

Q1 Prefix-free Definition

1 Point

When is an encoding $E : \mathcal{O} \rightarrow \{0, 1\}^*$ prefix-free? That is, write down the definition of prefix-free encoding.

Definition 2.17 — Prefix free encoding. For two strings y, y' , we say that y is a *prefix* of y' if $|y| \leq |y'|$ and for every $i < |y|$, $y'_i = y_i$.

Let \mathcal{O} be a non-empty set and $E : \mathcal{O} \rightarrow \{0, 1\}^*$ be a function. We say that E is *prefix-free* if $E(o)$ is non-empty for every $o \in \mathcal{O}$ and there does not exist a distinct pair of objects $o, o' \in \mathcal{O}$ such that $E(o)$ is a prefix of $E(o')$.

Q2 Converting to Prefix-Free

2 Points

In class we discussed a way to convert any encoding $E : \mathcal{O} \rightarrow \{0, 1\}^*$ to be a prefix-free encoding.

Describe a way to do this in two or three sentences. [It is recommended to type your response here.]

Given any element $a \in \mathcal{O}$

1. Compute $E(a)$
2. Convert all $1 \rightarrow 11$, $0 \rightarrow 00$
3. Add 01 to the end

Q3 Size?

1 Point

Is the following statement true or false: Every function $f : \{0, 1\}^n \rightarrow \{0, 1\}$ has a Boolean circuit of size at most $2^n/n^2$.

True

False.

We showed $\frac{2^n}{80n}$ false*, so $\frac{2^n}{n^2}$ should be False
as well (for large enough n)

*Theorem 5.3
Lower Bound

Q4 Equivalence?

1 Point

Which of the following circuit models (given by the gates specified in the choices) are equivalent to Boolean circuits?

Here, 0 denotes the constant 0 and 1 denotes the constant 1 function.

{NAND}

{AND, OR, 0, 1}

← We showed this in homework!

{AND, NOT}

{NAND, AND, OR, NOT}

Q5 Code as data

1 Point

In class, we saw a way to encode NAND circuit programs as elements of $\{0, 1\}^*$. What was the number of bits needed to encode a NAND circuit program of size s ?

[You can use big-Oh notation.]

$$\rightarrow O(s \log(s))$$

*Recall how this comes from encoding a list of tuples

Q6 Circuits

1 Point

Which of the following are true:

Every function $f : \{0, 1\}^n \rightarrow \{0, 1\}$ has a Boolean circuit of size at most $O(n2^n)$ computing it.

Theorem from
lecture 4

Every Boolean circuit of size s has an equivalent NAND circuit of size at most s .

Size at most
 $3s$

Every NAND circuit of size s has an equivalent Boolean circuit of size at most $2s$.

Not(AND)
Produces
NAND

There exists a NAND circuit for $Eval_{n,m,s}$ (the evaluation function as defined in class that takes the encoding of a NAND circuit, input and returns the evaluation of the circuit on the input) of size at most $O(s^2 \log s)$.

Universal
Circuit

Q7 Function Count

1 Point

How many functions are there from $f : \{0, 1\}^n \rightarrow \{0, 1, 2, \dots, n - 1\}$?

- We know # functions is same as # truth tables

→ e.g. $n=2$

The diagram shows three 2-variable truth tables labeled f_1 , f_2 , and f_3 . Each table has two input columns (i_1, i_2) and one output column (f_i). The rows represent all possible combinations of inputs: (0,0), (0,1), (1,0), and (1,1). The outputs for f_1 are 0, 0, 0, 0. The outputs for f_2 are 0, 1, 0, 0. The outputs for f_3 are 0, 0, 1, 0. An ellipsis indicates more tables follow.

i_1	i_2	f_1
0	0	0
0	1	0
1	0	0
1	1	0

i_1	i_2	f_2
0	0	0
0	1	1
1	0	0
1	1	0

i_1	i_2	f_3
0	0	0
0	1	1
1	0	0
1	1	0

⋮

...

- Each truth table has a unique combination of outputs

→ e.g. $n=2$

$$\begin{aligned}f_1 &\rightarrow 0000 \\f_2 &\rightarrow 1000 \\f_3 &\rightarrow 0100\end{aligned}$$

⋮

- Therefore our problem becomes counting # n -ary strings of length $L = 2^n$

$$\Rightarrow \# \text{ functions} = n^L = n^{2^n}$$

Q8 State PECTT

1 Point

State the Physical Extended Church-Turing Thesis as discussed in class.

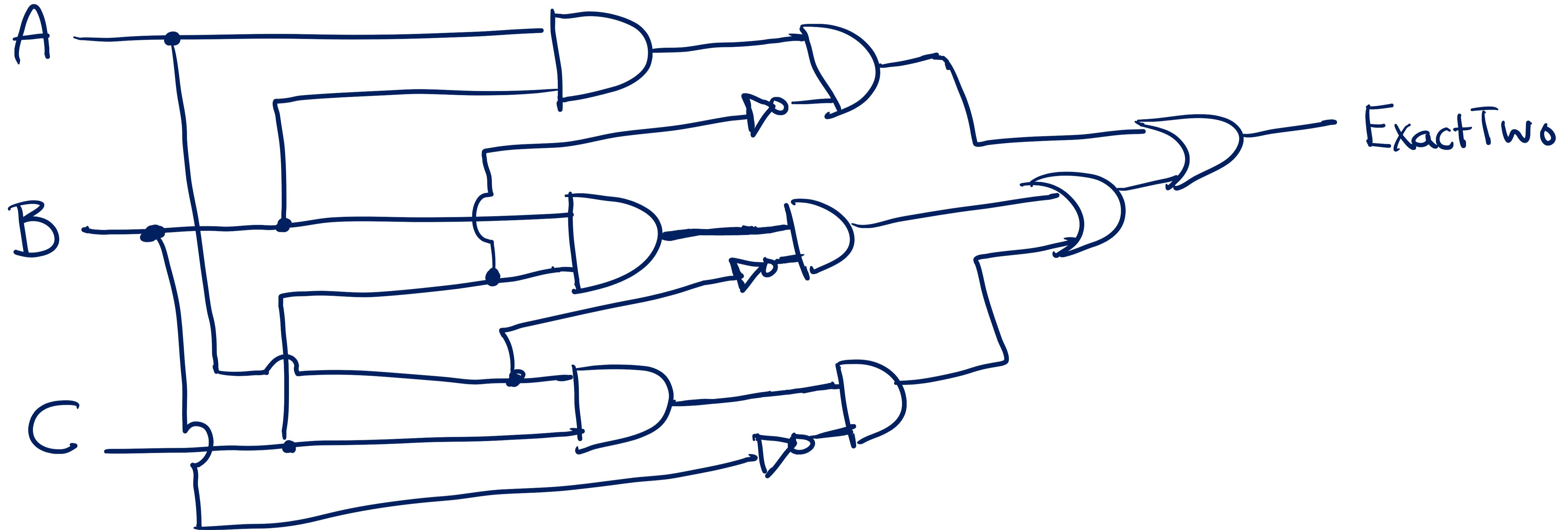
"Physical Extended Church-Turing Thesis" (PECTT): If a function $F : \{0, 1\}^n \rightarrow \{0, 1\}^m$ can be computed in the physical world using s amount of "physical resources" then it can be computed by a Boolean circuit program of roughly s gates.

Q9 Computing Exactly Two using AON

3 Points

Design a AND/OR/NOT circuit to compute the function $\text{ExactTwo} : \{0, 1\}^3 \rightarrow \{0, 1\}$ such that $\text{ExactTwo}(a, b, c) = 1$ if and only if exactly two of a, b, c are 1 and 0 otherwise. In other words, $\text{ExactTwo}(0, 1, 1) = \text{ExactTwo}(1, 0, 1) = \text{ExactTwo}(1, 1, 0) = 1$ and the function evaluates to 0 on all other inputs.

You can draw the circuit or write it as a AND/OR/NOT program.

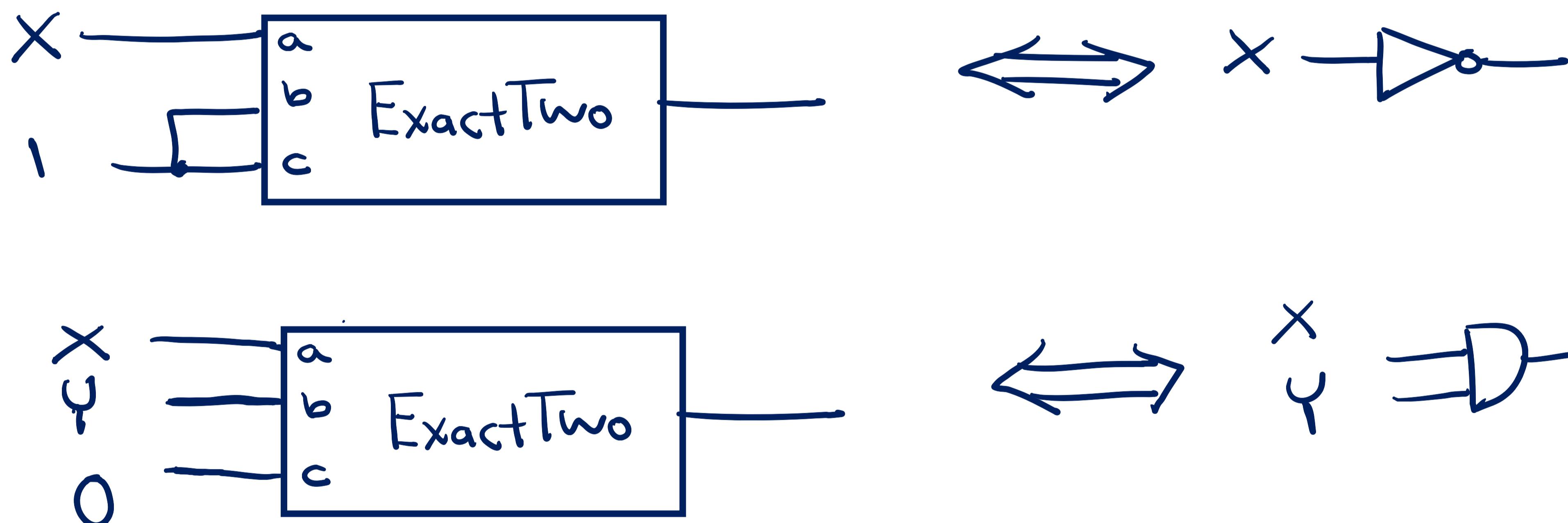


Q10 ExactTwo is Universal

3 Points

Show that $\{\text{ExactTwo}, 0, 1\}$ is a universal set of gates. Here, $\text{ExactTwo} : \{0, 1\}^3 \rightarrow \{0, 1\}$ is the function defined above and 0 is the constant zero function and 1 denotes the constant one function.

You can draw or write down the formulas for the computations as we did in class.

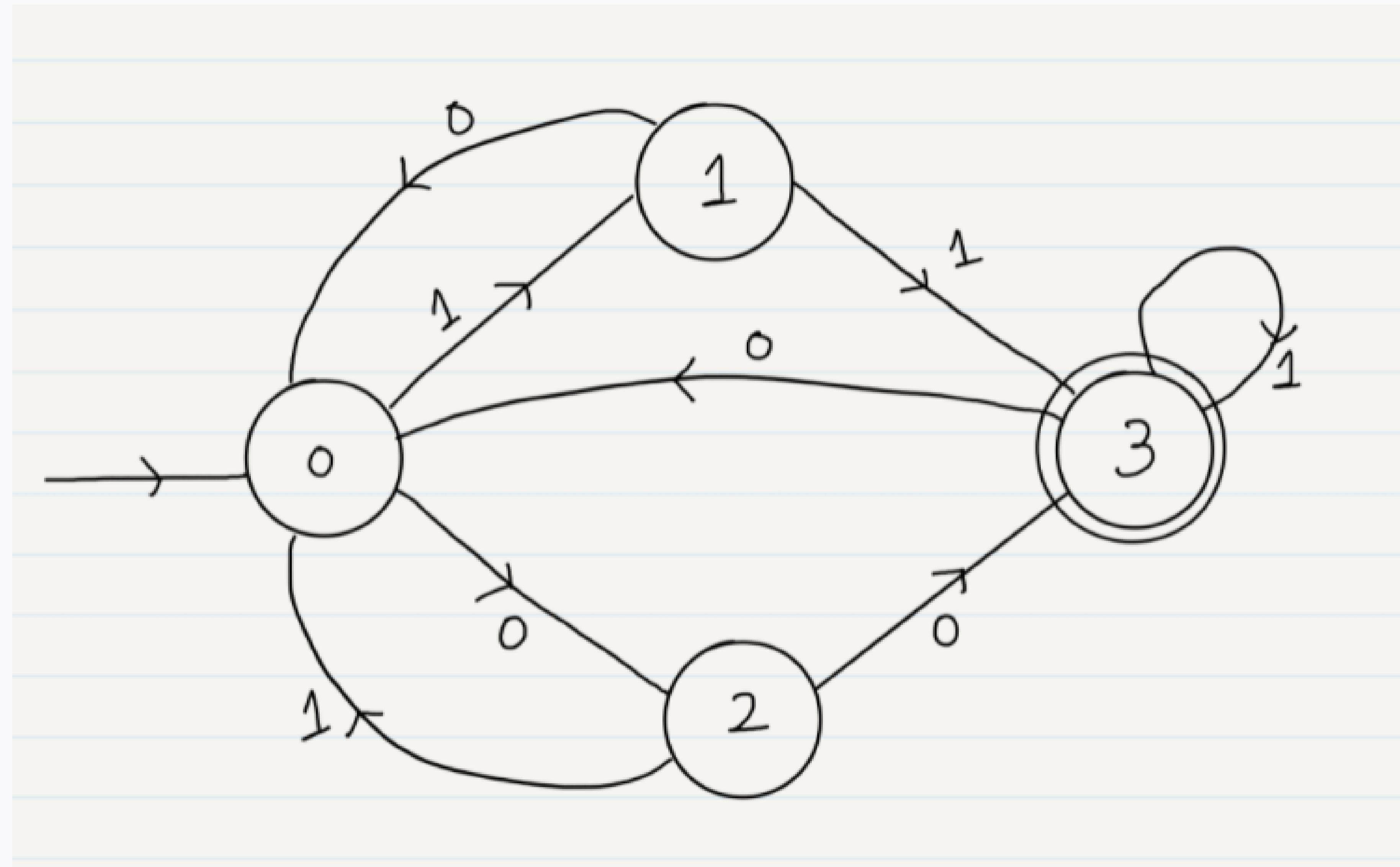


Chaining these together gives us $\text{NAND} \Rightarrow \text{universal circuit}$

Q11 Follow automata

2 Points

Write down the sequence of states visited by the automata D shown in the diagram on input $x = 0101001$. What is $D(x)$?



STATE	INPUT
0	0
2	1
0	0
2	1
0	0
2	0
3	1
3	

Because our last state = 3 \in Accept $\Rightarrow D(x) = 1$

Q12 Designing a DFA

3 Points

Draw a DFA that accepts strings that contain 101 as a substring. That is, draw a DFA D such that $D(x) = 1$ if and only if x contains 101 as a substring. You don't have to prove that your DFA works.

For instance, 1010, 10001010, 11101011, ... should be accepted by the DFA whereas strings such as 01001, 11000110, 11001000, ... should not lead to an accepting state.

It would be best to draw the DFA and upload the image.

