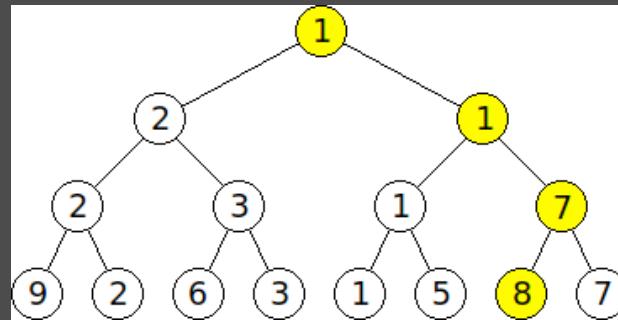
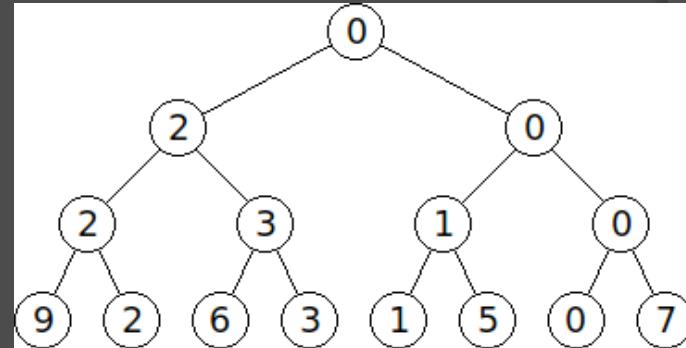


Advanced Computer Contest Preparation
Lecture 23

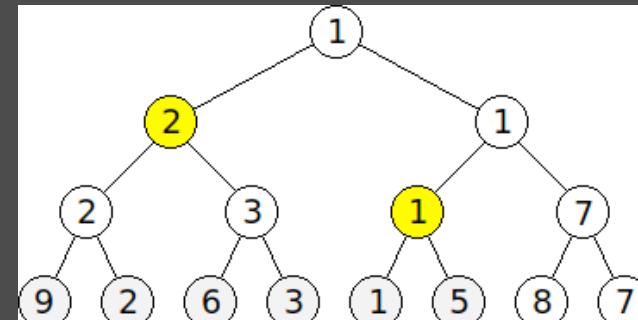
SEGMENT TREE LAZY PROPAGATION

Review

- Segment trees
 - Build
 - Point update
 - Point and range query



Update 7th element



Query [1,6]

Range Updates

- Range update means updating several consecutive elements at once
- Based on what we know, how do we update a range?
 - Call point update on each leaf node individually
 - Time complexity: $O(N \log N)$
 - Improvement: start at the leftmost node and follow a path similar to the one used for **build()** until the rightmost node
 - Time complexity: $O(N)$

Lazy Principle

- “We don’t need to do something unless it is absolutely necessary to do so”
- Example:
 - You have an array of N integers
 - 2 operations:
 - Add a constant to all elements
 - Get the x^{th} element
 - When updating, do not update each individual element
 - Instead, keep track of total updates and apply to individual element when query is required



Lazy Propagation

- The application of the lazy principle on segment trees
- We do not update nodes unless it is absolutely necessary to do so
- We mark nodes as lazy, meaning that it should be updated later, only if necessary

Lazy Value

- By setting a node's lazy value, we mark it as lazy
- Each answer in a node should have an associated lazy value
- ***Lazy value is generally what the node's answer should be after an actual update***
- Each lazy value should have an “unset” value, a value specifying that the node is not marked as lazy
 - Examples: 0, $+\infty$, $-\infty$
- This value should never be a possible update value

Lazy Propagation

- When is it absolutely necessary to update a node's value using its lazy value?
 - When a node is visited for any reason (update/query)
 - This includes after setting the initial lazy value!

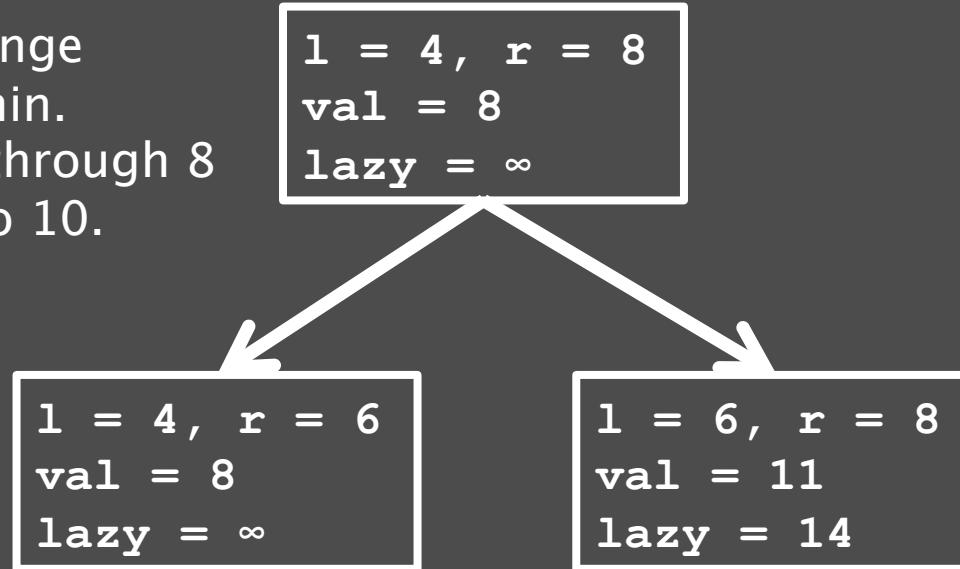
Lazy Propagation

- How do we update a node marked as lazy?
 - Node's values should reflect the answer as if the entire subtree was fully updated
 - Update children's lazy values, if the node has children
 - Method of updating node's values and children's lazy values using the current lazy value depends on the problem
 - Unset lazy values of current node
- We call this method **push_down()**

push_down

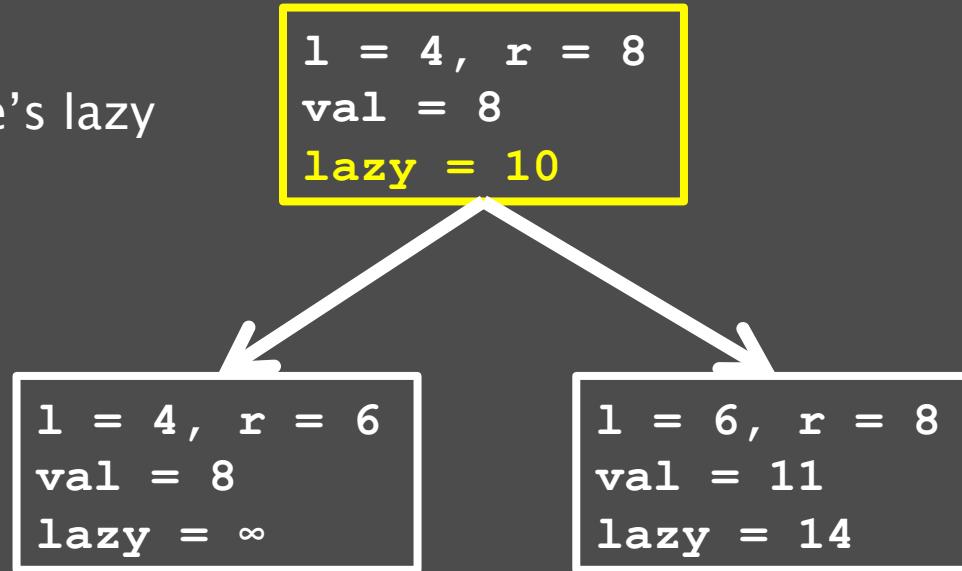
Example:

Segment tree has range
replace and query min.
Update elements 4 through 8
by changing them to 10.



push_down

Update node's lazy value to 10.



push_down

Begin `push_down` by updating node's `val`. The new value should reflect the answer if the entire range was updated. In this case, every number should be equal to 10.

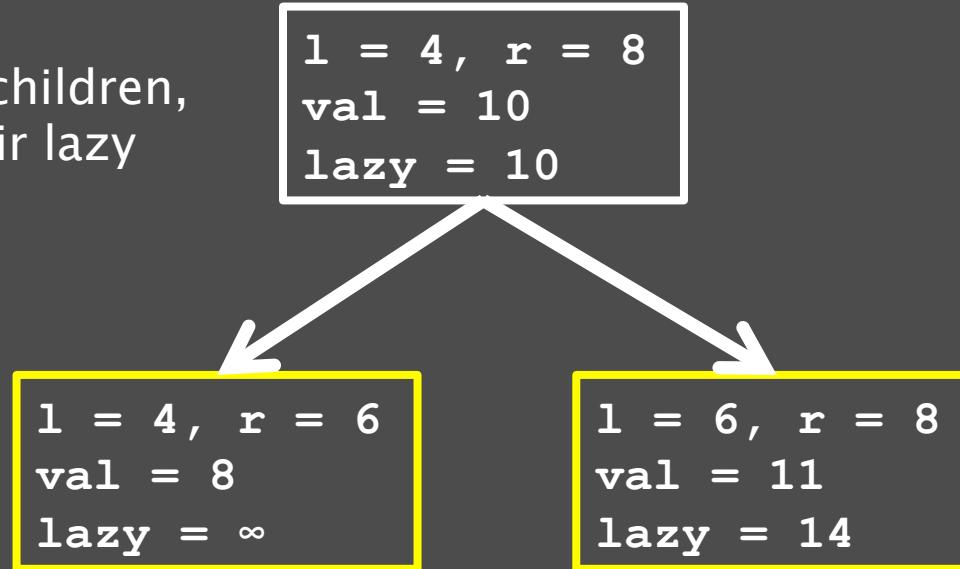
```
l = 4, r = 8  
val = 10  
lazy = 10
```

```
l = 4, r = 6  
val = 8  
lazy = ∞
```

```
l = 6, r = 8  
val = 11  
lazy = 14
```

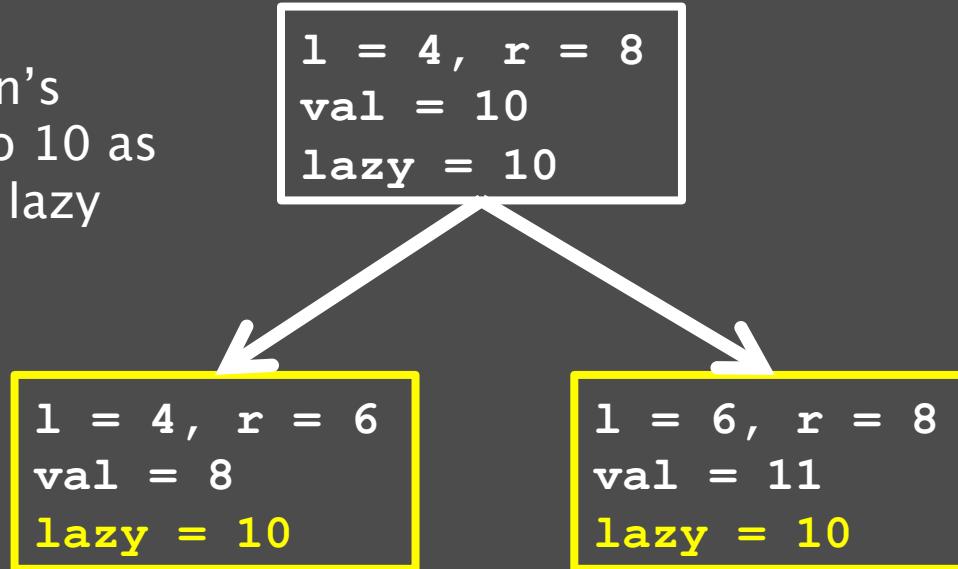
push_down

Since the node has children,
we must update their lazy
values.



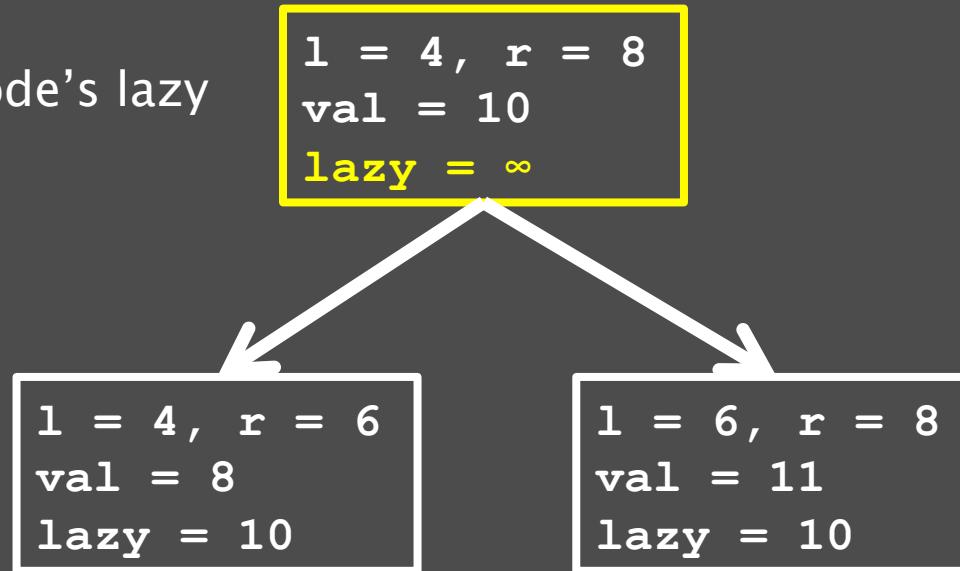
push_down

We want our children's values to be equal to 10 as well, so we set their lazy values.



push_down

Finally, unset the node's lazy value.



push_down Functionality Examples

- Replace, min:
 - `val = lazy`, replace children lazy
- Add, min:
 - `val += lazy`, add to children lazy
- Replace, max:
 - `val = lazy`, replace children lazy
- Add, max:
 - `val += lazy`, add to children lazy

push_down Functionality Examples

- Replace, sum:
 - `val = lazy*range_size`, replace children lazy
- Add, sum:
 - `val += lazy*range_size`, add to children lazy
- Replace, GCD:
 - `val = lazy`, replace children lazy

Lazy Propagation

- `push_down()` functionality is not always straightforward
- Example: add value to range, range query GCD
 - Possible to do, but not as simple as the other examples
 - Requires segment tree of a difference array

Lazy Propagation

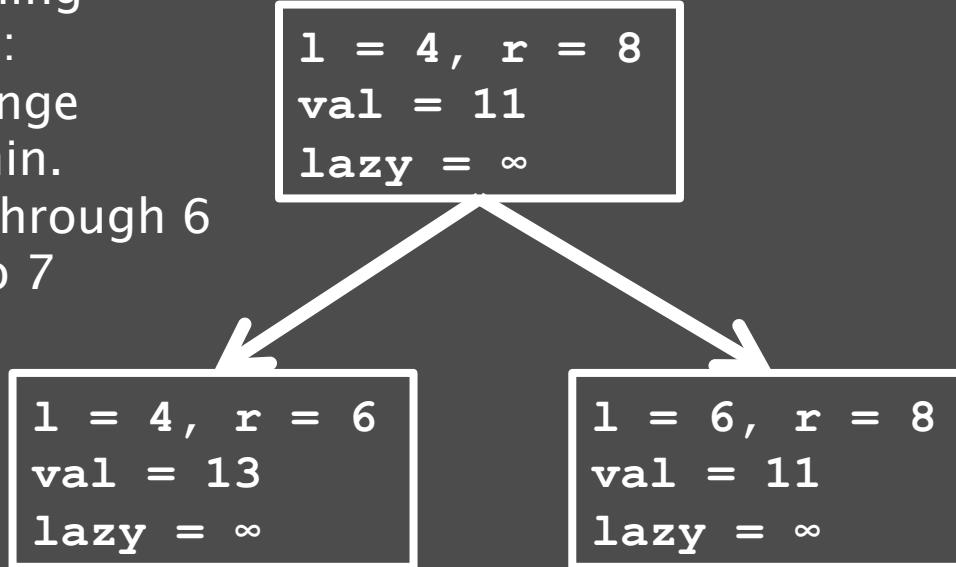
- Warning: in order for `push_up(nd)` to function as expected, `nd`'s children must be updated
 - Otherwise, `push_up()` will use `nd`'s children's pre-update values
- `push_down` must be called on `nd`'s children
- This is the reason why `push_down` is necessary after setting a node's lazy value – `push_up` will be called on the parent and it will use old values if we did not call `push_down`
- Also an issue if point functions used, since they do not typically visit (and thus `push_down`) both children

push_up/push_down Issue

Example of not pushing down after updating:

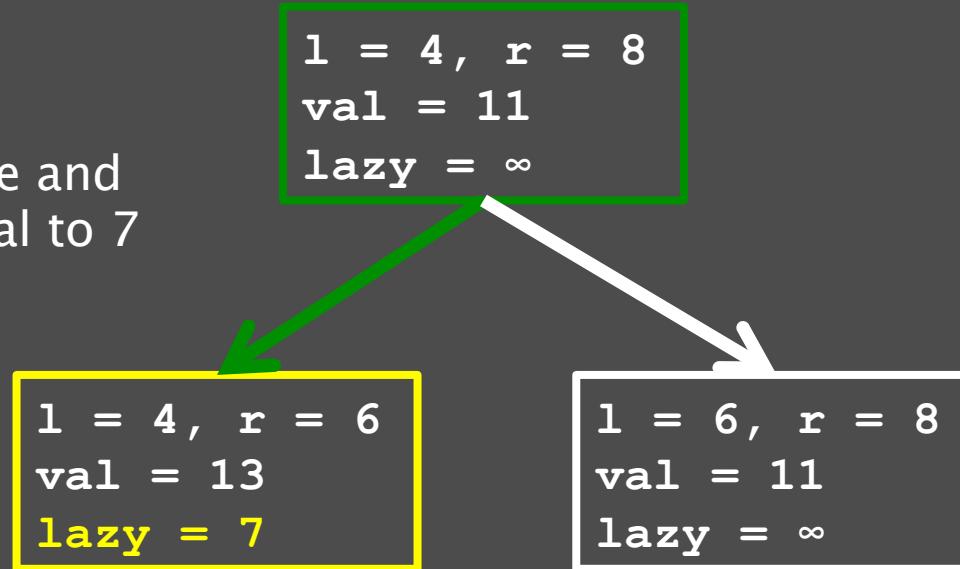
Segment tree has range replace and query min.

Update elements 4 through 6 by changing them to 7



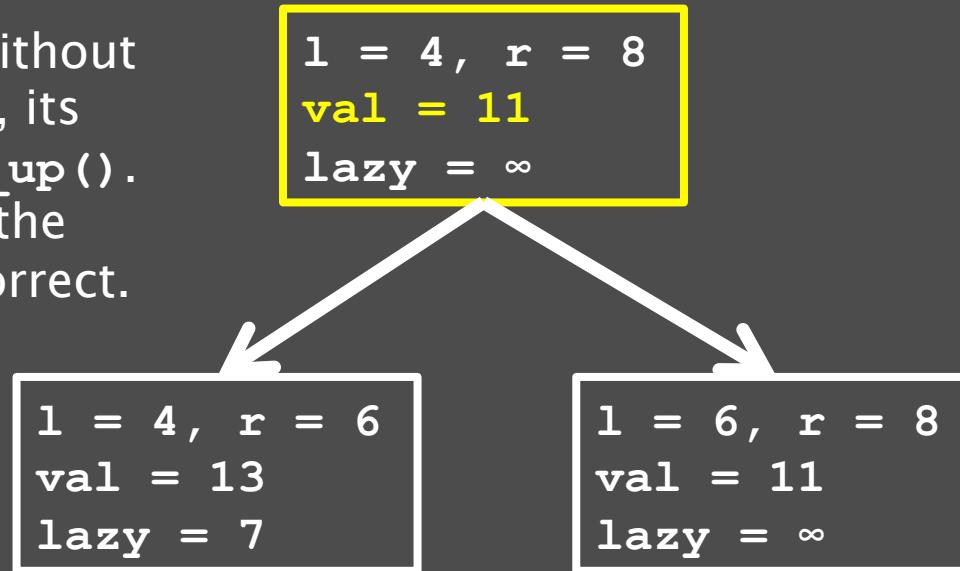
push_up/push_down Issue

Get to the target node and set its lazy value equal to 7



push_up/push_down Issue

If the node returns without calling `push_down()`, its parent will call `push_up()`. Note that the `val` in the parent node is not correct.



Lazy Propagation

- For each update operation, which nodes should be initially set as lazy?
 - Each of these nodes must be entirely within the update interval
 - Combination of all intervals of nodes chosen must equal the update interval
 - We use the same nodes as if we were doing a range query
 - Like range query, the time complexity must therefore be $O(\log N)$

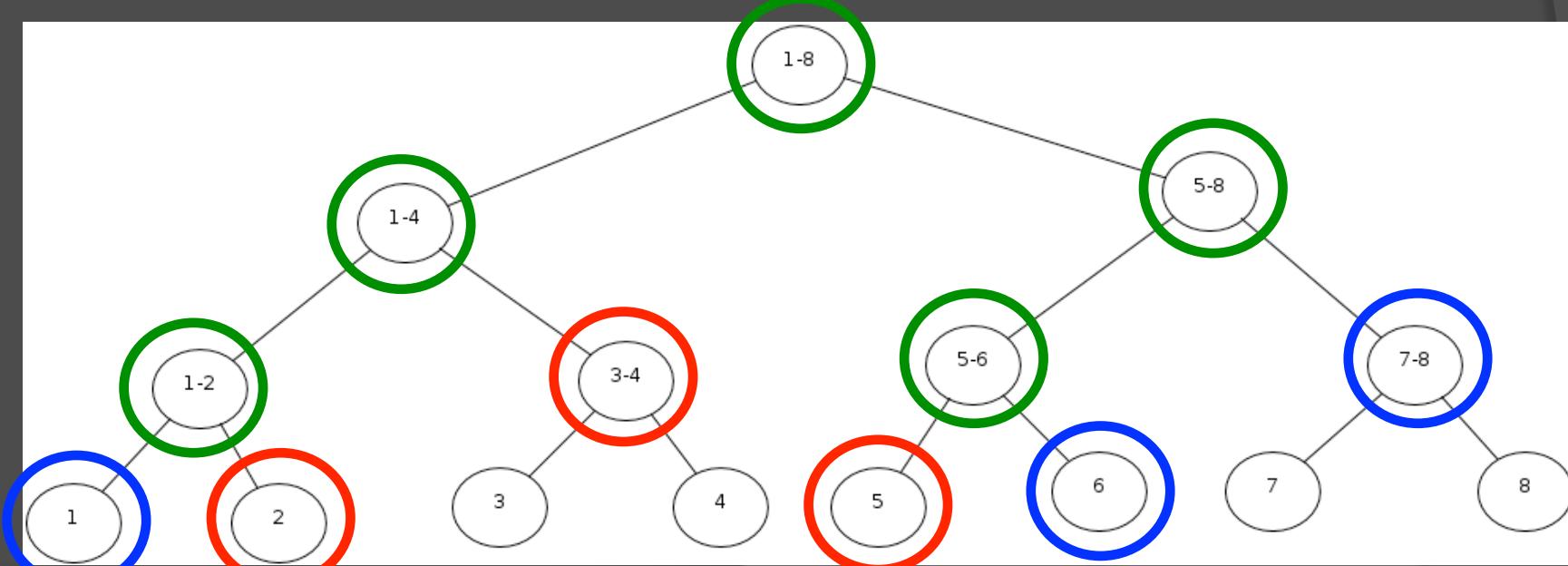
Range Update

Update: [2,5]

Green nodes specify visited nodes.

Red nodes specify nodes that we mark as lazy

Blue nodes are visited, but ignored



Range Update – Pseudocode (replace, min, elements > 0)

```
void update(int updL, int updR, int val, int idx) {
    push_down(idx); //necessary to update the node
    if (tree[idx].l >= updL && tree[idx].r <= updR) {
        //interval is entirely within range
        tree[idx].lazy = val; //mark as lazy
        push_down(idx);
        return;
    }
    if (tree[idx].l > updR || tree[idx].r < updL) {
        //interval is entire out of range
        return;
    }
    update(updL,updR,val,2*idx); //update left
    update(updL,updR,val,2*idx+1); //update right
    push_up(idx) // update answer
}
```

push_down - Pseudocode (replace, min, elements > 0)

```
void push_down(int idx) {
    if (tree[idx].lazy != 0){ //lazy is marked
        tree[idx].val = tree[idx].lazy; //update val
        if (tree[idx].l != tree[idx].r){
            //update child nodes
            tree[2*idx].lazy = tree[idx].lazy;
            tree[2*idx+1].lazy = tree[idx].lazy;
        }
        //unmark lazy
        tree[idx].lazy = 0;
    }
}
```

Query – Pseudocode

(replace, min, elements > 0)

```
int query(int queryL, int queryR, int idx){  
    push_down(idx); //query is same as before except this line  
    if (tree[idx].l >= queryL && tree[idx].r <= queryR){  
        //interval is completely within query range  
        return tree[idx].val;  
    }  
    if (tree[idx].l > queryR || tree[idx].r < queryL){  
        //interval is completely out of query range  
        return INF; //return some answer that will not affect  
        //final answer  
    }  
    int ansL = query(queryL,queryR,2*idx);  
    int ansR = query(queryL,queryR,2*idx+1);  
    return min(ansL,ansR); //choose the better answer  
}
```

THANK YOU!