

# Theory of Automata

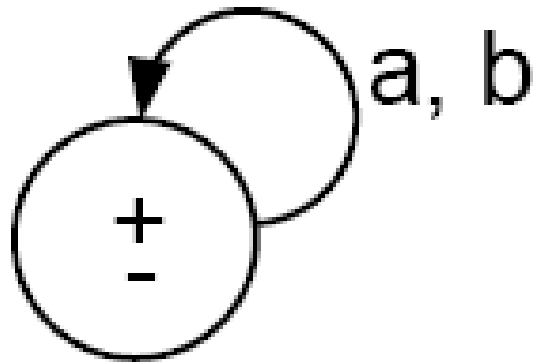
Shakir Ullah Shah

Lecture 5

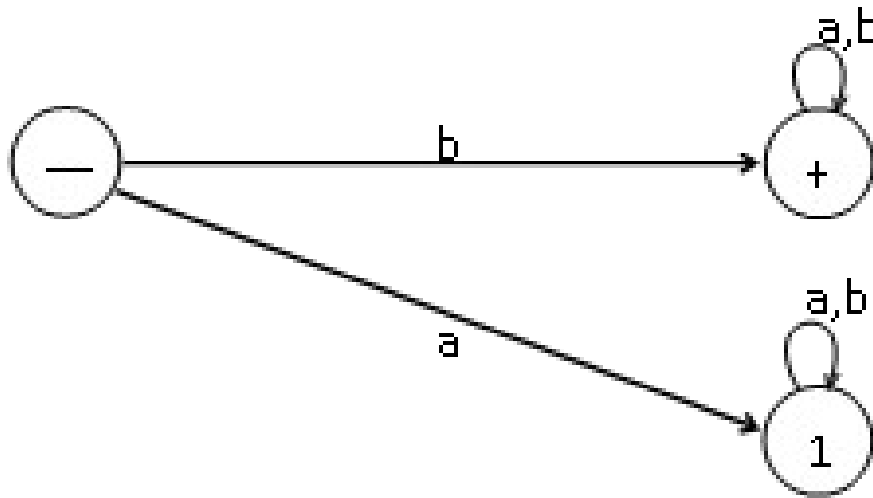
# FA and their Languages

- We will study FA from two different angles:
  1. Given a language, can we build a machine for it?
  2. Given a machine, can we deduce its language?
- **Note:**
  - **Every state has as many outgoing edges as there are letters in the alphabet.**
  - **It is possible for a state to have no incoming edges or to have many.**

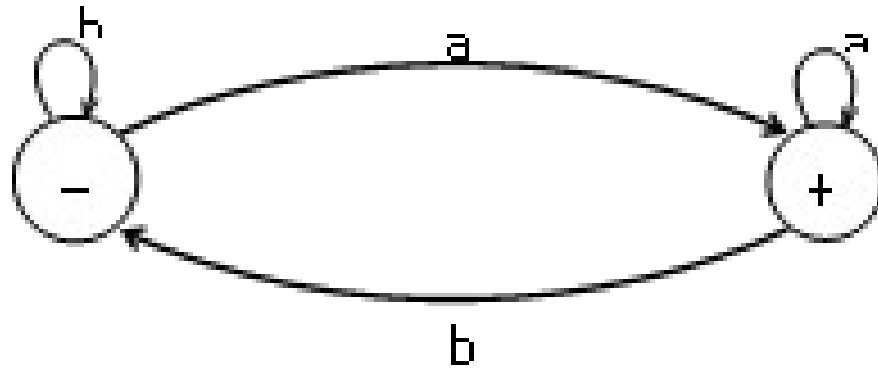
$(a + b)^*$



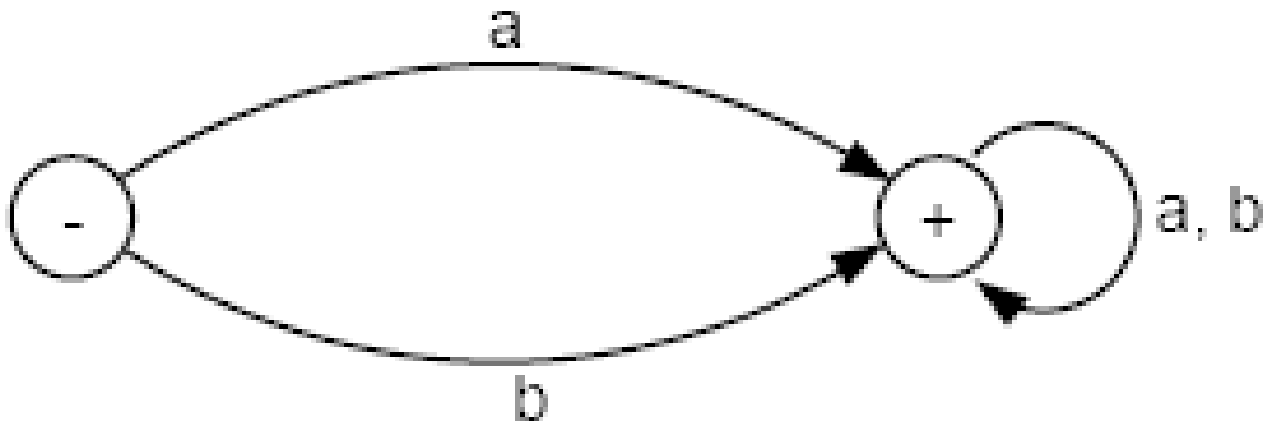
starting with  $b = b(a + b)^*$



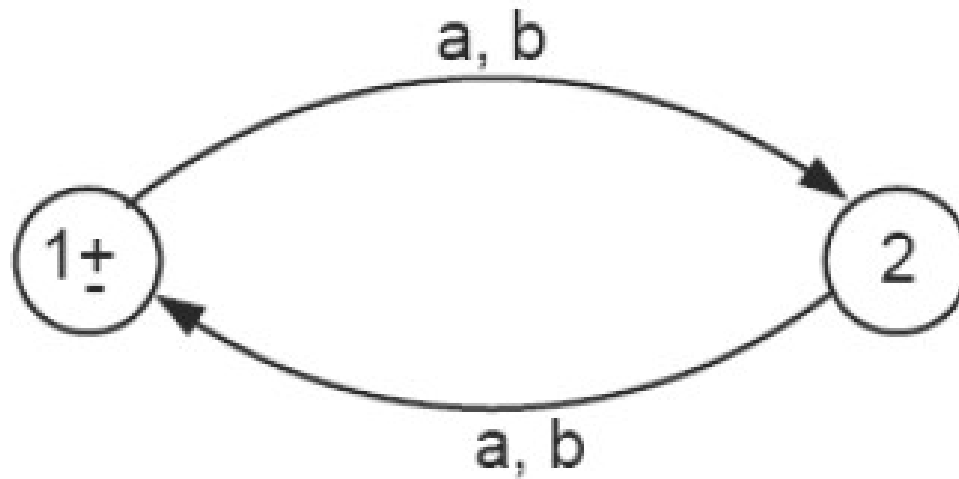
ending in  $a = (a + b)^* a$



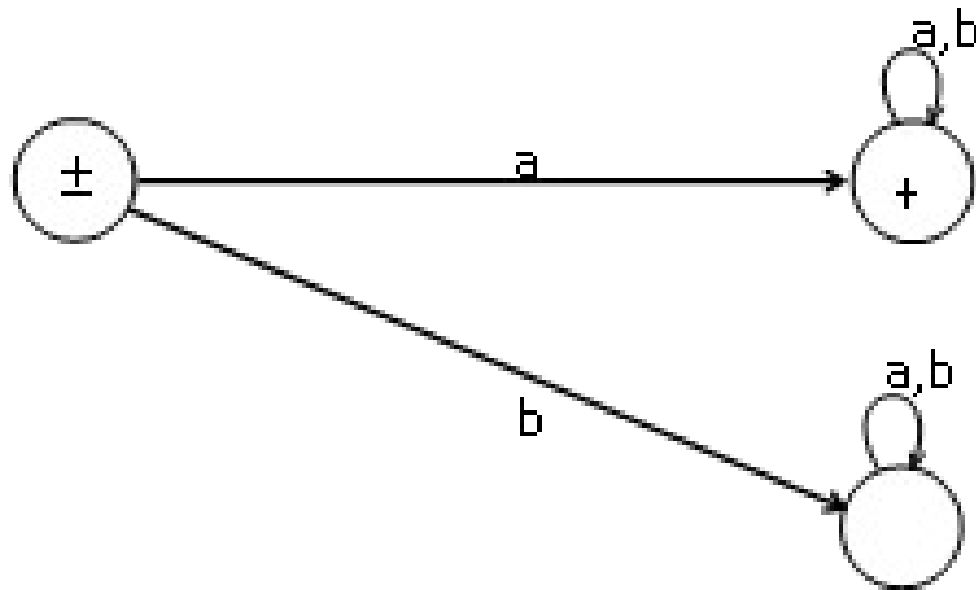
$$(a+b)(a+b)^*$$



# Even length

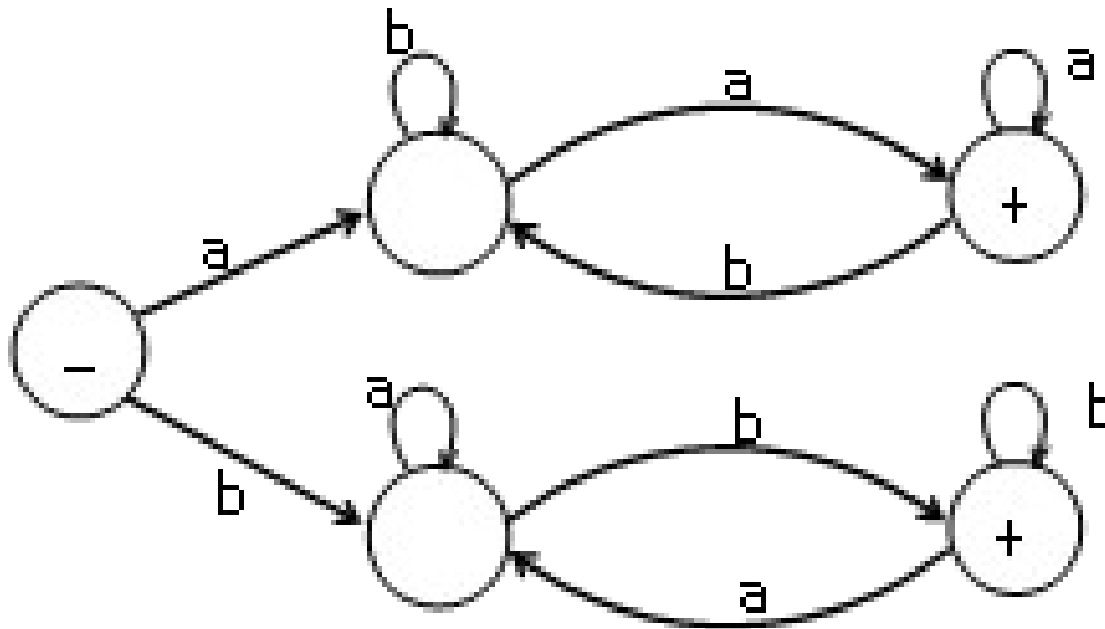


not beginning with  $(ab)^* + \Lambda$

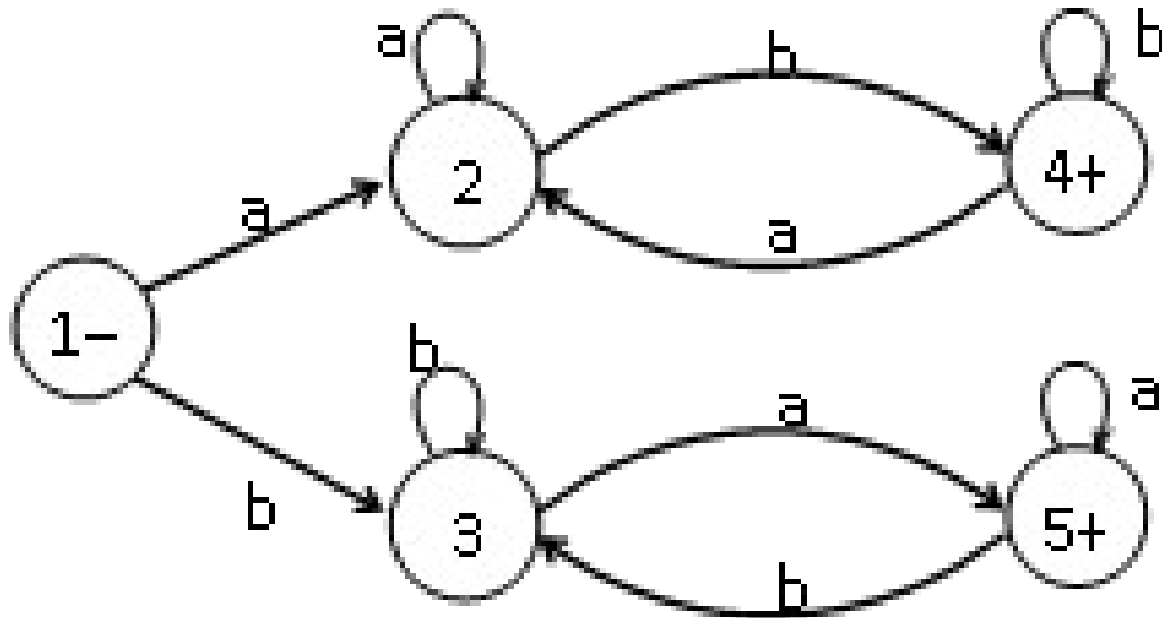




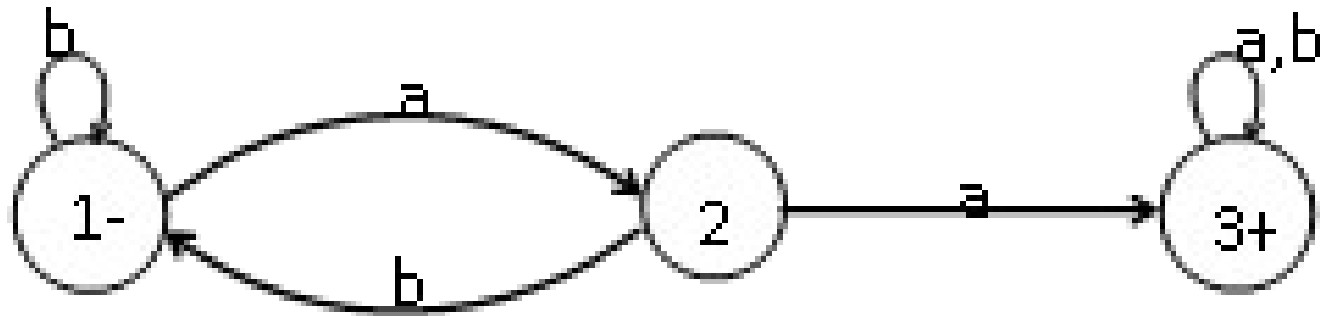
beginning with and ending in same  
letters.  $a(a + b)^*a + b(a + b)^*b$



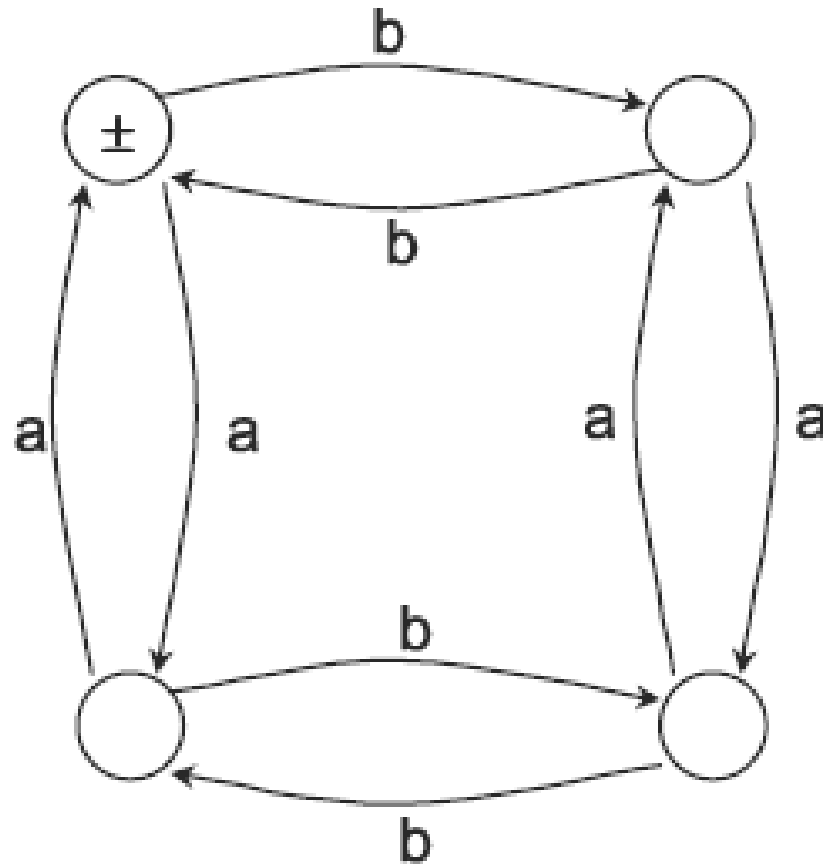
beginning with and ending in different letters.  $a(a + b)^*b + b(a + b)^*a$



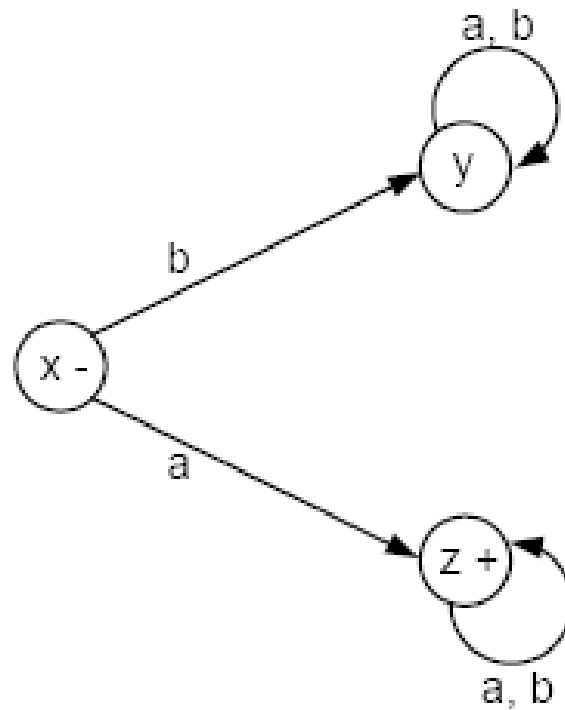
containing double  $a = (a+b)^* (aa)(a+b)^*$ .



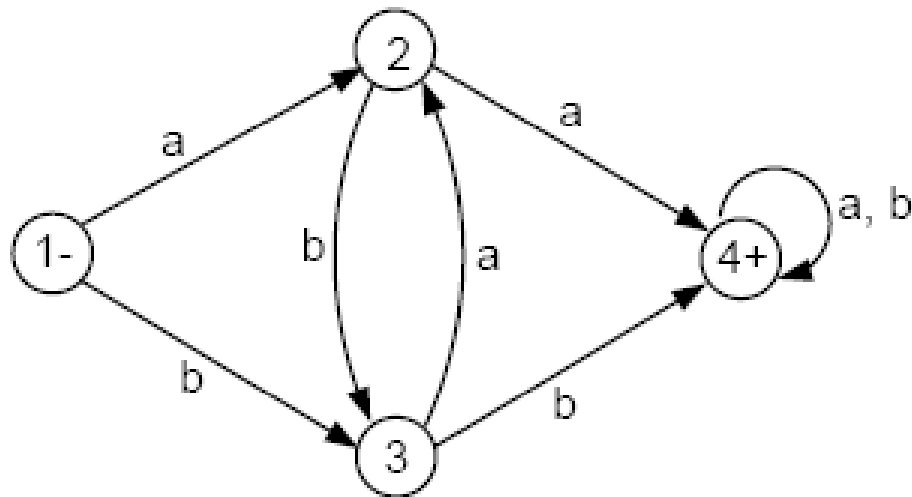
# ***EVEN-EVEN***



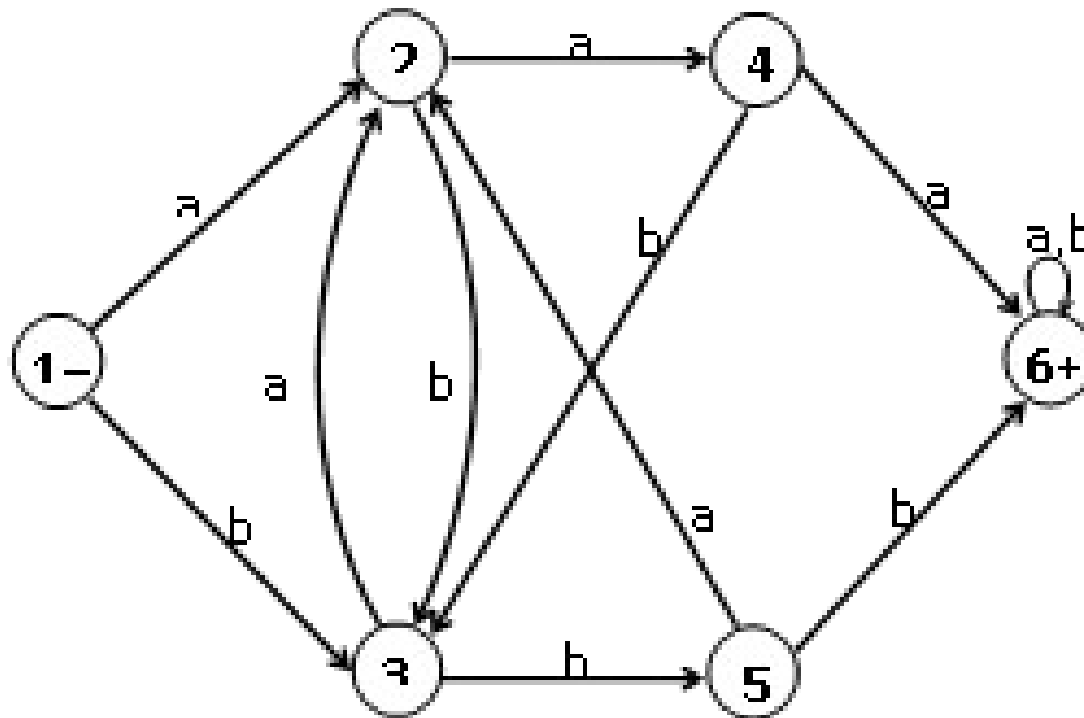
$a(a + b)^*$



$(a + b)^*(aa + bb)(a + b)^*$

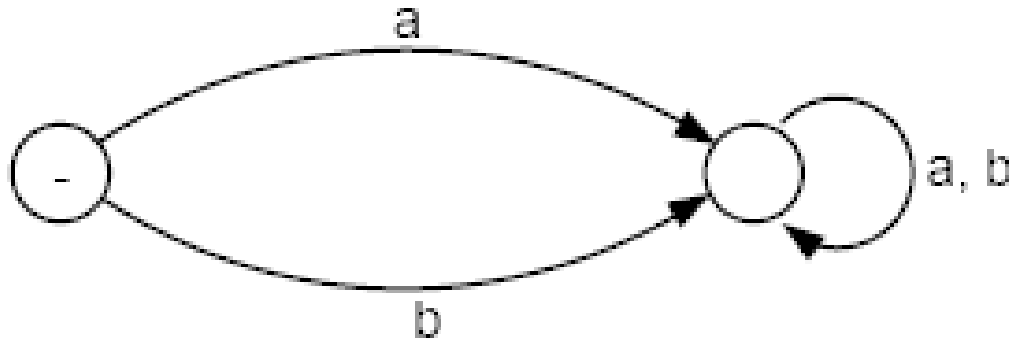


**containing a triple letter**, either  
*aaa* or *bbb*



# FAs that accept no language

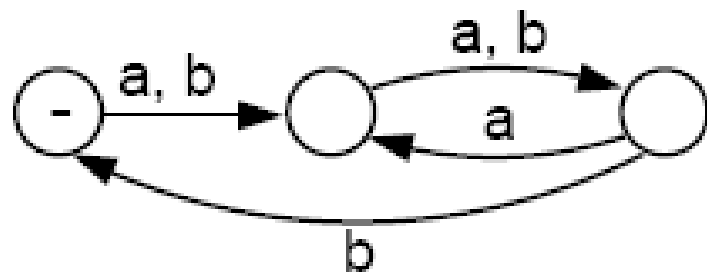
- There are FAs that accept no language. These are of two types:
- The first type includes FAs that have no final states, such as



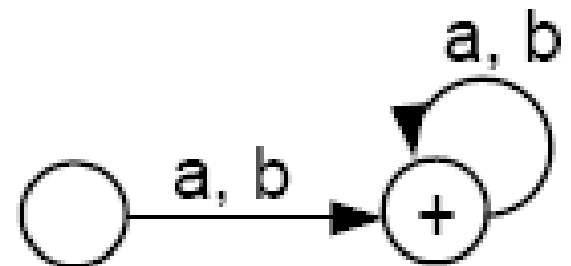


# FAs that accept no language

- The second type include FAs of which the final states can not be reached from the start state.
- This may be either because the diagram is in two separate components. In this case



is disconnected below:



- Or it is because the final state has no incoming edges, as shown below:

