

27.11.17 моно XII

Ускоренный оптимальный метод Хестерова

$$f(x) \rightarrow \min_{x \in \mathbb{R}^n}, \quad f \in C^1$$

$$\text{Фом: } x_k = x_0 - \sum_{i=0}^{k-1} \alpha_i \nabla f(x_i)$$

Утв $\forall 0 \leq k \leq \frac{n-1}{2} \exists f \in C^{2,1}_2$ и вып. (и сильно вып.):

$$\forall \text{ Фом } f(x_k) - f(x_{opt}) \geq \frac{1}{32(k+1)^2} \|x_0 - x_{opt}\|^2$$

$$\left(f(x_k) - f(x_{opt}) \geq \frac{\mu}{2} \left(\frac{\sqrt{\mu L} - 1}{\sqrt{\mu L} + 1} \right)^{2k} \|x_0 - x_{opt}\|^2 \right)$$

Угол гон-ва:

$$f(x) = \cos t \left(\frac{1}{3} \langle Ax, x \rangle - \langle x, e_2 \rangle \right)$$

$$e_1 = (1, 0, \dots, 0)^T, \quad A = \begin{bmatrix} 0 & & \\ & \ddots & \\ & & 0 \end{bmatrix}$$

$$\nabla f(x) = \cos t (Ax - e_2), \quad x_0 = 0$$

Ф-ия

Траг. спуск

Ускр. траг. спуск

$\in C^{2,1}_2$ и

$$O\left(\left(1 - \frac{\mu}{L}\right)^k\right)$$

$$O\left(\left(1 - \sqrt{\frac{\mu}{L}}\right)^k\right)$$

и сильно вып.

$\in C^{2,1}$ и

$$O(1/k)$$

$$O(1/k^2)$$

вып.

$$F(x) = f(x) + h(x) \rightarrow \min_x$$

$f \in C^{2,1}_2$ и вып., $h \in C, \notin C^2$, вып. простая

$$F(x) \leq m_k(x) = f(x_k) + \nabla f(x_k)^T (x - x_k) + \frac{1}{2} \|x - x_k\|^2 + h(x) \rightarrow \min_x$$

$$\text{prox}_h(x) = \arg\min_u \left(\frac{1}{2} \|u - x\|^2 + h(u) \right)$$

$$\text{prox-Grad} : x_{k+1} = \arg\min_x m_k(x) =$$

$$\begin{aligned} \text{градиентное} &= \text{prox}_{\frac{1}{2}h} \left(x_k - \frac{1}{2} \nabla f(x_k) \right) = \\ \text{отображение:} &= x_k - \frac{1}{2} G_{1/2}(x_k) \end{aligned}$$

$$G_{1/2}(x_k) = \frac{1}{2} (x_k - x_{k+1}) = \frac{1}{2} (x_k - \text{prox}_{\frac{1}{2}h} (x_k - \frac{1}{2} \nabla f(x_k)))$$

$$\text{критерий останова: } \|G_{1/2}(x_k)\|^2 \leq \varepsilon$$

$$\begin{aligned} \text{оценочная посл-ть: } \psi_k(x) &= \frac{1}{2} \|x - x_0\|^2 + \\ &+ \sum_{i=1}^k \alpha_i (f(x_i) + \langle \nabla f(x_i), x - x_i \rangle + h(x)) \end{aligned}$$

$$A_k = \sum_{i=1}^k \alpha_i, \quad A_0 = 0$$

$$\psi_k^* = \min_x \psi_k(x), \quad v_k = \arg\min_x \psi_k(x) = \text{prox}_{\frac{1}{A_k}h} \left(x_0 - \sum_{i=1}^k \alpha_i \nabla f(x_i) \right)$$

$$\{x_k\}_{k=0}^{\infty}, \quad \{v_k\}_{k=0}^{\infty}, \quad \{y_k\}_{k=0}^{\infty}$$

$$y_k = \frac{A_k x_k + \alpha_{k+1} v_k}{A_k + \alpha_{k+1}}, \quad x_{k+1} = \text{prox}_{1/2h} \left(y_k - \frac{1}{2} \nabla f(y_k) \right)$$

$$\exists (x_k, \psi_k(x)):$$

$$\textcircled{1} A_k F(x_k) \leq \psi_k^*$$

$$\textcircled{2} \psi_k(x) \leq \frac{1}{2} \|x - x_0\|^2 + A_k F(x) \quad \forall x, \text{ тогда}$$

$$F(x_k) - F(x_{opt}) \stackrel{\textcircled{1}}{\leq} \frac{\psi_k^* - F(x_{opt})}{A_k} \leq \frac{\psi_k(x_{opt}) - F(x_{opt})}{A_k} \stackrel{\textcircled{2}}{\leq}$$

$$\stackrel{\textcircled{2}}{\leq} \frac{1}{2 A_k} \|x_{opt} - x_0\|^2 + F(x_{opt}) - F(x_{opt}) = \frac{1}{2 A_k} \|x_0 - x_{opt}\|^2$$

$$\begin{aligned} \textcircled{2} \psi_k(x) &= \frac{1}{2} \|x - x_0\|^2 + \sum_{i=1}^k \alpha_i \underbrace{(f(x_i) + \langle \nabla f(x_i), x - x_i \rangle + h(x))}_{\leq f(x)} \leq \\ &\leq \frac{1}{2} \|x - x_0\|^2 + A_k F(x) \quad \forall \alpha_i \end{aligned}$$

но индукции

$$k=0 : \psi_0^* = 0; \quad A_0 = 0$$

Хотим: $\psi_{k+2}^* \geq A_{k+2} F(x_{k+2})$

$$\psi_{k+2}^* = \min_x \psi_{k+2}(x) = \min_x (\psi_k(x) + a_{k+2} (f(x_{k+2}) + \langle \nabla f(x_{k+2}), x - x_{k+2} \rangle + h(x))) \geq$$

$$\| \psi_k(x) \geq \psi_k^* + \frac{1}{2} \|x - v_k\|^2; \quad \psi_k^* \geq A_k F(x_k)$$

от сильной выпуклости, $k \geq 1$; индукционный переход

$$\geq \min_x (A_k F(x_k) + \frac{1}{2} \|x - v_k\|^2 + a_{k+2} (f(x_{k+2}) +$$

$$+ \langle \nabla f(x_{k+2}), x - x_{k+2} \rangle + h(x))) \geq$$

$$h(x_{k+2}) + \langle \partial h(x_{k+2}), x - x_{k+2} \rangle$$

$$\| A_k F(x_k) \geq F(x_{k+2}) + \langle \partial F(x_{k+2}), x_k - x_{k+2} \rangle$$

$$\geq \min_x (A_{k+2} F(x_{k+2}) + \langle \partial F(x_{k+2}), A_k x_k - A_k x_{k+2} \rangle + \frac{1}{2} \|x - v_k\|^2 + a_{k+2} \langle \partial F(x_{k+2}), x - x_{k+2} \rangle) =$$

$$\| A_k x_k = A_{k+2} y_k - a_{k+2} v_k$$

$$= \min_x (A_{k+2} F(x_{k+2}) + A_{k+2} \langle \partial F(x_{k+2}), y_k - x_{k+2} \rangle + a_{k+2} \langle \partial F(x_{k+2}), x - x_{k+2} \rangle + \frac{1}{2} \|x - v_k\|^2) =$$

$$= \int \arg \min_x = v_k - a_{k+2} \partial F(x_{k+2}) \} = A_{k+2} F(x_{k+2}) +$$

$$+ A_{k+2} \langle \partial F(x_{k+2}), y_k - x_{k+2} \rangle - \frac{a_{k+2}^2}{2} \|\partial F(x_{k+2})\|^2 =$$

$$= A_{k+2} F(x_{k+2}) + A_{k+2} (\langle \partial F(x_{k+2}), y_k - x_{k+2} \rangle - \frac{a_{k+2}^2}{2 A_{k+2}} \|\partial F(x_{k+2})\|^2) \geq A_{k+2} F(x_{k+2})$$

$$\langle \partial F(x_{k+2}), y_k - x_{k+2} \rangle \geq \frac{a_{k+2}^2}{2 A_{k+2}} \|\partial F(x_{k+2})\|^2$$

$$\underline{y_{mb}}. \quad f \in C_1^{2,2}, \quad x = \text{prox}_{\frac{1}{2}h} (y - \frac{1}{2} \nabla f(y)) \Rightarrow$$

$$\exists F'(x) \in \partial F(x) : \langle F'(x), y - x \rangle \geq \frac{1}{2L} \|F'(x)\|^2$$

$$\square \quad G_{1/2}(y) = L(y - x) ; \quad G_{1/2}(y) - \nabla f(y) \in \partial h(x)$$

$$\|\nabla f(y) - \nabla f(x)\|^2 \leq L^2 \|y - x\|^2$$

$$\partial F(x) \ni F'(x) = G_{1/2}(y) - \nabla f(y) + \nabla f(x)$$

$$\|L(y - x) - F'(x)\|^2 \leq L^2 \|y - x\|^2$$

$$L^2 \|y - x\|^2 - 2L \langle y - x, F'(x) \rangle + \|F'(x)\|^2 \leq L^2 \|y - x\|^2$$

$$\langle F'(x), y - x \rangle \geq \frac{1}{2L} \|F'(x)\|^2 \quad \square$$

$$\frac{a_{k+2}^2}{2A_{k+2}} = \frac{1}{2L} \Leftrightarrow \frac{a_{k+2}^2}{A_k + a_{k+2}} = \frac{1}{L}, \quad \text{откуда выразим } a_{k+2}$$

Схема проксимального ускоренного метода Флестерова

$$x_0 = 0, \quad A_0 = 0, \quad L, \quad v_0 = x_0$$

для $k = 0, 1, 2, \dots$

$$\text{Найти } a : \frac{a^2}{A_k + a} = \frac{1}{L} \quad (*)$$

$$y = \frac{A_k x_k + a \cdot v_k}{A_k + a} ; \quad A_{k+2} = A_k + a$$

$$x_{k+2} = \text{prox}_{\frac{1}{2}h} (y_k - \frac{1}{L} \nabla f(y_k))$$

$$\text{если } f(x_{k+2}) > f(y_k) + \langle \nabla f(y_k), x_{k+2} - y_k \rangle + \frac{1}{2} \|x_{k+2} - y_k\|^2, \text{ то } L \leftarrow L \cdot 2 \text{ то } (**)$$

$$\text{если } \|L(y_k - x_{k+2})\|^2 \leq \varepsilon, \text{ то выходи}$$

$$v_{k+2} = \text{prox}_{A_{k+2}h} (x_0 - \sum_{i=2}^{k+2} a_i \nabla f(x_i))$$

$$L \leftarrow L / 2$$

$$F(x_k) - F(x_{opt}) \leq \frac{1}{2A_k} \|x_0 - x_{opt}\|^2 \quad (\leq)$$

$$\frac{a_{k+2}^2}{a_{k+2} + A_k} = \frac{1}{2}, \quad \frac{(A_{k+2} - A_k)^2}{A_{k+2}} = \frac{1}{2},$$

$$A_{k+2} = 2(A_{k+2} - A_k)^2 = 2(\sqrt{A_{k+2}} - \sqrt{A_k})^2(\sqrt{A_{k+2}} + \sqrt{A_k})^2 \leq$$

$$\leq \{A_{k+2} > A_k\} \leq 2(\sqrt{A_{k+2}} - \sqrt{A_k})^2 4A_{k+2}$$

$$\sqrt{A_{k+2}} \geq \sqrt{A_k} + \frac{1}{2\sqrt{L}} \geq \sqrt{A_{k-1}} + \frac{2}{2\sqrt{L}} \geq \dots \geq \sqrt{A_0} + \frac{k+1}{2\sqrt{L}} = \frac{k+1}{2\sqrt{L}}$$

$$A_{k+2} \geq \frac{(k+1)^2}{4L}$$

$$(\leq) \frac{2L}{(k+1)^2} \|x_0 - x_{opt}\|^2$$

сходимость оптимальная
с точностью до фиксированной константы

$$F(\bar{x}_k) = \min_{0 \leq i \leq k} F(x_i)$$

монотонное уменьшение
функционала

семинар

$$\textcircled{2} \partial f_{\mathbb{R}_+^n} = ?$$

стандартные классы выпуклых задач

конические и полунормированные программирование

1 линейное программирование (LP)

$$\min_{x \in \mathbb{R}^n} \langle c, x \rangle \quad \Leftrightarrow \quad \begin{cases} \langle a_i, x \rangle \leq b_i \\ i=1..m \end{cases}$$

$$\text{s.t. } \boxed{Ax \leq b}, \quad Gx = h \quad \Leftrightarrow$$

2 кв. progr. (QP)

$$\min_{x \in \mathbb{R}^n} \frac{1}{2} \langle Qx, x \rangle + \langle p, x \rangle$$

$$Q \in \mathbb{S}_+^n$$

$$\text{s.t. } Ax \leq b, \quad Gx = h$$

③ кв. progr. с кв. op. ($Q(CQP)$)

$$\min \frac{1}{2} \langle Q_0 x, x \rangle + \langle p_0, x \rangle$$

$$\text{s.t. } \frac{1}{2} \langle Q_i x, x \rangle + \langle p_i, x \rangle \leq \hat{c}_i \quad \forall i=1 \dots m$$

$$Gx = h, \quad Q_0, Q_1, \dots, Q_m \in S_+^n$$

④ Конические progr. 2-го порядка ($SOCP$) second order conic programming

$$\min_{x \in \mathbb{R}^n} \langle c, x \rangle$$

$$\text{s.t. } \|A_i x + b_i\|_2 \leq \langle c_i, x \rangle + d_i$$

$$i=1 \dots m$$

$$Gx = h$$

$$x \in K \Rightarrow t \cdot x \in K \quad \forall t \geq 0, \quad K - \text{конус}$$

Конус Лоренца

$$(A_i x + b_i, \langle c_i, x \rangle + d_i) \in K_2 = \{(y, t) \in \mathbb{R}^m \times \mathbb{R} : \|y\|_2 \leq t\}$$

Конические программирование

$$\min_{x \in \mathbb{R}^n} \langle c, x \rangle$$

$$\text{s.t. } Gx = h, \quad x \in K, \quad \text{конус}$$

простые конусы:

$$① K = \mathbb{R}_+^n, \quad LP$$

$$② K = K_2, \quad \text{Лоренца}$$

$$③ K = S_+^n$$

⑤ Полуопределённое программирование (SDP)

$$\min_{x \in \mathbb{R}^n} \langle c, x \rangle \quad \text{матричное}$$

$$\text{s.t. } x_2 A_2^{(i)} + \dots + x_n A_n^{(i)} \leq A_0$$

$$i=1 \dots m, \quad Gx = h, \quad A_j^{(i)} \in S^n$$

$$\min_{x \in S_+^n} \langle c, x \rangle$$

эквибалентно

$$\text{s.t. } GX = H, \quad X \in S_+^n$$

$$LP \subset QP \subset QCQP \subset SOCP \subset SDP$$

Примеры:

$$\textcircled{1} \min_{x \in \mathbb{R}^n} \|Ax - b\|_2 \Leftrightarrow \min_{x \in \mathbb{R}^n} \|Ax - b\|_2^2$$

$$x \geq 0 \quad Q = A^T A \quad (QP)$$

$$\textcircled{2} \min_{x \in \mathbb{R}^n} \|Ax - b\|_2, \quad x \geq 0$$

$$\|Ax - b\|_2 = \sqrt{\sum_{i=1}^m (Ax - b)_i^2} = \sqrt{\sum_{i=1}^m t_i^2}$$

$$t_i \geq (Ax - b)_i, \quad t_i \leq -(Ax - b)_i$$

$$\min_{t \in \mathbb{R}^m} \sqrt{\sum_{i=1}^m t_i^2}$$

$$\| \cdot \|_2, \| \cdot \|_\infty \Rightarrow (LP)$$

$$\begin{cases} t_i \geq (Ax - b)_i \\ t_i \leq -(Ax - b)_i \end{cases} \quad \| \cdot \|_2: t_i \geq |Ax - b|_i \quad \text{epi } \sqrt{\cdot} \quad t \in \mathbb{R}_+^m, \min \sqrt{\sum t_i^2}$$

$$\| \cdot \|_\infty: t_i \geq |Ax - b|_i \quad \forall i \quad t \in \mathbb{R}_+, \min t$$

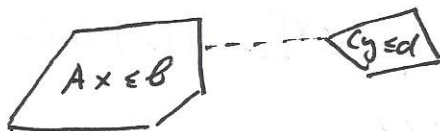
$$\textcircled{3} \min_{x, y \in \mathbb{R}^n} \|x - y\|_2$$

$$\min_{z \in \mathbb{R}^n} \|z\|_2$$

$$s.t. \quad Ax \leq b, \quad Cy \leq d \sim$$

$$T = [I, -I]$$

$$z = (x, y)$$



④

$$\frac{1}{2} \langle Q_i x, x \rangle + \langle p_i, x \rangle \leq z_i, \quad Q_i \in S_+^n \quad S^n$$

$$x_i \in [-1, 1] \Leftrightarrow x_i^2 \leq 1 \Leftrightarrow \begin{cases} x_i^2 \leq 1 & \text{всп.} \\ -x_i^2 \leq -1 & \text{всп.} \end{cases} \quad \underline{NP}$$

⑤

$$\min_t p(t)$$

$$\min_t a t^3 + b t^2 + c t + d$$

$$x_1 = t, \quad x_2 = t, \quad x_3 = x_1 \cdot x_2$$

$$x_3 \geq x_1 \cdot x_2$$

$$-x_3 \geq -x_1 \cdot x_2$$

невыпуклое
ограничение

минимизация произвольных
многочленов NP

$$\textcircled{6} \min_{x \in \mathbb{R}^n} \|Ax - b\|_2 + \|x\|_2$$

 t_2
 t_2

ионические

програм.

$$\min_{x, t_1, t_2} t_2 + t_2$$

$$\text{s.t. } \|Ax - b\|_2 \leq t_2, \|x\|_2 \leq t_2$$

неверно будет считать, что $\|Ax - b\|_2^2 \leq t_2^2, \|x\|_2^2 \leq t_2^2$

$$x^2 - t^2 \leq 0 \quad \begin{bmatrix} 1 & 0 \\ 0 & -1 \end{bmatrix} \notin S_+^n$$

$$\textcircled{7} \min_{x \in \mathbb{R}^n} \|A_0 - x_1 A_1 - \dots - x_n A_n\|_F, \text{ op.}$$

ив.и., полуопр.

$$A_0, A_1, \dots, A_n \in S^n$$

$$\min_{x \in \mathbb{R}^n} \lambda_{\max}(A_0 - x_1 A_1 - \dots - x_n A_n)$$

$$\min_{x, t} t, \text{ s.t. } \lambda_{\max}(A(x)) \geq t \Leftrightarrow A(x) \geq tI$$

$$\min_{x, t} t, \text{ s.t. } t \geq \sigma_{\max}(A(x))$$

$$t^2 \geq \lambda_{\max}(A(x)A(x)^T), t^2 I \geq A(x)A(x)^T$$

Лемма о выполнении Шура

$$M = \begin{bmatrix} A \rightarrow B \\ \uparrow \quad \downarrow \\ B^T \leftarrow C \end{bmatrix} \quad A \in S^n, C \in S^m, B \in \mathbb{R}^{n \times m}$$

Пусть $A > 0$. Тогда $M \geq 0 \Leftrightarrow \underbrace{C - B^T A^{-1} B}_{\text{гон-е Шура}} \geq 0$

$$\square \quad z = (x, y) \in \mathbb{R}^n \times \mathbb{R}^m$$

$$0 \leq \langle Mz, z \rangle = \langle Ax, x \rangle + 2\langle x, By \rangle + \langle Cy, y \rangle \rightarrow \min_x$$

$$Ax^* + By = 0, x^* = -A^{-1}By, \langle (C - B^T A^{-1} B)y, y \rangle \geq 0 \quad \square$$

тогда

переход от невыпуклых к линейным

$$\begin{bmatrix} tI \rightarrow A(x)^T \\ \uparrow \\ A(x) \leftarrow tI \end{bmatrix} \geq 0 \Leftrightarrow tI - \frac{A(x)A(x)^T}{t} \geq 0$$

$$t^2 I \geq A(x)A(x)^T$$

① L D

$$Ax \leq b \Leftrightarrow \text{Diag} \{b - Ax\} \geq 0 \quad (\text{SDP})$$

② Q D

$$\frac{1}{2} \langle Qx, x \rangle + \langle p, x \rangle \leq 2$$

$$\frac{1}{2} x^T Q x \leq 2 - p^T x$$

$$Q = L^T L \Rightarrow \frac{1}{2} x^T L^T L x \leq 2 - p^T x$$

lemma Угpa: ↗

$$\begin{bmatrix} 2I & Lx \\ (Lx)^T & (2 - p^T x)I \end{bmatrix} \geq 0$$

$$\begin{bmatrix} tI & x \\ x^T & t \end{bmatrix} \geq 0$$

↕

③ SDCP

$$\|x\|_2 \leq t \Leftrightarrow \|x\|_2^2 \leq t^2 \Leftrightarrow t^2 - x^T x \geq 0$$