# **HPC Concepts**

Queues Schedulers, and workflow design for shared systems.

#### Overview

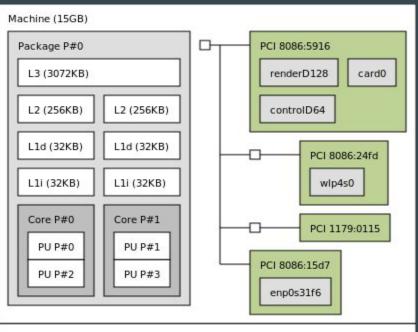
- Types
  - Of Computing systems
  - o Of Jobs
- Limitations
  - Compute
  - o I/O
- Cluster components
- Seperation
- Queuing
  - Queue Partitioning
  - o Queue Filling
- Workflow organization as solution
- Examples

## What is a HPC?

# What is a HPC? a device for turning compute-bound problems into I/O-bound problems.

Ken Batcher

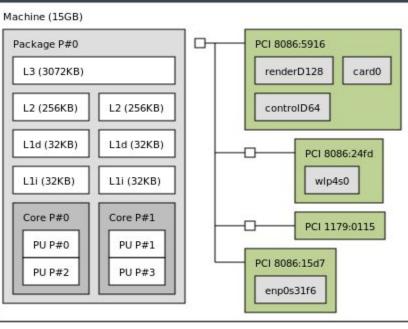
## More powerful CPU's, More Chips Per board

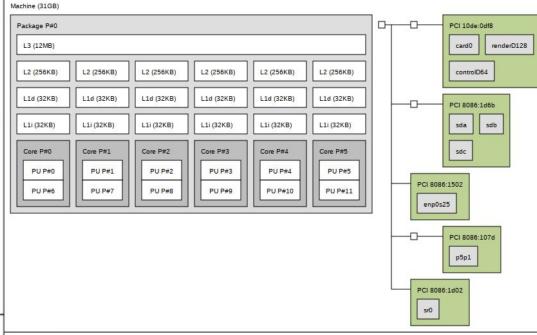


Host: prospero Indexes: physical

Date: Mon 22 Oct 2018 11:03:41 AM EDT

## More powerful CPU's, More Chips Per board





Host: prospero

Indexes: physical

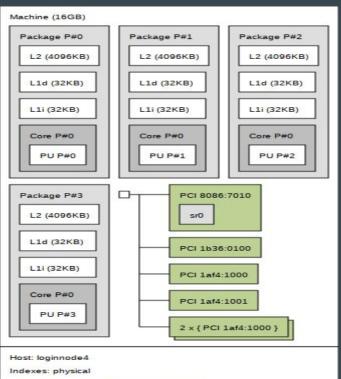
Date: Mon 22 Oct 2018 11:03:41 AM EDT

Host: oberon Indexes: physical

Date: Mon 22 Oct 2018 11:00:45 AM EDT

## More powerful CPU's, More Chips Per board - Reorganized

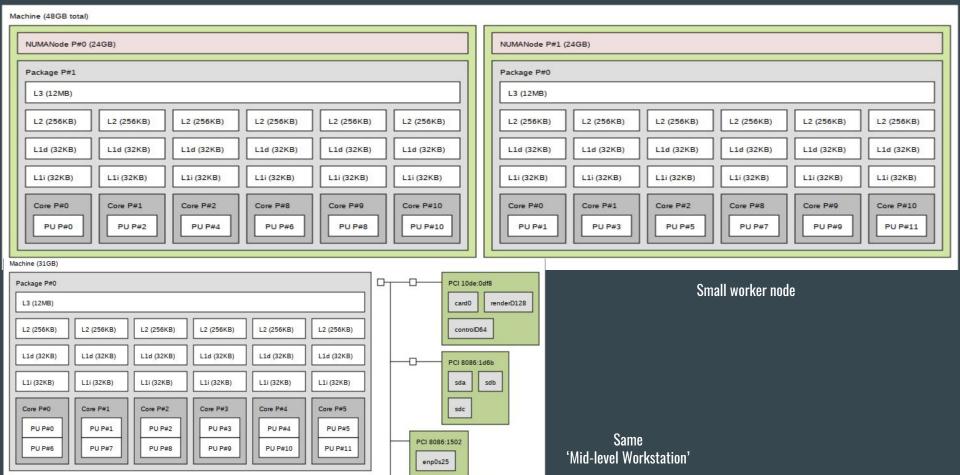
Machine (31GB)



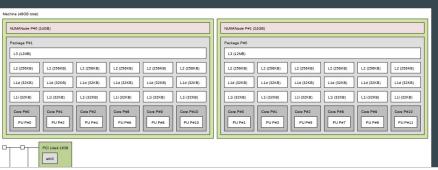
Package P#0 PCI 10de:0df8 L3 (12MB) card0 renderD128 L2 (256KB) L2 (256KB) L2 (256KB) L2 (256KB) L2 (256KB) L2 (256KB) controlD64 L1d (32KB) L1d (32KB) L1d (32KB) L1d (32KB) L1d (32KB) L1d (32KB) PCI 8086:1d6b L1i (32KB) L1i (32KB) L1i (32KB) L1i (32KB) L1i (32KB) L1i (32KB) sda sdb Core P#0 Core P#1 Core P#2 Core P#3 Core P#4 Core P#5 sdc PU P#0 PU P#1 PU P#2 PU P#3 PU P#4 PU P#5 PCI 8086-1502 PU P#6 PU P#7 PU P#8 PU P#9 PU P#10 PU P#11 enp0s25 PCI 8086:107d p5p1 PCI 8086:1d02 sr0 Host: oberon Date: Mon 22 Oct 2018 11:00:45 AM EDT

Date: Mon 22 Oct 2018 11:10:50 AM EDT

## More powerful CPU's, More Chips Per board - Reorganized



## More powerful CPU's, More Chips Per board - Reorganized



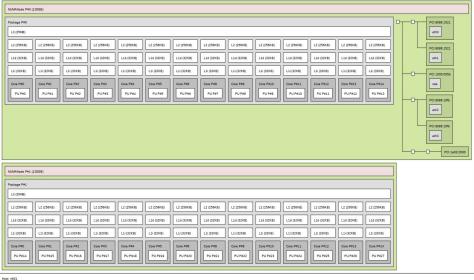
Small worker node 2x 6 cores 2x 24Gb RAM

Mid-size worker node 2x 12 cores 2x 128 Gb RAM

Marking (256GR seral)

Big-compute worker node 2x 14 cores 2x 128 Gb RAM



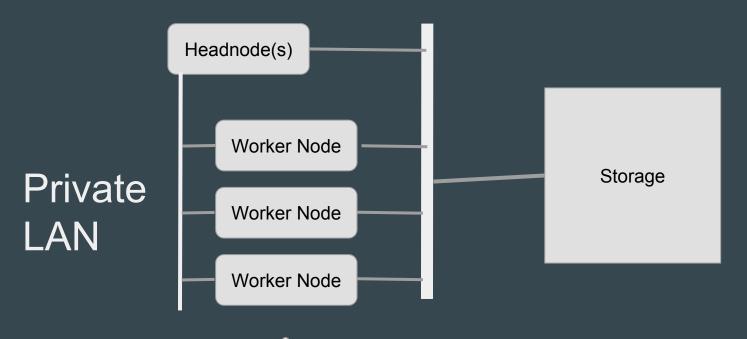


Indexes: physical

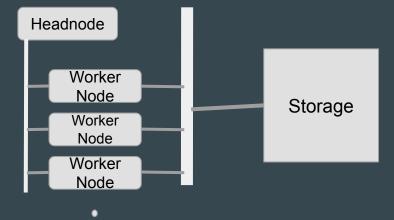
Date: Mon 22 Oct 2018 11:10:50 AM EDT

Indexes: physical
Date: Mon 22 Oct 2018 11:10:50 AM EDT

## **More Interconnects**

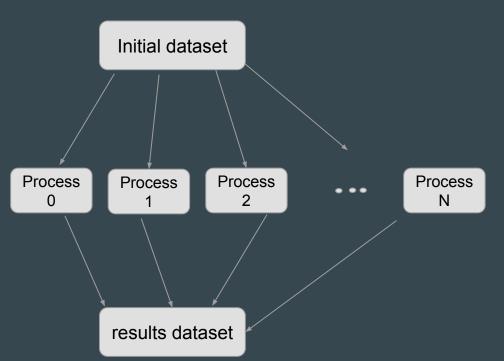


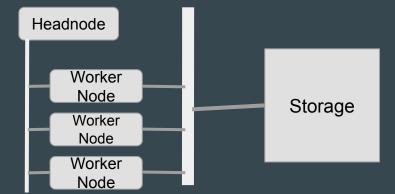
## **More Interconnects**



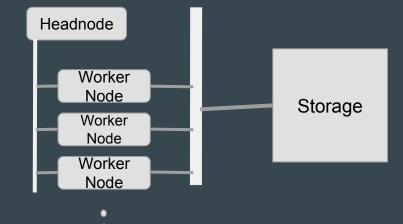
More Interconnects

"Pleasingly Parallel" Minimal complexity





Types of HPC **More Interconnects** Loosely coupled workflows Initial dataset **Process Process** Process Intermediate dataset **Process** Process 3b 3a Process **Process** 3c

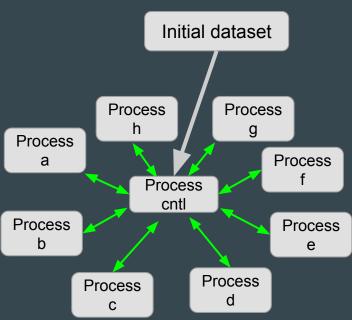


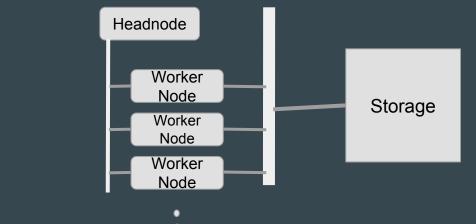
results dataset

Types of HPC

More Interconnects

Loosely coupled workflows





Loosely coupled jobs have dependent processes, but distribute the data according to the task the are

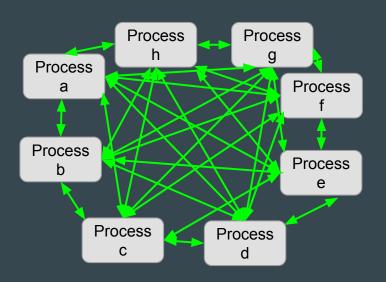
Dependent in terms of task order and

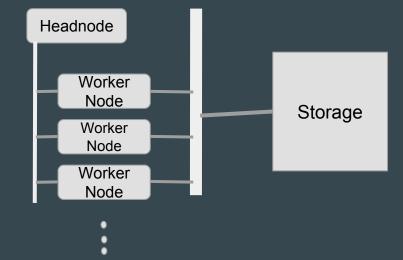
Independent internal to the task and its individual data flow

Special setup can improve speed but is not strictly required

#### **More Interconnects**

Tightly coupled workflows

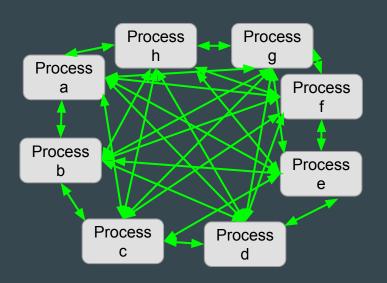


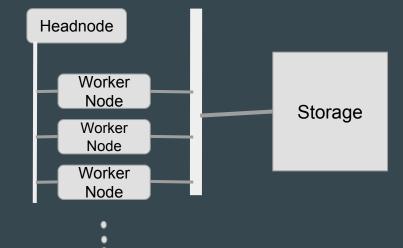


Tightly coupled jobs have interdependent processes, and Shared memory.

#### **More Interconnects**

Tightly coupled workflows





Tightly coupled jobs have interdependent processes, and Shared memory.

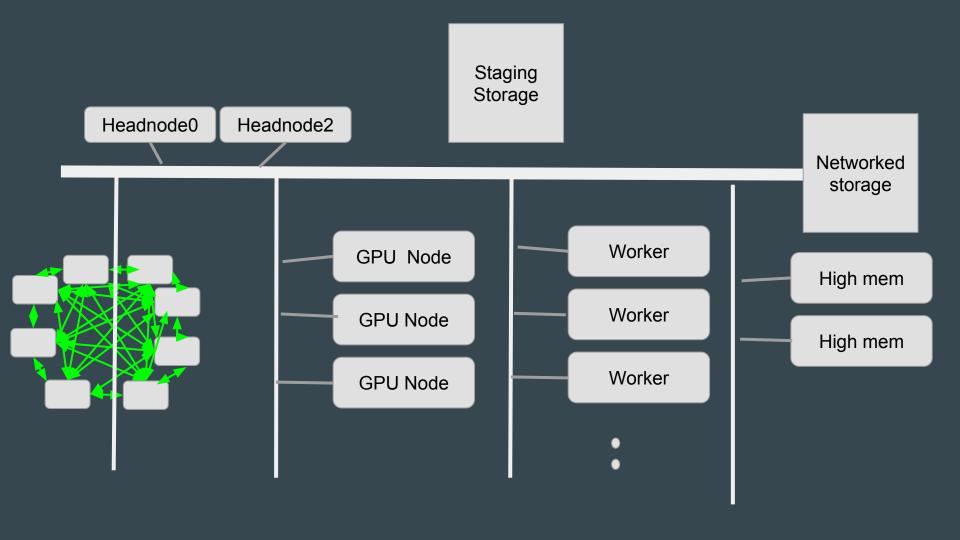
Requires special physical set up and programing (PVM or MPI)

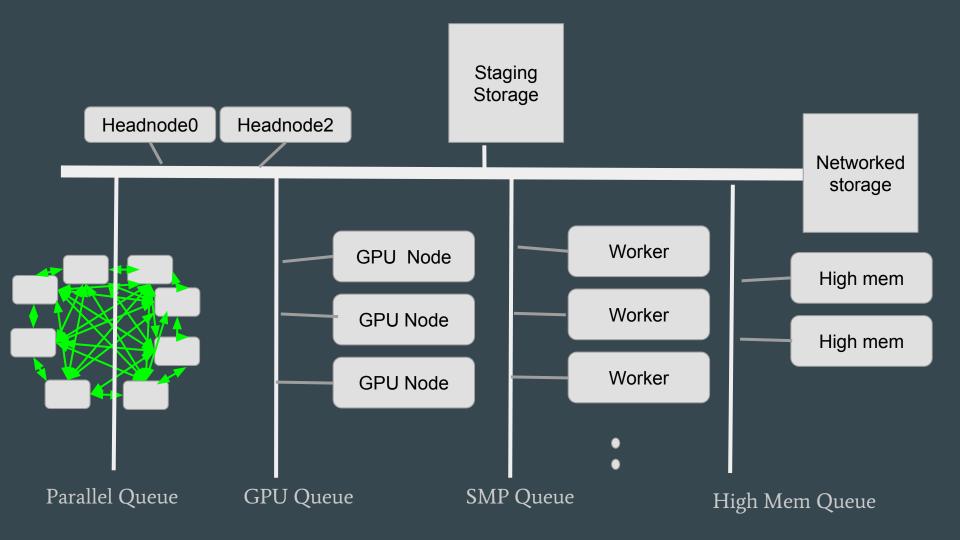
#### **Acceleration**

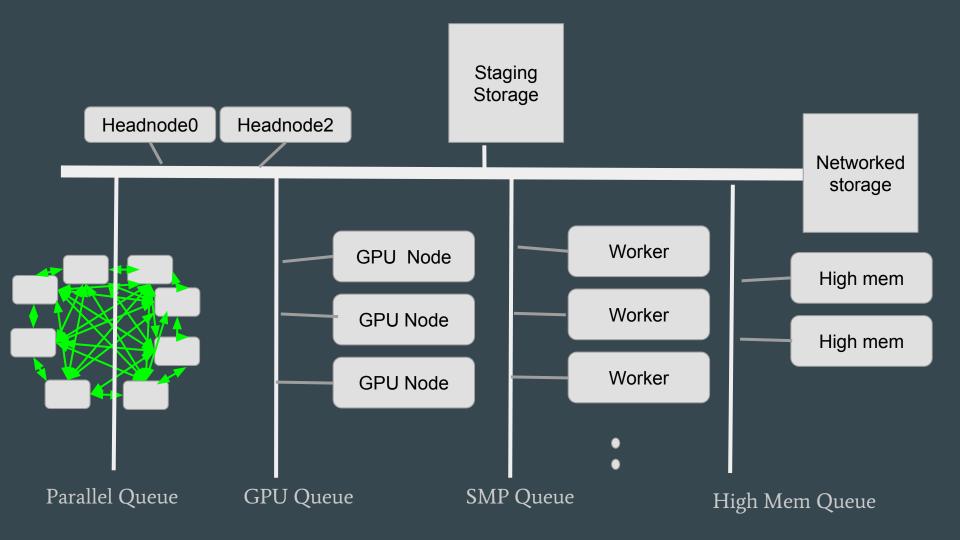
GPU's: cssentially a dense tightly bound system.

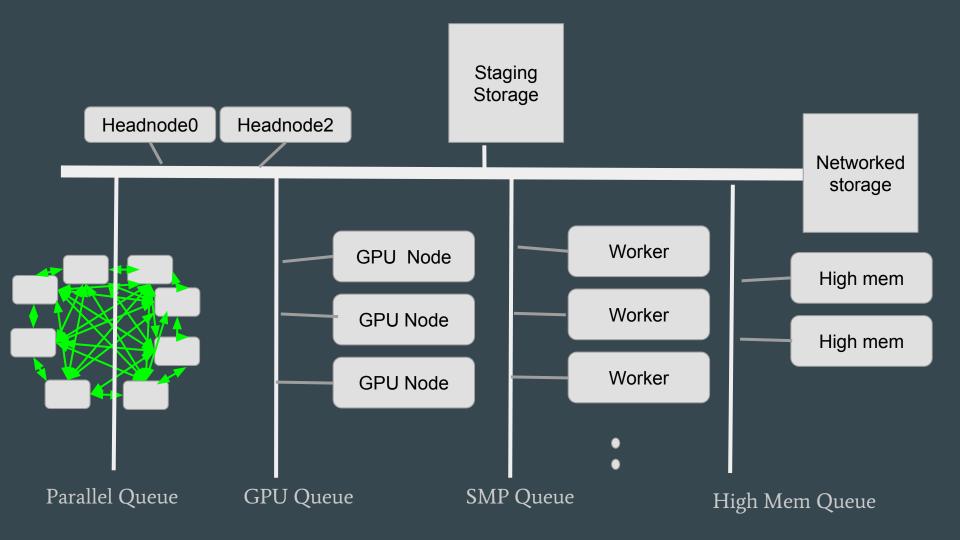
ARM: Single Instruction Multiple Data (SIMD) vector processors

FPGA: Field Programmable Gate arrays Customizing the physical structure of the chip for specific problem sets.









# A mile high View

- Solutions are Organizational
- Matches the hardware with
  - Data flow
  - o Computational patterns.

## Common Queue Scheduler systems

### For distributed work

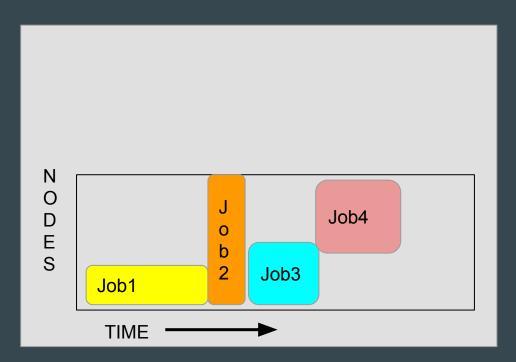
- HTCondor
- Yarn

## For localized systems

- Grid Engine (SGE, UGE)
- PBS/torque
- Slurm

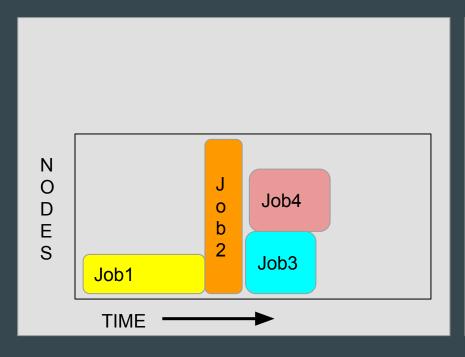
- First in First Out FIFO
- Priorities (Backfill scheduling)
- Preemptions

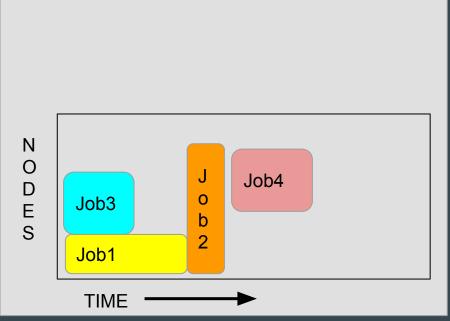
#### First in First Out FIFO



First in First Out FIFO

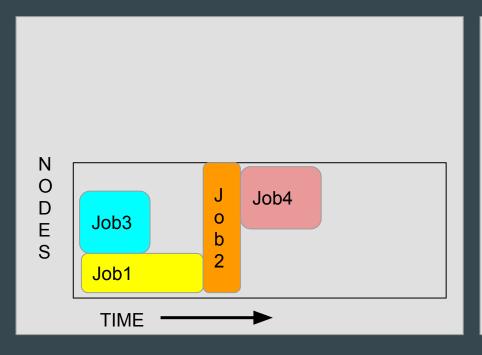
BackFill

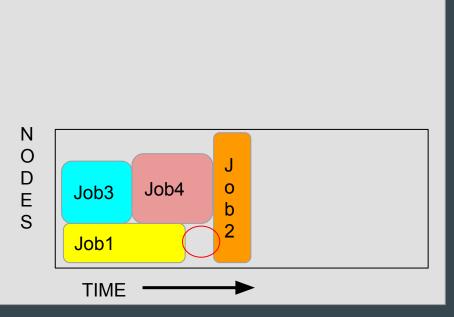




BackFill

Optimizing backfill





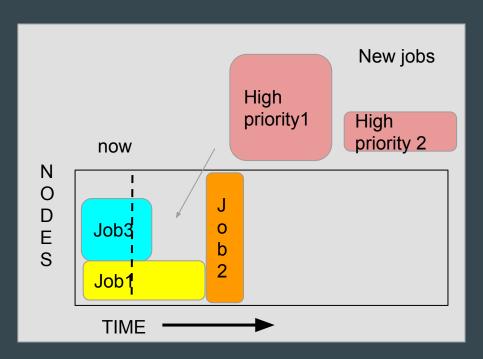
#### Fair share

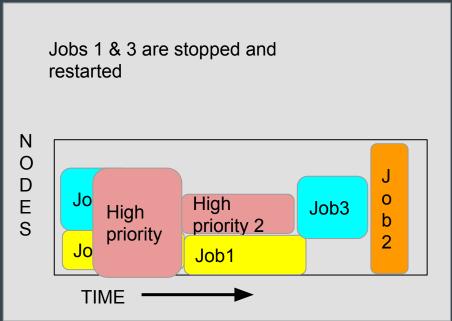
## Adds Priorities to backfill scheduling

Weight several factors to generate a priority level for each job

- How long it has been waiting
- Job Size
- Where the job is being run
- User's base priority level

## Preemption





## What is Checkpointing?

- Saves a representation of the current state at regular intervals
- Allows Automation
  - Scheduler with automatically restarts
  - Add a function at the beginning to check for the save state file
    - If a save state file exists reads it in and starts from there.

```
If statefile exits:
        Startjobat(statefile)
else:
        Startjobat(beginning)
```

## When to use Checkpointing

- Long running jobs
  - To not start over if there's a "hiccup"
    - Balancing between speed and loss
      - Writing to a file is much much slower than anything else.
      - Longer period between checkpoint means faster runs
      - Potential lost time with a restart is bigger.
  - Runtime is longer than maximum job time

# Planning workflows

- Seperate jobs
- Identify dependencies
- Decided what (if any)
  intermediate data to keep
- Match your pattern to available resources and limitations.

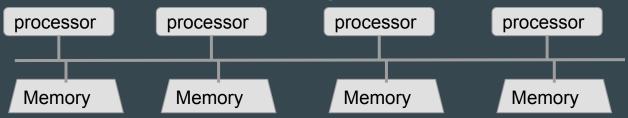
# Turning compute-bound problems into I/O-bound problems.

#### **SMP and NUMA**

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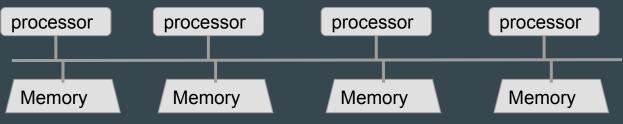
#### Shared Memory Processing *Uniform access*



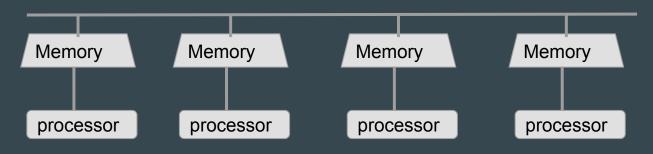
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#### **SMP and NUMA**

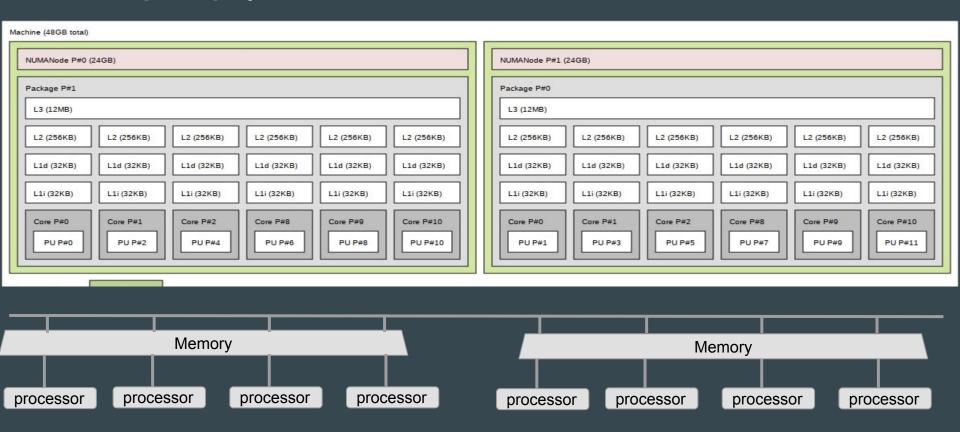
#### Shared Memory *Uniform access*



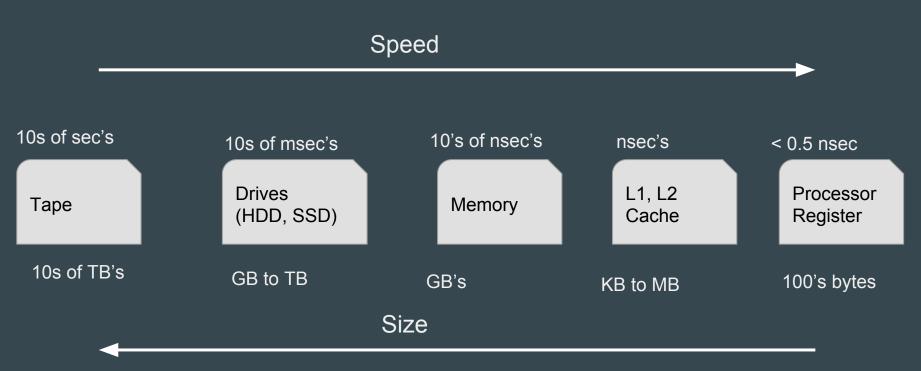
#### **Shared Memory Non-Uniform Memory Access**



# Turning compute-bound problems into I/O-bound problems. Everything is a graph



# Turning compute-bound problems into I/O-bound problems.urning compute-bound problems into I/O-bound problems.



## Distance = 1/Speed

Remote networked storage	Slow - not for computing
Nearby networked/warm storage	Pre-staging Completed data
'Fast scratch' (or /tmp if not available)	Staging, interprocess holding, bigger than RAM for for worker I/o
Tmpfs (memory as filesystem)	For faster I/o

#### Distance = 1/Speed

#### Caution every system is different

Remote networked storage	Slow - not for computing
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'Fast scratch' (or /tmp if not available)	Staging, interprocess holding, bigger than RAM for for worker I/o
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Check with your local systems documentation

Or ask the system folks

## Modules

#### Modules

modules allow us to keep the working space clean, when there are a lot of installed programs.

#### \$module avail

Will list the currently available modules

#### \$module load <modulename>

Will make the module available Be explicit!

## \$module unload <modulename> Will unload the module

## \$module purge Will unload all modules

## \$module whatis <modulename> Will show the module information

# **Examples**

## SGE basics

```
# bigio_example.sh
# jobname
#$ -N example
# use this shell
#$ -S /bin/bash
# how much memory per processor
#$ -1 h_vmem=1g
# maximum run time
#$ -1 h_rt=00:05:00
# parallel environment SMP, 4 processes
#$ -pe smp 4
# email user when done
#$ -m me@myemail.com
runscript "$mytmp"
```

#!/bin/bash

# "Big" I/O

- a. mkdir mytemp @
  - i. /tmp, or /dev/shm
- b. Copy files into and untar
  - i. cp remote/myfiles.tar mytempdir/
  - ii. tar xf myfiles.tar
- c. Run jobscript
  - i. with working output pointed to mytempdir
- d. Tar contents of mytempdir
  - i. tar cf myresults.tar myfiles
- e. Copy the results back to storage
  - i. cp myresults.tar remote/myresults
- f. (check the move with checksum)
- g. Cleanup (some systems do this for you)
  - i. rm -r mytempdir

# SGE Using the /tmp space

```
#!/bin/bash
#$ -N bigIO_example
#$ -S /bin/bash
#$ -1 h_vmem=1g
#$ -1 h_rt=02:00:00
#$ -pe smp 4
```startupscript
# mydir="$2" tmploc="$1"
starttar=$3 #
mytmp="$1"/"$2"
mkdir "$mytmp"
cp "$starttar" "$mytmp"
cd $mytmp
tar xf "$startar"
runscript "$mytmp"
tar cf results.tar $mytmp
cp results.tar <remote system>
rm -r "$mytemp"
```

# SGE Array jobs

```
#!/bin/bash
#$ -N bigIO_example
#$ -S /bin/bash
#$ -1 h_vmem=500m
#$ -1 h_rt=00:015:00
#$ -pe smp 1
#$ -t 1-5
e
# don't do this
```

# SGE Array jobs

- Split the files into batched directories
- 2. Make a file listing
   those directories
   a. ls -1 \* > list.txt
- 3. Get a count of files a. wc -l list.txt
- 4. Use sed or awk to pop from list.txt and use bigIO pattern

```
#!/bin/bash
#$ -N bigIO_example
#$ -S /bin/bash
#$ -1 h_vmem=500m
#$ -1 h_rt=00:015:00
#$ -pe smp 1
\#$ -t 1-<number from "wc -1">
infile=$(sed -n -e "$SGE_TASK_ID p" list.txt)
starttar="$infile"
mytmp=/tmp/"$2"
mkdir "$mytmp"
cp "$starttar" "$mytmp"
cd $mytmp
tar xf "Sstartar"
runscript "$mytmp"
tar cf "$infile"_results.tar $mytmp
cp "$infile"_results.tar <remote system>
rm -r "$mytemp"
```

## **HPC** resources

- With ITsupport at Einstein
- XSEDE
- Open Science Grid
- "The Cloud"
  - Amazon, Google, Microsoft and others