(44)

Remark 1: If f(p) = 0 and  $f'(p) \neq 0$ ,

then for any Po sufficiently close to (s.t. 9 & C'(16), 9 (x) & Ca,6), 19 (x) & k)

Newton's method will converge at least quadratically.

Remark 2: If f(p) =0, then for po close to p secont method converges to p with order

N5 +1 ≈ 1.618

Order & Method

Bisection

1 it g'(p) \$ 0 ≥ 2 if g'(p)=0

Fived Pt. Iteration Newton ≥2 if f'(p) ≠0

51 if f'(p)=0 ≈ 1.611 if f'(p) \$0 Secant

if f'(p) =0

2.4 Multiple Roots

45

Def: let  $f(x) = (x-p)^m q(x)$  and  $\lim_{x\to p} q(x) \neq 0$ 

Then p is a zero of multiplicity m (multiple root) of f.

m=1 => p is a simple root.

Thm 1: A function  $f \in C'(a,b)$  has a simple root at  $p \in (a,b)$  if and only if f(p) = 0 and  $f'(p) \neq 0$ 

OPf: (=>) If f has a simple root at  $\rho$  in (a,b),
then  $\int f(\rho) = 0$ 2)  $f(x) = (x-\rho) \cdot g(x)$  where  $\lim_{x \to \rho} g(x) \neq 0$ 

Since  $f \in C'(a,b)$ ,  $f'(p) = \lim_{x \to p} f'(x) = \lim_{x \to p} (g(x) + (x-p)g'(x))$   $= \lim_{x \to p} g(x) \neq 0$ 

(=) If f(q) = 0 and  $f'(p) \neq 0$ ,  $f(x) = f(p) + f'(s(x))(x-p) = f'(s(x)) \cdot (x-p)$ So between  $f(x) = f'(s(x)) \cdot (x-p)$ Taylor

5 between x and p

(=) If f(1)=0 and f'(p) =0,

 $f(x) = f'(p) + f'(s(x))(x-p) = f'(s(x)) \cdot (x-p)$ 

 $0 = f(\rho) = f'(\rho) = f''(\rho) = -$ 

 $f^{(m)}(\rho) \neq 0$ 

Ex: let  $f(x) = e^x - x - 1$ .

a) Show that x = 0 is a zero of mult. z = 0 f(x)b) Show that Newton's Method does not converge guardrestically

Ex: let  $f(x) = e^x - x - 1$ .

a) Show that x = 0 is a zero of mult. z = 0 f(x)b) Show that Newton's Method does not converge guardrestically

Sol: a) Compute f(0), f'(0), f''(0)

Sol: 6) Compute f(0), f'(0), f''(0)  $f(x) = e' - x - 1 \implies f(0) = 0$   $f'(x) = e^{x} - 1 \implies f'(0) = 0$ 

 $f'(x) = e^{x} - 1 \implies f'(0) = 0$   $f''(x) = e^{x} \implies f''(0) = 1 \neq 0$ By Thm 2, p = 0 is a zero of mult. 2

By Thm 2, p=0 is a zero of mult. 2
b) Since f'(0) = 0, Newton's method does not converge quadratically.

2.4 Modified Newton's Method

(48)

Recall: we lose quadratic convergence when f'(p) = 0, i.e.,

when multiplicity of p is mal

Let  $M(x) = \frac{f(x)}{f'(x)} =$  $(x-p)^m q(x)$  $(x-p)^{m}q'(x) + m(x-p)^{m-1}q(x)$ 

 $(x-p)^{m-1}$  (x-p)q(x) $(x-p)^{m-1}$   $(x-p)q^{3}(x) + m q(x)$ 

Note that p is a simple root of M(x) since

 $\lim_{x\to p} \frac{q(x)}{(x-p)q'(x)+mq(x)} = \frac{1}{m} \neq 0$ 

I dea: Apply Newton's Method to M(x) rather than f(x) if f'(p) = 0. Since p is simple root => quadratic convergence to p.

Modified Newton: f(P.) · f'(P.)

 $\rho_{n+1} = \rho_n - \frac{\mathcal{M}(\rho_n)}{\mathcal{M}'(\rho_n)} = \rho_n - \frac{f(\rho_n) \cdot f'(\rho_n)}{\left[f'(\rho_n)\right]^2 - f(\rho_n) f''(\rho_n)}$ Remarks: 1) Modified Newton converges quadratically regardless of multiplicity of p.

2) Require second derivative information (expensive!)

## 2.5 Accelarating Convergence

(49)

Last time: Modified Newton improved linearly convergent to quadratically convergent scheme, but required higher derivatives (expensive!)

Q: Given a linearly convergent sequence, how to modify it to achieve faster convergence?

Def: Assume  $\{P_n\}_{n=1}^{\infty}$  Converge, to p and  $\lim_{n\to\infty} \frac{|P_{n+1}-p|}{|P_n-p|} = \lambda$ 

a) It  $\lambda = 0$ ,  $\{\rho_n\}_{n=1}^{\infty}$  converges superlinearly to  $\rho$ 

b) If  $0 < \lambda < 1$ ,  $\{ f_n \}_{n=1}^{\infty}$  converges linearly to pc) If  $\lambda = 1$ ,  $\{ f_n \}_{n=1}^{\infty}$  converges sublinearly to p

It N=), ((n) n=, converges sublinearly to p

Ex: Show that the seq.  $p_n = \frac{1}{n+1}$  converges sublinearly to 0.

 $\lim_{n\to\infty} \frac{|P_{n+1}-0|}{|P_n-0|} = \lim_{n\to\infty} \frac{\left|\frac{1}{n+2}\right|}{\left|\frac{1}{n+1}\right|} = \lim_{n\to\infty} \frac{|\eta+1|}{|\eta+2|} = 1$ 

Aithen's D' Method

I deq: speed up convergence of linearly convergent seq.

Let { p. } converge linearly to P)

i.e., lim IPn-Pl =  $\lambda \in (0,1)$ 

Assume that Pn-P, Pn+1-P, Pn+2-P have the same sign.

Then  $\frac{\rho_{n+2}-\rho}{\rho_{n+1}-\rho} \approx \frac{\rho_{n+1}-\rho}{\rho_{n}-\rho} \quad \text{for sufficiently large } n.$   $\approx \lambda \qquad \approx \lambda$ 

Isolating P:

 $\frac{\left(\rho_{n+1}-\rho_{n}\right)^{2}}{\rho_{n+2}-2\rho_{n+1}+\rho_{n}}$  $\rho \approx \frac{\rho_{n+2} - 2\rho_n - \rho_{n+1}^2}{\rho_{n+2} - 2\rho_{n+1} + \rho} =$ 

Define new sequence

 $\hat{\beta}_n = \beta_n - \frac{(\beta_{n+1} - \beta_n)^2}{\beta_{n+2} - 2\beta_{n+1} + \beta_n}$ 

Remark: need to compute two sequences:

P. P. ( P. ) → P.

Def: Giren {Pa}n=1, the first order forward difference is defined as △ Pn = Pn+1 - Pn two diff: \$2(Pn) = \$ (\$Pn) = 0 Pn+1 - 0 Pn = PA12 - 2 PA11 +PA 3rd order two diff: 13(Pn) = Pn+3 - 3Pn+2 +3Pn+1 - Pn kt order fud diff: \$\( \rangle (Pn) = \D (\Delta (Pn)), k=2. Aithen's 2 Method: Thm (Aithen's Method) converges to p linearly, and Assume that {Pn} = 1  $\lim_{n\to\infty} \frac{\rho_{n+1}-\rho}{\rho_n-\rho} \leq 1$ Then Aithons 22 sequence EP, 3 no converges to p faster than  $\{P_n\}$  in the sense that  $\lim_{n\to\infty} \frac{\hat{P}_n - P}{P_n - P} = 0$ 

Thm (Aithers Method) converges to p linearly, and Assume that {Pn}\_n=  $\lim_{n\to\infty}\frac{\rho_{n+1}-\rho}{\rho_n-\rho} \leq 1$ Then Aithons 22 sequence EP, 3 nz converges to p faster than (Pn) in the sense that  $\lim_{n\to\infty}\frac{\rho_n-\rho}{\rho_n-\rho}=0$  $\frac{\text{Dff}}{\text{let}}: \text{ Let } \lim_{n\to\infty} \frac{P_{n+1}-P}{P_n-P} = \lambda,$ 

Then we have

 $P_{n+1} - P_n = (P_{n+1} - P) - (P_n - P) = (S_n - I)(P_n - P)$ (\*) Pn+2-2Pn+1+Pn=(Pn+2-P)-2(Pn+1-P)+ Pn-p

 $= (\delta_{n+1} - \delta_n - 2\delta_n + 1) (\rho_n - \rho)$ 

 $\lim_{h\to\infty} \frac{\rho_n - \rho}{\rho_n - \rho} = \lim_{n\to\infty} \frac{1}{\rho_n - \rho} \left( \rho_n - \frac{(\rho_{n+1} - \rho_n)^2}{\rho_{n+2} - 2\rho_{n+1} + \rho_n} \right)$ 

 $=\lim_{n\to \infty}\frac{1}{\rho_{n}-\rho}\left(\rho_{n}-\rho\right)-\left[\frac{(\rho_{n+1}-\rho_{n})^{2}}{(\rho_{n}-\rho)\left(\rho_{n+2}-2\rho_{n+1}+\rho_{n}\right)}\right]$ 

 $= \lim_{n\to\infty} \left| -\frac{(\delta_{n-1})^{2}}{\delta_{n+1}\delta_{n}-2\delta_{n}+1} \right| = \left| -\frac{(\lambda-1)^{2}}{\lambda^{2}-2\lambda+1} \right|$ from (\*), (\*\*)

Aithers Method for Fixed Pt. Iteration

$$P_{0}, P_{1} = g(P_{0}), P_{2} = g(P_{1}) \longrightarrow \hat{P}_{0} = P_{0} - \frac{(P_{1} - P_{0})^{2}}{P_{2} - 2P_{1} + P_{0}}$$

$$P_{3} = g(P_{2}) \longrightarrow \hat{P}_{1} = P_{1} - \frac{(P_{2} - P_{1})^{2}}{P_{3} - 2P_{2} + P_{1}}$$

Steffenson's Method. Let 
$$\{\Delta^{2}\}(\rho_{n}) = \ell_{n} - \frac{(\rho_{n+1} - \rho_{n})^{2}}{\rho_{n+2} - 2\rho_{n+1} + \rho_{n}}$$

O)  $\rho_{0}^{(0)}$ ,  $\rho_{1}^{(0)} = g(\rho_{0}^{(0)})$ ,  $\rho_{2}^{(0)} = g(\rho_{1}^{(0)})$ 

P<sub>1</sub> =  $g(\rho_{0}^{(0)})$ ,  $\rho_{1}^{(0)} = g(\rho_{0}^{(0)})$ 

P<sub>2</sub> =  $g(\rho_{1}^{(0)})$ 

P<sub>3</sub> =  $g(\rho_{1}^{(0)})$ 

P<sub>4</sub> =  $g(\rho_{0}^{(0)})$ 

P<sub>5</sub> =  $g(\rho_{1}^{(0)})$ 

P<sub>6</sub> =  $g(\rho_{1}^{(0)})$ 

P<sub>6</sub> =  $g(\rho_{1}^{(0)})$ 

P<sub>1</sub> =  $g(\rho_{0}^{(0)})$ 

P<sub>2</sub> =  $g(\rho_{1}^{(0)})$ 

P<sub>2</sub> =  $g(\rho_{1}^{(0)})$ 

P<sub>3</sub> =  $g(\rho_{1}^{(0)})$ 

$$P_{0}^{(n)} = P_{0}^{(n)} - P_{0}^{(n)} + P_{0}^{(n)}$$

$$P_{0}^{(n)} = \{\Delta^{2}\} (P_{0}^{(n-1)}) \qquad P_{1}^{(n)} = g(P_{0}^{(n)}) \qquad P_{2}^{(n)} = g(P_{1}^{(n)})$$

$$P_{0}^{(n)} = \{\Delta^{2}\} (P_{0}^{(n-1)}) \qquad P_{1}^{(n)} = g(P_{0}^{(n)}) \qquad P_{2}^{(n)} = g(P_{1}^{(n)})$$

Steffenson's Method. Let {\Displays } (Pn) =  $P_{1}^{(0)}=g\left(P_{0}^{(0)}\right)$ 0)  $f_{(a)}^{0}$  $\rho_{2}^{(0)} = g(\rho_{1}^{(0)})$  $P_{i}^{(1)} = g(P_{o}^{(1)})$  $p_{2}^{(1)} = g(p_{1}^{(1)})$ 

 $P_{0}^{(1)} = \{ \triangle^{2} \} (P_{0}^{(0)})$   $= P_{0}^{(0)} - (P_{0}^{(0)} - P_{0}^{(0)})$   $= P_{0}^{(0)} - 2P_{0}^{(0)} + P_{0}^{(0)}$  $\rho_{1}^{(2)} = g(\rho_{0}^{(2)})$   $\rho_{2}^{(2)} = g(\rho_{1}^{(2)})$ 2) [ (2) = { 2} (10)

 $h \rho_0^{(n)} = \{ \Delta^2 \} \left( \rho_0^{(n-1)} \right)$  $f_{1}^{(n)} = g(f_{0}^{(n)})$   $f_{2}^{(n)} = g(f_{1}^{(n)})$ Thm (Steffensen's): Suppose g(x)=x has sol. p

with  $g'(p) \neq 1$ . If there exists d>0 s.t. g ∈ C³[P-S,P+S), then Steffensen's gives quadratic convergence for  $\rho_0 \in [\rho - \delta, \rho + \delta]$ 

Def: A polynomial of degree n has the form  $P(x) = a_n x^n + a_{n-1} x^{n-1} + \cdots + a_1 x + q_0, \quad q_n \neq 0$ 

Remarks: •  $q_i^2 S$  are coefficients of  $\rho$ •  $\rho(x) = 0$  is a polynomial of degree zero.

Q: Does Pn have any zeros? How many?

Thm: (Fundamental Thm of Algebra): If p(x) is a polynomial of degree  $n \ge 1$  with  $q_i \in C$ , i = 0, 1, ..., n. Then p(x) = 0 has at least one complex root

Cor 1: If  $\rho(x)$  is a polynomial of degree  $n \ge 1$  with complex coefficients, then there exist unique  $x_1, x_2, ..., x_k$  and unique  $m_1, m_2, ..., m_k$  satisfying  $\lim_{i \ge 1} m_i = n$  s.t. integers  $\lim_{i \ge 1} \rho(x) = a_n (x-x_1)^m (x-x_2)^{m_1} ... (x-x_k)^m k$ 

Remark: Cor 1 => collection of zeros of pn are unique, and it each zero xi counted as many times as its

multiplicity, mi, then Pn has exactly in zeros.

Cor 2: Let P(x) and Q(x) be polynomials of degree at most n

It x1, x2,..., Xk with k>n are distinct numbers

 $P(x_i) = Q(x_i) , i = 1,...,k$ Thun P(x) = Q(x) for all values of X.

Remark: 7) To show two polynomials of deg. < n are the same, we only need to show that they agree on n+1 values (will use in Chpt 3)

2) Proof sketch of Cor?:

R(x) = P(x) - Q(x), deg (R) & h, R hos n+1 100ts.

R(x) = P(x) - Q(x), deg (R) & h, R hos n+1 100ts.

 $\Rightarrow R(x) \equiv 0 \Rightarrow P(x) = Q(x)$ 

 $\rho(x_i) = x_i^n$  for  $x_1 = 1$ ,  $x_2 = 2$ , ...,  $x_n = n$ ,  $x_{n+1} = n+1$ then p(x) = X" (By Corz).

Ex: If P(x) with deg (P(x)) = n, and

(57)

Horner's Method:

Newton's Method to find zeros of PCX) => need to evaluate p(x) and p'(x) repeatedly

Note: p and p' are both polynomials

Q: How can we evaluate polynomials efficiently? Idea: nested evaluation / synthetic division

Recall from Chyt 1:

Ex: Evaluate  $f(x) = x^3 - 6.1x^2 + 3.2x + 1.5$  at x=4.71

-mult. - 2 - - - 2 - - - 1 - - -

Now consider nested formulation:

 $f(x) = (x^3 - 6.1x^2 + 3.2x) + 1.5$ 

 $= (x^2 - 6.1x + 3.2) \times + 1.5$ 

 $= ((x - 6.1) \times + 3.2) \times + 1.5$ 2 multiplications

3 additions

bk= 9k + bk+1 Xo,

More over, it Q(x) = 6, x -1 + 6, -, x - $\rho(x) = (x - x_0) Q(x) + 6.$ 

60 = P(X0)