



# CS 540 Introduction to Artificial Intelligence Games I

University of Wisconsin-Madison  
**Spring 2023**

# Outline

## Homeworks:

- Homework 9 due Thursday April 27
- Homework 10 due Thursday May 4

## Class roadmap:

Tuesday, April 18	Games I
Thursday, April 20	Games II
Tuesday, April 25	Reinforcement Learning I
Thursday, April 27	Reinforcement Learning I
Tuesday, May 2	Review of RL + Games
Thursday, May 4	Ethics and Trust in AI

# Outline

- Introduction to game theory
  - Properties of games, mathematical formulation
- Simultaneous-Move Games
  - Normal form, strategies, dominance, Nash equilibrium

# So Far in The Course

We looked at techniques:

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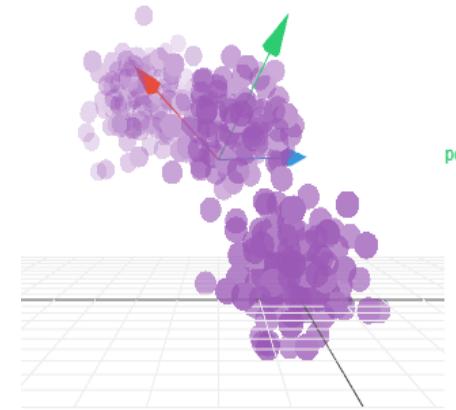
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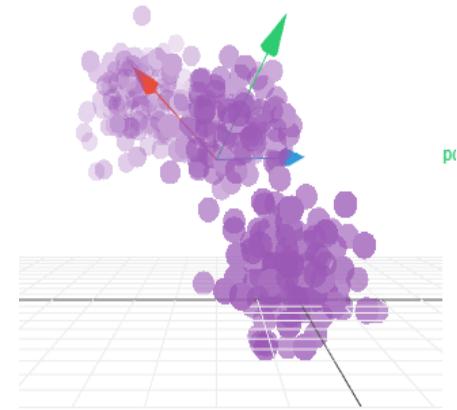


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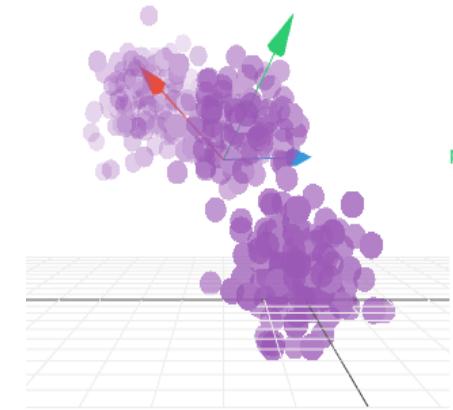


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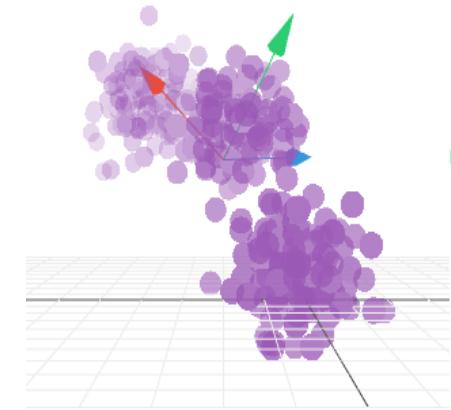


outdoor

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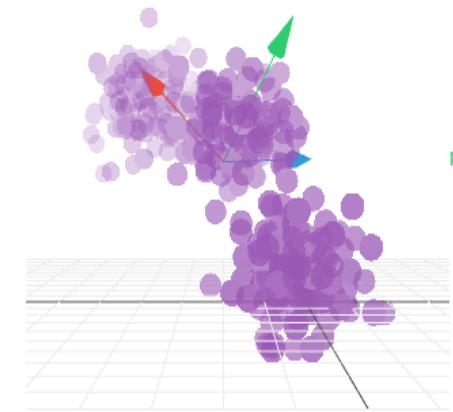


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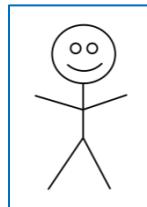


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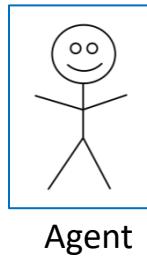
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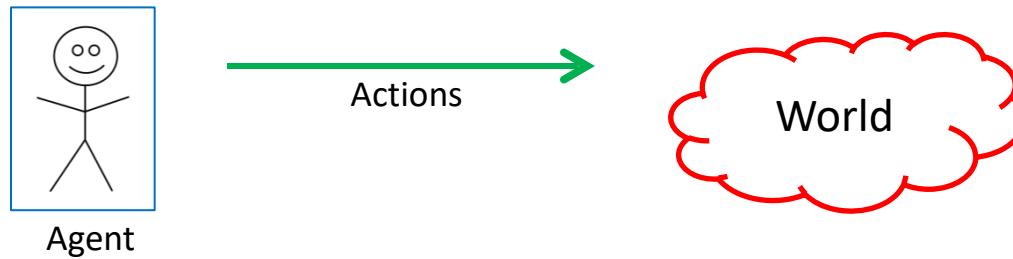
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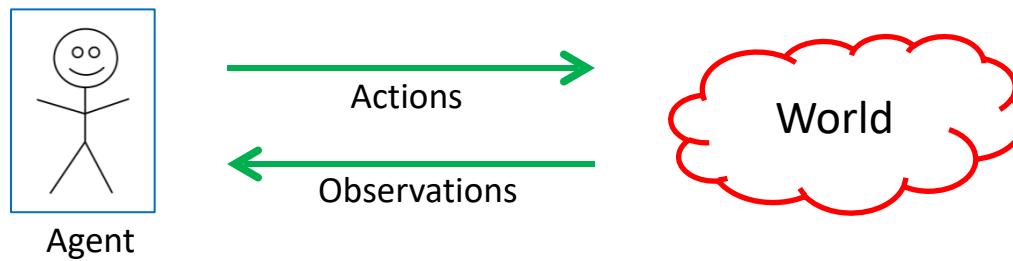
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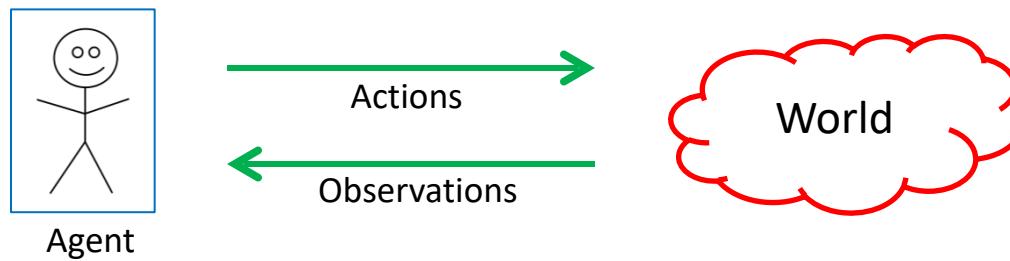
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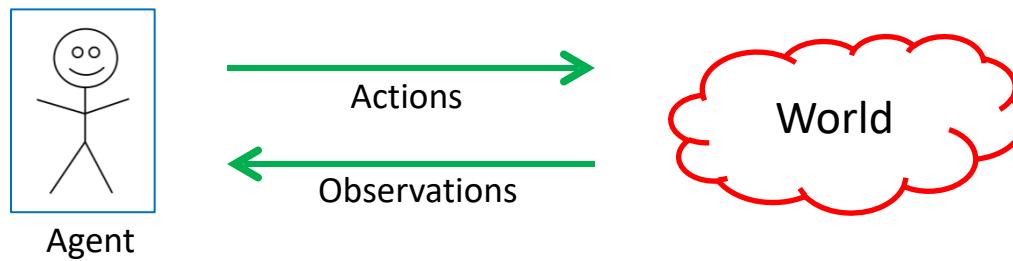
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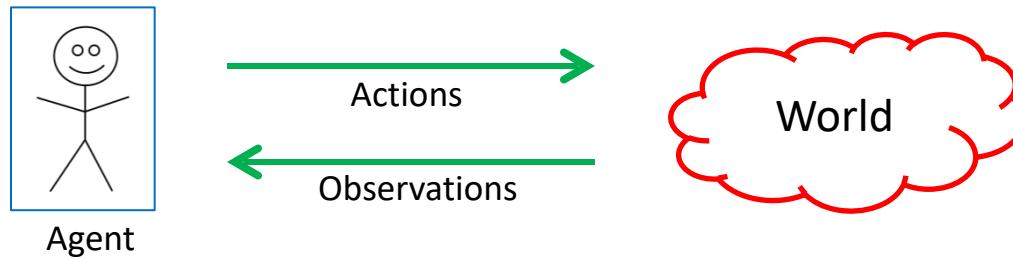
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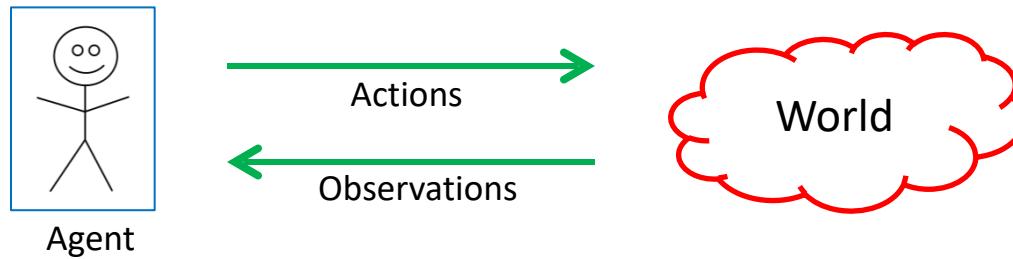
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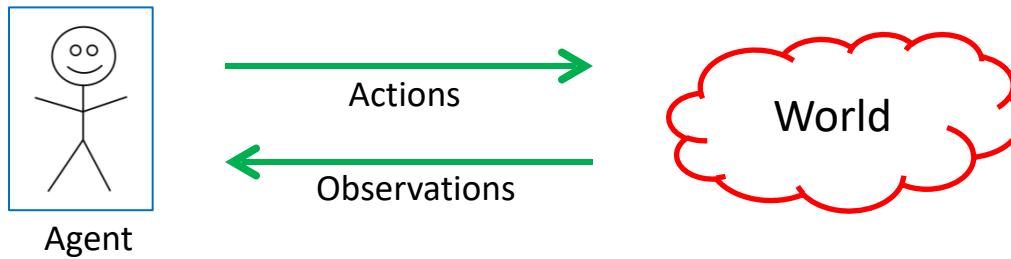
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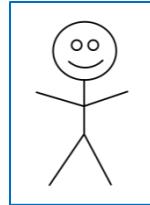
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  - Setup for decision theory, reinforcement learning, planning

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Games setup: **multiple** agents

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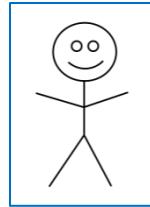
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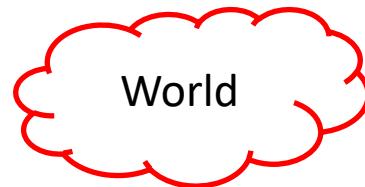
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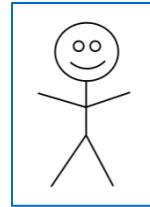


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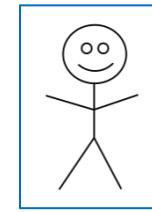
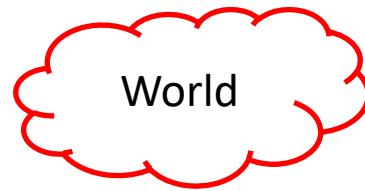


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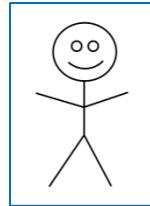
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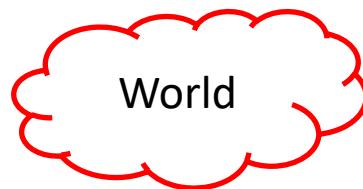
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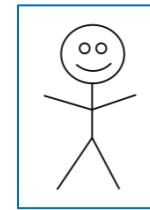
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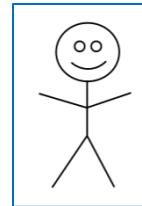
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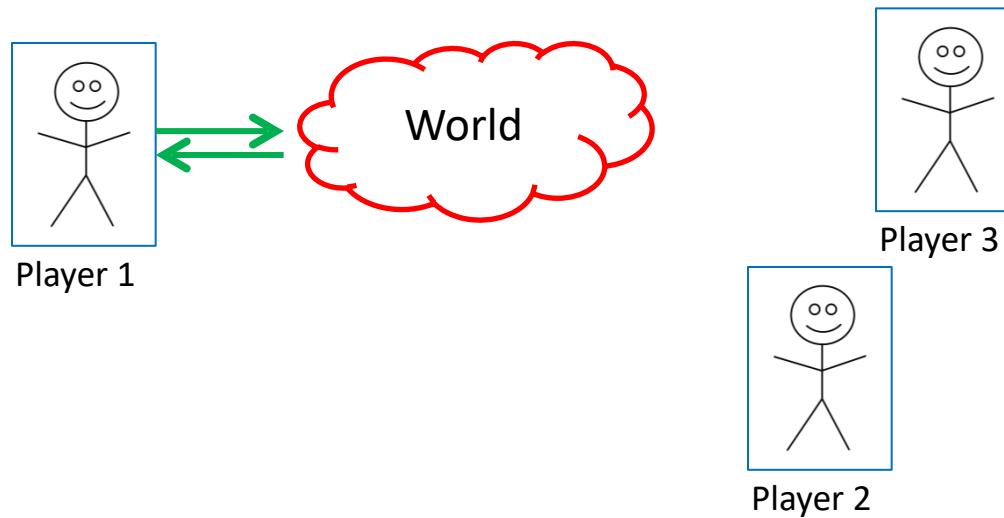
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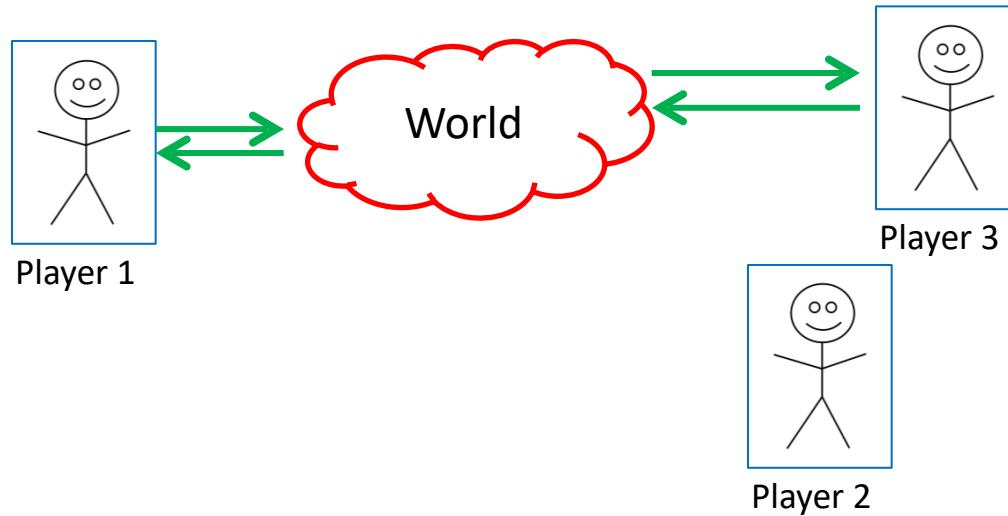
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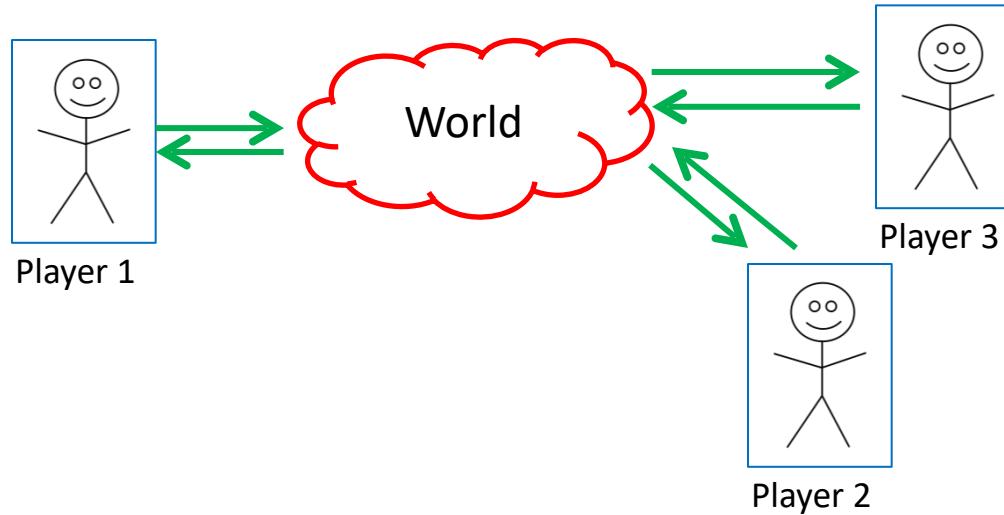
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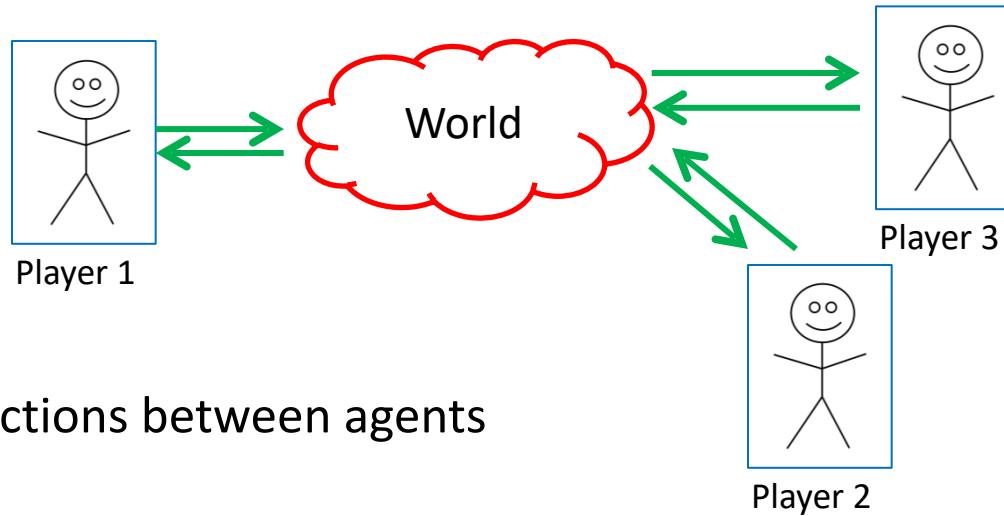
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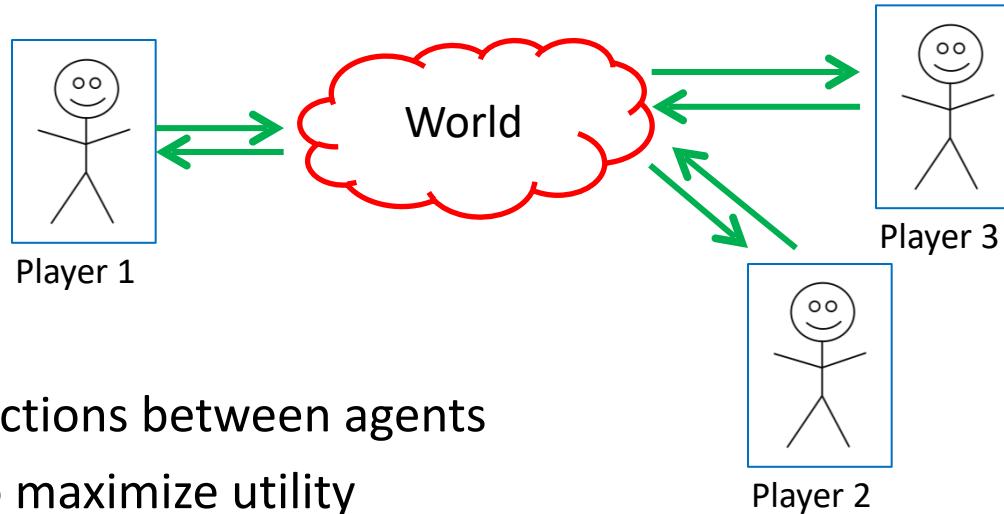
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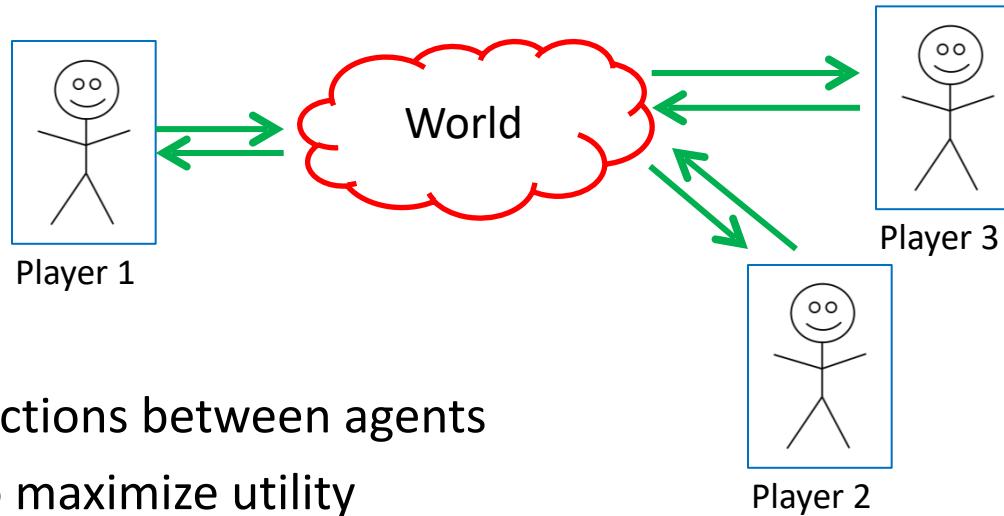
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Wiki

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# Property 2: Action Space

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Action space: set of possible actions an agent can choose from.

Can be finite or infinite.

Examples:

- Rock-paper-scissors
- Tennis

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  - Example: driving to work, prisoner's dilemma

## **Property 5: Sequential or Simultaneous Moves**

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- Sequential: take turns (but payoff only revealed at end of game)

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- We consider the simple case where all reward functions are common knowledge.

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- Sometimes a dominant strategy does not exist!

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<i>T</i>		2, 1	0, 0
<i>B</i>		0, 0	1, 2

# Nash Equilibrium

$a^*$  is a Nash equilibrium if no player has an incentive to **unilaterally deviate**

$$u_i(a_i^*, a_{-i}^*) \geq u_i(a_i, a_{-i}^*) \quad \forall a_i \in A_i$$

		Player 2		
Player 1			<i>L</i>	<i>R</i>
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- Pure Nash equilibrium:
  - A **pure strategy** is a deterministic choice (no randomness).
  - Later: we will consider **mixed** strategies
  - In pure Nash equilibrium, players can only play pure strategies.

# Finding (pure) Nash Equilibria by hand

- As player 1: For each column, find the best response, underscore it.

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# Finding (pure) Nash Equilibria by hand

- As player 2: for each row, find the best response, upper-score it.

Player 2		<i>L</i>	<i>R</i>
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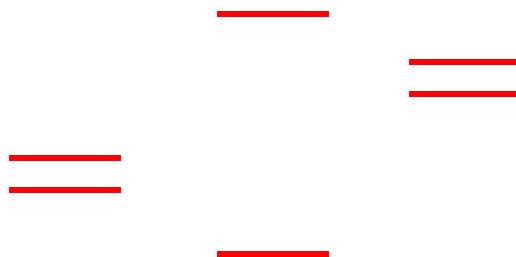
# Finding (pure) Nash Equilibria by hand

- Entries with both lower and upper bars are pure NEs.

		Player 2	
		<i>L</i>	<i>R</i>
Player 1			
<i>T</i>		2, 1 _____	0, 0
<i>B</i>		0, 0	1, 2 _____

# Pure Nash Equilibrium may not exist

So far, pure strategy: each player picks a deterministic strategy. But:



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		Player 2	rock	paper	scissors
		Player 1			
rock	0, 0	-1, 1	1, -1		
paper	1, -1	0, 0	-1, 1		
scissors	-1, 1	1, -1	0, 0		

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- Now consider **expected rewards**

$$u_i(x_i, x_{-i}) = E_{a_i \sim x_i, a_{-i} \sim x_{-i}} u_i(a_i, a_{-i}) = \sum_{a_i} \sum_{a_{-i}} x_i(a_i) x_{-i}(a_{-i}) u_i(a_i, a_{-i})$$

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- Intuition: nobody can **increase expected reward** by changing only their own strategy.

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Example:  $x_1(\cdot) = x_2(\cdot) = (\frac{1}{3}, \frac{1}{3}, \frac{1}{3})$

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Player 2		<i>rock</i>	<i>paper</i>	<i>scissors</i>
Player 1	<i>rock</i>	0, 0	-1, 1	1, -1
<i>paper</i>	1, -1	0, 0	-1, 1	
<i>scissors</i>	-1, 1	1, -1	0, 0	

# Finding Mixed NE in 2-Player Zero-Sum Game

Example: Two Finger Morra. Show 1 or 2 fingers. The “even player” wins the sum if the sum is even, and vice versa.

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	odd	
even	$f_1$	$f_2$
$f_1$	2, -2	-3, 3
$f_2$	-3, 3	4, -4

# Finding Mixed NE in 2-Player 2-action Zero-Sum Game

Two Finger Morra. Two-player zero-sum game. No pure NE:

	odd		
	<i>f1</i>		<i>f2</i>
even			
<i>f1</i>	2, -2 <u>  </u>	-3, 3 <u>  </u>	
<i>f2</i>	-3, 3 <u>  </u>	4, -4 <u>  </u>	

# Finding Mixed NE in 2-Player 2-action Zero-Sum Game

		q	1-q
		odd	
		even	
p	$f_1$	$2, -2$	<u>-3, 3</u>
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Suppose odd's mixed strategy at NE is  $(q, 1-q)$ , and even's  $(p, 1-p)$

		$q$	$1-q$
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even		$f_1$	$f_2$
$p$			
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Average is no greater than components

$$\rightarrow u_1(p, q) = u_1(f_1, q) = u_1(f_2, q)$$

We want to find  $q$  such that equality holds.

Then even has no incentive to change strategy.

		q	1-q
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even		$f_1$	$f_2$
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$$u_1(f_1, q) = u_1(f_2, q)$$

$$2q + (-3)(1 - q) = (-3)q + 4(1 - q)$$

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$$u_1(f_1, q) = u_1(f_2, q)$$

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$$q = \frac{7}{12}$$

		q	1-q
		odd	
		even	
p	f1	2, -2	-3, 3
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1-p			

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At this NE, even gets  $-1/12$ , odd gets  $1/12$ .

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- Every **finite** (players, actions) game has at least one Nash equilibrium
  - But not necessarily **pure** (i.e., deterministic strategy)
- Could be more than one
- Searching for Nash equilibria: computationally **hard**.
  - Exception: two-player zero-sum games (can be found with linear programming).

# Break & Quiz

**Q 2.1:** Which of the following is false?

- (i) Rock/paper/scissors has a dominant pure strategy
  - (ii) There is a Nash equilibrium for rock/paper/scissors
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- A. Neither
  - B. (i) but not (ii)
  - C. (ii) but not (i)
  - D. Both

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  - D. Both (There is a mixed strategy Nash Eq.)

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  - (ii) Nash equilibria require rational play
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- A. Neither (See below)
  - B. (i) but not (ii) (Rational play required: i.e., what if prisoners desire longer jail sentences?)
  - C. (ii) but not (i) (The basic assumption of Nash equilibria is knowing all of the strategies involved)
  - D. Both

# Summary

- Intro to game theory
  - Characterize games by various properties
- Mathematical formulation for simultaneous games
  - Normal form, dominance, Nash equilibria, mixed vs pure