPDATE b

UPDATE a

Transaction 1

Rolled back

Transaction 2

UPDATE b

UPDATE a

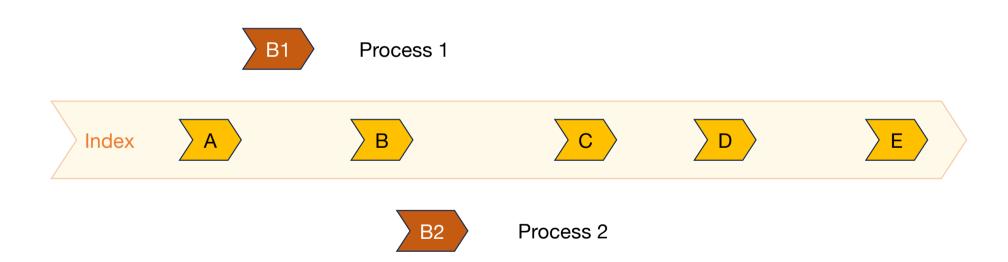
COMMIT

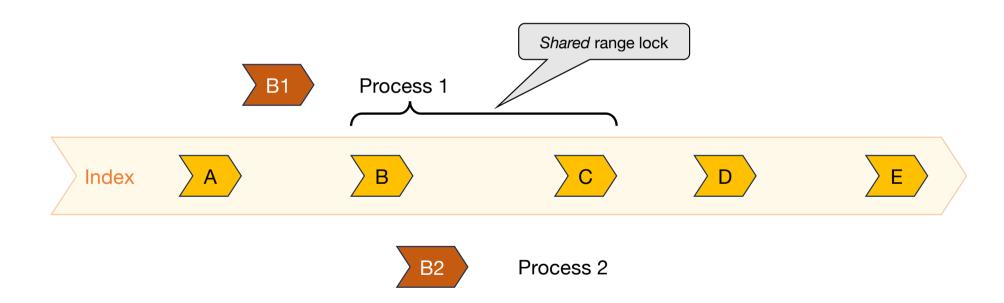


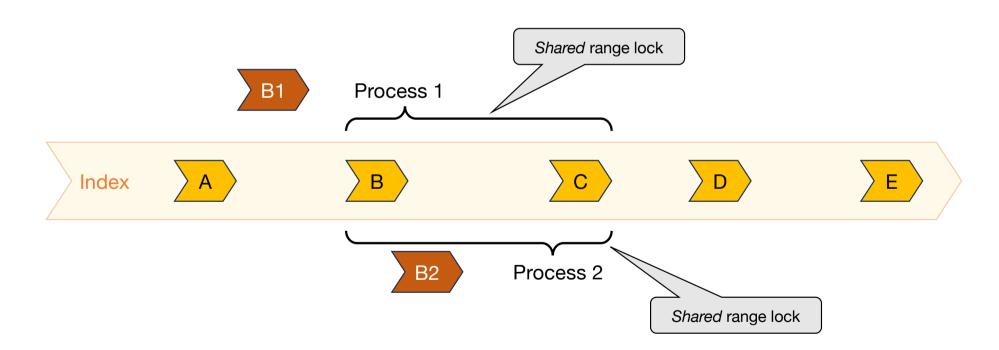


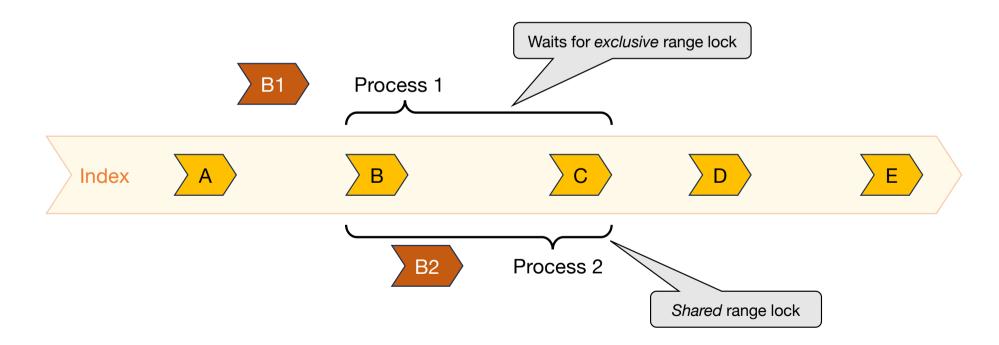


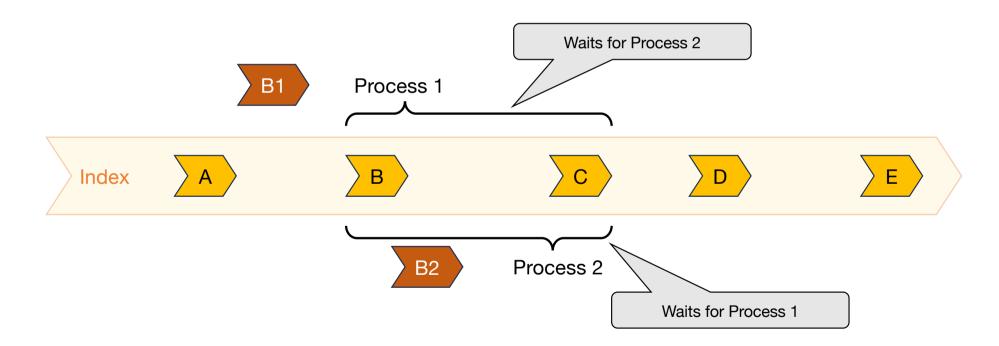




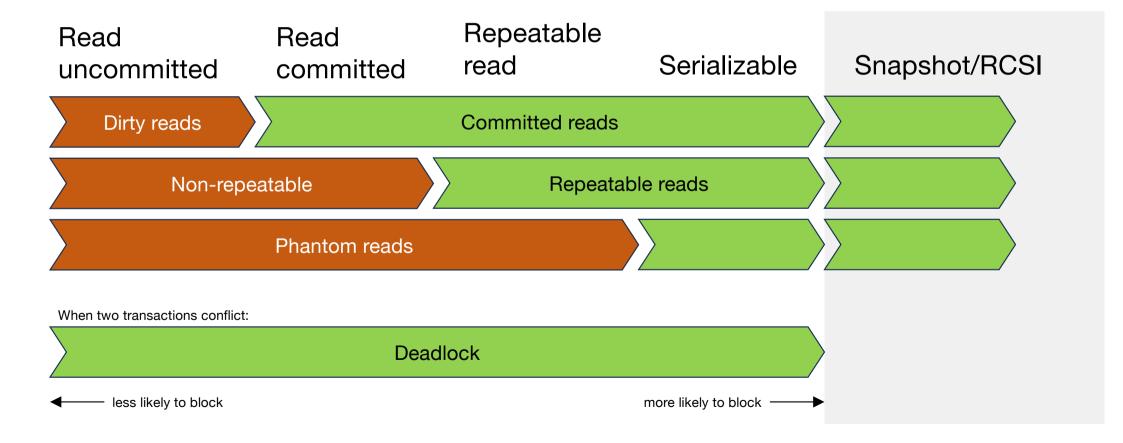




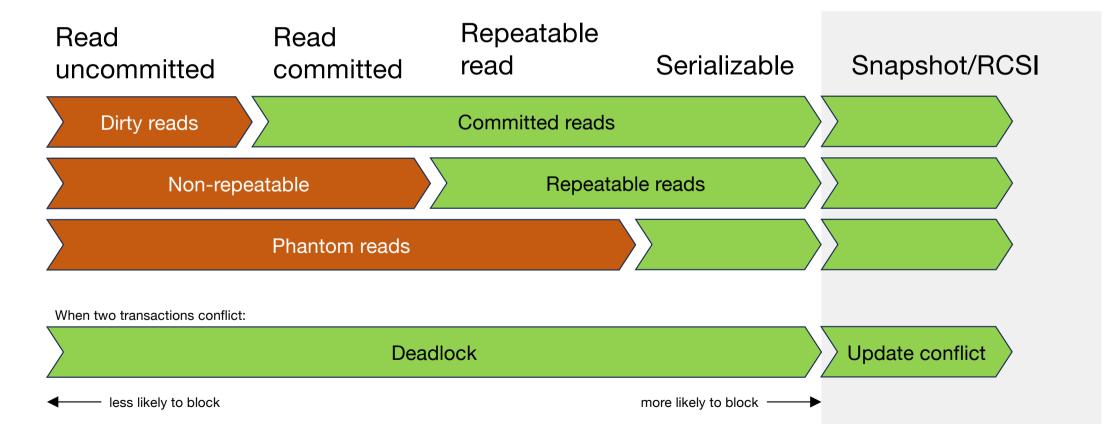




Snapshot

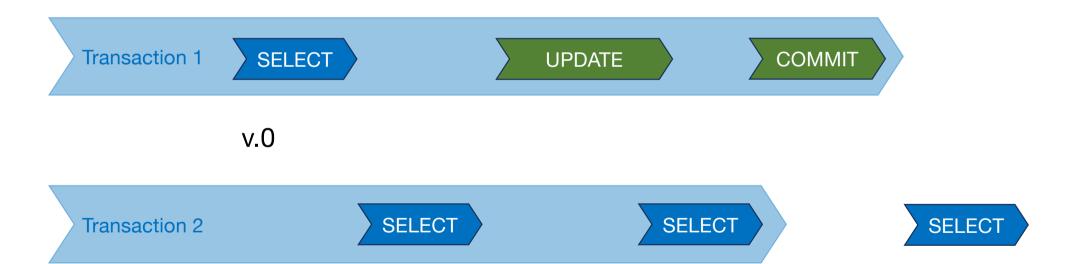


Snapshot



Transaction 1 SELECT UPDATE COMMIT

Transaction 2 SELECT SELECT SELECT







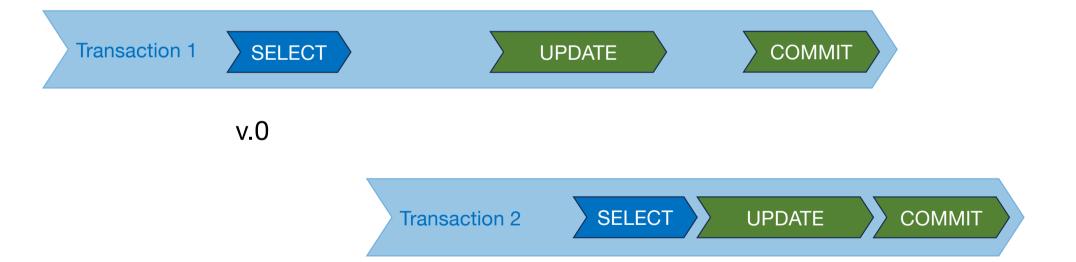








Transaction 2 SELECT UPDATE COMMIT

















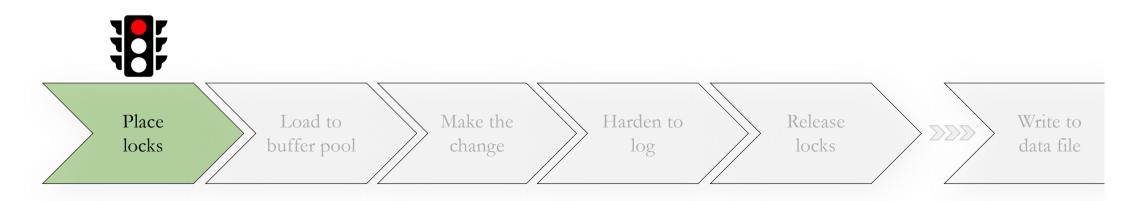
Read Committed Snapshot Isolation

vs. Snapshot Isolation

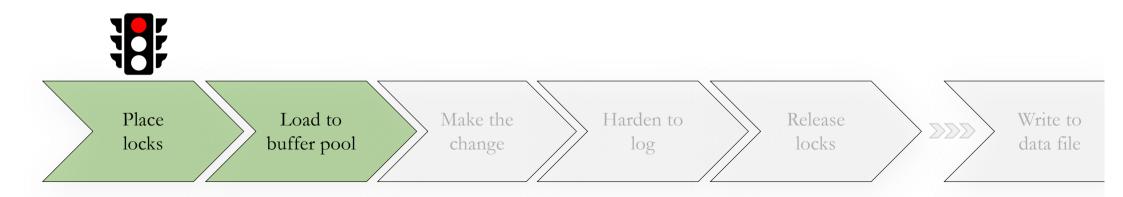
- Snapshot Isolation is set per transaction.
- Snapshot Isolation protects the transaction.
- Snapshot Isolation requires code change.
- Read Committed Snapshot Isolation is a database setting.
- Read Committed Snapshot Isolation protects the statement.
- Both require testing, because they behave differently.

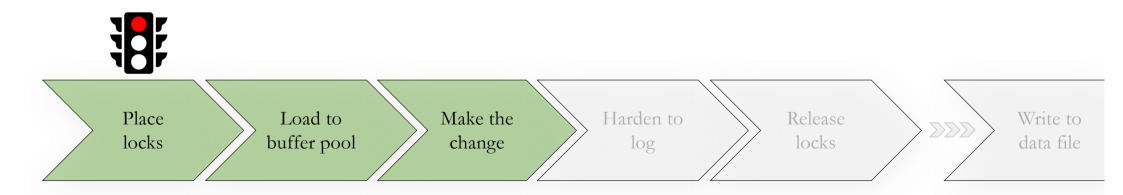
More detail: https://brentozar.com/go/rcsi

How the transaction log works



How the transaction log works







Place locks

Load to buffer pool

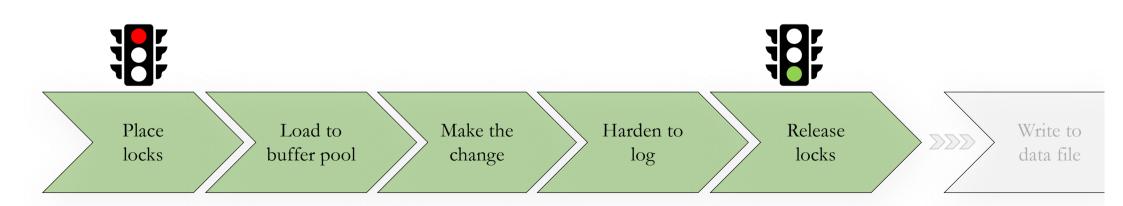
Make the change

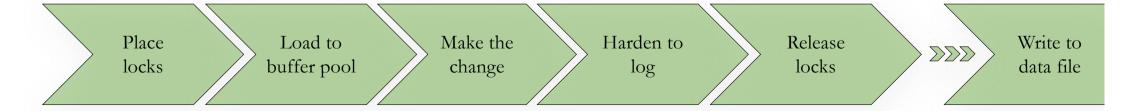
Harden to log

Release locks



Write to data file







Place locks

Load to buffer pool

Make the change

Harden to log

Release locks

Write to data file



Place locks

Load to buffer pool

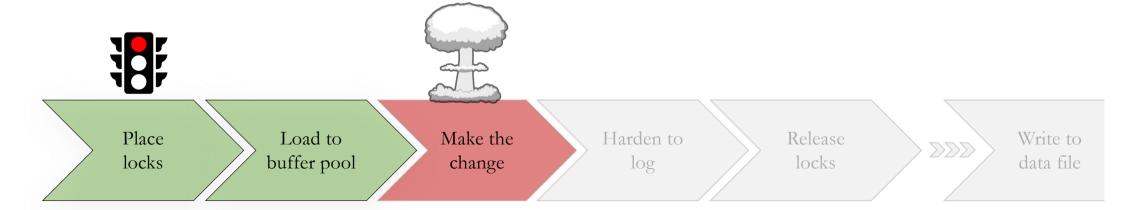
Make the change

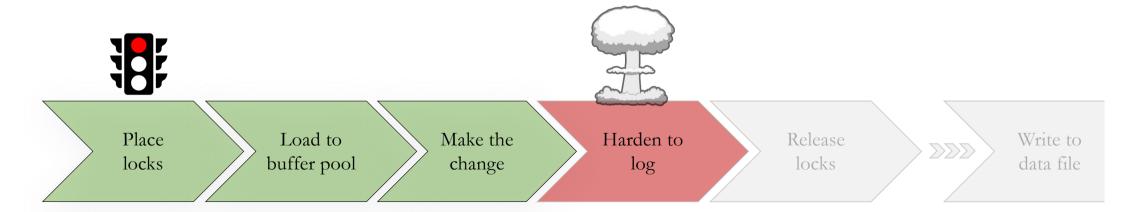
Harden to log

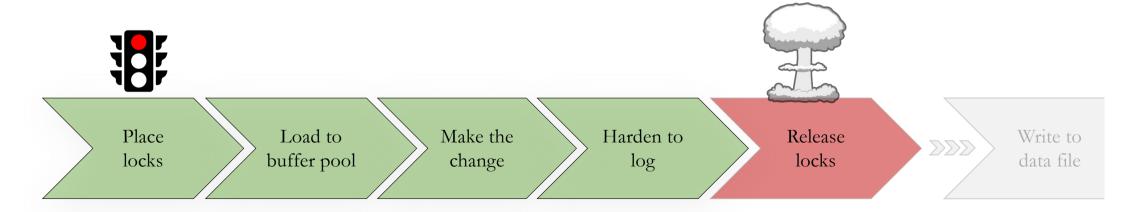
Release locks

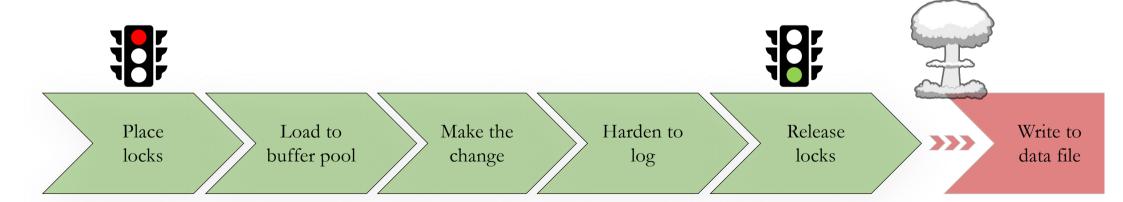


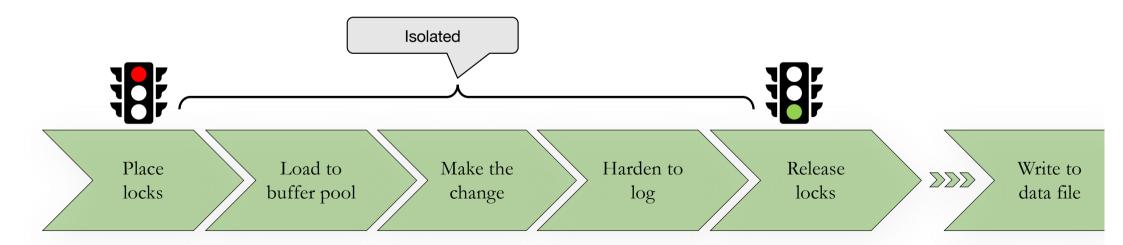
Write to data file

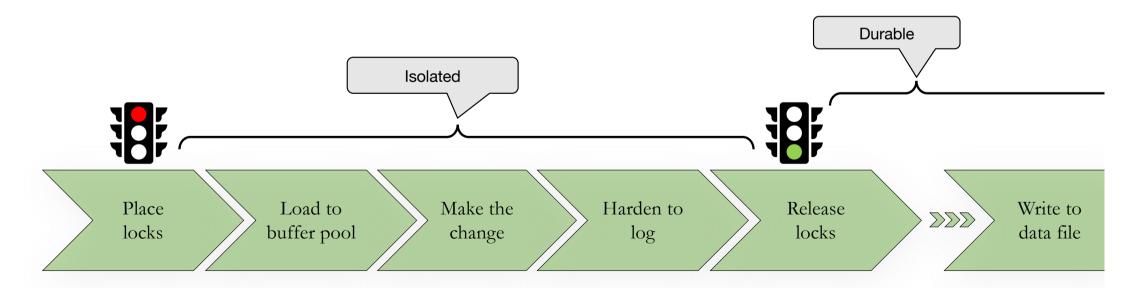


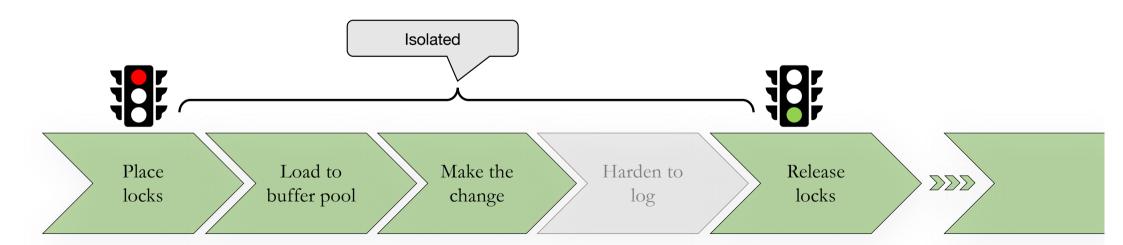


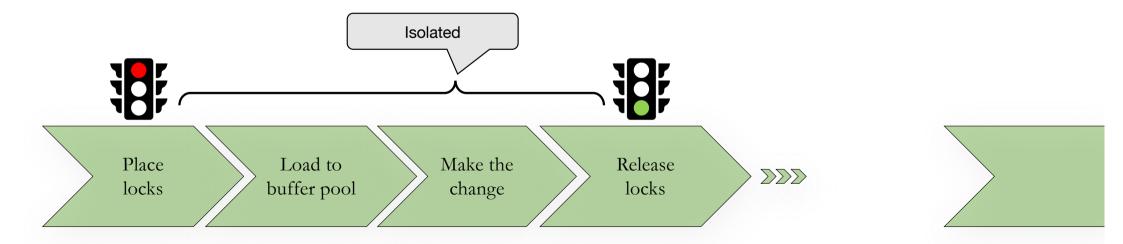


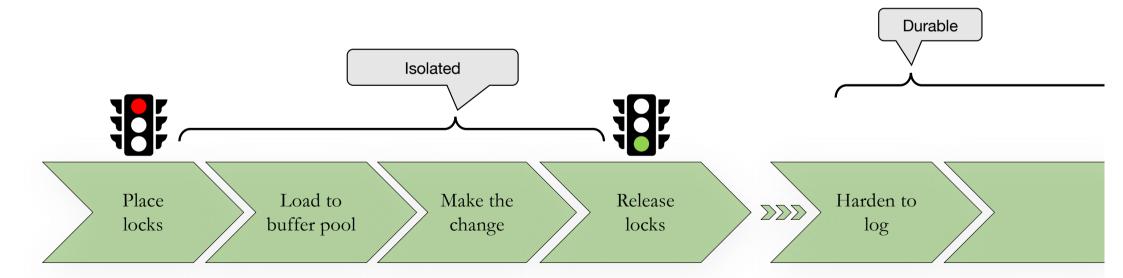


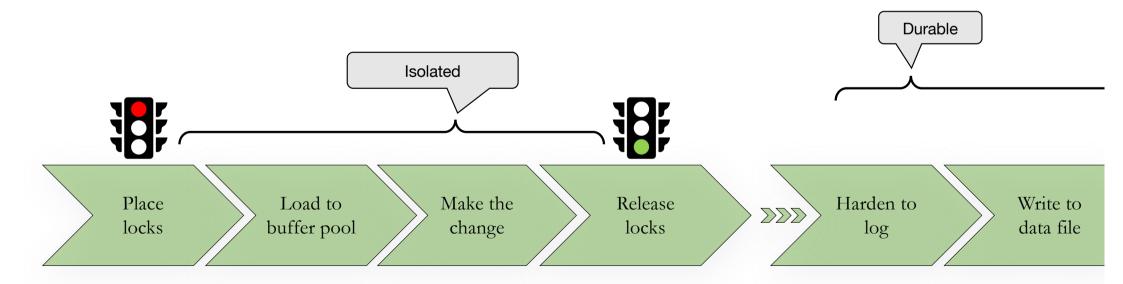




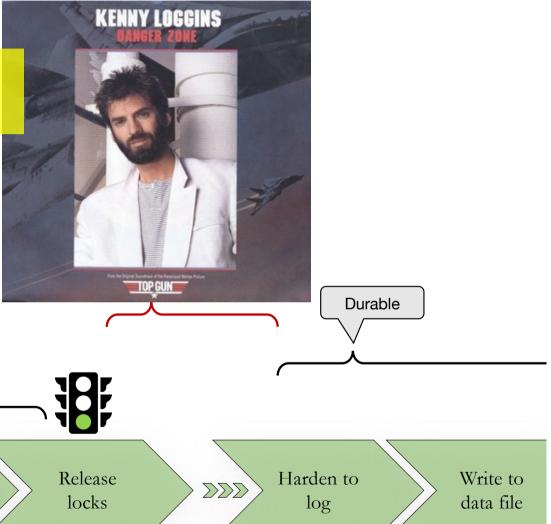














Place locks

Load to buffer pool

Make the change

Isolated

- Writes are batched into the transaction log
- Significantly reduced latency for tiny transactions
- Great if you don't mind losing the data like a staging environment
- tempdb uses a form of Delayed Durability under the hood

Key take-aways

• I was going to do this slide later, but here we are.

Give me feedback. Please?

