

# Advanced Computer Networking (ACN) IN2097

Prof. Dr.-Ing. Georg Carle

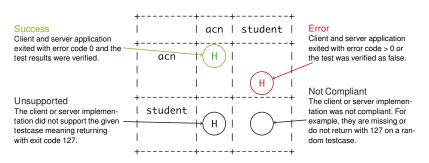
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#### General Remarks



- · Use the Moodle forum for discussion
- Compliance check of your implementations
  - If given a random test case, the client and server need to terminate with code 127
  - and print "exited with code 127"
- · Binaries need to be executable outside the CI as well
  - · We will test your implementations against each other and on real hardware
- The handshake for the first test case needs to be without a retry (see Moodle for more information)
- We use HTTP/3 and therefore h3 as ALPN
- For Problem2, all tests in the CI should be successful:

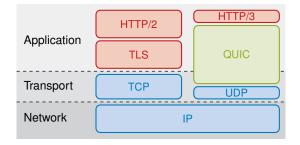




#### 1a)

- · Introduced by Google and Jim Roskind
- First draft from 08.07.2016<sup>1</sup>
- While QUIC was an acronym initially, it is not used as such anymore

# 1b)

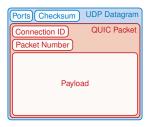


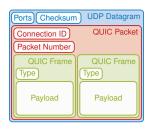
 $<sup>\</sup>mathbf{1}_{\text{https://datatracker.ietf.org/doc/html/draft-hamilton-quic-transport-protocol-00}}$ 



#### 1c)







# 1d)

- Unilaterally declared parameters by each endpoint
- 17 different parameters exist in the specification
- Peers can e.g., set the initial size of the flow control window for the connection or streams, the maximum number of allowed streams, and options regarding connection migration
- Transport parameters can be send during the TLS handshake to authenticate values but can be set or updated later as well



#### 1f)

#### Client:

- The client can send data via a QUIC stream (similar to TCP: in order, reliable, and stream-oriented)
- QUIC splits the stream into chunks which are encoded as frames
- The frames are combined into QUIC packets, multiple frames can be included in one packet
- The whole QUIC packet is encrypted and header protection is applied
- The encrypted QUIC PDU is passed to the UDP interface, e.g. using send()

#### Server:

- A UDP datagram arrives on the bound UDP socket
- The connection ID is parsed and the packet can be assigned to a connection
- The header protection is removed and the payload is decrypted using the connection's session keys
- The frames are parsed and the data stream is reconstructed from the data chunks using the offset field
- The server application can read from that stream interface, as with TCP, and, for example, respond to the received data in the same stream

# Project - Problem 1 Selected Libraries



# 1e,g,h)

Library	Versions	CC Algorithms	Students
aioquic	v1, draft-[29-32] v1, draft-[29,27]	NewReno	8
Isquic quic-go	v1, draft-29	Cubic, BBR, adaptive Cubic	4 3
quinn quiche	v1, draft-[29-34] v1, draft-[27-29]	Reno, Cubic, BBR Reno, Cubic	3 1
haskell quic	v1, draft-29	RFC9002 pseudo code	1



We will publish the description and additional information as Moodle post next Tuesday

- Start: December 21, 2021, 4:00 PM
- End: January 18, 2022, 4:00 PM
- Tasks: Implement additional test cases, e.g.:
  - Conduct multiple handshakes
  - Support retries
  - · Configure transport parameters
- We will test implementations against each other and verify test cases