

Simple Relational Correctness Proofs for Static Analyses and Program Transformations

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Soundness of Program optimization

Lot of work on functional languages especially in

- formalization
- validation

Few work on imperative programming languages

- seems trivial
- ... but it's not



Some notations

- $\mathbb{V} = \{X, Y, \dots\}$ a set of variables
- $n \in \mathbb{Z}$ a number, $b \in \mathbb{B}$ a boolean literal
- $iop \in \{+, -, \times, \dots\} \subseteq \mathbb{Z} \times \mathbb{Z} \rightarrow \mathbb{Z}$ an integer operation
- $bop \in \{<, =, \dots\} \subseteq \mathbb{Z} \times \mathbb{Z} \rightarrow \mathbb{B}$ an integer to boolean operation
- $lop \in \{\wedge, \vee, \dots\} \subseteq \mathbb{B} \times \mathbb{B} \rightarrow \mathbb{B}$ a logical operation
- $E := n | X | E \ iop \ E$ integer expressions
- $B := b | E \ bop \ E | \text{not } B | B \ lop \ B$ boolean expressions
- $C := \text{skip} | X := E | C; C | \text{if } B \text{ then } C \text{ else } C | \text{while } B \text{ do } C$ commands
- \supset
- $\llbracket E \rrbracket_S$ evaluation
- ddd

ss

(1)

(PE-COND)

$$\frac{\sqsubseteq \rho}{\hat{v}\rho}$$