

5.3 Code Generation Rules

c	<pre> dest <- fresh_tmp() code.add("li dest, c") return dest </pre>
x	<pre> # get the temporary associated to x. reg <- symbol_table[x] return reg </pre>
$e_1 + e_2$	<pre> t1 <- GenCodeExpr(e_1) t2 <- GenCodeExpr(e_2) dest <- fresh_tmp() code.add("add dest, t1, t2") return dest </pre>
$e_1 - e_2$	<pre> t1 <- GenCodeExpr(e_1) t2 <- GenCodeExpr(e_2) dest <- fresh_tmp() code.add("sub dest, t1, t2") return dest </pre>
true	<pre> dest <- fresh_tmp() code.add("li dest, 1") return dest </pre>
$e_1 < e_2$	<pre> dest <- fresh_tmp() t1 <- GenCodeExpr(e1) t2 <- GenCodeExpr(e2) endrel <- new_label() code.add("li dest, 0") # if t1>=t2 jump to endrel code.add("bge endrel, t1, t2") code.add("li dest, 1") code.addLabel(endrel) return dest </pre>

Figure 5.1: 3@ Code generation for numerical or Boolean expressions

<code>x = e</code>	<pre> dest <- GenCodeExpr(e) loc <- symbol_table[x] code.add("mv loc, dest") </pre>
<code>S1; S2</code>	<pre> # Just concatenate codes GenCodeSmt(S1) GenCodeSmt(S2) </pre>
<code>if b then S1 else S2</code>	<pre> lelse <- new_label() lendif <- new_label() t1 <- GenCodeExpr(b) #if the condition is false, jump to else code.add("beq lelse, t1, 0") GenCodeSmt(S1) # then code.add("j lendif") code.addLabel(lelse) GenCodeSmt(S2) # else code.addLabel(lendif) </pre>
<code>while(b){ S }</code>	<pre> ltest <- new_label() lendwhile <- new_label() code.addLabel(ltest) t1 <- GenCodeExpr(b) code.add("beq lendwhile, t1, 0") GenCodeSmt(S) # execute S code.add("j ltest") # and jump to the test code.addLabel(lendwhile) # else it is done. </pre>

Figure 5.2: 3@ Code generation for Statements

5.4 Allocations

EXERCISE #5 ► Prepare the lab: allocations

After code generation, we obtain the following code:

```

li temp_0, 42
li temp_1, 1
add temp_2, temp_1, temp_0

```

- Compute the naive allocation and rewrite the code accordingly.
- Compute the all-in-mem allocation and rewrite the code accordingly.