Homework 5

The logic behind this statement is simple. If you double the number of edges, you get the sum of all the degrees. The sim of degrees is even. The sum of the degrees of even-degree vertices is also even. Because an even number subtracted from an even number yields another even number, we can conclude that the sum of the degrees of odd degree vertices must also be even. Because the sum of the degrees of odd-degree vertices is even, we can further conclude that there will be an even amount of these vertices in the graph. D There is no Hamiltonian cycle in this graph. The fact that a and b are the only entry points to the inner part of the graph makes it impossible for a Manifornan cycle to be found. No matter which vertex you start from, it is inevitable that a and for b will show up moltiple times when aftempting to create a flamiltonian cycle. Maybe if there were more entry points to the inner part of the graph there would be a Hamitonian Cycle. For the complete bipartite graph Km, n to contain an Euleman cycle, m and n must both be even. This is because an Eulerian cycle only forms when each of the vertices has an even degree. For example:

m = n = 7

Each vertex has an even degree (deg(2)) Cycle goes from A -> D -> B -> C -> A



Each vertex has an even degree of either deg (4) or deg (2), so there's an Eulerian cycle. V

Each vertex does not have an even degree so there's no Eulerian cycle, X

a 2.	G=(V, E) a graph
	V~W if and only if there's or path from v to w
	Show this relation is on equivalence relation
	equivalence relation: a reflexive, symmetric, and transitive relation
0	Reflexive: This relation ~ is reflexive because if we
	let v & V, there will be a path v-V. This is
	essentially an empty path for a path from a vertex
	to itself. This means the relation ~ is reflexive,
0	Symmetric: This relation ~ is symmetric because if
	V, w E V and there is a path v-w, there
	will also be a path w-v. A real life example
	would be this: if you walked a path in the
	park that started at the playgrand (v) and
	followed the path to the tennis court (w), you
	can follow the path from the tennis court back
	to the playground. This shows that the relation ~
	is symmetric.
(3)	Transitive: This relation ~ is transitive becourse of
	a similar logic to how the relation is symmetric.
	I will build upon the example with the park.
	X (factball field)
	(playgrund)
	In this dragram there is a path from v to w
~	and another path from w to x Because
	v is connected to w and w is connected to x,
	this shows that v is connected (or related) to
	X. This shows the relation ~ is transitive.
-	Because the relation ~ is reflexive, symmetric, and
	transitive, we can conclude the relation is indeed
	an equivalence relation.

Vi. .. Vm the blocks in the partition of V giren by equivalence relation in Problem # 2 (a. = (Vi, Ei) induced subgraph of (on Vi for i=1, ..., m is Ei=E - Gi consists of all vertices in V. and the edges of G that exist within paths of the verkes in V: This problem is saying to prove the sum of the edges of all the induced subgraphs links have formed by the blocks of the partition) is equivalent to the number of edges E from the complete graph. We know the graph is symmetric, reflexive, transitive. A partition will form in a graph when there are vertices that aren't connected, or non-overlapping (because of eg. relation) For example, EX Graph G: (4 total edges) ledge 3 edges partition because neither D or E are connected to / neighbors with A, B, or C. In the above example, we see that there is one partition in graph G, whice creates 2 blocks. The first block has 3 edges and the second block has I edge. When adding them together, 3-1=4= total number of edges of G. This is one example that satisfies the slatement , D. E .: E. V Every vertex in a block Vi is connected to each other, so every vertex in Vi is only incident on edges that're incident on other vertices in Vi. All of these edges combine to form Ei. This means Ei consists of the e E E when the vertex X E V:, and X is incident on the edge e Every edge e EE is made of lor 2 vertices in V: These two vertices are adjacent since they reside in the same block. If an edge only has I vertex, that just means the block only has one vertex in it. Every edge e & E belongs in an E:, the same E: that corresponds with V; where its incident vertices reside. Every edge ett belongs to a set E; of E, ..., Em, This means EI, ... Em must be a partition of E, proving that U E; = E. D