



Calculating the future

How to model PostgreSQL performance

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14-03-2024

- Why is it relevant for you?
- Back of the envelope calculations
- Approximation
- Simulation

Why bother?

Benchmarking instead?

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- Resource intensive

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- Hard to get full coverage

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Enhance benchmarking!

It's easy

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→ just bump max_wal_size?

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- just bump max_wal_size?
- just increase shared_buffers?

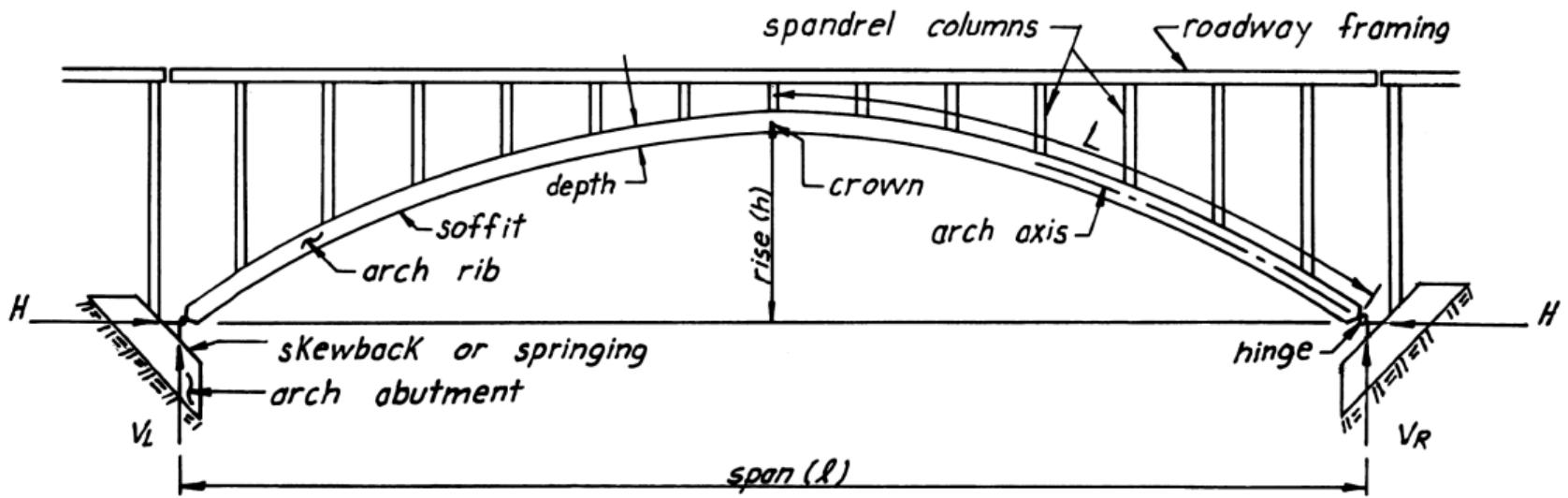
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- just configure autovacuum?

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Well...



Douglas A. Nettleton, John S. Torkelson Department of Transportation, Federal Highway Administration, Office of Engineering, Bridge Division



OSM



Red Hat

Target of the research

```
create table test(a int);  
create index on test(a);
```

```
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create index on test(a);
```

Back of the envelope calculations

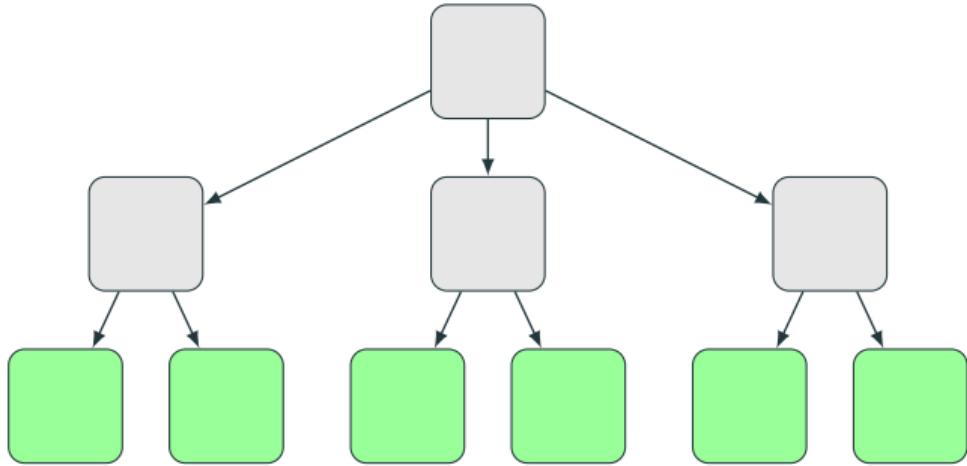
[...] we present a minimal, parsimonious model to account for the cost of maintaining a high body temperature [...]. A body temperature of 36.7°C maximizes fitness by restricting the growth of most fungal species relative to its metabolic cost.

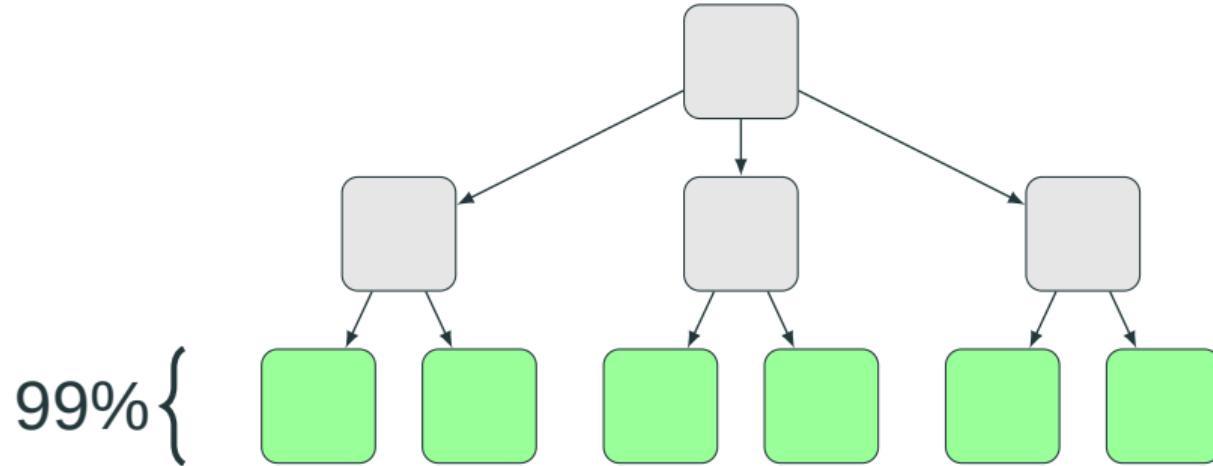
Bergman A, Casadevall A. 2010. Mammalian Endothermy Optimally Restricts Fungi and Metabolic Costs. *mBio* 1:10.1128/mbio.00212-10. <https://doi.org/10.1128/mbio.00212-10>

Assuming we know the schema,
how to approximate space usage?

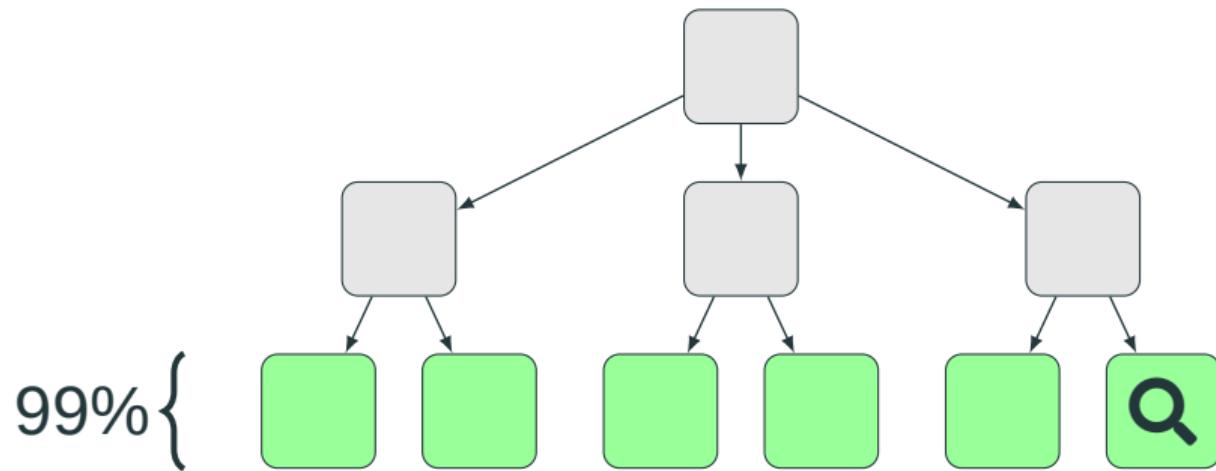
```
create table test(a int);  
create index on test(a);
```

```
create table test(a int);
create index on test(a)
    with (fillfactor = 100);
```



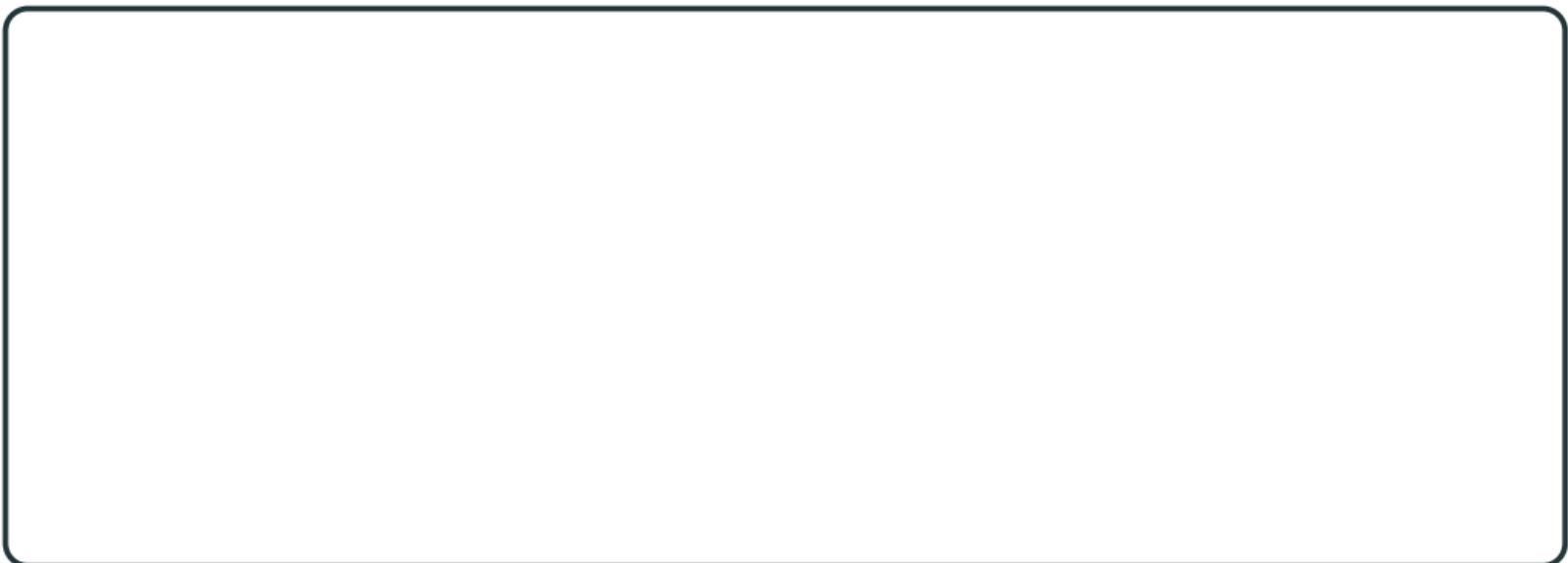


Goetz Graefe. "Modern B-Tree Techniques." Foundations and Trends in Databases 3.4 (2010) 203-402



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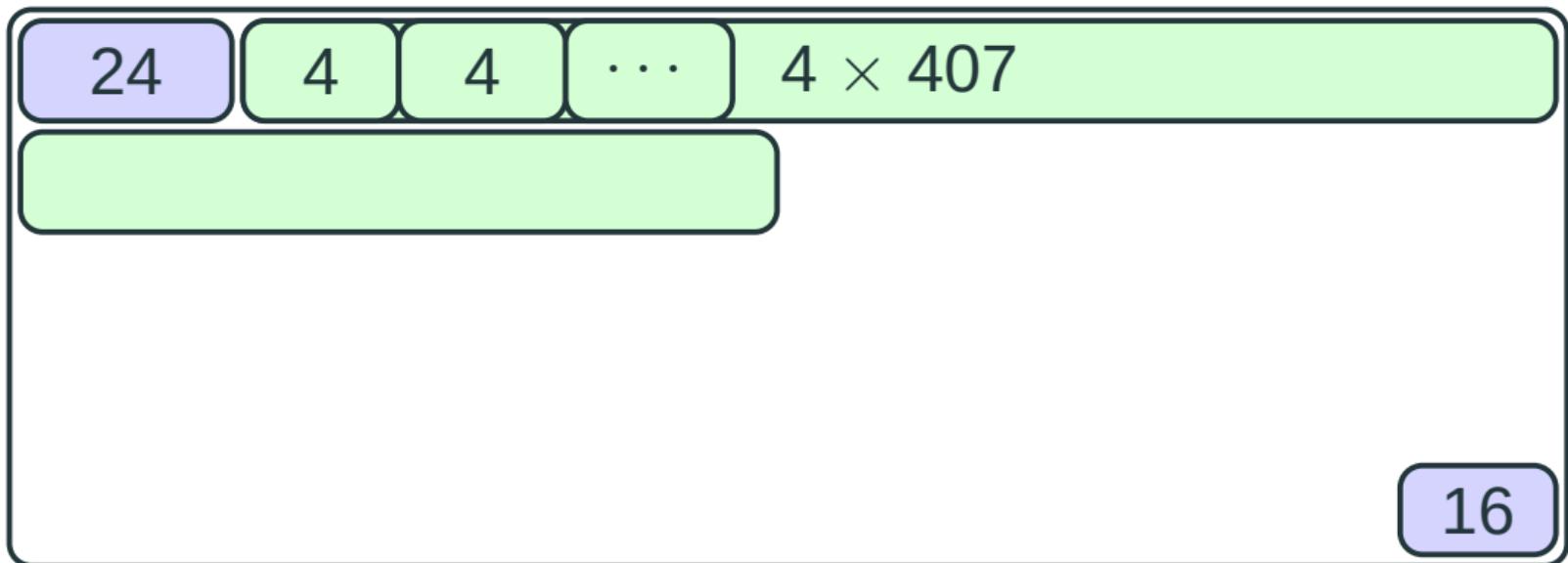


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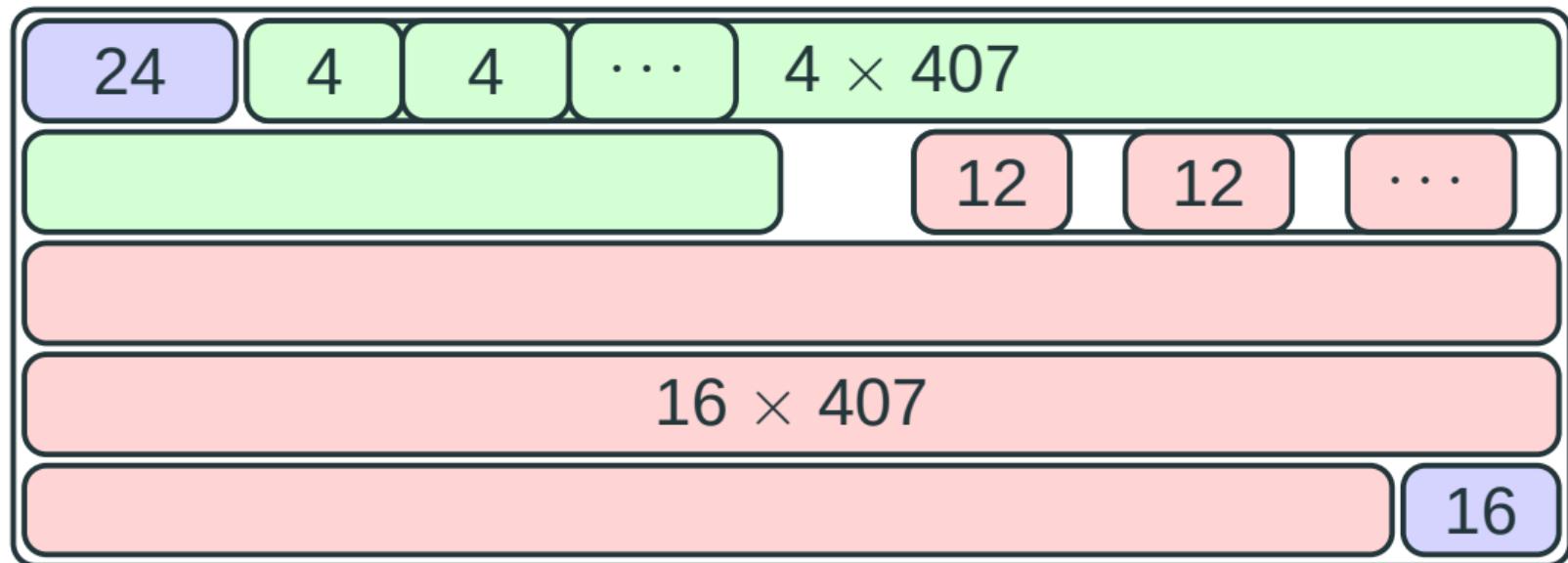
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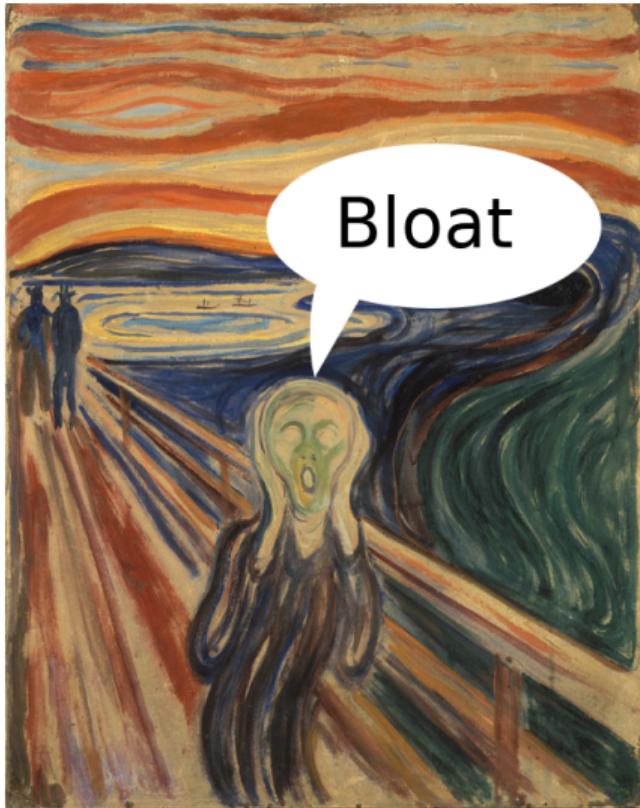
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Edvard Munch – The Scream (about a bloated index)

Assuming we know the workload,
how to approximate free space or bloat?

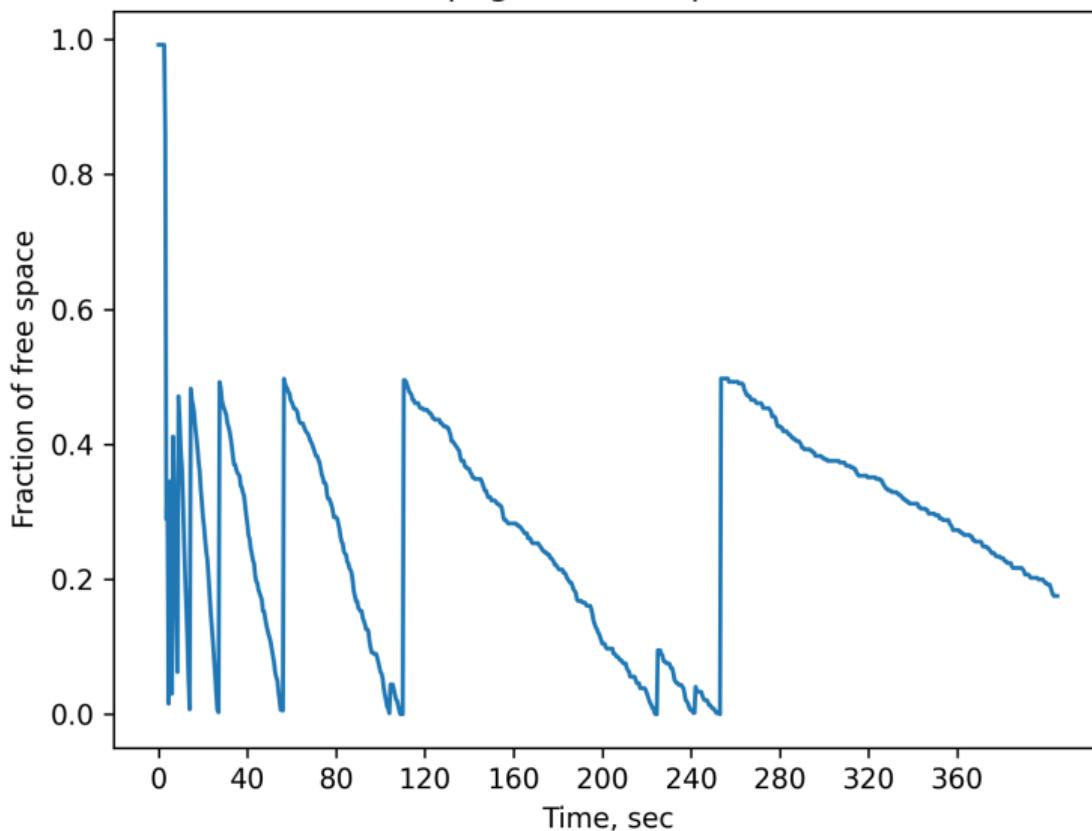
```
-- -M prepared --rate=max-rate
\set aid random(0, N)
\set bid random(0, N)

update test set a = :aid where a = :bid;
```

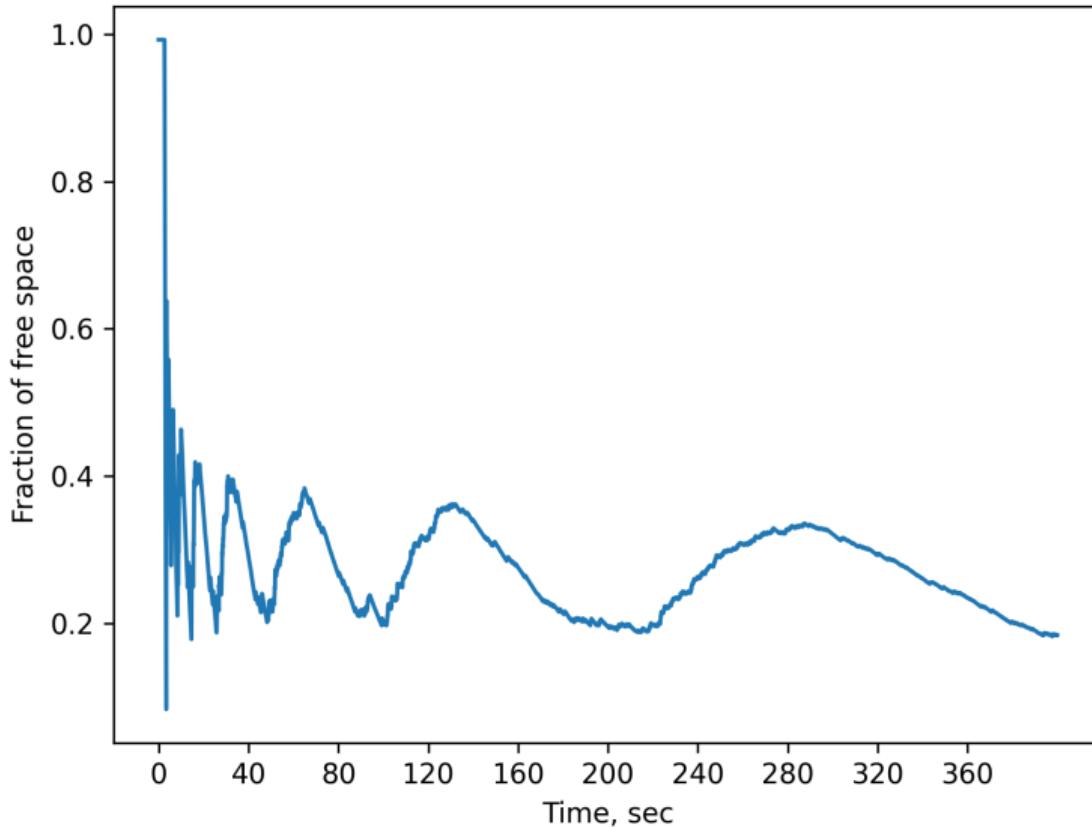


Portrait de l'artiste sous les traits d'un moqueur

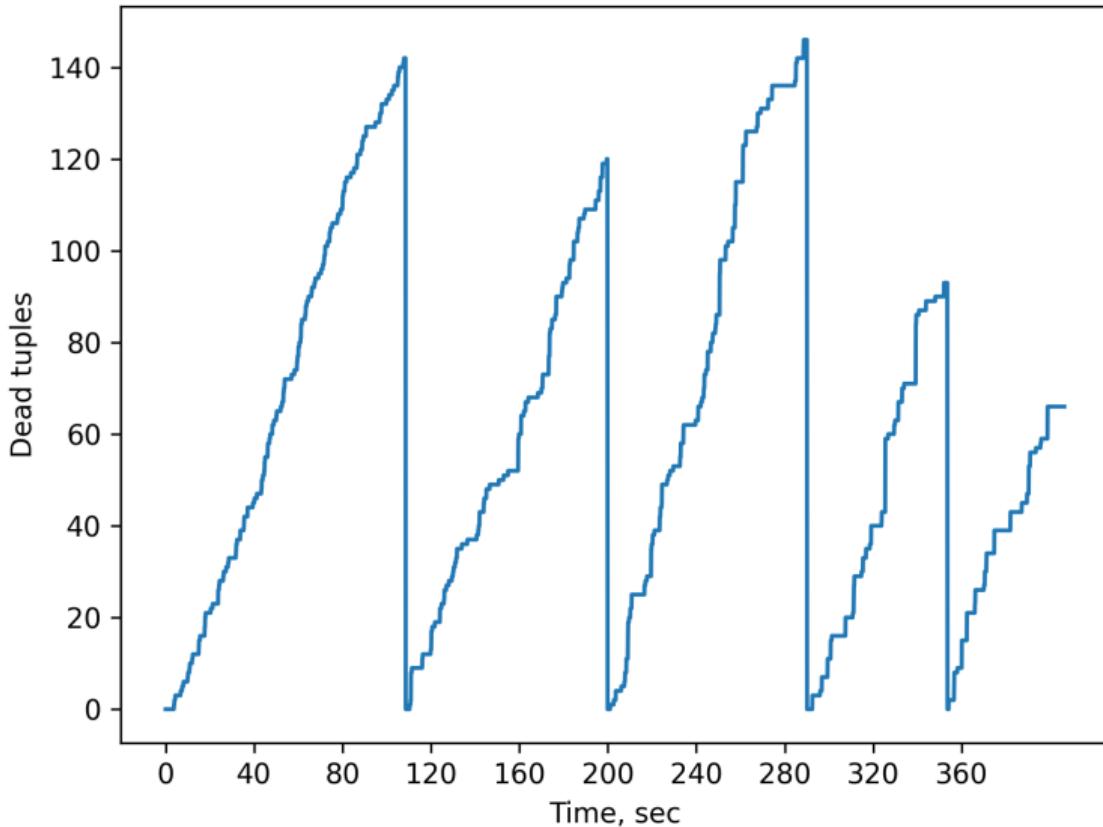
B-Tree page, uniform point insert



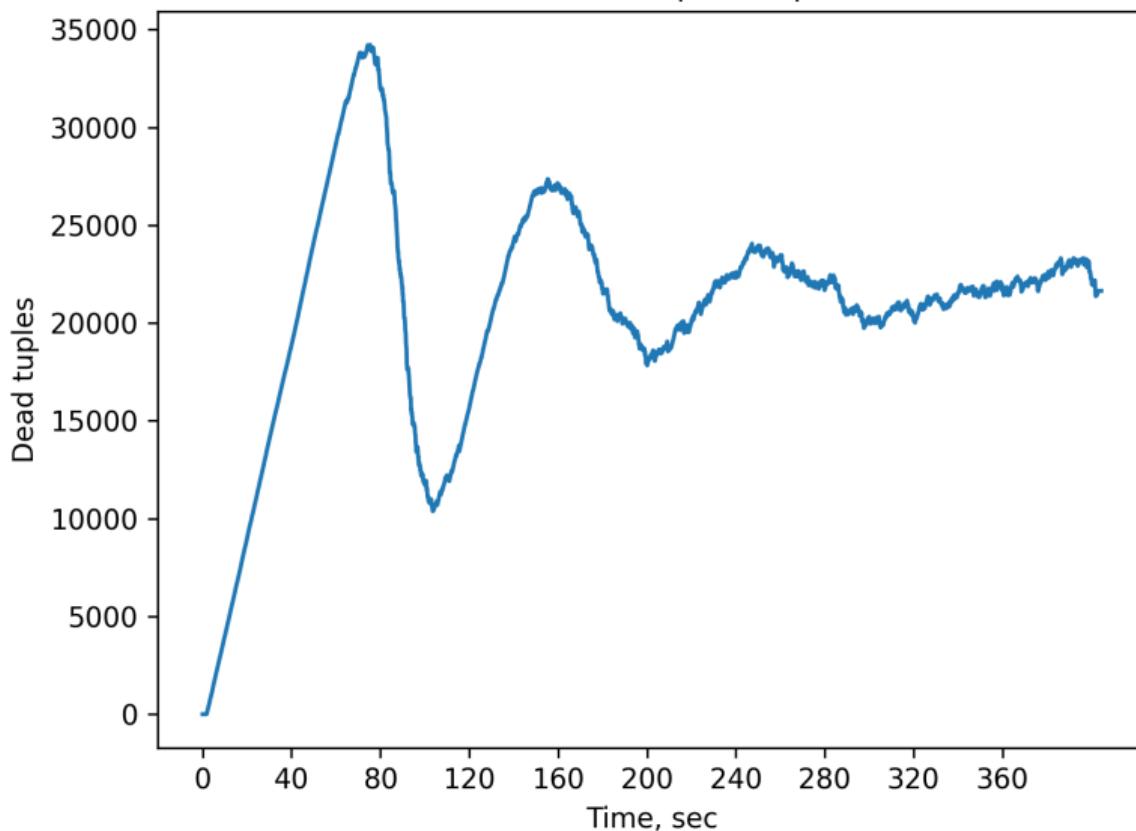
B-Tree, avg free space per page, uniform point insert



B-Tree page, uniform point update



B-Tree, uniform point update



Approximation

```
create table test(a int);
create index on test(a)
    with (fillfactor = 100);
```

```
create unlogged table test(a int);
create index on test(a)
    with (fillfactor = 100);
```

```
create unlogged table test(a int);
create index on test(a)
    with (fillfactor = 100);
```

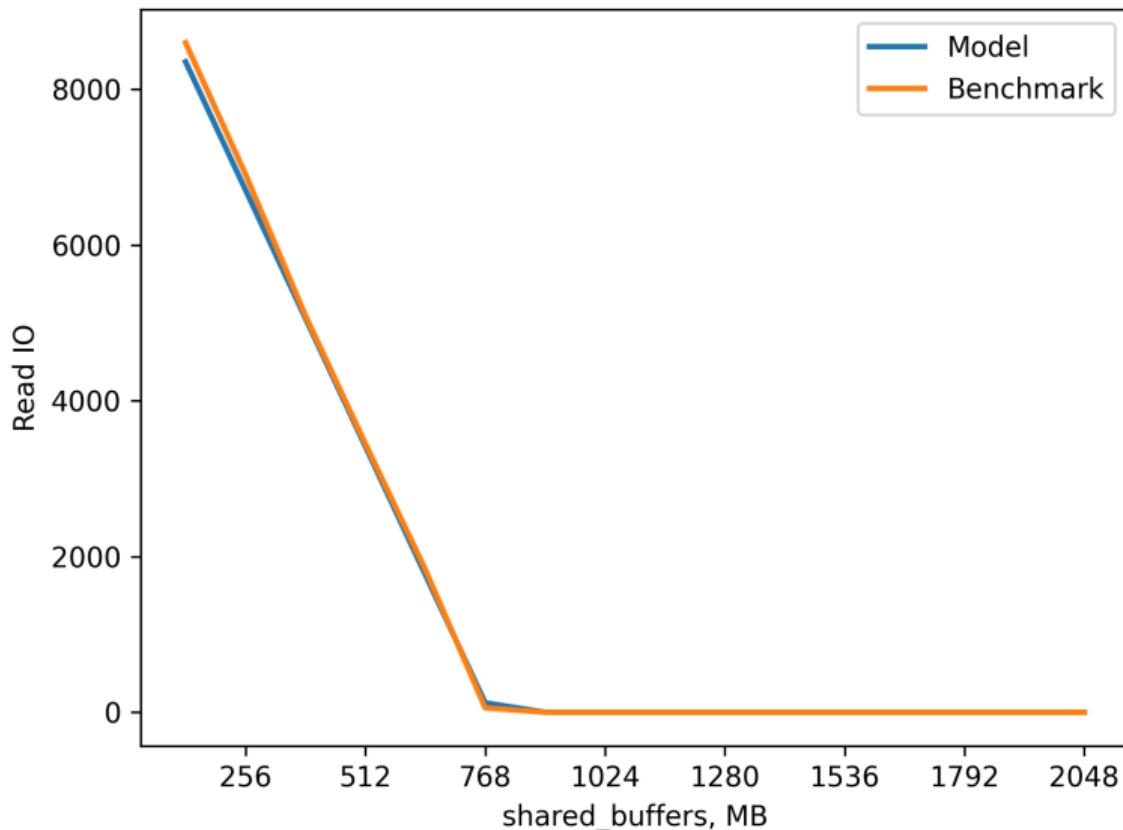
```
# fsync = off
# autovacuum = off
# etc
```

Assuming we know the workload,
how to approximate amount of IO?

```
-- -M prepared --rate=max-rate
\set aid random(0, N)

-- a pre-populated table
select * from test where a = :aid;
```

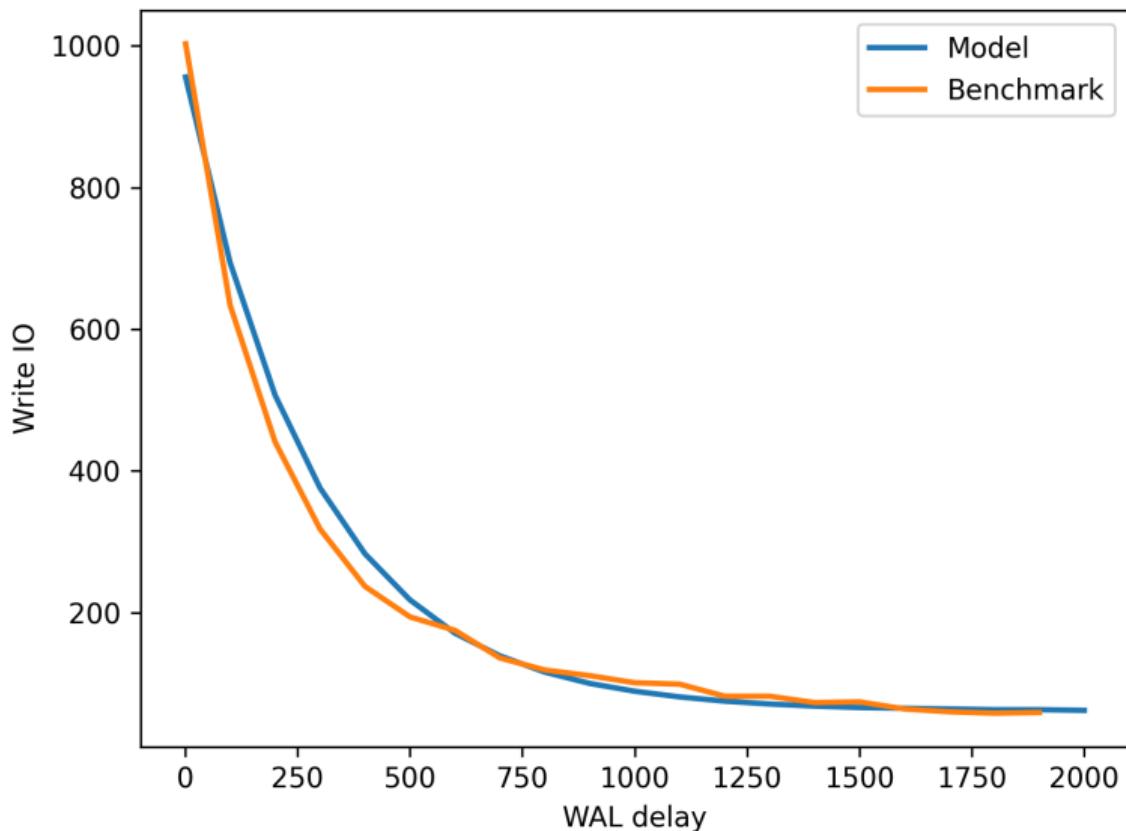
B-Tree, uniform read



```
-- -M prepared --rate=max-rate
\set aid random(0, N)

-- an empty table
insert into test values(:aid);
```

B-Tree, uniform insert



$$IO = Q \cdot M_{press} \sum_{l=1}^L \frac{N_l}{N} + Q_w \cdot W_g + \frac{Q_w}{I}$$

$$IO = Q \cdot M_{press} \sum_{l=1}^L \frac{N_l}{N} + Q_w \cdot W_g + \frac{Q_w}{I}$$

Read Split Insert

The diagram illustrates the components of IO as defined by the equation. It features three colored boxes: a pink box for $Q \cdot M_{press} \sum_{l=1}^L \frac{N_l}{N}$, a green box for $Q_w \cdot W_g$, and a purple box for $\frac{Q_w}{I}$. Above the pink box is the word "Read" with a red arrow pointing down to it. Above the purple box is the word "Split" with a blue arrow pointing down to it. Above the green box is the word "Insert" with a green arrow pointing up to it.

Modeling the Linux page cache for accurate simulation of data-intensive applications

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arXiv:2101.01335v1 [cs.DC] 5 Jan 2021

Abstract—The emergence of Big Data in recent years has resulted in a growing need for efficient data processing solutions. While infrastructures with sufficient compute power are available, the I/O bottleneck remains. The Linux page cache is an efficient approach to reduce I/O overheads, but few experimental studies of its interactions with Big Data applications exist, partly due to limitations of real-world experiments. Simulation is a popular approach to address these issues, however, existing simulation frameworks do not simulate page caching fully, or even at all. As a result, simulation-based performance studies of data-intensive applications lead to inaccurate results.

In this paper, we propose an I/O simulation model that includes the key features of the Linux page cache. We have implemented this model as part of the WRENCH workflow simulation framework, which itself builds on the popular SimGrid distributed systems simulation framework. Our model and its implementation enable the simulation of both single-threaded and multithreaded applications, and of both writeback and writethrough caches for local or network-based filesystems. We evaluate the accuracy of our model in different conditions, including sequential and concurrent applications, as well as local and remote I/Os. We find that our page cache model reduces the simulation error by up to an order of magnitude when compared to state-of-the-art, cacheless simulations.

I. INTRODUCTION

The Linux page cache plays an important role in reducing filesystem data transfer times. With the page cache, previously read data can be re-read directly from memory, and written data can be written directly to memory, bypassing the disk.

type of hardware/software stacks are best suited to different application classes, as well as understanding the limitations of current algorithms, designs and technologies. Unfortunately, performance studies relying on real-world experiments on compute platforms face several difficulties (high operational costs, labor-intensive experimental setups, shared platforms with dynamic loads that hinder reproducibility of results) and shortcomings (experiments are limited to the available platform/software configurations, which precludes the exploration of hypothetical scenarios). Simulations address these concerns by providing models and abstractions for the performance of computer hardware, such as CPU, network and storage. As a result, simulations provide a cost-effective, fast, easy and reproducible way to evaluate application performance on arbitrary platform configurations. It thus comes as no surprise that a large number of simulation frameworks have been developed and used for research and development [1], [2], [3], [4], [5], [6], [7], [8], [9], [10], [11], [12], [13].

Page caching is an ubiquitous technique for mitigating the I/O bottleneck. As such, it is necessary to model it when simulating data-intensive applications. While existing simulation frameworks of parallel and distributed computing systems capture many relevant features of hardware/software stacks, they lack the ability to simulate page cache with enough details to capture key features such as dirty data and cache eviction policies [5], [6]. Some simulators, such as the one

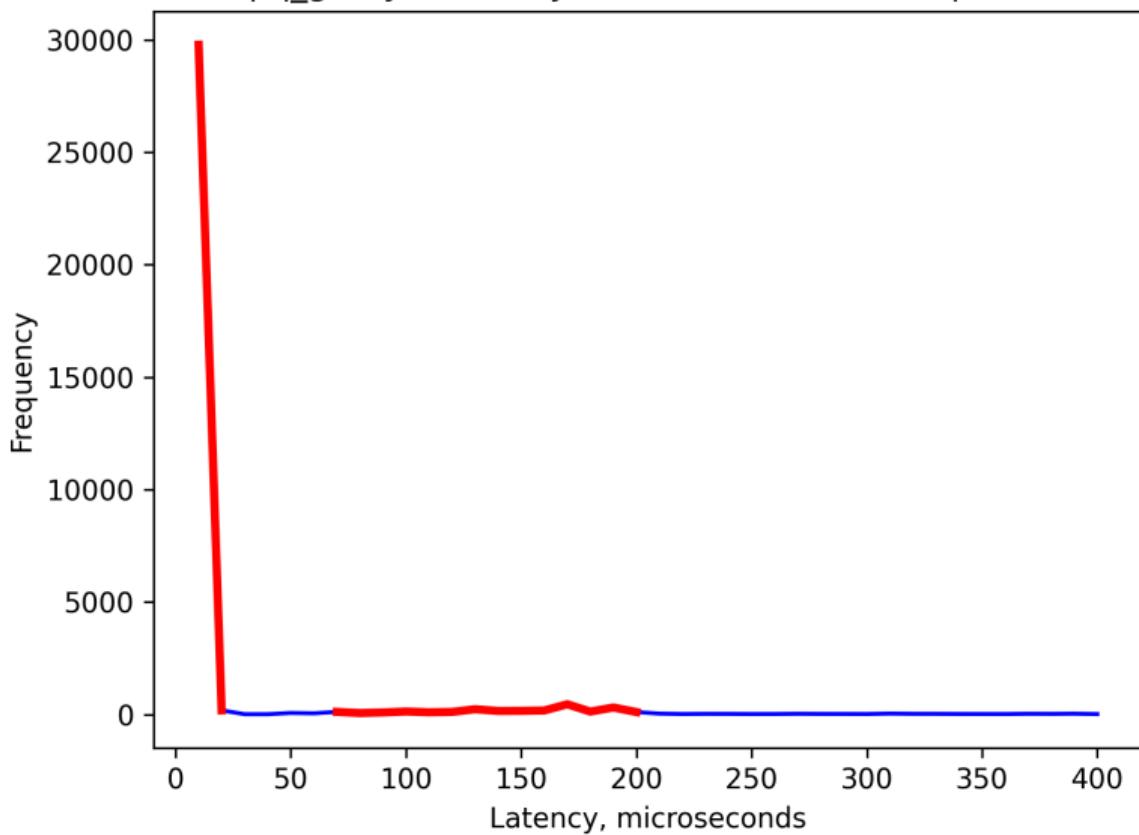


Simulation

Assuming we know the workload,
how to approximate query latency?

```
data Event = GetCachedPlan TxLatency
            | PgGetByte TxLatency
            | BtGetTuple TxLatency
            | BtInsert TxLatency
            | HeapPagePrune TxLatency
            | HeapUpdate TxLatency
            | CommitTx TxLatency
            | SocketFlush TxLatency
```

pq_getbyte latency distribution, uniform update



Questions?

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