

Is Seven Shuffles Enough?

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- ▶ Riffle shuffle again

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How?

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Consider the following trick:

- ▶ Start with an ordered deck **1 rising sequence**
- ▶ Riffle shuffle once
- ▶ Riffle shuffle again
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How?

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Consider the following trick:

- ▶ Start with an ordered deck **1 rising sequence**
- ▶ Riffle shuffle once **2 rising sequences**
- ▶ Riffle shuffle again
- ▶ Pick a card and remember it
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- ▶ Riffle shuffle again **4 rising sequences**
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- ▶ Riffle shuffle one last time **$8 + 1$ rising sequences**
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How?

Main Theorem

After shuffling n cards for $\frac{3}{2} \log_2 n + \theta$ times, the distance to uniform is

$$1 - 2\Phi\left(\frac{-2^{-\theta}}{4\sqrt{3}}\right) + O\left(\frac{1}{n^{1/4}}\right)$$

where

$$\Phi(x) = \frac{1}{\sqrt{2\pi}} \int_{-\infty}^x e^{-t^2/2} dt.$$

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$$\left[1 - 2\Phi\left(\frac{-1}{4c\sqrt{3}}\right) \xrightarrow{c \rightarrow \infty} \frac{1}{2c\sqrt{6\pi}}\right]$$

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Some Data

TABLE 4
Shuffles needed to mix 25, 32, 52, 78, 104, 208 or 312 cards

n	25	32	52	78	104	208	312
$\frac{3}{2} \log_2 n$	6.97	7.50	8.55	9.43	10.05	11.55	12.43

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TABLE 3
Total variation distance for m shuffles of 25, 32, 52, 78, 104, 208 or 312 distinct cards

m	1	2	3	4	5	6	7	8	9	10
25	1.000	1.000	0.999	0.775	0.437	0.231	0.114	0.056	0.028	0.014
32	1.000	1.000	1.000	0.929	0.597	0.322	0.164	0.084	0.042	0.021
52	1.000	1.000	1.000	1.000	0.924	0.614	0.334	0.167	0.085	0.043
78	1.000	1.000	1.000	1.000	1.000	0.893	0.571	0.307	0.153	0.078
104	1.000	1.000	1.000	1.000	1.000	0.988	0.772	0.454	0.237	0.119
208	1.000	1.000	1.000	1.000	1.000	1.000	1.000	0.914	0.603	0.329
312	1.000	1.000	1.000	1.000	1.000	1.000	1.000	0.999	0.883	0.565

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 - ▶ Analysis of *smooshing* shuffle

Another Way to Riffle Shuffle

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- ▶ Split the deck into a packets using the multinomial distribution
- ▶ Iteratively riffle the packets together

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3. If π has r rising sequences, then the probability of obtaining π in an a -shuffle is $\frac{\binom{a+n-r}{n}}{a^n}$.

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Goal: upper bound the distance

$$\|Q^m - U\| = \sum_{r: Q^m(r) \geq \frac{1}{n!}} A_{n,r} \left(Q^m(r) - \frac{1}{n!} \right)$$

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$$\begin{aligned} \|Q^m - U\| &= \sum_{r: Q^m(r) \geq \frac{1}{n!}} A_{n,r} \left(Q^m(r) - \frac{1}{n!} \right) \\ &= \sum_{r=0}^{r^*} A_{n,r} \left(Q^m(r) - \frac{1}{n!} \right) \end{aligned}$$

Bounding the First Term

Let $c = 2^\theta$, then

$$\sum_{r=0}^{r^*} A_{n,r} \cdot Q^m(r) \approx \sum_{r=r_0}^{r^*} A_{n,r} \cdot Q^m(r)$$

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Putting Everything Together

$$\sum_{r=0}^{r^*} A_{n,r} \left(Q^m(r) - \frac{1}{n!} \right) = \Phi \left(\frac{1}{4c\sqrt{3}} \right) - \frac{1}{n!} \sum_{r=0}^{r^*} A_{n,r}$$

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Putting Everything Together

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