

CS 202 - Fundamental Structures of Computer Science 2

Homework 1 – Algorithm Efficiency and Sorting

Student Name: Esad İsmail Tök

Student ID: 21801679

Section: 2

Question 1

Part A

- In order to prove that $f(x) = 5n^3 + 4n^2 + 10$ is $O(n^4)$, we need to be able to find two numbers, c and n_0 , such that $0 \leq 5n^3 + 4n^2 + 10 \leq c \times n^4$ when $n \geq n_0$.
- By finding such two integers, c and n_0 , we will be able to prove that $c \times n^4$ constitutes an upper bound on $f(x) = 5n^3 + 4n^2 + 10$.
- When we choose $c = 5$ and $n_0 = 2$, the equation that must be satisfied becomes $0 \leq 5n^3 + 4n^2 + 10 \leq 5 \times n^4$ when $n \geq 2$. This equation is always true and it can be seen that $5 \times n^4$ is always an upper bound for the equation $f(x) = 5n^3 + 4n^2 + 10$ when we substitute any value for the input size, n , that is greater than or equal to 2.
- Thus we have proved that there is such numbers, c and n_0 , that holds the equation $0 \leq 5n^3 + 4n^2 + 10 \leq c \times n^4$ when $n \geq n_0$, with one of the selections as $c = 5$ and $n_0 = 2$.
- Therefore the equation $f(x) = 5n^3 + 4n^2 + 10$ is order of n^4 , namely $O(n^4)$.

Question 1

Part B

- Tracing the array [24, 8, 51, 28, 20, 29, 21, 17, 38, 27] using the insertion sort algorithm:

24	8	51	28	20	29	21	17	38	27
----	---	----	----	----	----	----	----	----	----

Copy 8

24	24	51	28	20	29	21	17	38	27
----	----	----	----	----	----	----	----	----	----

Shift 24

8	24	51	28	20	29	21	17	38	27
---	----	----	----	----	----	----	----	----	----

Insert 8; copy 51, insert 51 on top of itself

8	24	51	28	20	29	21	17	38	27
---	----	----	----	----	----	----	----	----	----

Copy 28

8	24	51	51	20	29	21	17	38	27
---	----	----	----	----	----	----	----	----	----

Shift 51

8	24	28	51	20	29	21	17	38	27
---	----	----	----	----	----	----	----	----	----

Insert 28; copy 20

8	24	24	28	51	29	21	17	38	27
---	----	----	----	----	----	----	----	----	----

Shift 51, 28, 24

8	20	24	28	51	29	21	17	38	27
---	----	----	----	----	----	----	----	----	----

Insert 20; copy 29

8	20	24	28	51	51	21	17	38	27
---	----	----	----	----	----	----	----	----	----

Shift 51

8	20	24	28	29	51	21	17	38	27
---	----	----	----	----	----	----	----	----	----

Insert 29; copy 21

8	20	24	24	28	29	51	17	38	27
---	----	----	----	----	----	----	----	----	----

Shift 51, 29, 28, 24

8	20	21	24	28	29	51	17	38	27
---	----	----	----	----	----	----	----	----	----

Insert 21; copy 17

8	20	20	21	24	28	29	51	38	27
---	----	----	----	----	----	----	----	----	----

Shift 51, 29, 28, 24, 21, 20

8	17	20	21	24	28	29	51	38	27
---	----	----	----	----	----	----	----	----	----

Insert 17; copy 38

8	17	20	21	24	28	29	51	51	27
---	----	----	----	----	----	----	----	----	----

Shift 51

8	17	20	21	24	28	29	38	51	27
---	----	----	----	----	----	----	----	----	----

Insert 38; copy 27

8	17	20	21	24	28	28	29	38	51
---	----	----	----	----	----	----	----	----	----

Shift 51, 38, 29, 28

8	17	20	21	24	27	28	29	38	51
---	----	----	----	----	----	----	----	----	----

Insert 27

- Tracing the array [24, 8, 51, 28, 20, 29, 21, 17, 38, 27] using the bubble sort algorithm:

24	8	51	28	20	29	21	17	38	27
8	24	51	28	20	29	21	17	38	27
8	24	51	28	20	29	21	17	38	27
8	24	28	51	20	29	21	17	38	27
8	24	28	20	51	29	21	17	38	27
8	24	28	20	29	51	21	17	38	27
8	24	28	20	29	21	51	17	38	27
8	24	28	20	29	21	17	51	38	27
8	24	28	20	29	21	17	38	51	27
8	24	28	20	29	21	17	38	27	51

Pass 1

- In the *Pass 1* figure above, the algorithm compares the array items pairwise and changes the order if the first element is greater than the second. At the end of *Pass 1*, the greatest element is bubbled.

8	24	28	20	29	21	17	38	27	51
8	24	28	20	29	21	17	38	27	51
8	24	28	20	29	21	17	38	27	51
8	24	20	28	29	21	17	38	27	51
8	24	20	28	29	21	17	38	27	51
8	24	20	28	21	29	17	38	27	51
8	24	20	28	21	17	29	38	27	51
8	24	20	28	21	17	29	38	27	51
8	24	20	28	21	17	29	27	38	51

Pass 2

- In the *Pass 2* figure above, the algorithm compares the array items pairwise until the last bubbled item and changes the order if the first element is greater than the second. At the end of *Pass 2*, the second greatest element is bubbled.

8	24	20	28	21	17	29	27	38	51
8	24	20	28	21	17	29	27	38	51
8	20	24	28	21	17	29	27	38	51
8	20	24	28	21	17	29	27	38	51
8	20	24	21	28	17	29	27	38	51
8	20	24	21	17	28	29	27	38	51
8	20	24	21	17	28	29	27	38	51
8	20	24	21	17	28	27	29	38	51

Pass 3

- In the *Pass 3* figure above, the algorithm compares the array items pairwise until the last bubbled item and changes the order if the first element is greater than the second. At the end of *Pass 3*, the third greatest element is bubbled.

8	20	24	21	17	28	27	29	38	51
8	20	24	21	17	28	27	29	38	51
8	20	24	21	17	28	27	29	38	51
8	20	21	24	17	28	27	29	38	51
8	20	21	17	24	28	27	29	38	51
8	20	21	17	24	28	27	29	38	51
8	20	21	17	24	27	28	29	38	51

Pass 4

- In the *Pass 4* figure above, the algorithm compares the array items pairwise until the last bubbled item and changes the order if the first element is greater than the second. At the end of *Pass 4*, the fourth greatest element is bubbled.

8	20	21	17	24	27	28	29	38	51
8	20	21	17	24	27	28	29	38	51
8	20	21	17	24	27	28	29	38	51
8	20	17	21	24	27	28	29	38	51
8	20	17	21	24	27	28	29	38	51
8	20	17	21	24	27	28	29	38	51

Pass 5

- In the *Pass 5* figure above, the algorithm compares the array items pairwise until the last bubbled item and changes the order if the first element is greater than the second. At the end of *Pass 5*, the fifth greatest element is bubbled.

8	20	17	21	24	27	28	29	38	51
8	20	17	21	24	27	28	29	38	51
8	17	20	21	24	27	28	29	38	51
8	17	20	21	24	27	28	29	38	51
8	17	20	21	24	27	28	29	38	51

Pass 6

- In the *Pass 6* figure above, the algorithm compares the array items pairwise until the last bubbled item and changes the order if the first element is greater than the second. At the end of *Pass 6*, the sixth greatest element is bubbled.

8	17	20	21	24	27	28	29	38	51
8	17	20	21	24	27	28	29	38	51
8	17	20	21	24	27	28	29	38	51
8	17	20	21	24	27	28	29	38	51
8	17	20	21	24	27	28	29	38	51

Pass 7

- In the *Pass 7* figure above, the algorithm compares the array items pairwise until the last bubbled array item and there is no data move happens between the pairwise array elements in this pass.
- Therefore bubble sort algorithm decides that the array is sorted using its boolean flag. And since there is no need to go along the other passes and the array becomes sorted and the algorithm ends with the *Pass 7*.

Question 2

Part C

```
---Selection Sort---  
  
The array before selection sort  
[12, 7, 11, 18, 19, 9, 6, 14, 21, 3, 17, 20, 5, 12, 14, 8]  
  
Number of key comparisons: 120  
Number of data moves: 45  
  
The array after selection sort  
[3, 5, 6, 7, 8, 9, 11, 12, 12, 14, 14, 17, 18, 19, 20, 21]  
  
---Merge Sort---  
  
The array before merge sort  
[12, 7, 11, 18, 19, 9, 6, 14, 21, 3, 17, 20, 5, 12, 14, 8]  
  
Number of key comparisons: 46  
Number of data moves: 128  
  
The array after merge sort  
[3, 5, 6, 7, 8, 9, 11, 12, 12, 14, 14, 17, 18, 19, 20, 21]  
  
---Quick Sort---  
  
The array before quick sort  
[12, 7, 11, 18, 19, 9, 6, 14, 21, 3, 17, 20, 5, 12, 14, 8]  
  
Number of key comparisons: 45  
Number of data moves: 93  
  
The array after quick sort  
[3, 5, 6, 7, 8, 9, 11, 12, 12, 14, 14, 17, 18, 19, 20, 21]  
  
---Radix Sort---  
  
The array before radix sort  
[12, 7, 11, 18, 19, 9, 6, 14, 21, 3, 17, 20, 5, 12, 14, 8]  
  
The array after radix sort  
[3, 5, 6, 7, 8, 9, 11, 12, 12, 14, 14, 17, 18, 19, 20, 21]
```

- The above figure is the console output of the Question 2 Part C. For the selection sort, merge sort, and quick sort algorithms; the program displays the array before the sorting operation, then displays the number of key comparisons and data moves to sort the array, and lastly displays the sorted form of the array. For the radix sort algorithm, the program displays the array before the sorting operation and the array after the sorting operation.

Question 2

Part D

```
-----Performing Analysis on Randomly Ordered Arrays-----
```

Analysis of Selection Sort				
Array Size	Elapsed Time		compCount	moveCount
6000	34	ms	17997000	17997
10000	94	ms	49995000	29997
14000	181	ms	97993000	41997
18000	296	ms	161991000	53997
22000	440	ms	241989000	65997
26000	615	ms	337987000	77997
30000	817	ms	449985000	89997

Analysis of Merge Sort				
Array Size	Elapsed Time		compCount	moveCount
6000	2	ms	67729	151616
10000	2	ms	120233	267232
14000	3	ms	174990	387232
18000	4	ms	231522	510464
22000	5	ms	289475	638464
26000	5	ms	348075	766464
30000	6	ms	407612	894464

Analysis of Quick Sort				
Array Size	Elapsed Time		compCount	moveCount
6000	0	ms	225923	86823
10000	2	ms	584190	173112
14000	3	ms	1112980	259110
18000	4	ms	1771131	306582
22000	5	ms	2584627	289272
26000	8	ms	3610581	361692
30000	10	ms	4753363	464052

Analysis of Radix Sort		
Array Size	Elapsed Time	
6000	1	ms
10000	0	ms
14000	0	ms
18000	0	ms
22000	1	ms
26000	1	ms
30000	0	ms

- The above figure is the output for the performing analysis of the sorting random order array scenario from Question 2 Part D.
- For selection sort, merge sort, and quick sort; the program displays the elapsed times (ms), key comparison count, and data move count for 7 different input sizes.
- For the radix sort, the program displays the elapsed time (ms) for 7 different input sizes.

Question 2

Part D

```
-----Performing Analysis on Ascendingly Ordered Arrays-----  
  
-----  
Analysis of Selection Sort  
Array Size    Elapsed Time    compCount    moveCount  
6000          34    ms        17997000     17997  
10000         102    ms        49995000     29997  
14000         180    ms        97993000     41997  
18000         324    ms        161991000    53997  
22000         452    ms        241989000    65997  
26000         616    ms        337987000    77997  
30000         817    ms        449985000    89997  
  
-----  
Analysis of Merge Sort  
Array Size    Elapsed Time    compCount    moveCount  
6000          0    ms        42538        151616  
10000         1    ms        74599        267232  
14000         2    ms        108369       387232  
18000         2    ms        142209       510464  
22000         3    ms        176166       638464  
26000         4    ms        211081       766464  
30000         5    ms        250701       894464  
  
-----  
Analysis of Quick Sort  
Array Size    Elapsed Time    compCount    moveCount  
6000          36    ms        17997000     17997  
10000         99    ms        49995000     29997  
14000         178    ms        97993000     41997  
18000         305    ms        161991000    53997  
22000         427    ms        241989000    65997  
26000         592    ms        337987000    77997  
30000         816    ms        449985000    89997  
  
-----  
Analysis of Radix Sort  
Array Size    Elapsed Time  
6000          0    ms  
10000         0    ms  
14000         0    ms  
18000         0    ms  
22000         1    ms  
26000         1    ms  
30000         1    ms
```

- The above figure is the output for the performing analysis of the sorting ascending order array scenario from Question 2 Part D.
- For selection sort, merge sort, and quick sort; the program displays the elapsed times (ms), key comparison count, and data move count for 7 different input sizes.
- For the radix sort, the program displays the elapsed time (ms) for 7 different input sizes.

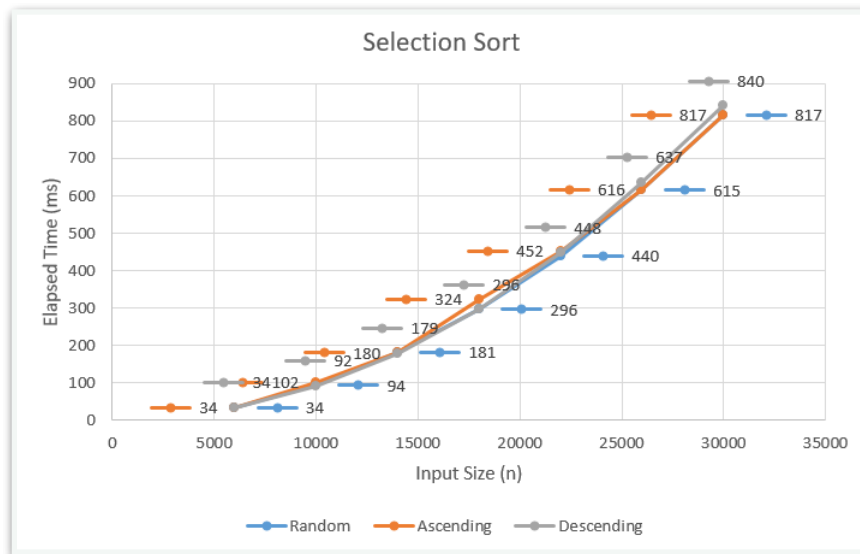
Question 2

Part D

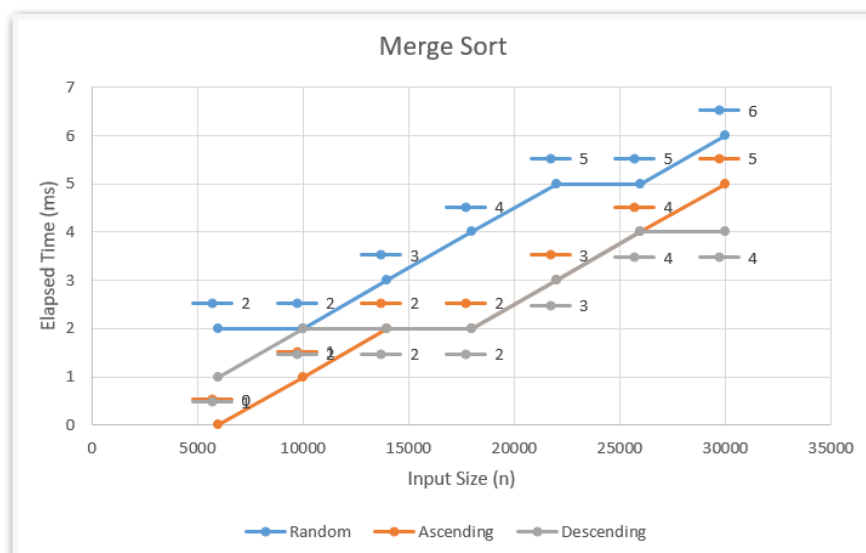
```
-----Performing Analysis on Descendingly Ordered Arrays-----  
  
-----  
Analysis of Selection Sort  
Array Size    Elapsed Time    compCount    moveCount  
6000          34 ms          17997000     17997  
10000         92 ms          49995000     29997  
14000         179 ms         97993000     41997  
18000         296 ms         161991000    53997  
22000         448 ms         241989000    65997  
26000         637 ms         337987000    77997  
30000         840 ms         449985000    89997  
  
-----  
Analysis of Merge Sort  
Array Size    Elapsed Time    compCount    moveCount  
6000          1 ms           36656        151616  
10000         2 ms           64608        267232  
14000         2 ms           94256        387232  
18000         2 ms           124640       510464  
22000         3 ms           154208       638464  
26000         4 ms           186160       766464  
30000         4 ms           219504       894464  
  
-----  
Analysis of Quick Sort  
Array Size    Elapsed Time    compCount    moveCount  
6000          2 ms           764245       898224  
10000         3 ms           1460716      1477608  
14000         5 ms           2345797      2100615  
18000         8 ms           3380398      2714685  
22000         10 ms          4564966      3308643  
26000         14 ms          5918483      3920526  
30000         16 ms          7430242      4521792  
  
-----  
Analysis of Radix Sort  
Array Size    Elapsed Time  
6000          0 ms  
10000         0 ms  
14000         1 ms  
18000         1 ms  
22000         1 ms  
26000         1 ms  
30000         1 ms
```

- The above figure is the output for the performing analysis of the sorting descending order array scenario from Question 2 Part D.
- For selection sort, merge sort, and quick sort; the program displays the elapsed times (ms), key comparison count, and data move count for 7 different input sizes.
- For the radix sort, the program displays the elapsed time (ms) for 7 different input sizes.

Question 3



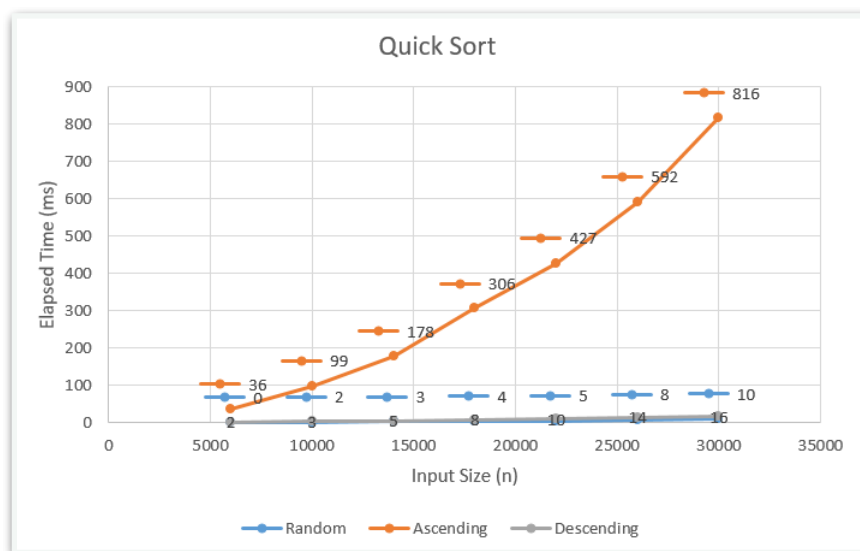
- Selection sort algorithm is $O(n^2)$ for random, ascending, and descending order arrays and its growth rate is independent of the order of the input array since the algorithm needs to find the greatest element in the unsorted sublist in each iteration. As it can be seen in the graph the records are considerably similar for all scenarios.
- When the experimental results are compared with theoretical expectations, it can be seen that the experimental pattern is significantly similar with the theoretical ones. However there occurs some minor elapsed time differences for some executions that can be caused by the amount of CPU usage of the computer at the run time or machine dependent factors if the executions are done in different computers



- Merge sort algorithm is $O(n \cdot \log n)$ for random, ascending, and descending order arrays and its growth rate is independent of the order of the input array. The elapsed time records of the merge sort is considerably smaller than the selection sort.

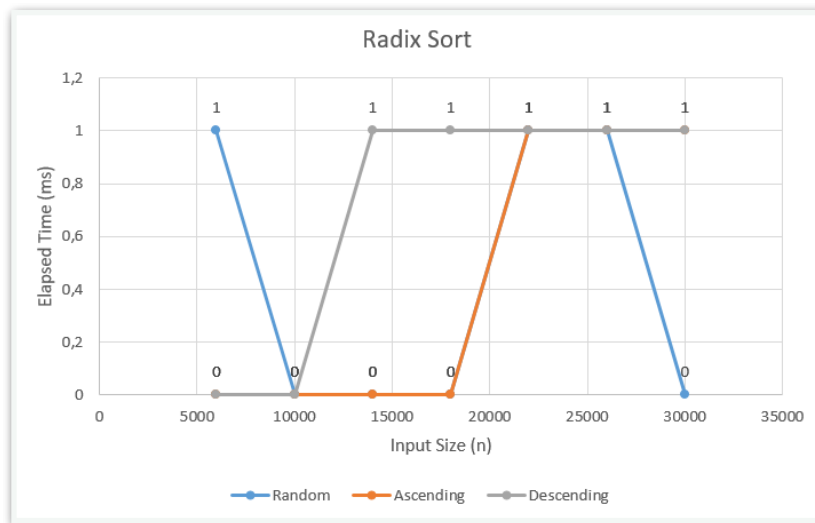
Question 3

- Although the algorithm is always $O(n \cdot \log n)$, it can be seen that for randomly ordered array scenario, the elapsed time records are greater than the other scenarios. So there is a contrast between the theoretical expectations.
- It is because even though the number of data moves does not change across the scenarios and decides the growth rate, there is less number of key comparisons in ascending and descending order arrays since one half is always smaller or greater than the other half. Therefore there can happen minor elapsed time differences among different scenarios.



- Quick sort algorithm is $O(n \cdot \log n)$ for random order input array but $O(n^2)$ for ascending and descending order input arrays.
- The random input array is the average case of the quick sort algorithm. Both ascending and descending input arrays are the worst cases of quick sort algorithm because in those cases the pivot value, which is the first element of the input array, partitions the array into 2 subarrays with sizes 0 and $(n - 1)$. Therefore bad partition causes the algorithm to grow as fast as selection sort algorithm which is also $O(n^2)$.
- According to the theoretical results, the descending performance of the algorithm also needs to have a similar pattern with ascending scenario. However in the graph, the elapsed times of the ascending array scenario dominates the descending one in contrast to theory. The causes of this situation may be the compiler or any other intermediate step between the compilation step and generating the executable.

Question 3



- Radix sort algorithm is $O(n)$ for random, ascending, and descending order arrays and its growth rate is independent of the order of the input array. It is a significantly faster algorithm compared to other ones.
- The reason why this algorithm is so fast and does not depend on the input array arrangement is, because it does not use key comparisons when sorting the array, so does not waste the time with comparisons.
- At first the patterns of the graph seems not to be fitting the theoretical $O(n)$ pattern of the algorithm. However it is because the elapsed time values are too small and only changes between 0 and 1. Therefore the input sizes up to 30000 are not enough to observe the complete pattern of the radix sort algorithm and larger input sizes are needed to fit the algorithm into theoretical pattern.