## 0.1 Esercizi finali

- 1) Sia  $T_1$  la teoria dei tipi definita del tipo singoletto con le regole strutturali, definite in questo capitolo, incluse quelle di sostituzione e indebolimento. Allora stabilire se i seguenti termini sono tipabili come termini del tipo singoletto, secondo  $T_1$  e quali sono uguali definizionalmente.
  - $El_{N1}(*, *)$
  - $El_{N1}(x, *)$
  - $El_{N1}(*, y)$
  - $\mathrm{El}_{N1}(\mathbf{x},\,\mathbf{y})$
  - $\mathrm{El}_{N1}(\mathrm{El}_{N1}(*, y), \mathrm{El}_{N1}(x, *))$

## Soluzione

 $Per\ una\ maggiore\ comprensione\ delle\ derivazioni,\ ho\ ritenuto\ opportuno,\ ove\ necessario,\ spezzare\ l'albero\ in\ più\ parti.$ 

1

$$\begin{array}{c} \text{I-S} & \underbrace{ \left[ \right] \text{ cont} }_{\text{E-S}} & \underbrace{ \begin{array}{c} \text{F-S} \\ \text{N}_1 \text{ type} \left[ \right] \\ \text{z} \in \text{N}_1 \text{ cont} \end{array} }_{\text{F-S} \text{cont}} (z \in \text{N}_1) \notin \left[ \right] \\ \text{F-S} & \underbrace{ \begin{array}{c} \text{F-S} \\ \text{N}_1 \text{ type} \left[ \text{z} \in \text{N}_1 \right] \end{array} }_{\text{El}_{N_1}(*, *) \in N_1 \left[ \right]} \\ \text{I-S} & \underbrace{ \begin{array}{c} \left[ \right] \text{ cont} \\ \text{*} \in \text{N}_1 \left[ \right] \end{array} }_{\text{*}} \end{array}$$

Applicando la  $\beta$   $N_1$ -red allora  $\text{El}_{N_1}(*,*) \to_1 * \text{El}_{N_1}(*,*)$  è uguale definizionalmente.

 $\mathbf{2}$ 

$$\begin{array}{c} F\text{-S} & \frac{ \text{[] cont}}{N_1 \text{ type}[]} \\ \text{F-S} & \frac{\text{F-C} & \frac{N_1 \text{ type}[]}{N_1 \text{ type}[]}}{\frac{1}{x \in N_1 \text{ cont}}} (x \in N_1) \notin [] \\ \text{var} & \frac{F\text{-C} & \frac{N_1 \text{ type}[]}{x \in N_1 \text{ cont}}}{x \in N_1 \text{ (x \in N_1)}} (x \in N_1) \notin [] \\ \text{E-S} & \frac{F\text{-C} & \frac{N_1 \text{ type}[x \in N_1]}{x \in N_1 \text{ cont}}}{\frac{1}{x \in N_1 \text{ cont}}} (z \in N_1) \notin x \in N_1 \\ \text{E-I}_{N_1} & \text{type}[x \in N_1, z \in N_1] \\ & \text{E-I}_{N_1} & \text{type}[x \in N_1] \\ \end{array}$$

Applicando la  $\beta$   $N_1$ -red allora  $\text{El}_{N_1}(\mathbf{x}, *) \rightarrow_1 \text{El}_{N_1}(\mathbf{x}, *)$  non è uguale definizionalmente.

3

$$\begin{array}{c} F\text{-S} & \frac{ \left[ \; \right] \; \text{cont} }{N_1 \; \text{type}[\;]} \\ F\text{-C} & \frac{ \left[ \; \right] \; \text{cont} }{N_1 \; \text{type}[\;]} \\ F\text{-C} & \frac{ \left[ \; \right] \; \text{cont} }{y \in \; N_1 \; \text{cont}} \\ \text{I-S} & \frac{ \left[ \; \right] \; \text{cont} }{y \in \; N_1 \; \text{cont}} \\ \text{E-S} & \frac{ \left[ \; \right] \; \text{cont} }{y \in \; N_1 \; \text{cont}} \\ \text{E-S} & \frac{ \left[ \; \right] \; \text{cont} }{y \in \; N_1 \; \text{cont}} \\ \text{E-S} & \frac{ \left[ \; \right] \; \text{cont} }{N_1 \; \text{type}[y \in \; N_1]} \\ \text{E-S} & \frac{ \left[ \; \right] \; \text{cont} }{N_1 \; \text{type}[y \in \; N_1]} \\ \text{E-S} & \frac{ \left[ \; \right] \; \text{cont} }{N_1 \; \text{type}[y \in \; N_1]} \\ \text{E-S} & \frac{ \left[ \; \right] \; \text{cont} }{N_1 \; \text{type}[y \in \; N_1]} \\ \text{E-S} & \frac{ \left[ \; \right] \; \text{cont} }{N_1 \; \text{type}[y \in \; N_1]} \\ \text{E-S} & \frac{ \left[ \; \right] \; \text{cont} }{N_1 \; \text{type}[y \in \; N_1]} \\ \text{E-S} & \frac{ \left[ \; \right] \; \text{cont} }{N_1 \; \text{type}[y \in \; N_1]} \\ \text{E-S} & \frac{ \left[ \; \right] \; \text{cont} }{N_1 \; \text{type}[y \in \; N_1]} \\ \text{E-S} & \frac{ \left[ \; \right] \; \text{cont} }{N_1 \; \text{type}[y \in \; N_1]} \\ \text{E-S} & \frac{ \left[ \; \right] \; \text{cont} }{N_1 \; \text{type}[y \in \; N_1]} \\ \text{E-S} & \frac{ \left[ \; \right] \; \text{cont} }{N_1 \; \text{type}[y \in \; N_1]} \\ \text{E-S} & \frac{ \left[ \; \right] \; \text{cont} }{N_1 \; \text{type}[y \in \; N_1]} \\ \text{E-S} & \frac{ \left[ \; \right] \; \text{cont} }{N_1 \; \text{type}[y \in \; N_1]} \\ \text{E-S} & \frac{ \left[ \; \right] \; \text{cont} }{N_1 \; \text{type}[y \in \; N_1]} \\ \text{E-S} & \frac{ \left[ \; \right] \; \text{cont} }{N_1 \; \text{type}[y \in \; N_1]} \\ \text{E-S} & \frac{ \left[ \; \right] \; \text{cont} }{N_1 \; \text{type}[y \in \; N_1]} \\ \text{E-S} & \frac{ \left[ \; \right] \; \text{cont} }{N_1 \; \text{type}[y \in \; N_1]} \\ \text{E-S} & \frac{ \left[ \; \right] \; \text{cont} }{N_1 \; \text{type}[y \in \; N_1]} \\ \text{E-S} & \frac{ \left[ \; \right] \; \text{cont} }{N_1 \; \text{type}[y \in \; N_1]} \\ \text{E-S} & \frac{ \left[ \; \right] \; \text{cont} }{N_1 \; \text{type}[y \in \; N_1]} \\ \text{E-S} & \frac{ \left[ \; \right] \; \text{cont} }{N_1 \; \text{type}[y \in \; N_1]} \\ \text{E-S} & \frac{ \left[ \; \right] \; \text{cont} }{N_1 \; \text{type}[y \in \; N_1]} \\ \text{E-S} & \frac{ \left[ \; \right] \; \text{cont} }{N_1 \; \text{type}[y \in \; N_1]} \\ \text{E-S} & \frac{ \left[ \; \right] \; \text{cont} }{N_1 \; \text{type}[y \in \; N_1]} \\ \text{E-S} & \frac{ \left[ \; \right] \; \text{cont} }{N_1 \; \text{type}[y \in \; N_1]} \\ \text{E-S} & \frac{ \left[ \; \right] \; \text{cont} }{N_1 \; \text{type}[y \in \; N_1]} \\ \text{E-S} & \frac{ \left[ \; \right] \; \text{cont} }{N_1 \; \text{type}[y \in \; N_1]} \\ \text{E-S} & \frac{ \left[ \; \right] \; \text{cont} }{N_1 \; \text{type}[y \in \; N_1]} \\ \text{E-S} & \frac{ \left[ \; \right] \; \text{cont} }{N_$$

Applicando la  $\beta$   $N_1$ -red allora  $\text{El}_{N_1}(*, y) \rightarrow_1 y$   $\text{El}_{N_1}(*, y)$  è uguale definizionalmente.

4

Applicando la  $\beta$   $N_1$ -red allora  $\text{El}_{N_1}(\mathbf{x}, \mathbf{y}) \nrightarrow_1$   $\text{El}_{N_1}(\mathbf{x}, \mathbf{y})$  non è uguale definizionalmente.

5

$$\text{E-S} \frac{\mathbf{5}_{A}}{\text{El}_{N1}(*,\,\mathbf{y}) \in N_{1}[\mathbf{y} \in N_{1},\,\mathbf{x} \in N_{1}]} \frac{\mathbf{5}_{B}}{N_{1} \text{ type}[\mathbf{y} \in N_{1},\,\mathbf{x} \in N_{1},\,\mathbf{z} \in N_{1}]} \frac{\mathbf{5}_{C}}{\text{El}_{N1}(\mathbf{x},\,*) \in N_{1}[\mathbf{y} \in N_{1},\,\mathbf{x} \in N_{1}]}$$

$$= \frac{\mathbf{5}_{A}}{N_{1} \text{ type}[\mathbf{y} \in N_{1},\,\mathbf{x} \in N_{1},\,\mathbf{z} \in N_{1}]} \frac{\mathbf{5}_{C}}{\text{El}_{N1}(\mathbf{x},\,*) \in N_{1}[\mathbf{y} \in N_{1},\,\mathbf{x} \in N_{1}]}$$

 $\mathbf{5}_{A}$ 

$$\inf_{\text{ind-te}} \frac{F-S}{\frac{F-S}{N_1 \text{ type}[]}} \frac{[] \text{ cont}}{y \in N_1 \text{ type}[]}}{F-S} \frac{[] \text{ cont}}{y \in N_1 \text{ cont}} (y \in N_1) \notin []}{F-S} \frac{[] \text{ cont}}{N_1 \text{ type}[y \in N_1]}}{F-S} \frac{[] \text{ cont}}{N_1 \text{ type}[y \in N_1]}}{[] \text{ el}_{N_1}(*, y) \in N_1[y \in N_1]} (x \in N_1) \notin [] \text{ cont}}{[] \text{ el}_{N_1}(*, y) \in N_1[y \in N_1, x \in N_1]}}$$

 $\mathbf{5}_{B}$ 

$$F-S \frac{ \left[ \right] \operatorname{cont}}{N_1 \operatorname{type}[\right]} \\ F-C \frac{ \left[ \left[ \right] }{y \in N_1 \operatorname{cont}} \left( y \in N_1 \right) \notin \left[ \right] \\ F-S \frac{ \left[ \left[ \right] }{N_1 \operatorname{type}[y \in N_1]} \right] }{ \left[ \left[ \left[ \right] \right] } \\ F-S \frac{ \left[ \left[ \left[ \right] \right] }{y \in N_1, \ x \in N_1 \operatorname{cont}} \left( x \in N_1 \right) \notin y \in N_1 \right] } \\ F-C \frac{ \left[ \left[ \left[ \right] \right] }{N_1 \operatorname{type}[y \in N_1, \ x \in N_1]} \left( z \in N_1 \right) \notin \left( y \in N_1, \ x \in N_1 \right) \right] } \\ F-S \frac{ \left[ \left[ \left[ \right] \right] }{N_1 \operatorname{type}[y \in N_1, \ x \in N_1, \ z \in N_1]} \right] } \\$$

3

 $\mathbf{5}_C$ 

$$\inf \text{det} \frac{ \begin{bmatrix} F-S & \frac{[\ ] \ \text{cont}}{N_1 \ \text{type} \ [\ ]} \\ F-C & \frac{N_1 \ \text{type} \ [\ ]}{x \in N_1 \ \text{cont}} \\ F-S & \frac{N_1 \ \text{type} \ [\ ]}{x \in N_1 \ \text{cont}} \\ \hline \\ \text{ind-te} & \frac{El_{N1}(x,\,*) \in N_1[x \in N_1]}{\text{ex-te}} \underbrace{ \begin{bmatrix} \ ] \ \text{cont}}_{F-S} & \frac{F-S & \frac{[\ ] \ \text{cont}}{N_1 \ \text{type} \ [x \in N_1]}}_{F-S \ \text{toperator}} \\ & F-S & \frac{N_1 \ \text{type} \ [\ ]}{y \in N_1 \ \text{count}} \\ \hline \\ F-C & \frac{N_1 \ \text{type} \ [\ ]}{y \in N_1 \ \text{count}} \\ & F-C & \frac{N_1 \ \text{type} \ [\ ]}{y \in N_1 \ \text{count}} \\ \hline \\ \text{for } & \frac{N_1 \ \text{type} \ [\ ]}{y \in N_1 \ \text{toperator}} \\ \hline \\ \text{for } & \frac{N_1 \ \text{type} \ [\ ]}{y \in N_1 \ \text{toperator}} \\ \hline \\ \text{for } & \frac{N_1 \ \text{type} \ [\ ]}{y \in N_1 \ \text{toperator}} \\ \hline \\ \text{for } & \frac{N_1 \ \text{type} \ [\ ]}{y \in N_1 \ \text{toperator}} \\ \hline \\ \text{for } & \frac{N_1 \ \text{type} \ [\ ]}{y \in N_1 \ \text{toperator}} \\ \hline \\ \text{for } & \frac{N_1 \ \text{type} \ [\ ]}{y \in N_1 \ \text{toperator}} \\ \hline \\ \text{for } & \frac{N_1 \ \text{type} \ [\ ]}{y \in N_1 \ \text{toperator}} \\ \hline \\ \text{for } & \frac{N_1 \ \text{type} \ [\ ]}{y \in N_1 \ \text{toperator}} \\ \hline \\ \text{for } & \frac{N_1 \ \text{type} \ [\ ]}{y \in N_1 \ \text{toperator}} \\ \hline \\ \text{for } & \frac{N_1 \ \text{type} \ [\ ]}{y \in N_1 \ \text{toperator}} \\ \hline \\ \text{for } & \frac{N_1 \ \text{type} \ [\ ]}{y \in N_1 \ \text{toperator}} \\ \hline \\ \text{for } & \frac{N_1 \ \text{type} \ [\ ]}{y \in N_1 \ \text{toperator}} \\ \hline \\ \text{for } & \frac{N_1 \ \text{type} \ [\ ]}{y \in N_1 \ \text{toperator}} \\ \hline \\ \text{for } & \frac{N_1 \ \text{type} \ [\ ]}{y \in N_1 \ \text{toperator}} \\ \hline \\ \text{for } & \frac{N_1 \ \text{type} \ [\ ]}{y \in N_1 \ \text{toperator}} \\ \hline \\ \text{for } & \frac{N_1 \ \text{type} \ [\ ]}{y \in N_1 \ \text{toperator}} \\ \hline \\ \text{for } & \frac{N_1 \ \text{type} \ [\ ]}{y \in N_1 \ \text{toperator}} \\ \hline \\ \text{for } & \frac{N_1 \ \text{type} \ [\ ]}{y \in N_1 \ \text{toperator}} \\ \hline \\ \text{for } & \frac{N_1 \ \text{type} \ [\ ]}{y \in N_1 \ \text{toperator}} \\ \hline \\ \text{for } & \frac{N_1 \ \text{type} \ [\ ]}{y \in N_1 \ \text{toperator}} \\ \hline \\ \text{for } & \frac{N_1 \ \text{type} \ [\ ]}{y \in N_1 \ \text{toperator}} \\ \hline \\ \text{for } & \frac{N_1 \ \text{type} \ [\ ]}{y \in N_1 \ \text{toperator}} \\ \hline \\ \text{for } & \frac{N_1 \ \text{type} \ [\ ]}{y \in N_1 \ \text{toperator}} \\ \hline \\ \text{for } & \frac{N_1 \ \text{type} \ [\ ]}{y \in N_1 \ \text{toperator}} \\ \hline \\ \text{for } & \frac{N$$

Per le premesse  $El_{N1}(*, y) \in N_1[y \in N_1]$  e  $El_{N1}(x, *) \in N_1[x \in N_1]$  ho già dimostrato sopra (in 3 e 2) la loro tipabilità per il tipo singoletto. Applicando la  $red_I$  e  $\beta$   $N_1$ -red allora  $El_{N1}(El_{N1}(*, y), El_{N1}(x, *)) \rightarrow_1 El_{N1}(y, y)$ 

Più nel dettaglio la riduzione è la seguente

$$\operatorname{red}_{I} \frac{\beta \operatorname{N}_{1}\operatorname{-red} \overline{\operatorname{El}_{N1}(*, y) \to_{1} y}}{\operatorname{El}_{N1}(\operatorname{El}_{N1}(*, y), \operatorname{El}_{N1}(x, *)) \to_{1} \operatorname{El}_{N1}(y, \operatorname{El}_{N1}(x, *))}$$

 $\mathrm{El}_{N1}(\mathrm{El}_{N1}(*,\,\mathrm{y}),\,\mathrm{El}_{N1}(\mathrm{x},\,*))$  è uguale definizionalmente.

2) Definire  $w + 2 \in Nat[w \in Nat]$ , ove 2 è l'abbreviazione del termine ottenuto applicando  $2 \equiv succ(succ(0))$ .

## Soluzione

La ricorsione la faccio su w, usando lo schema di ricorsione primitiva, vale che w  $+2 \equiv \text{El}_{Nat}(w, 2, (x,y).\text{succ}(y)) \in \text{Nat}[w \in \text{Nat}].$ 

$$\begin{array}{c} F\text{-Nat} \xrightarrow{ \left[ \begin{array}{c} \text{Cont} \\ \text{Nat type} \left[ \begin{array}{c} \text{I} \end{array} \right] } \text{Cont} \\ F\text{-Nat} \xrightarrow{ \left[ \begin{array}{c} \text{Nat type} \left[ \begin{array}{c} \text{I} \end{array} \right] } \text{Cont} \\ \hline \text{F-Nat} \xrightarrow{ \left[ \begin{array}{c} \text{Nat type} \left[ \begin{array}{c} \text{I} \end{array} \right] } \text{Cont} \\ \hline \text{I}_2\text{-Nat} \xrightarrow{ \left[ \begin{array}{c} \text{I} \end{array} \right] } \hline \text{I}_2\text{-Nat} \left[ \begin{array}{c} \text{I} \end{array} \right] \\ \hline \text{II}_2\text{-Nat} \left[ \begin{array}{c} \text{II} \end{array} \right] \\ \hline \text{II}_2\text{-Nat} \left[ \begin{array}{c} \text{II$$

Dimostrazione di correttezza di  $El_{Nat}(w, 2, (x,y).succ(y)) \in Nat[w \in Nat]$ 

- $\mathrm{El}_{Nat}(0, 2, (\mathbf{x}, \mathbf{y}).\mathrm{succ}(\mathbf{y})) \rightarrow_1 2 \ \mathrm{per} \ \beta_{1Nat}\text{-red}$
- $\text{El}_{Nat}(\text{succ}(\text{m}), 2, (\text{x,y}).\text{succ}(\text{y})) \rightarrow_1 \text{succ}(\text{El}_{Nat}(\text{m}, 2, (\text{x,y}).\text{succ}(\text{y}))) \text{ per } \beta_{2Nat}\text{-}red \Rightarrow \text{per } \text{m} = 0 \equiv \text{succ}(\text{El}_{Nat}(0, 2, (\text{x,y}).\text{succ}(\text{y}))) \rightarrow_1 \text{succ}(2) \in \text{Nat} = 3 \text{ (dal punto precedente)}.$
- 4) Si dimostri che esiste un proof-term pf del tipo

$$ext{pf} \in ext{Id}( ext{N}_1, ext{x}, st)[ ext{x} \in ext{N}_1]$$

## Soluzione

$$z \equiv x$$
  
 $El_{N1}(z,c) \equiv pf \equiv El_{N1}(x,x.id(*))$ 

$$\begin{array}{l} M(z) \equiv \operatorname{Id}(N_1,\!x,\!*) \\ c \in M(*)[\Gamma] \equiv \operatorname{id}(*) \in \operatorname{Id}(N_1,\!*,\!*)[x \in N_1] \end{array}$$

$$F-Id = \frac{\frac{1}{N_1 \text{ type}[x \in N_1]} \frac{1}{x \in N_1[x \in N_1]} \frac{1}{x \in N_1[x \in N_1]} \frac{1}{x \in N_1[x \in N_1]} = \frac{1-S \frac{[] \text{ cont}}{x \in N_1[]}}{1-Id \frac{id(x) \in Id(N_1, x, x)[]}{id(x) \in Id(N_1, x, x)[]}}$$

$$= \frac{E-S_{dip}}{El_{N_1}(x, x.id(x)) \in Id(N_1, x, x)[x \in N_1]}$$