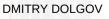




NOSQL FOR POSTGRESQL

BEST PRACTICES



11-07-2017





→ Jsonb internals and performance-related factors



- → Jsonb internals and performance-related factors
- → Tricky queries



- → Jsonb internals and performance-related factors
- → Tricky queries
- → Benchmarks



- → Jsonb internals and performance-related factors
- → Tricky queries
- → Benchmarks
- → How to shoot yourself in the foot



Live Long, and Prosper

- Han Solo



AWS EC2

m4.large instance separate instance (database and generator) 16GB memory, 4 core 2.3GHz Ubuntu 16.04 Same VPC and placement group AMI that supports HVM virtualization type at least 4 rounds of benchmark



PostgreSQL 9.6.3/10

MySQL 5.7.9/8.0

MongoDB 3.4.4

YCSB 0.13

 10^6 rows and operations



Configuration

```
shared buffers
effective cache size
max wal size
innodb buffer pool size
innodb_log_file_size
write concern level (journaled or transaction sync)
checkpoint
eviction
```

Document types

"simple" document 10 key/value pairs (100 characters)

"large" document 100 key/value pairs (200 characters)

"complex" document 100 keys, 3 nesting levels (100 characters)





→ On-disk representation



- → On-disk representation
- → In-memory representation



- → On-disk representation
- → In-memory representation
- → Indexing support



Indexing support

- → Postgresql single path, multiple paths, entire document
- → MongoDB single path, multiple paths
- → MySQL virtual columns, single path, multiple paths



PG indexing details

- → jsonb_path
- → jsonb_path_ops

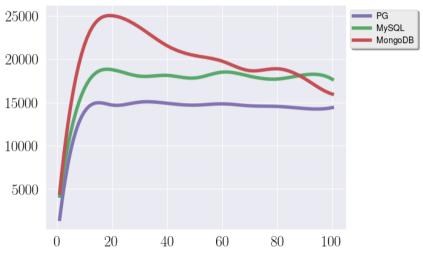


Select, GIN

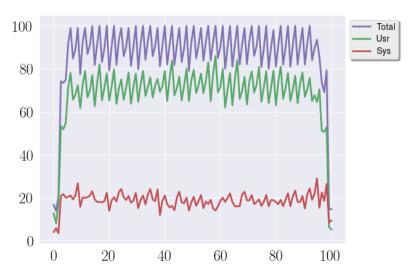
"simple" document jsonb_path_ops



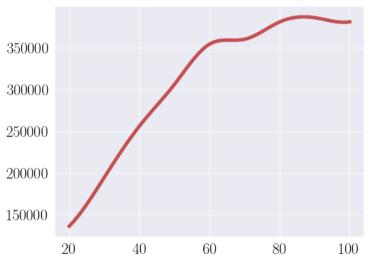
Throughput (ops/sec)



CPU%



CPU migrations (MongoDB)



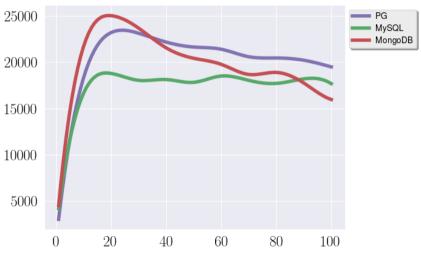


Select, BTree

"simple" document btree

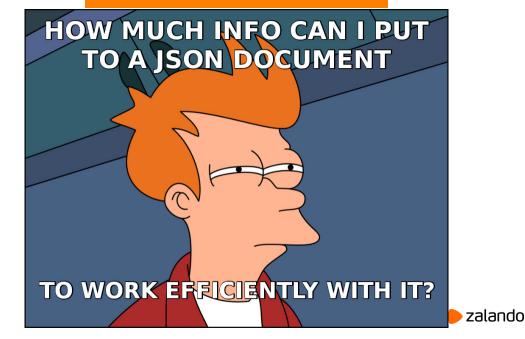


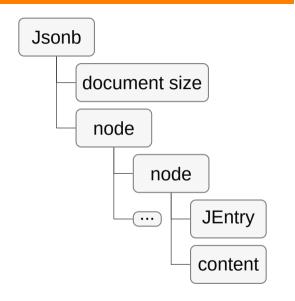
Throughput (ops/sec)

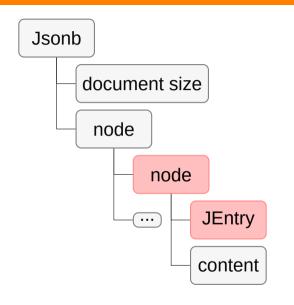


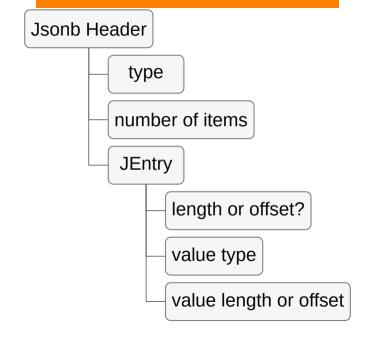
Internals







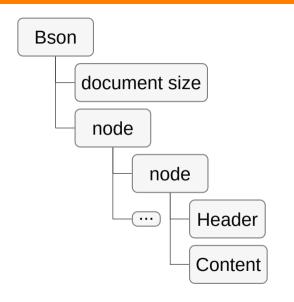


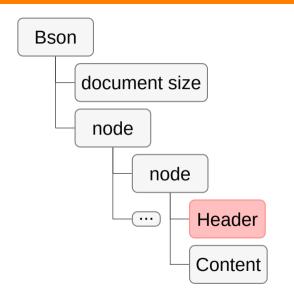


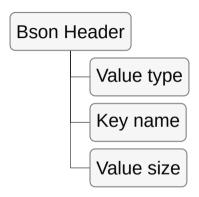
JB_OFFSET_STRIDE

- → JEntry may contains a value length or offset
- → Offset = access speed
- → Length = compressibility
- → Every **JB_OFFSET_STRIDE**'th JEntry contains an offset
- → Rest of them contain length

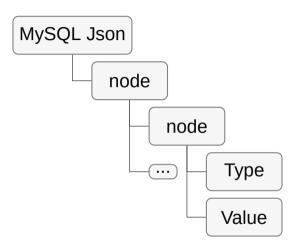


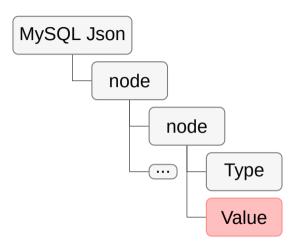


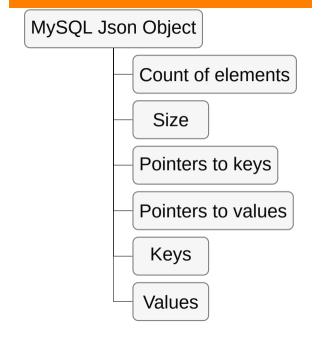


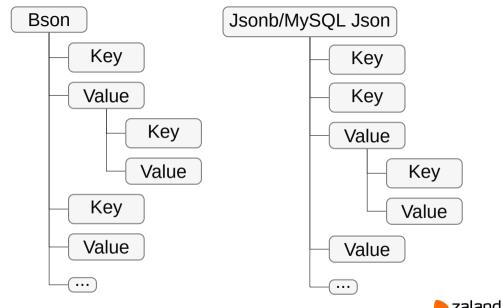












```
{"a": 3, "b": "xyz"}
```



```
select pg relation filepath(oid),
relpages from pg class
where relname = 'table name':
 pg relation filepath | relpages
 base/40960/325477
(1 \text{ row})
```



```
bson.dumps({"a": 3, "b": u"xyz"})
```



\$ hexdump -C database/table.ibd

\x00\x02\x00\x18\x00\x12\x00\x01\x00\x13\x00\x01\x00\x05\x03\x00\x0c\x14\x00ab\x03xyz\x00



Alignment

Variable-length portion is aligned to a 4-byte

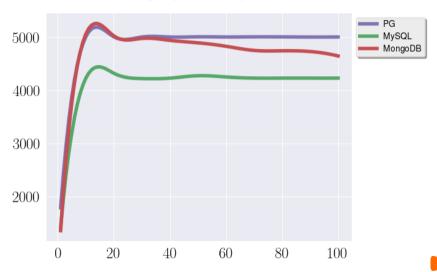
```
insert into test
values('{"a": "aa", "b": 1}');
abaa\x20\x00\x00\x00\x00\x80\x01\x00
insert into test
values('{"a": 1, "b": "aa"}');
```



Select, BTree

"complex" document btree



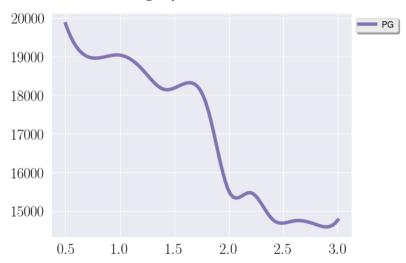


TOAST_TUPLE_THRESHOLD

"simple" document 40 threads different document size select



Throughput, 40 clients





TOAST



- → TOAST_TUPLE_THRESHOLD bytes (normally 2 kB)
- → PostgreSQL and MySQL use LZ variation
- → MongoDB uses snappy block compression



In-memory representation

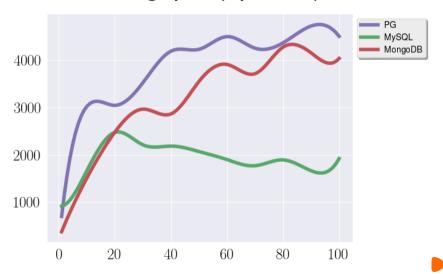
- → Tree-like representation (JsonbValue, Document, Json_dom)
- → Little bit more expensive but more convenient to work with
- → Mostly in use to modify data (except MySQL)
- → Most of the read operations use on-disk representation



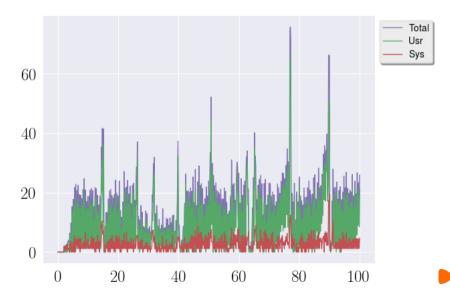
Insert

"simple" document journaled

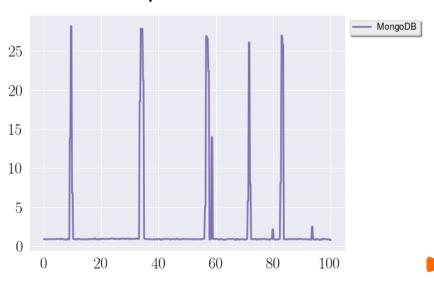


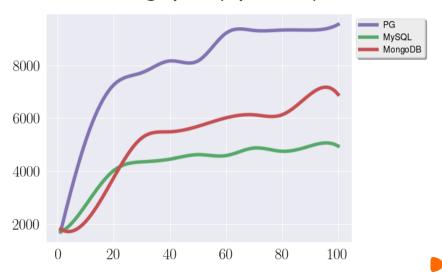


CPU%



IO queue size







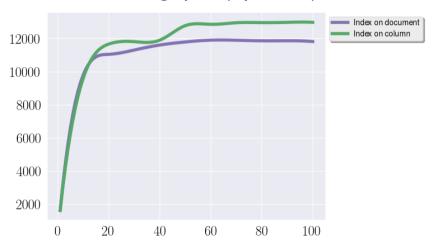
- → Update one field of a document
- → DETOAST of a document (select, constraints, procedures etc.)
- → Reindex of an entire document



Index update

"simple" document Update one field



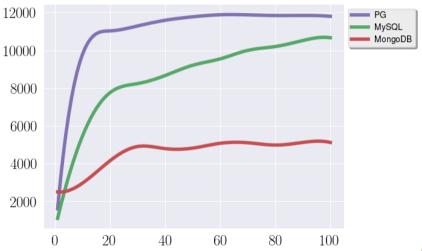




Update 50%, Select 50%

"simple" document
Update one field
journaled
max wal size 5GB

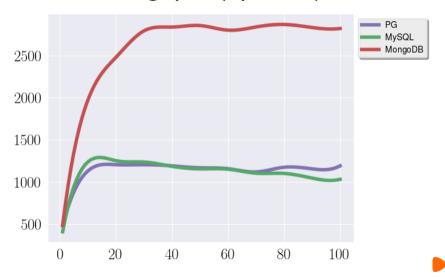




Update 50%, Select 50%

"large" document Update one field





Journal



Postgresql

 $\\ \verb|\x00| x00| aaabbbcccdddbbbccc| x00| x00| x00| x00| x02| x00| x00| \\$

Mysql

MongoDB

 $\xspace{0.05} \xspace{0.05} \xspace{0.05}$



Document slice

"large" document

One field from a document



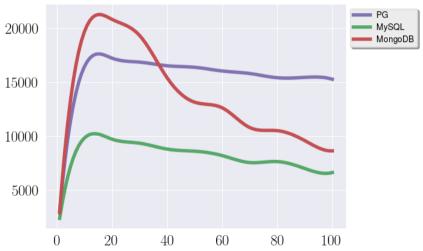
```
select data->'key1'->'key2' from table;
select data->'key1', data->'key2' from table;
```



```
select data->'key1'->'key2' from table;
select data->'key1', data->'key2' from table;
```



```
select data->'key1'->'key2' from table;
select data->'key1', data->'key2' from table;
```

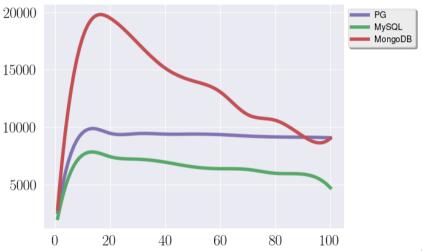




Document slice

"large" document 10 fields from a document





Solutions?

- → set storage external
- → different query



Document slice

```
create type test as ("a" text, "b" text);
insert into test jsonb
values('{"a": 1, "b": 2, "c": 3}');
select q.* from test jsonb,
isonb populate record(NULL::test, data) as q;
a b
1 2
(1 \text{ row})
```

Document slice

```
create type test as ("a" text, "b" text):
insert into test jsonb
values('{"a": 1, "b": 2, "c": 3}');
select q.* from test jsonb,
jsonb_populate_record (NULL::test, data) as q;
a b
1 2
(1 \text{ row})
```

Queries

Pitfalls

- → No Json path out of the box (jsquery, SQL/JSON)
- → Queries with an array somewhere in the middle
- → Iterating through document
- → Update inside document



```
"items": [
    {"id": 1, "value": "aaa"},
    {"id": 2. "value": "bbb"}
"items": [
    {"id": 3, "value": "aaa"},
    {"id": 4, "value": "bbb"}
```

zalando

```
WITH items AS (
    SELECT jsonb_array elements(data->'items')
    AS item FROM test
SELECT * FROM items
WHERE item->>'value' = 'aaa';
item
 {"id": 1, "value": "aaa"}
 {"id": 3, "value": "aaa"}
(2 rows)
```

```
WITH items AS (
    SELECT jsonb_array_elements (data->'items')
    AS item FROM test
SELECT * FROM items
WHERE item->>'value' = 'aaa';
item
 {"id": 1, "value": "aaa"}
 {"id": 3, "value": "aaa"}
(2 rows)
```

```
"items": {
    "item1": {"status": true},
    "item2": {"status": true},
    "item3": {"status": false}
```

```
WITH items AS (
    SELECT jsonb each(data->'items')
    AS item FROM test
SELECT (item).key FROM items
WHERE (item).value->>'status' = 'true';
kev
item1
item2
(2 rows)
```

```
WITH items AS (
    SELECT jsonb_each (data->'items')
    AS item FROM test
SELECT (item).key FROM items
WHERE (item).value->>'status' = 'true';
kev
item1
item2
(2 rows)
```



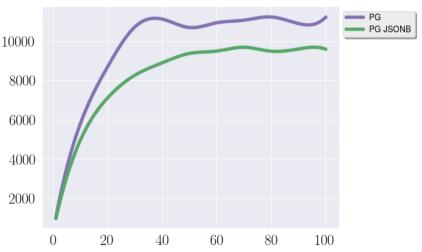
SQL vs JSONB

"simple" document

btree

insert





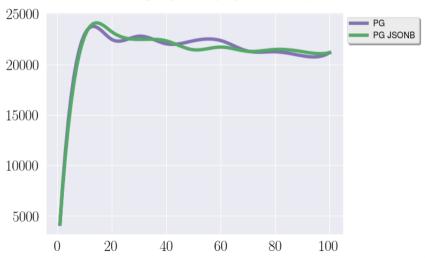


SQL vs JSONB

"simple" document btree

select





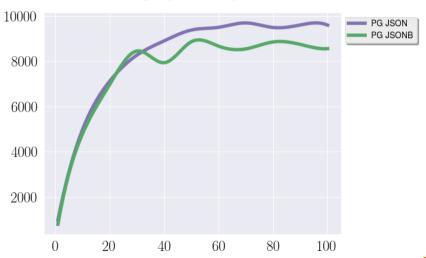
JSON vs JSONB

"simple" document

btree

insert





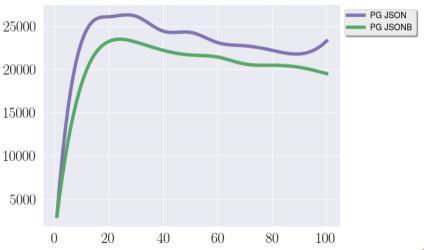
JSON vs JSONB

"simple" document

btree

select







→ Jsonb is more that good for many use cases



- → Jsonb is more that good for many use cases
- → Reasons for performance difference is mostly database itself

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- → Benchmarks above are only "hints"



- → Jsonb is more that good for many use cases
- → Reasons for performance difference is mostly database itself
- → Benchmarks above are only "hints"
- → You need your own tests

Questions?

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- ≥ 9erthalion6 at gmail dot com

