# STUDY AND APPLICATION OF ANT COLONY SYSTEM ON TRAVELLING SALESMAN PROBLEM

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# OVERVIEW

- Travelling salesman problem (**TSP**)
- ANT COLONY SYSTEM(ACS)
- Description of algorithms
- Rules applied
- Data sets used
- Results
- Summary

#### TRAVELLING SALESMAN PROBLEM

• Minimum length Hamiltonian circuit.

• Shortest possible route that visits each city once and returns to the origin city.

• NP-HARD problem.



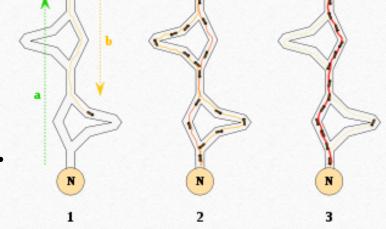
# Graph Theory

- Each node is referred to as a cities
- A path between two cities is denoted by a sequence of edge pairs or a sequence of nodes
- For study purposes: simple and fully connected graphs
- Large TSP data sets available online

#### ANT COLONY SYSTEM

• Ants deposit pheromones while searching for food.

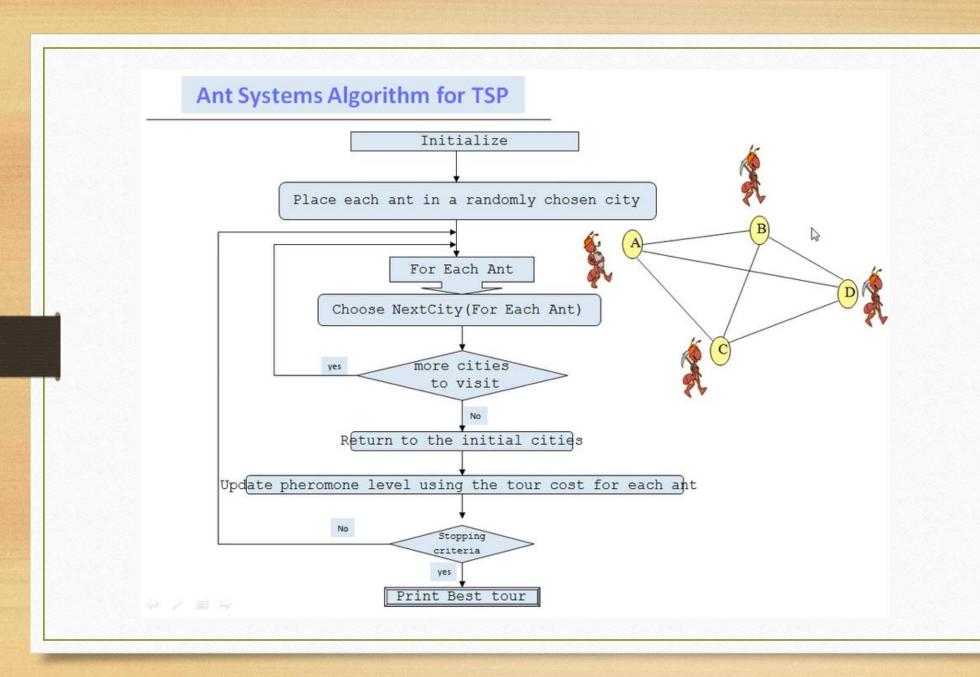
• A pheromone evaporates with time.



• Ants choose a path depending upon the amount of pheromones.

### TOUR CONSTRUCTION IN ACS

- 1. An ant starts at a some city.
- 2. Use pheromone and heuristic (educated or trial and error) values to probabilistically construct a tour by iteratively adding cities that the ant has not visited yet, until all cities have been visited.
- 3. Go back to the initial city.



# PHEROMONE AS A HEURISTIC IN ACS

- Initialize edge weights
- Update edge weights (add the pheromone to traversed edges)
- Decrement edge weights (evaporate pheromone)

#### INITIALIZATION OF PHEROMONE IN ACS

- Set equal amount of pheromone for all edges
  - value > expected amount of pheromone deposited by an ant during an iteration.
- If set too low, the search is quickly biased by the first tours generated by the ants.
- If set too high, many iterations are lost waiting until pheromone evaporates enough.

#### SIGNIFICANT FACTORS IN ACS ALGORITHM

- Pheromone  $\tau$  remaining information on the path the ants pass through
  - $\tau_{ij}$  = weightof the edge between nodes i and j
- Visibility  $\eta$  reciprocal distance between nodes
  - $\eta_{ij} = \frac{1}{d_{ij}}$ , d is distance between nodes i and j

#### SIGNIFICANT FORMULA IN ACS ALGORITHM

- Probability P At each construction step of the shortest path, ant<sub>k</sub> applies a probabilistic rule to decide which city to visit next.
- The shortest travelled path so far, will have a high  $\tau$  and a high  $\eta$  making P (probability) stronger for selection purposes.

# RANDOM PROPORTIONAL RULE

• The rule is given as follows:

$$p_{ij}^{k} = \frac{\tau_{ij}^{\alpha} \, \eta_{ij}^{\beta}}{\sum_{z \in i} \tau_{iz}^{\alpha} \, \eta_{iz}^{\beta}}, \quad 0 \text{ if no next city}$$

k is an ant.

 $p_{ij}^k$  is probability of choosing the next city j when an ant is at i.

 $\tau_{ij}$  is pheromone quantity between the edge between x and y.

 $\eta_{ij} = \frac{1}{d_{ij}}$ , where d is the distance between i and j.

 $\alpha$  is a parameter to control the influence of pheromone.

 $\beta$  is a parameter to control the influence of distance.

#### ACS TOUR CONSTRUCTION

- With a probability of P the ant makes the best possible move.
- With a probability of (1-P), the ant performs a biased exploration of edges.
- Tuning P allows modulation of the degree of exploration and the choice to explore other tours.
- The ant which has the best tour will now add pheromone after each iteration according to Global Pheromone Trail Update.

# Global Pheromone Trail Update

In pheromone trail update, both evaporation and new pheromone deposit only apply to the arcs that belong to the best so far solution.

- $\Delta \tau_{ij} = Q/Lk$  is quantity of pheromone laid on the path (i,j) by ant k, where Q is a constant value of total pheromone and  $L_k$  is the tour length of the kth ant.
- Update  $\tau_{ij} = (1 \rho) \tau_{ij} + \sum_{k=1}^{m} \Delta \tau_{ij}$ , where  $\rho$  is the constant evaporation factor and m is the number of ants in the colony

# Improvement in Efficiency

• The pheromone update at each iteration is reduced from O(n^2) to O(n).

# Local Pheromone Trail Update

• Ants use a local pheromone update rule that they apply immediately after having crossed an arc (I,j) during their tour.

$$\tau_{ij} = (1 - \rho) \tau_{ij} + \rho \tau_0$$

• Each time an ant uses an arc its pheromone level is reduced, so that the arc becomes less desirable for the following ants.

#### PYTHON CODE IMPLEMENTATION

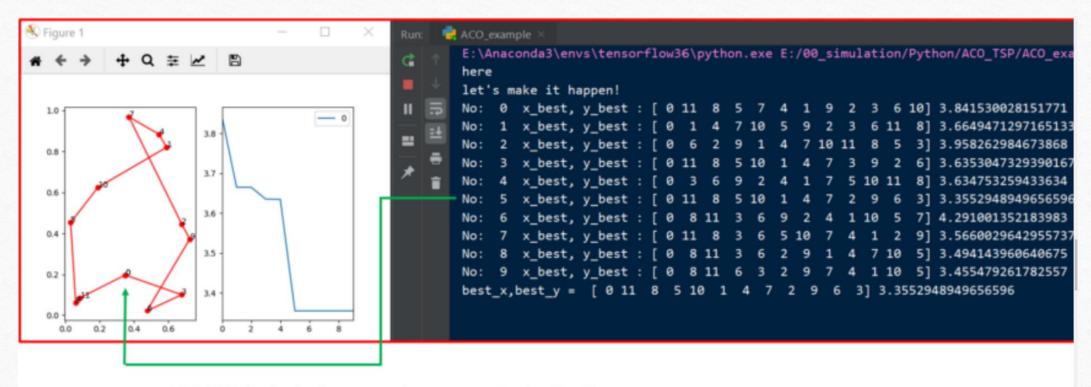
• The SKO package consists of optimization algorithms and sample code.

from sko.ACA import ACA\_TSP

- In the following example code, size\_pop is set to 3 ants to find the shortest path among 12 cities (points).
- Function cal\_total\_distance() refers to a path length computation and is the objective function.
- Function ACA\_TSP() finds the best path given the data.

```
from __future__ import division
import numpy as np
from scipy import spatial
import pandas as pd
import matplotlib.pyplot as plt
from sko.ACA import ACA TSP
# generate coordinate points
num points = 12
points_coordinate = np.random.rand(num_points, 2)
distance_matrix = spatial.distance.cdist(points_coordinate, points_coordinate, \
                       metric='euclidean')
def cal total distance(routine):
    num points, = routine.shape
    return sum([distance_matrix[routine[i % num_points], \
                 routine[(i + 1) % num points]] \
                      for i in range(num_points)])
```

```
def main():
    #find shortest path
    aca = ACA TSP(func=cal total distance, n dim=num points,
                  size pop=3, max iter=10,
                  distance matrix=distance matrix)
    best x, best y = aca.run()
    # Plot the result
    fig, ax = plt.subplots(1, 2)
    best_points_ = np.concatenate([best_x, [best_x[0]]])
    best_points_coordinate = points_coordinate[best_points_, :]
    for index in range(0, len(best points )):
        ax[0].annotate(best points [index], (best points coordinate[index, 0], \
                                              best_points_coordinate[index, 1]))
    ax[0].plot(best_points_coordinate[:, 0], best_points_coordinate[:, 1], 'o-r')
    pd.DataFrame(aca.y best history).cummin().plot(ax=ax[1])
    plt.show()
```



ACO-TSP finds the best travel-sequence in the No. 5

# INSIGHT INTO THE ACO\_TSP RUN() FUNCTION

ACO\_TSP problem (1997)

 TSP: travel sales problem.
 Find a shortest route of visiting all red-dot points

$$\eta_{ij} = \frac{1}{d_{ij}}$$
 ,  $d$  is distance

$$P_{ij}^{k} = egin{cases} rac{ au_{ij}^{lpha} \eta_{ij}^{eta}}{\sum_{Z \in i}^{\square} au_{iz}^{lpha} \eta_{iz}^{eta}} \ 0 \ , otherwise \end{cases}$$

$$\Delta \tau_{ij}^k = \begin{cases} \frac{Q}{L_k} \\ 0 \text{ , otherwise} \end{cases}$$

$$\tau_{ij}(t+1) = (1-\rho)\tau_{ij}(t) + \sum_{k=1}^{m} \Delta \tau_{ij}^{k}$$

```
run(self, max iter=None):
self.max_iter = max_iter or self.max_iter
for i in range(self.max_iter): # 附海文學:
   prob matrix = (self.Tau ** self.alpha) * ((self.prob matrix distance) ** self.beta) # 转移概率 王原知一代
    for j in range(self.size_pop): # T開个例数
       self.Table[j, 0] = 0 # start point,其实可以随机,但没什么区别
       for k in range(self.n_dim - 1): # 對权限法的每个节点
           allow_list * list(set(range(self.n_dim)) - taboo_set) # 存取機点中歐速排
           prob = prob_matrix[self.Table[j, k], allow_list]
           prob = prob / prob.sum() # EST
          next point = np.random.choice(allow list, size=1, p=prob)[0]
           self.Table[j, k + 1] = next_point
   y = np.array([self.func(i) for i in self.Table])
    index_best = y.argmin()
   x_best, y_best = self.Table[index_best, :].copy(), y[index_best].copy()
    self.x_best_history.append(x_best)
   self.y_best_history.append(y_best)
   delta_tau = np.zeros((self.n_dim, self.n_dim))
   for j in range(self.size_pop): # 問个對权
       for k in range(self, n_dim - 1): # # 170
           delta_tau[n1, n2] +* 1 / y[j] # 涂抹的信息素
       n1, n2 = self.Table[j, self.n_dim - 1], self.Table[j, 0] # 解权从最后一个节点使回到第一个节点
       delta_tau[n1, n2] += 1 / y[j] # 涂抹信息素
  self.Tau = (1 - self.rho) * self.Tau + delta_tau
best_generation = np.array(self.y_best_history).argmin()
self.best x = self.x best history[best generation]
self.best_y = self.y_best_history[best_generation]
return self.best x, self.best y
```

# **OTHER VIDEOS**

For a trace of the algorithm:

https://www.youtube.com/watch?v=783ZtAF4j5g

# NEAREST NEIGHBOR (NN) ALGORITHM

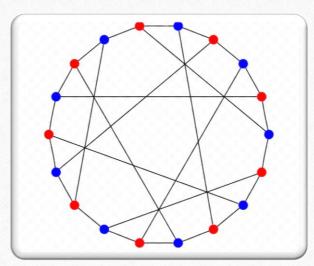
- 1. Initialize all vertices as unvisited.
- 2. Select an arbitrary vertex, set it as the current vertex U. Mark U as visited.
- 3. Find the shortest edge from U to one of the unvisited vertex V.
- 4. Mark V as the current vertex and mark it as visited.
- 5. If all the vertices are visited, then terminate.
- 6. Otherwise go to step 3.

#### **DATA SETS**

# All points on a circle

ACS and Nearest Neighbour algorithm will work perfectly as the

nearest node will always be on the circle.



#### Data sets

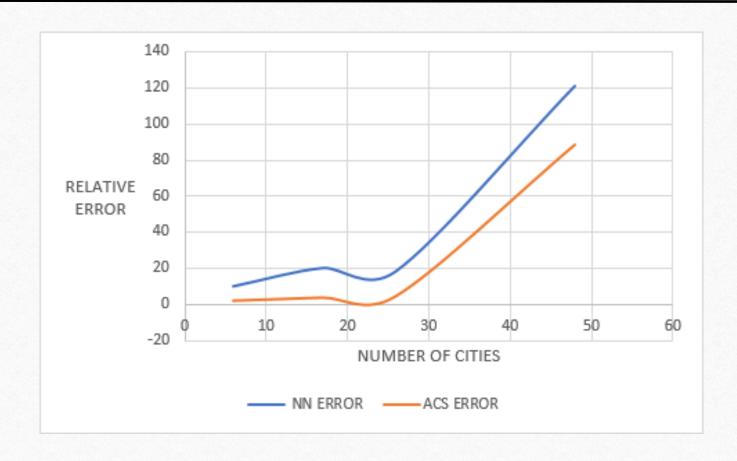
Random Nodes: As the number of Nodes increase the ACS performance decreases as the number of solutions increase.

# Some NN vs ACS Results

Relative errors obtained using NN and ACS on Travelling Sales Man Problem (compared to known optimal solution).

NO OF NODES	<b>NN</b> RELATIVE ERROR	ACS RELATIVE ERROR
6	10.37	2.37
17	20.47	4.05
26	18.67	4.85
48	121.07	88.96

# **RESULTS**



#### **SUMMARY**

- Tough ACS does not yield an optimal solution it gives much better results than many other algorithms.
- More number of paths are explored in ACS as the individual ants take individual routes while constructing tour.
- Tuning parameters will improve the results obtained by ACS.