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CS 2341

Prof. Fontenot

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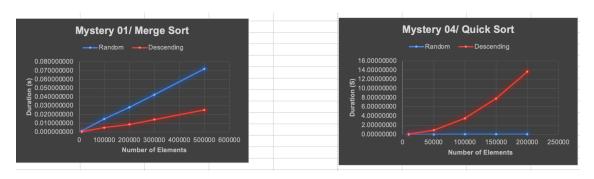
Lab Report of Project Assignment 03: "Mystery Sort"

As an introduction, my name is Ethan Le, and I am a student at SMU studying for a bachelor's degree in computer science. My goal is to delve into the field of software engineering, because I want to explore the design and implementation of computer algorithms. This semester I took a course on Data Structure and came across project 3 "Mystery Sort". The project assignment requires to determine the type of sorting in each mystery function that has been implemented. I would first separate the sorting function into two groups based on average run time then further study them when running in their best and worst case scenario.

I started the distinguishing process by separating the mystery function into two groups using run time of O(nlogn) and $O(n^2)$ in the average case. Merge sort and quicksort would be O(nlogn) as they are efficient sorting algorithms while selection sort, bubble sort, and insertion sort would be in place of $O(n^2)$. I would first create a randomized array of integers with size ranging from 1000 to 500000 and then run those mystery functions on these arrays. To increase the accuracy of my test, I run the sorting function three times and take the average time in case there are variances in run time. Based on what I observe, the run time of mystery 1 and 4 are 0.028106s and 0.016059s, which are both less than 1 second in the case of 500000. Compared to the run time of the other mystery sorts, which are all several minutes long, I can easily identify that mystery sort 1 and 4 must be the group with run time of O(nlogn).

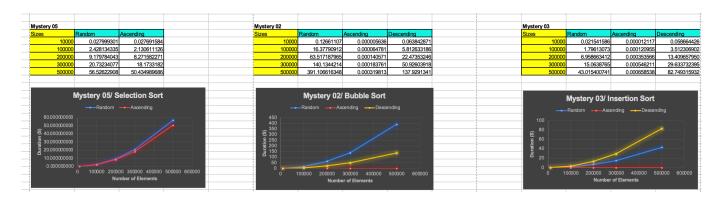
| Mystery 01 | | | Mystery 04 | | |
|------------|-------------|-------------|------------|------------|-------------|
| Sizes | Random | Descending | Sizes | Random | Descending |
| 10000 | 0.001294315 | 0.000411183 | 10000 | 0.00056808 | 0.04106702 |
| 100000 | 0.014985683 | 0.005031366 | 50000 | 0.00358493 | 0.89904065 |
| 200000 | 0.028106890 | 0.008590262 | 100000 | 0.00795571 | 3.49296730 |
| 300000 | 0.042574468 | 0.014167441 | 150000 | 0.01208633 | 7.82219026 |
| 500000 | 0.072072206 | 0.025143856 | 200000 | 0.01605936 | 13.64281474 |
| | | | | | |

The next procedure was to determine the type of sorts between mystery 01 and mystery 04. In the worst case run time of quicksort, where the list is sorted in reverse order, the run time of the algorithm is very close to O(n^2). Therefore, I create an array of integers that is in reverse order and run these two mystery functions on this array. The Mystery 04 function is recognized as a quick sort because it has the slowest running time when sorting in this type of array. The data plot graph of Mystery 04 shows that it takes way more time to build and run through the number of elements than the Mystery 01. At a size of 200 000, mystery 04 had a significant amount of 13.64 seconds in descending order while the mystery 01 function took about 0.0086 seconds. Therefore, the mystery 01 function is identified as a merge sort Hence, mystery 01 represents the merge sort and mystery 04 represents the quick sort.



At this point, mystery 02, mystery 03, and mystery 05 had the slower run time. For the group O(n^2), there are two procedures to find the sorting algorithms. The first procedure is to determine the selection sort by running the number of elements in a sorted order. This is because selection sort is a consistent sorting algorithm, which means that even in the best case scenario of a sorted array, the run time is still O(n^2). Based on the observations of the graph, the ascending column of Mystery 05 shows a higher upward trend than the other two graphs. Consequently, mystery 05 function is the selection sort since it takes more time to run through the number of elements. Within almost sorted data, bubble sort and insertion sort need very few swaps. However, selection sort requires the same amount of search process even in almost sorted data.

Quan 3

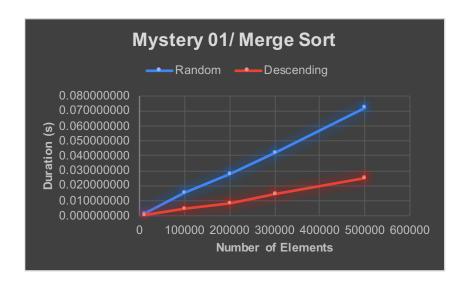


The second procedure is to determine the bubble sort and insertion sort by running the number of elements in reverse order. The chart of Mystery 02 shows that during the descent process, the running time is longer than that of Mystery 03. In addition, bubble sort performs more swap operations than the insertion sort. The higher number of swaps leads to a higher runtime for the bubble sort algorithm. Thus, the difference in runtime grows as the number of elements to be sorted increases on a random list.

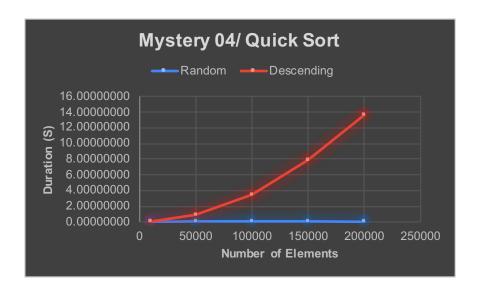
Therefore, the mystery 02 function is recognized as a bubble sort because it has the slower running time when sorted in the reverse order.

In conclusion, after, separate the sorting function into two groups based on average run time, then further study them when running in their best and worst case scenario. For O(nlogn), mystery 1 is identified as merge sort and mystery 4 is quick sort. Hence, for $O(n^2)$, mystery 5 is identified as selection sort, mystery 2 is bubble sort and mystery 3 is insertion sort.

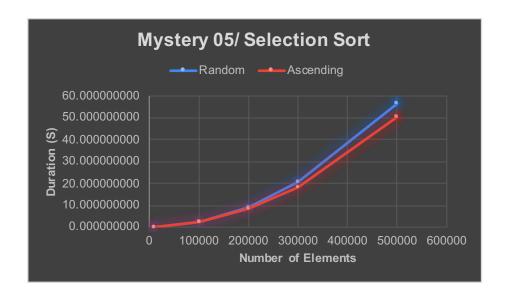
| Sizes | Random | Descending |
|--------|-------------|-------------|
| 10000 | 0.001294315 | 0.000411183 |
| 100000 | 0.014985683 | 0.005031366 |
| 200000 | 0.028106890 | 0.008590262 |
| 300000 | 0.042574468 | 0.014167441 |
| 500000 | 0.072072206 | 0.025143856 |



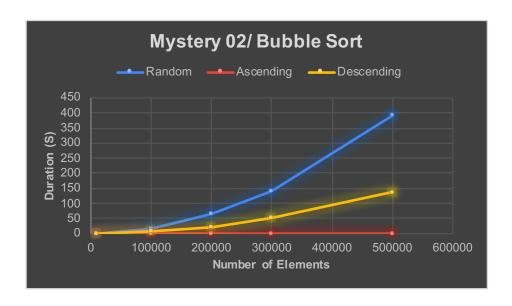
| Sizes | Random | Descending |
|--------|------------|--------------|
| 10000 | 0.00056808 | 0.041067020 |
| 50000 | 0.00358493 | 0.899040658 |
| 100000 | 0.00795571 | 3.492967303 |
| 150000 | 0.01208633 | 7.822190262 |
| 200000 | 0.01605936 | 13.642814743 |



| Sizes | Random | Ascending |
|--------|-------------|--------------|
| 10000 | 0.027999301 | 0.027691584 |
| 100000 | 2.428134335 | 2.130611126 |
| 200000 | 9.179784043 | 8.271582271 |
| 300000 | 20.73234077 | 18.1733182 |
| 500000 | 56.52622908 | 50.434989686 |



| Sizes | | Random | Ascending | Descending | |
|-------|----|---------------|-------------|-------------|--|
| 1000 | 00 | 0.12661107 | 0.000005636 | 0.063842871 | |
| 1000 | 00 | 16.37790912 | 0.000064781 | 5.812633186 | |
| 2000 | 00 | 63.517187965 | 0.000140571 | 22.47353246 | |
| 3000 | 00 | 140.1344214 | 0.000183761 | 50.92603918 | |
| 5000 | 00 | 391.106616348 | 0.000319813 | 137.9291341 | |



| Sizes | Random | Ascending | Descending | |
|--------|--------------|-------------|--------------|--|
| 10000 | 0.021541586 | 0.000012117 | 0.058664426 | |
| 100000 | 1.79613073 | 0.000120955 | 3.512306902 | |
| 200000 | 6.958663412 | 0.000353566 | 13.409657950 | |
| 300000 | 15.0638765 | 0.000546211 | 29.633732395 | |
| 500000 | 43.015400741 | 0.000658538 | 82.749315932 | |

