Eric Schulte

Department of Computer Science University of New Mexico Albuquerque, NM

2014

Extant Software

The existing Software ecosystem.

- applications
- libraries
- compilers
- operating systems
- architectures

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The existing Software ecosystem.

- applications
- libraries
- compilers
- operating systems
- architectures

Evolution

Evolved product of evolutionary forces

Evolvable amenable to automated improvement

Automatically Fix Bugs in C Software

Automatically Fix Bugs in C Software

Collaboration between UNM and UVA

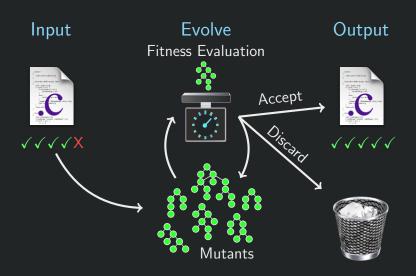


Dr. Stephanie Forrest



Dr. Westley Weimer

Automatically Fix Bugs in C Software



Automatically Fix Bugs in C Software

Strengths

Effective Repaired 55/105 bugs for \$8 each
General Multiple classes of bugs and security defects
Best Papers ICSE 2009, GECCO 2009, SBST 2009
Humies Gold 2009, Bronze 2012

Automatically Fix Bugs in C Software

But

Why doesn't this break my software?

Outline

Introduction

Software Mutational Robustness

Program Representations

Embedded Systems

NETGEAR Exploit

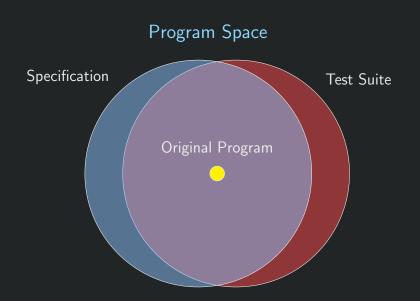
Program Optimization

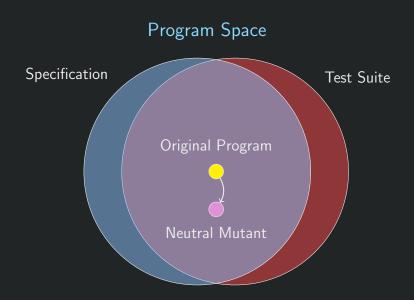
Future Work

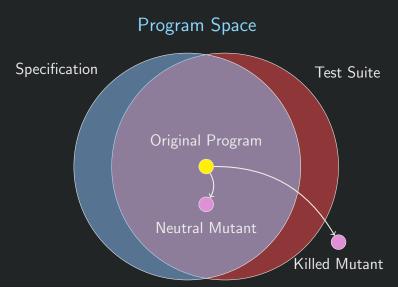
Conclusion

[Schulte, GPEM 2013]

percentage of mutants which are functional





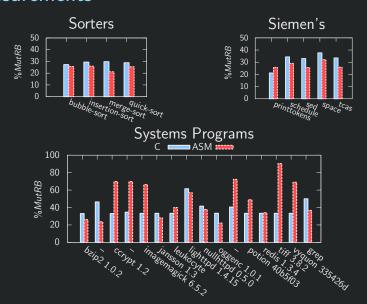


Definition

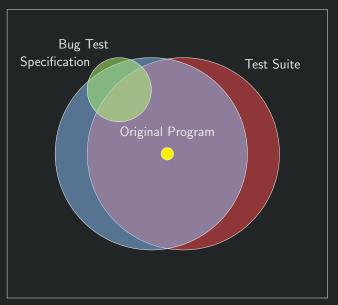
```
MutRB(P, T, M)
P \mid program
T \mid test suite
M \mid mutation operators
```

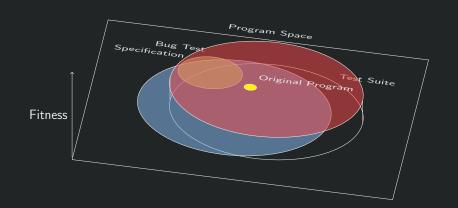
$$MutRB(P, T, M) = \frac{|\{P' \mid m \in M. \ P' \leftarrow m(P) \ \land \ T(P')\}|}{|\{P' \mid m \in M. \ P' \leftarrow m(P)\}|}$$

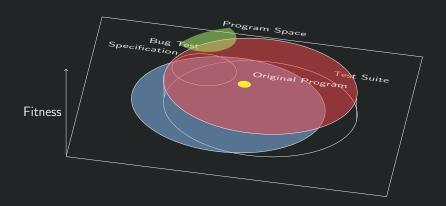
Measurements

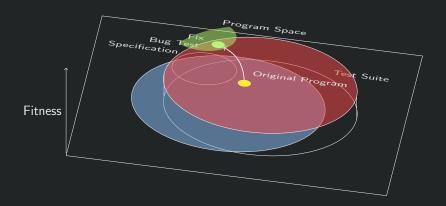


Program Space





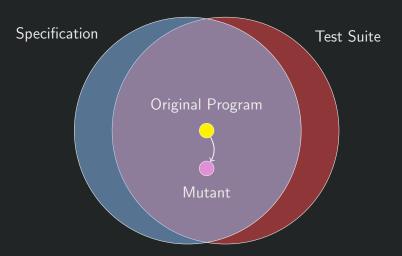




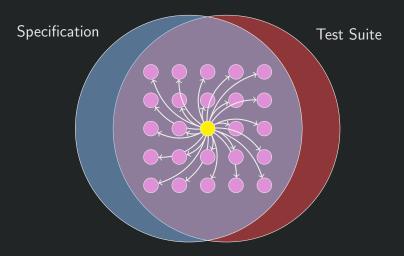




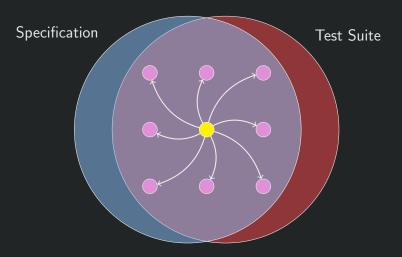
Program Space



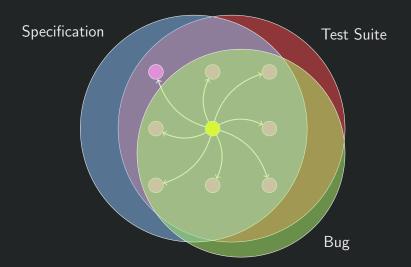
Program Space



Program Space



Automated Diversity and Proactive Bug Repair Program Space



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Source

```
if (a==0){
  printf("%g\n", b); }
else {
  while (b!=0){
    if (a>b){ a=a-b; }
    else { b=b-a; } }
printf("%g\n", a);
```

Source

AST

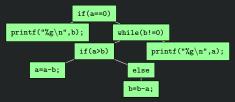
Source

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if (a==0){
  printf("%g\n", b); }
else {
  while (b!=0){
    if (a>b){ a=a-b; }
    else { b=b-a; } }
printf("%g\n", a);
```

ASM [Schulte, ASE 2010]

```
.file "gcd.c"
.globl main
.type main, @function
main:
.cfi.startproc
pushq %rbp
.cfi.def.cfa.offset 16
.cfi.offset 6, -16
movq %rsp, %rbp
.cfi.def.cfa.register 6
subq $48, %rsp
```

AST



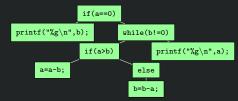
Source

ASM [Schulte, ASE 2010]

```
.file "gcd.c"
.globl main
.type main, @function
main:
.cfi.startproc
pushq %rbp
.cfi.def_cfa_offset 16
.cfi.offset 6, -16
movq %rsp, %rbp
.cfi.def_cfa_register 6
subq $48, %rsp
```

The Evolution of Extant Software

AST



ELF [Schulte, ASPLOS 2013]

```
ELF\?

ELF header

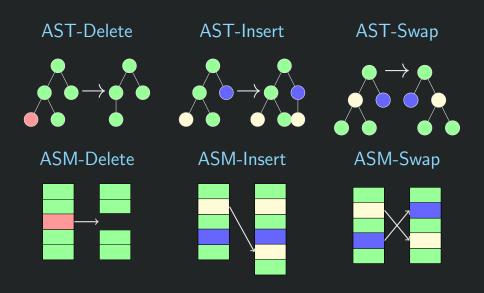
program header table

section 1
...
.text section
[55] [48 89 e5] [48 83 ec 20] [48 89 7d e8] [89 75 e4] [83 7d e4 01] [7e 60] ...

section n

section header table
```

Program Mutation Operations



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Program Repair in Embedded Devices

[Schulte, ASPLOS 2013]

Resource Constraints

- Small disks
- Less memory
- Slow processors
- ► Slow, costly comm.



Light Weight Fault Localization

- 1. Sample program counter
- 2. Translate memory addresses to program offsets.
- 3. Smooth sample with Gaussian convolution.



movq 8(%rdx), %rdi xorl %eax, %eax movl %eax, (%r15) addl \$1, %r14d call atoi movq -80(%rbp), %rdx movq %rdx, -80(%rbp) addq \$4, %r15 movq 8(%rdx), %rdi xorl %eax, %eax movl %eax, (%r15)

Machine-code Instructions

Light Weight Fault Localization

- 1. Sample program counter.
- 2. Translate memory addresses to program offsets.
- 3. Smooth sample with Gaussian convolution.

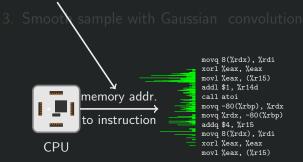


movq 8(%rdx), %rdi xorl %eax, %eax movl %eax, (%r15) addl \$1, %r14d call atoi movq -80(%rbp), %rdx movq %rdx, -80(%rbp) addq \$4, %r15 movq 8(%rdx), %rdi xorl %eax, %eax movl %eax, (%r15)

Machine-code Instructions

Light Weight Fault Localization

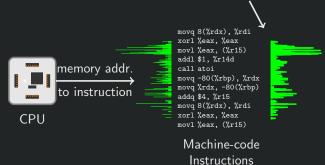
- 1. Sample program counter.
- 2. Translate memory addresses to program offsets.



Machine-code Instructions

Light Weight Fault Localization

- 1. Sample program counter.
- 2. Translate memory addresses to program offsets.
- 3. Smooth sample with Gaussian convolution.



Genprog

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Limitations

- Requires source code
- Requires build tool chain
- ► Requires program instrumentation
- Expensive fitness function (compilation, test execution)

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Embedded Repair Benchmark Programs

Program	Program Description	Bug	
atris	graphical tetris game	stack buffer exploit	
ccrypt	encryption utility	segfault	
deroff	document processing	segfault	
flex	lexical analyzer generator	segfault	
indent	source code processing	infinite loop	
look svr4	dictionary lookup	infinite loop	
look ultrix	dictionary lookup	infinite loop	
merge	merge sort	duplicate inputs	
merge-cpp	merge sort (in $C{++})$	duplicate inputs	
s3	sendmail utility	buffer overflow	
uniq	duplicate text processing	segfault	
units	metric conversion	segfault	
zune	embedded media player	infinite loop	

- Effective
- ▶ 62% faster runtime
- 95% smaller disk footprint
- ▶ 86% less memory

Total bugs repaired				
Rep.	Num. Bugs			
AST	13			
ASM	12			
ELF	11			

Average success rate			
100 runs per bug			
Rep.	Success Rate		
AST	78.17%		
ASM	70.75%		
ELF	65.83%		

- Effective
- ▶ 62% faster runtime
- ▶ 95% smaller disk footprint
- 86% less memory

Expected fitness evaluations Rep. Evaluations AST 583.98 ASM 188.38 ELF 207.15

Total runtime		
Rep.	Sec.	
AST	229.50	
ASM	278.30	
ELF	74.20	

- Effective
- ▶ 62% faster runtime
- ▶ 95% smaller disk footprint
- ▶ 86% less memory

Example: Merge Sort Repair by representation

- ▶ AST, 2 of 4900 Swaps
- ► ASM, 1 of 280 Deletes

merge.c

merge.s

```
cmpl %eax, %edx ; fix: del.
jg .L12
movq -72(%rbp), %rax
```

Disk size

- Effective
- ▶ 62% faster runtime
- ▶ 95% smaller disk footprint
- ▶ 86% less memory

	Requirements
AST	Source code
	& build toolchain
ASM	Assembly code
	& linker
ELF	Compiled executable

Working memory

Rep.	MB	
AST	1402	
ASM	756	
ELF	200	

- Effective
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NETGEAR Exploit

[Schulte, unpublished]

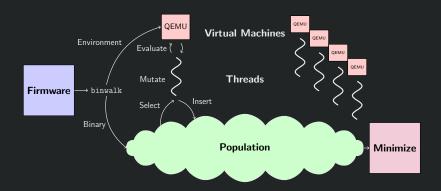


- 1. URI starting with BRS bypasses authentication
- 2. URI including unauth.cgi or securityquestions.cgi bypass authentication
- 3. unprotected page removes authentication for every page http://router/BRS_02_genieHelp.html

Common Vulnerability



NETGEAR Repair Overview



NETGEAR Repair Results

Results

- repair without regression tests
- repair of real un-patched exploit
- multiple concurrent repairs

Repair Runtime

```
evaluations \sim 36,000
runtime 86.6 minutes
threads 32
```

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Optimization Problem

[Schulte, ASPLOS 2014]

Optimizing complex non-functional properties

 $properties \times hardware \times environment$ $properties \ memory, \ network, \ energy, \ etc. \ . \ .$ $hardware \ architectures, \ processors, \ memory \ stack, \ etc. \ . \ .$

environment variables, load, etc. . .

Every program transformation requires

- a-priori reasoning
- manual implementation
- © guaranteed correctness

Our Solution

Genetic Optimization Algorithm (GOA)

- empirically guided (guess and check)
- automated evolutionary search
- relaxed semantics

Applied to PARSEC benchmarks

- ▶ reduces energy consumption by 20% on average
- maintain functionality on withheld tests

Technique

Post-compiler, test-driven, Genetic Optimization Algorithm

Post-compiler



Technique

Post-compiler, test-driven, Genetic Optimization Algorithm

Test driven

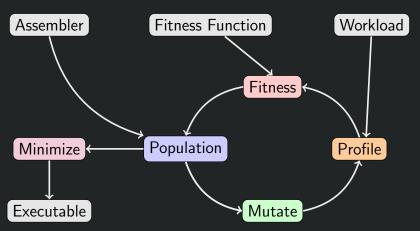
Use test cases to exercise program

- evaluate functionality
- measure runtime properties

Technique

Post-compiler, test-driven, Genetic Optimization Algorithm

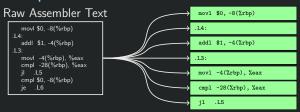
Genetic Optimization Algorithm (GOA)



Program Mutation



Software Representation

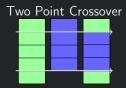


Mutation Operations









Profiling





Hardware
Performance
Counters

Profiling



\$ perf stat -- ./blackscholes 1 input /tmp/output

6,864,315,342 cycles

5,062,293,918 instructions

2,944,060,039 r533f00

1,113,084,780 cache-references

1,122,960 cache-misses

3.227585368 seconds time elapsed

Fitness Function



$$\frac{\textit{energy}}{\textit{time}} = \textit{C}_{\textit{const}} + \textit{C}_{\textit{ins}} \frac{\textit{ins}}{\textit{cycle}} + \textit{C}_{\textit{flops}} \frac{\textit{flops}}{\textit{cycle}} + \textit{C}_{\textit{tca}} \frac{\textit{tca}}{\textit{cycle}} + \textit{C}_{\textit{mem}} \frac{\textit{mem}}{\textit{cycle}}$$

Steady State Genetic Algorithm



Details

- ▶ population size: 2¹⁰
- ▶ 2¹⁸ fitness evaluations
- $ightharpoonup \sim 16$ hour runtime per optimization

```
5358c5358
< .L808:
> addl %ebx. %ecx
5416c5416
< addl %ebx. %ecx
> .L808:
5463c5463
< .L970:
> .byte 0x33
5651d5650
< .loc 1 457 0 is_stmt 0 discriminator 2
5841d5839
< addg %rdx, %r14
6309c6307
< xorpd %xmm1, %xmm7
> cmpq %r13, %rdi
6413a6412
> cmpl %ecx, %esi
```



```
5358c5358
< .L808:
> addl %ebx. %ecx
5416c5416
< addl %ebx. %ecx
> .L808:
5463c5463
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```



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```







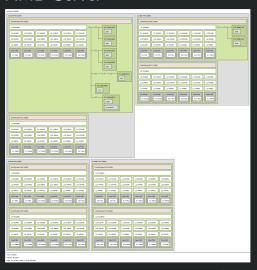
Benchmark Applications

	C/C++	ASM	
Program	Lines	of Code	Description
blackscholes	510	7,932	Finance modeling
bodytrack	14,513	955,888	Human video tracking
facesim			no alternate inputs
ferret	15,188	288,981	Image search engine
fluidanimate	11,424	44,681	Fluid dynamics animation
freqmine	2,710	104,722	Frequent itemset mining
raytrace			no testable output
swaptions	1,649	61,134	Portfolio pricing
vips	142,019	132,012	Image transformation
×264	37,454	111,718	MPEG-4 video encoder
total	225,467	1,707,068	

Table: PARSEC benchmark applications.

Hardware Platforms

AMD Server



Intel Desktop



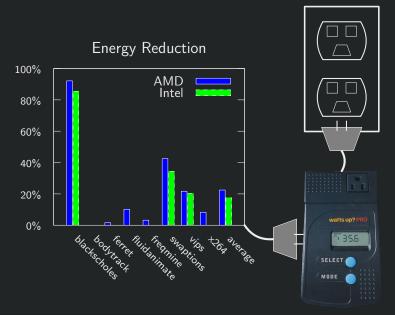
Energy Model

$$\frac{\textit{energy}}{\textit{time}} = \textit{C}_{\textit{const}} + \textit{C}_{\textit{ins}} \frac{\textit{ins}}{\textit{cycle}} + \textit{C}_{\textit{flops}} \frac{\textit{flops}}{\textit{cycle}} + \textit{C}_{\textit{tca}} \frac{\textit{tca}}{\textit{cycle}} + \textit{C}_{\textit{mem}} \frac{\textit{mem}}{\textit{cycle}}$$

		Intel	AMD
Coefficient	Description	(4-core)	(48-core)
C_{const}	constant power draw	31.530	394.74
C_{ins}	instructions	20.490	-83.68
C_{flops}	floating point ops.	9.838	60.23
C_tca	cache accesses	-4.102	-16.38
C_{mem}	cache misses	2962.678	-4209.09

Table: Energy model coefficients.

Results: Energy Reduction



Functionality on Withheld Tests

Program	AMD	Intel
blackscholes	100%	100%
bodytrack	92%	100%
ferret	100%	100%
fluidanimate	6%	31%
freqmine	100%	100%
swaptions	100%	100%
vips	100%	100%
×264	27%	100%

Anecdotes

Blackscholes

- ▶ 90% less energy
- removed redundant outer loop
- modified semantics

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Swaptions

- ▶ 42% less energy
- improved branch prediction
- hardware specific

Anecdotes

Blackscholes

- ▶ 90% less energy
- removed redundant outer loop
- modified semantics

Swaptions

- ▶ 42% less energy
- improved branch prediction
- hardware specific

Vips

- ▶ 20% less energy
- substitution of memory access for calculation
- resource trade-off

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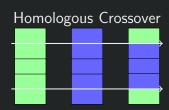
Conclusion

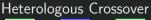
Future Work

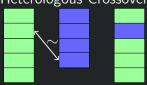
Heterologous crossover, Verification, Static, Hardening

Crossover between

- different compilation flags
- different versions
- different implementations
- different programs







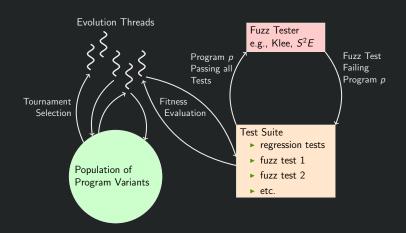
Heterologous crossover, Verification, Static, Hardening

- Assembler diff chunks formally equivalent
- Invariant preservation e.g., Daikon

Heterologous crossover, Verification, Static, Hardening

- Static analysis tools
- Annotations
 e.g., Eiffel contracts, Meta-level compilation

Heterologous crossover, Verification, Static, Hardening



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Conclusion

Extant software is

1. Mutationally robust

Conclusion

Extant software is

- 1. Mutationally robust
- 2. A product of evolution

Conclusion

Extant software is

- 1. Mutationally robust
- 2. A product of evolution
- 3. Amenable to automated evolutionary improvement

Thank You

Eric Schulte

email eschulte@cs.unm.edu

homepage https://cs.unm.edu/~eschulte

Backup Slides

Slides

- Benefits of ASM
- ► Embedded Repair Search Space
- ► Many Bugs
- ► LLVM Representation
- ► Optimization: Runtime and Energy Reduction

Benefits of ASM

Back

The ASM representation performs well.

Avoids

- direct addresses
- argumented instructions

Similarities to DNA

- Linear vector genome
- read sequentially
- reading frames
- start and stop codes
- padding
- bootstrapped compilers

Embedded Repair Search Space

Back

```
Program size
```

pprox 3 imes more assembly instructions than C statements

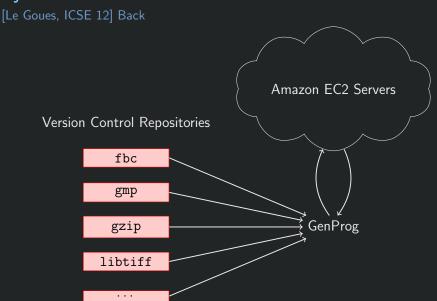
Search space size

 $= |alphabet|^{program \ size}$

Possible program coverage



Systematic Evaluation



Fixing Bugs for \$8 a Bug

Back

	Defects	Avg. Cost per Non-Repair		Avg. Cost Per Repair	
Program	Repaired	Hours	US\$	Hours	US\$
fbc	1 / 3	8.52	5.56	6.52	4.08
gmp	1 / 2	9.93	6.61	1.60	0.44
gzip	1 / 5	5.11	3.04	1.41	0.30
libtiff	17 / 24	7.81	5.04	1.05	0.04
lighttpd	5 / 9	10.79	7.25	1.34	0.25
php	28 / 44	13.00	8.80	1.84	0.62
python	$1 \ / \ 11$	13.00	8.80	1.22	0.16
wireshark	1 / 7	13.00	8.80	1.23	0.17
total	55 / 105	11.22h		1.60h	

LLVM Representations

Back

%x = alloca i32, align 4
store i32 2, i32* %x, align 4
%0 = load i32* %x, align 4
%add = add nsw i32 %0, 3
store i32 %add, i32* %x, align 4
%1 = load i32* %x, align 4
%2 = load i32* %x, align 4
%mul = mul nsw i32 %1, %2
store i32 %mul, i32* %x, align 4
%3 = load i32* %x, align 4
%call = call ... @printf(@.str %3)



LLVM Delete

Back

delete 4

%x = alloca i32, align 4
store i32 2, i32* %x, align 4
%0 = load i32* %x, align 4
%add = add nsw i32 %0, 3
store i32 %0, i32* %x, align 4
%1 = load i32* %x, align 4
%2 = load i32* %x, align 4
%mul = mul nsw i32 %1, %2
store i32 %mul, i32* %x, align 4
%3 = load i32* %x, align 4
%3 = load i32* %x, align 4
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LLVM Delete

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delete 4

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%0 = load i32* %x, align 4
%add = add nsw i32 %0, 3
store i32 %0, i32* %x, align 4
%1 = load i32* %x, align 4
%2 = load i32* %x, align 4
%mul = mul nsw i32 %1, %2
store i32 %mul, i32* %x, align 4
%3 = load i32* %x, align 4
%3 = load i32* %x, align 4
%call = call ... @printf(@.str %3)



LLVM Insert

Back

insert 9 4

%x = alloca i32, align 4
store i32 2, i32* %x, align 4
%0 = load i32* %x, align 4
%add = add nsw i32 %0, 3
store i32 %add, i32* %x, align 4
%1 = load i32* %x, align 4
%2 = load i32* %x, align 4
%mul = mul nsw i32 %1, %2

store i32 %mul, i32* %x, align 4
%3 = load i32* %x, align 4
%call = call ... @printf(@.str %3)

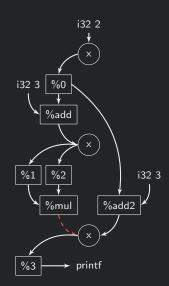


LLVM Insert

Back

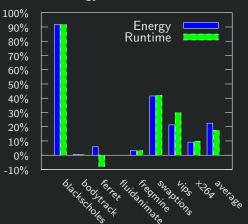
insert 9 4

%x = alloca i32, align 4
store i32 2, i32* %x, align 4
%0 = load i32* %x, align 4
%add = add nsw i32 %0, 3
store i32 %add, i32* %x, align 4
%1 = load i32* %x, align 4
%2 = load i32* %x, align 4
%mul = mul nsw i32 %1, %2
%add.insert = add nsw i32 %0, 3
store i32 %mul, i32* %x, align 4
%3 = load i32* %x, align 4
%call = call ... @printf(@.str %3)



Results: Runtime and Energy Reduction

AMD Energy and Runtime Reduction



The End

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Eric Schulte

email eschulte@cs.unm.edu
homepage https://cs.unm.edu/~eschulte