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March 22, 2021

Integrity

We, Ryan Mao and Ethan Febinger, certify that the code and report for this assignment are our work alone.

Work Division

Ryan Mao: Representation, Play-by-play, Performance graphs

Ethan Febinger: Inference, Decision, Efficiency

Contents

1 Writeup
2 Non-Standard Libraries Used
18

1 Writeup

1. Representation: How did you represent the board in your program, and how did you represent the information/knowledge that clue cells reveal? How could you represent inferred relationships between cells?

The board was simply represented as a 2D matrix of cells. In this, each cell can be one of three things. If the cell is covered (i.e. the user/agent has not decided to query this cell), there will be a '?' at the corresponding location. If the cell has been queried before and is a mine, there will be a 'M' present. If the cell has been queried and is safe (i.e. the cell itself is NOT a mine), there will be the appropriate clue value which describes the number of mines directly surrounding the current cell.

The knowledgebase consists of four 2D matrices. The first 2D matrix is the current board using the same representation methods discussed above. The second 2D matrix is the number of safe (non-mine) squares around a given cell. For example, if there are 2 mines around a cell with indices (i, j), then at indices (i, j) for this 2D matrix will be 6, as there are 6 cells around the current cell that are not mines. Our third 2D matrix in our knowledgebase is precisely the negation of the second 2D matrix. Examining the same cell as in the previous example, at indices (i, j) for the third 2D matrix would be a 2. Lastly, the fourth 2D matrix stores the amount of covered/unqueried cells around a current cell.

The inferred relationships between cells is represented as sets. For example, if a cell has a clue value of 2, and none of the uncovered cells next to it are mines, then we know that the set of cells around the current cell has 2 mines in it. So, we say that $C = \{\text{number of nearby covered cells}\}$ and this has a value of 2. We then know any subset, A, of this set has to have a value say v, which is ≤ 2 . Thus, the C - A = 2 - v. From this, we can perform a sort of induction to aid our agent in the minesweeper program.

2. Inference: When you collect a new clue, how do you model/process/compute the information you gain from it? In other words, how do you update your current state of knowledge based on that clue? Does your program deduce everything it can from a given clue before continuing? If so, how can you be sure of this, and if not, how could you consider improving it?

Each new clue is stored a pair of a set and an integer. This set-integer pair represents an equation. For example, $(\{A, B, C\}, 2)$ represents the equation A+B+C=2 where A, B, and C, are the three covered squares surrounding the clue. In addition to adding a new equation, old equations are also updated in the knowledge base. This means that if the location of a revealed clue is a variable of another equation, that variable is removed from the set. Also, if a mine is revealed, the corresponding variable is removed from the set and the integer value is reduced by one.

First, the program looks for trivial equations to solve. If an equation sums to zero, all variables represent clear locations. If an equation's sum is equal to the number of variables it contains, then every variable represents a mine.

The program also makes inferences by looking if any set is a strict subset of another, i.e. for equations A and B then $A \subseteq B$. If this is the case, B is replaced with B - A in the knowledge base. From this new equation, the basic rules from above can then be checked again to look for additional knowledge.

This program does not deduce everything possible for a given clue. To improve this, the program could iterate through all possible permutations of the variable set and check which permutations satisfy every equation. If any variables were the same for every satisfiable permutation, then we can be sure of whether that variable is a mine. This, however, would be very computationally expensive.

3. Decisions: Given a current state of the board, and a state of knowledge about the board, how does your program decide which cell to search next?

When a new cell is revealed, that information is added to the knowledge base. Then, the program infers as much as possible via the inferences described in the previous question. The program then and adds the locations of those cells to a queue and stores whether any should be flagged as a mine. The program then searches through each location in the queue. If nothing can be inferred, a random covered location is added to the queue.

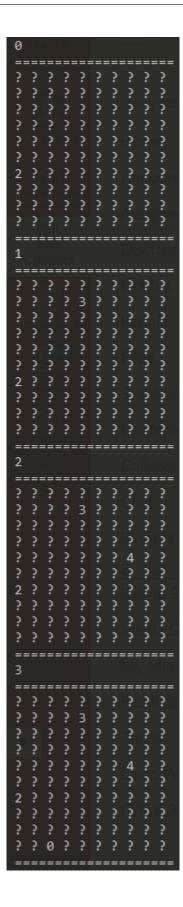
4. Performance: For a reasonably-sized board and a reasonable number of mines, include a play-by-play progression to completion or loss. Are there any points where your program makes a decision that you don't agree with? Are there any points where your program made a decision that surprised you? Why was your program able to make that decision?

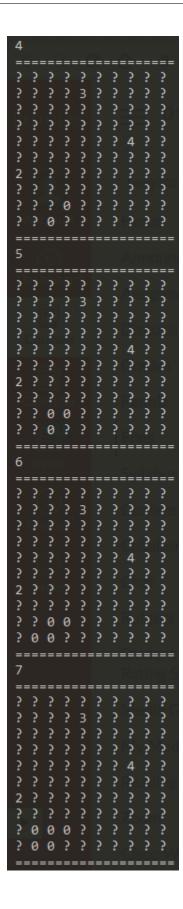
We have included a play_by_play.py script that uses our algorithm methods and prints what the algorithm sees as it attempts to solve the minesweeper game.

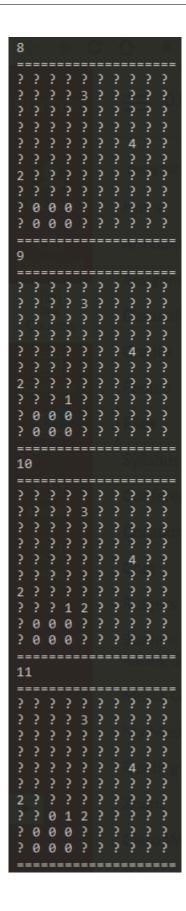
Here is a play by play that made 31 decisions before making an incorrect one and failing on the 32nd.

play by play begins on the next page (apologize for the inconvenience, but it could not fit without resizing)

There is also a description of when our advanced algorithm is able to make a decision the basic algorithm could not. Located after all the screenshots of the play by play.







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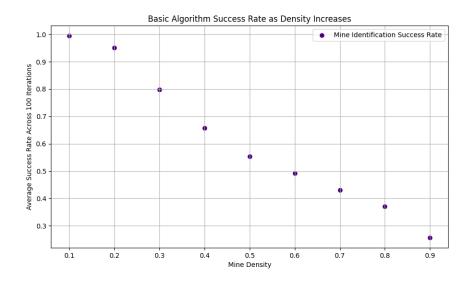
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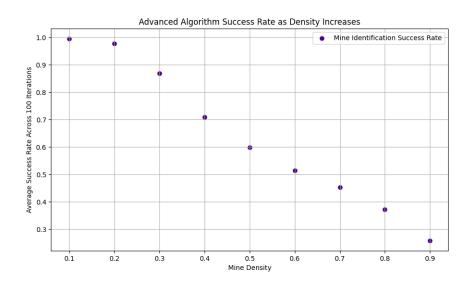


Focusing on step 25, this is a step that the original basic algorithm could not have made. We see that we have three '2s' stacked on top of each other. The bottom most 2 has, correctly, a mine identified next to it. The basic algorithm would not be able to deduce anything else, since the middle 2 has 2 undiscovered squares but only 1 undiscovered mine. Using the logic of the top most 2 and the 1 to the left of it, our advanced algorithm correctly marks the second mine, and also safely discovered the cell to the right of the top most 2.

5. Performance: For a fixed, reasonable size of board, plot as a function of mine density the average final score (safely identified mines / total mines) for the simple baseline algorithm and your algorithm for comparison. This will require solving multiple random boards at a given density of mines to get good average score results. Does the graph make sense / agree with your intuition? When does minesweeper become 'hard'? When does your algorithm beat the simple algorithm, and when is the simple algorithm better? Why? How frequently is your algorithm able to work out things that the basic agent cannot?

For both the basic algorithm and the advanced algorithm, we use a dimension size of 25, so the minesweeper board is a 25 by 25 square board.





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We can see that the advanced algorithm slightly outperforms the basic algorithm up until around a mine density of 0.8, where both algorithms obtain the same success results.

This makes sense for the advanced algorithm to do slightly better than the basic algorithm, as it is performing more induction, and attempting to 'reason' using logic and finding contradictions.

Minesweeper becomes 'hard' for both algorithms at around a mine density value of 0.6 where both algorithms perform practically the same.

There is not a time where the simple algorithm is better. The advanced algorithm is either better by a small margin, or the same. This liens up with our prediction/algorithm since the advanced algorithm performs the basic algorithm logic until it has to guess. Then it attempts to deduce something by assuming certain cells to be a mine. Depending on whether or not the algorithm then reaches a contradiction, the algorithm can slowly rule out certain cells as mines. If it is unable to deduce anything more, however, it will then make another guess.

6. Efficiency: What are some of the space or time constraints you run into in implementing this program? Are these problem specific constraints, or implementation specific constraints? In the case of implementation constraints, what could you improve on?

We ran into both space and time constraints for this program. We store an equation for every revealed clue. For a board of dimension d there are then $O(d^2)$ equations that need to be stored in the knowledge base at any given time. This is a problem specific constraint.

For every inference, the program iterates through every pair of constraints to see if they are subsets of each other. Given n equations, this takes $O(n^2)$ time. This constraint is implementation specific, and certain approaches could be used to improve upon this. For our implementation, all equations are stored in a singular array; however, not every variable appears in every equation. When the board is not dense there is very little overlap between equations. Because of this, the set of all equations can be separated into multiple disjoint sets of equations. Then, the time to iterate over all pairs of equations could be reduced to $O(m^2)$ where m is the size of the largest of these sets.

2 Non-Standard Libraries Used

- 1. matplotlib Generating and saving graphs.
- 2. numpy arange function to generate equally separated float values between two floats.