Answer the questions in the boxes provided on the question sheets. If you run out of room for an answer, add a page to the end of the document.

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More Network Flow

1. *Kleinberg, Jon. Algorithm Design (p.416 q.6).* Suppose you're a consultant for the Ergonomic Architecture Commission, and they come to you with the following problem.

   They're really concerned about designing houses that are "user-friendly", and they've been having a lot of trouble with the setup of light fixtures and switches in newly designed houses. Consider, for example, a one-floor house with n light fixtures and n locations for light switches mounted in the wall. You'd like to be able to wire up one switch to control each light fixture, in such a way that a person at the switch can see the light fixture being controlled.

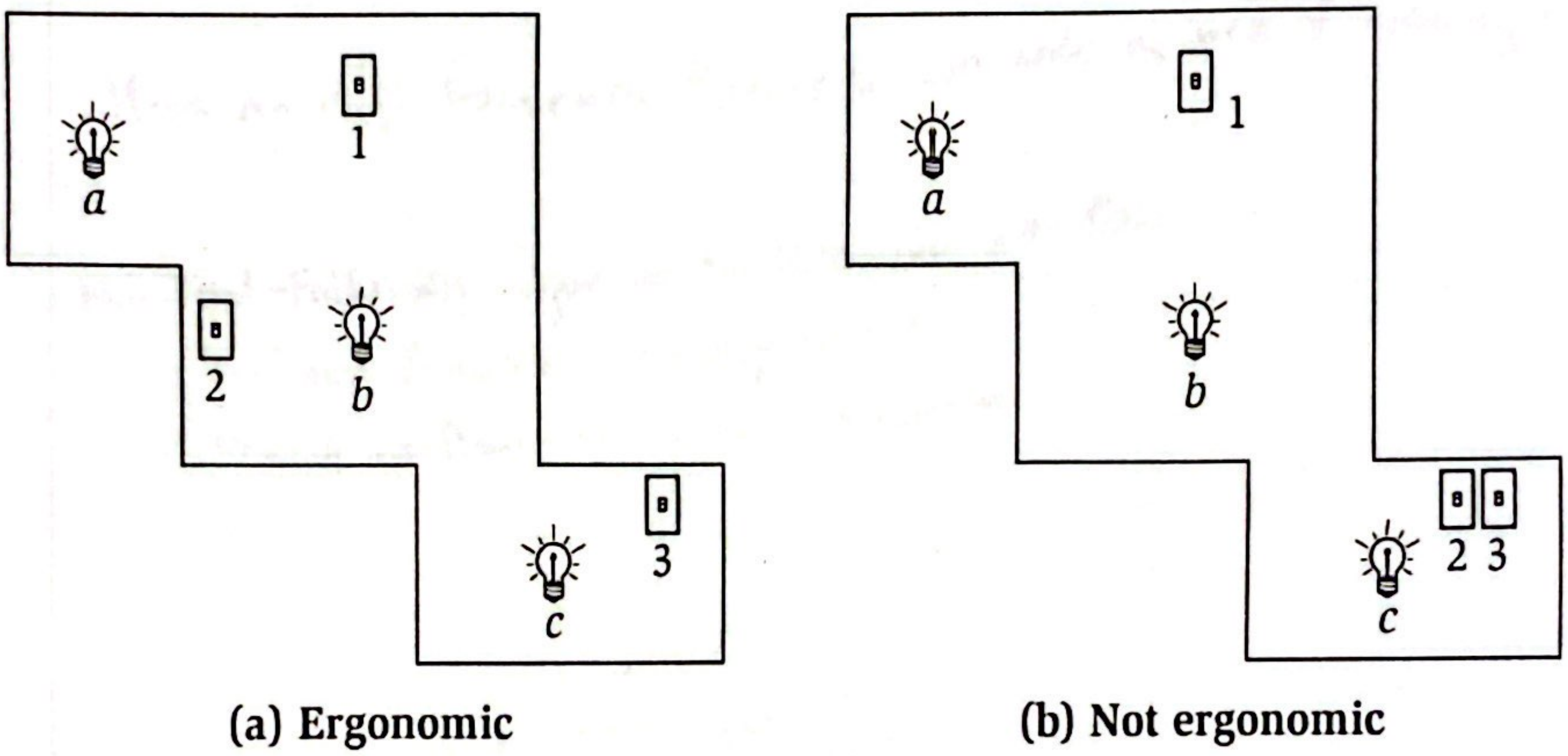


Figure 1: The floor plan in (a) is ergonomic, because we can wire switches to fixtures in such a way that each fixture is visible from the switch that controls it. (This can be done by wiring switch 1 to a, switch 2 to b, and switch 3 to c.) The floor plan in (b) is not ergonomic, because no such wiring is possible.

   Sometimes this is possible and sometimes it isn't. Consider the two simple floor plans for houses in Figure 1. There are three light fixtures (labelled a, b, c) and three switches (labelled 1, 2, 3). It is possible to wire switches to fixtures in Figure 1(a) so that every switch has a line of sight to the fixture, but this is not possible in Figure 1(b).

   Let's call a floor plan, together with n light fixture locations and n switch locations, ergonomic if it's possible to wire one switch to each fixture so that every fixture is visible from the switch that controls it. A floor plan will be represented by a set of m horizontal or vertical line segments in the plane (the walls), where the i-th wall has endpoints $(x_i, y_i), (x'_i, y'_i)$. Each of the n switches and each of the n fixtures is given by its coordinates in the plane. A fixture is visible from a switch if the line segment joining them does not cross any of the walls.

   Give an algorithm to decide if a given floor plan is ergonomic. The running time should be polynomial in m and n. You may assume that you have a subroutine with $O(1)$ running time that takes two line segments as input and decides whether or not they cross in the plane.

Solution:

Have fixtures as nodes & switches as nodes.

Have a source node with edge to each switch of capacity 1.

~~A source~~ edge from each switch to fixture if they are in each other's line of sight, pass in both x & y coordinates to the O(1) function.

n switches, n fixtures, m walls, so this takes $O(n^2m)$.

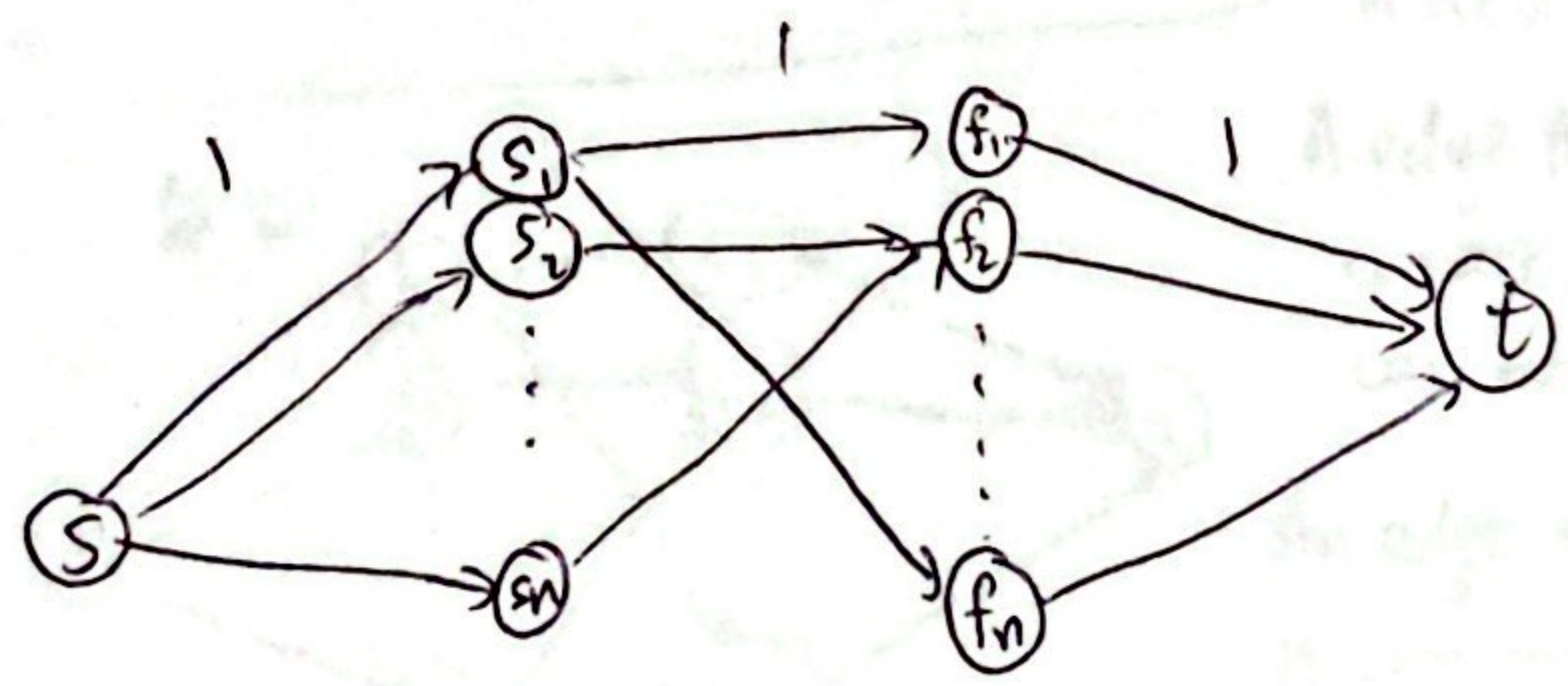
capacity 1.

Have an edge from each fixture to sink node as well of capacity 1.

Run Ford-Fulkerson algorithm to determine max flow.

↳ if max flow = n ⇒ ergonomic

↳ else max flow < n ⇒ non-ergonomic.



2. *Kleinberg, Jon. Algorithm Design (p.426 q.20).*

Your friends are involved in a large-scale atmospheric science experiment. They need to get good measurements on a set S of n different conditions in the atmosphere (such as the ozone level at various places), and they have a set of m balloons that they plan to send up to make these measurements. Each balloon can make at most two measurements. Unfortunately, not all balloons are capable of measuring all conditions, so for each balloon $i = 1, \ldots, m$, they have a set S_i of conditions that balloon i can measure. Finally, to make the results more reliable, they plan to take each measurement from at least k different balloons. (Note that a single balloon should not measure the same condition twice.) They are having trouble figuring out which conditions to measure on which balloon.

Example. Suppose that $k = 2$, there are $n = 4$ conditions labelled c_1, c_2, c_3, c_4, and there are $m = 4$ balloons that can measure conditions, subject to the limitation that $S_1 = S_2 = c_1, c_2, c_3$, and $S_3 = S_4 = c_1, c_3, c_4$. Then one possible way to make sure that each condition is measured at least $k = 2$ times is to have

- balloon 1 measure conditions c_1, c_2,
- balloon 2 measure conditions c_2, c_3,
- balloon 3 measure conditions c_3, c_4, and
- balloon 4 measure conditions c_1, c_4.

(a) Give a polynomial-time algorithm that takes the input to an instance of this problem (the n conditions, the sets S_i for each of the m balloons, and the parameter k) and decides whether there is a way to measure each condition by k different balloons, while each balloon only measures at most two conditions.

Solution:

1 k

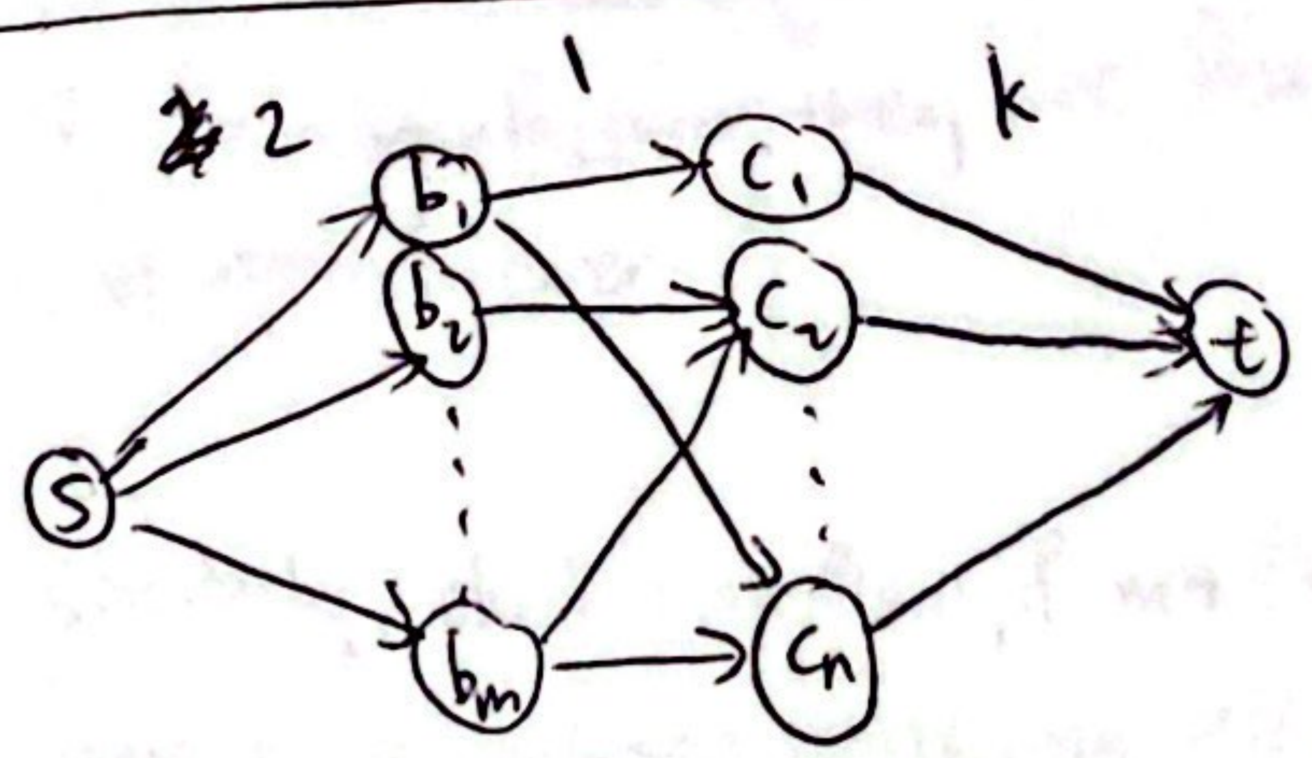
2 ballons conditions

sink.

source

if $|S| \times k$ then yes.

Each ~~node~~ ballon, condition is a node.



A edge from source to balloon with capacity ~~2~~ 2, since each balloon can measure at most 2 conditions.

An edge from balloon to each condition it can measure, capacity 1.

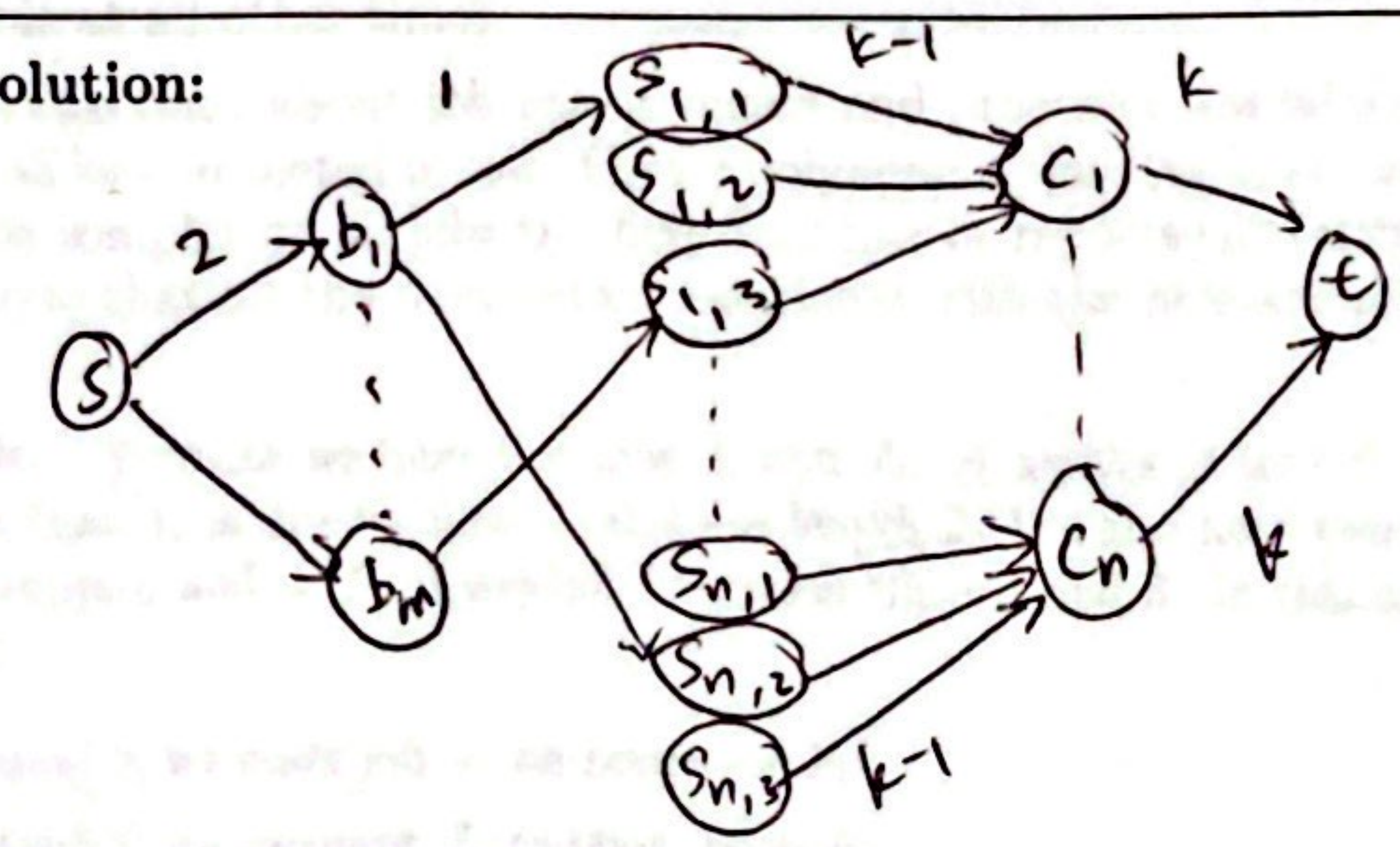
An each from each condition to sink, capacity k, the number of times each condition must be measured.

If Max Flow $= |S| \times k$ then there is a way to measure each condition by k different balloons, where each balloon only measures at most 2 conditions. $\Rightarrow$ Use Ford Fulkerson $O((mn)^2 (m+n))$ in this case

edges nodes.

(b) You show your friends a solution computed by your algorithm from (a), and to your surprise they reply, "This won't do at all—one of the conditions is only being measured by balloons from a single subcontractor." You hadn't heard anything about subcontractors before; it turns out there's an extra wrinkle they forgot to mention...
Each of the balloons is produced by one of three different subcontractors involved in the experiment. A requirement of the experiment is that there be no condition for which all k measurements come from balloons produced by a single subcontractor.
Explain how to modify your polynomial-time algorithm for part (a) into a new algorithm that decides whether there exists a solution satisfying all the conditions from (a), plus the new requirement about subcontractors.

Solution:



3n more nodes.
$|S| \times k$.

Add middle nodes, 3 nodes for each manufacturer per condition with capacity $k-1$, to ensure each condition is measured by at least balloons of 2 manufacturers. There will be 3n more nodes.

Each balloon has edge to the respective condition they can measure of the manufacturer they are from, so $S_{n,1}$ means nth condition, 1st manufacturer.

Similarly, check max flow, if max flow = $|S| \times k$, then we know there is a solution satisfying all conditions & requirements.

3. *Kleinberg, Jon. Algorithm Design (p.442, q.41).*

Suppose you're managing a collection of k processors and must schedule a sequence of m jobs over n time steps.

The jobs have the following characteristics. Each job j has an arrival time a_j when it is first available for processing, a length ℓ_j which indicates how much processing time it needs, and a deadline d_j by which it must be finished. (We'll assume $0 < \ell_j \leq d_j - a_j$.) Each job can be run on any of the processors, but only on one at a time; it can also be preempted and resumed from where it left off (possibly after a delay) on another processor.

Moreover, the collection of processors is not entirely static either: You have an overall pool of k possible processors; but for each processor i, there is an interval of time $[t_i, t_i']$ during which it is available; it is unavailable at all other times.

Given all this data about job requirements and processor availability, you'd like to decide whether the jobs can all be completed or not. Give a polynomial-time (in k, m, and n) algorithm that either produces a schedule completing all jobs by their deadlines or reports (correctly) that no such schedule exists. You may assume that all the parameters associated with the problem are integers.

Example. Suppose we have two jobs J_1 and J_2. J_1 arrives at time 0, is due at time 4, and has length 3. J_2 arrives at time 1, is due at time 3, and has length 2. We also have two processors P_1 and P_2. P_1 is available between times 0 and 4; P_2 is available between times 2 and 3. In this case, there is a schedule that gets both jobs done.

- At time 0, we start job J_1 on processor P_1.
- At time 1, we preempt J_1 to start J_2 on P_1.
- At time 2, we resume J_1 on P_2. (J_2 continues processing on P_1.)
- At time 3, J_2 completes by its deadline. P_2 ceases to be available, so we move J_1 back to P_1 to finish its remaining one unit of processing there.
- At time 4, J_1 completes its processing on P_1. Notice that there is no solution that does not involve preemption and moving of jobs.

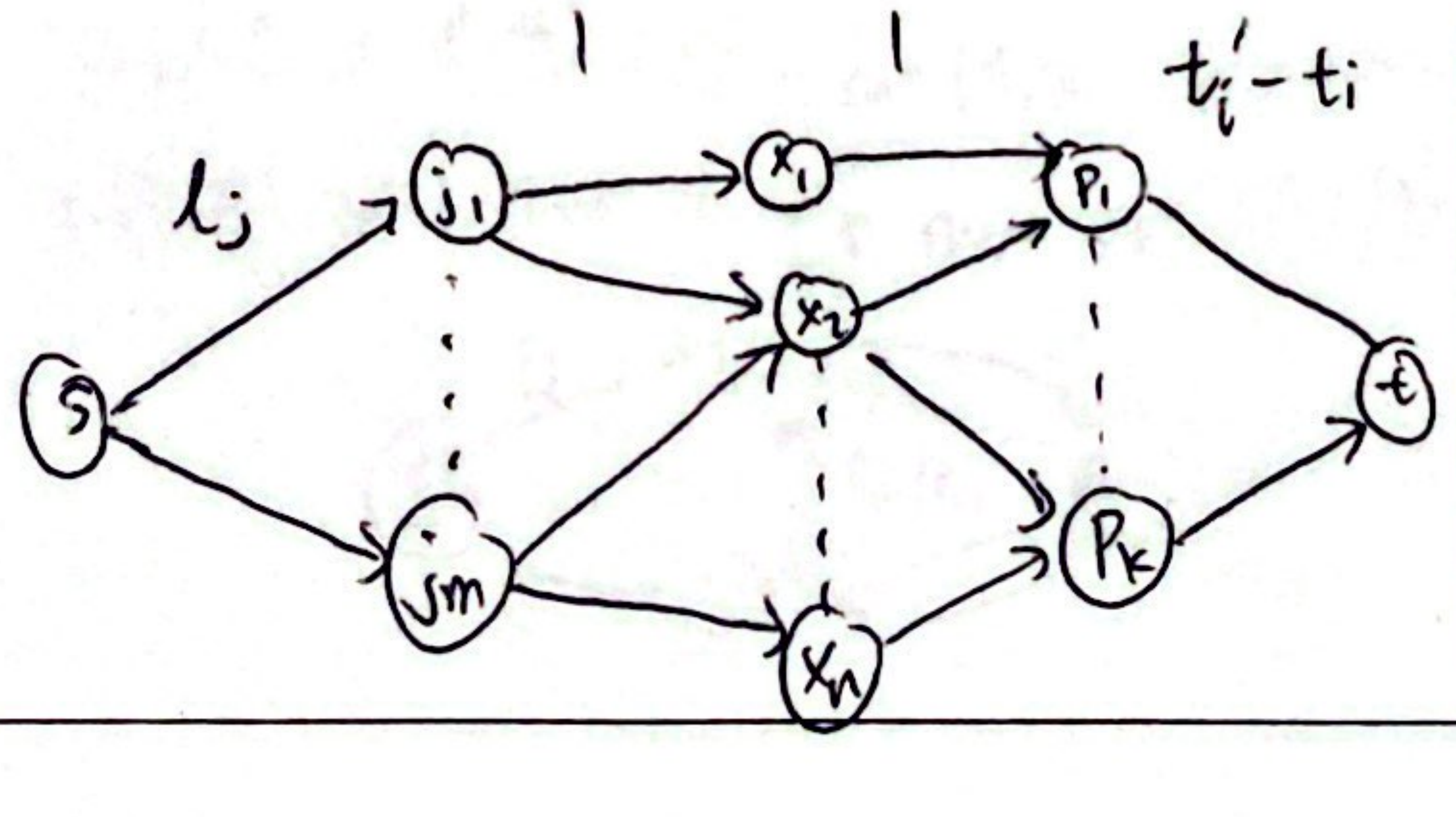
Solution:

Each job have a node with source node s, with capacity l_j

Each time step is a node with edge from job if job j is available at time x.

Each processor p will have node with an edge of capacity 1 from time x if p is available at time x.

sink node t with an edge from p with capacity $t_i' - t_i$. ~~representing~~ ~~max flow if~~ $= \sum_{j=1}^{m} l_j$



Check if max flow $= \sum_{j=1}^{m} l_j$

↳ if yes ⇒ there is a schedule

↳ else no.

Can use Ford Fulkerson.

$O(\text{edges}^2 \text{nodes}) = O((mn+nk)^2 (m+n+k)))$

4. *Kleinberg, Jon. Algorithm Design (p.444, q.45).*

Consider the following definition. We are given a set of n countries that are engaged in trade with one another. For each country i, we have the value s_i of its budget surplus; this number may be positive or negative, with a negative number indicating a deficit. For each pair of countries i, j, we have the total value e_{ij} of all exports from i to j; this number is always nonnegative. We say that a subset S of the countries is *free-standing* if the sum of the budget surpluses of the countries in S, minus the total value of all exports from countries in S to countries not in S, is nonnegative. Give a polynomial-time algorithm that takes this data for a set of n countries and decides whether it contains a nonempty free-standing subset.

Solution:

Let $S^+ = \Sigma_{s_j>0} s_j$

free-standing sum $f(A)$ can be calculated

$$f(A) = \sum_{j\in A} s_j - \sum_{i\in A, j\notin A} e_{ij} = S^+ + \sum_{j\in A, s_j<0} s_j - \sum_{j\notin A, s_j>0} s_j - \sum_{i\in A, j\notin A} e_{ij}$$

- Node for each country j
- For each pair of countries i, j, there is an edge (i,j) with capacity e_{ij} and edge (j,i) with capacity e_{ji}.
- Node s, source, with an edge with capacity s_j to each country j with $s_j>0$.
- Node t, with edge, capacity of $-s_j$ from each country j with $s_j<0$.
- consider cut (A,B)
  - Edge $i\in A$ to t contributes $-s_j$
  - Edge s to $j\in B$ contributes s_j
  - Edge from $i\in A$ to $j\in B$ contributes e_{ij}
  - Hence $\text{cut}(A,B) = -\sum_{j\in A, s_j<0} s_j + \sum_{j\notin A, s_j>0} s_j + \sum_{i\in A, j\notin A} e_{ij}$ which corresp to all but constant S^+ in $f(A)$.

Therefore $f(A) = S^+ - \text{cut}(A,B)$, so minimizing $\text{cut}(A,B)$ maximizes $f(A)$.

If $f(A)$ for min cut is >0 then set $A \setminus \{s\}$ is free-standing.