VALIDATOR FOR FLAMBDA2 SIMPLIFIER

1. Reduction strategy

This language has a call-by-value style reduction strategy. Notice the unusual [LetR] rule—the expression N refers to an expression in the normal form, which may refer to a normalized effectful expression. This rule is *not* analogous to the [ApplyContR] rule, since the lambda abstraction is always implicit in let expressions, ensuring that the "lefthand-side" of the application is always a value. This is necessary because for the case of several effectful expressions (such as a print statement), inlining the occurence of the expression multiple times will be behaviorally different from the original expression.

The [Apply-β] rule describes the case when the callee is a lambda expression, and the argument is fully evaluated. The expression is beta-reduced, then the resulting value get passed on as an argument to either the return or exception continuation, depending on whether or not the expression throws an exception.

LET-β
$\mathsf{let}\ x\ =\ v\ \mathsf{in}\ e_2 \longrightarrow e_2\ [x \setminus v]$

LETL
$$\frac{e_1 \longrightarrow e_1'}{\mathsf{let}\ x\ =\ e_1\ \mathsf{in}\ e_2 \longrightarrow \mathsf{let}\ x\ =\ e_1'\ \mathsf{in}\ e_2}$$

LETR
$$\frac{e_2 \longrightarrow e_2'}{\mathsf{let}\ x\ =\ N\ \mathsf{in}\ e_2 \longrightarrow \mathsf{let}\ x\ =\ N\ \mathsf{in}\ e_2'}$$

LETCONT-β
$e_1\ \mathsf{where}\ k\ \overrightarrow{args}\ =\ e_2 \longrightarrow e_1\ [k \setminus \lambda\ \overrightarrow{args}.\ e_2]$

APPLYCONTR
$$\frac{\overrightarrow{args} \longrightarrow \overrightarrow{args}'}{v\ \overrightarrow{args} \longrightarrow v\ \overrightarrow{args}'}$$

APPLYCONTL
$$\frac{k \longrightarrow k'}{k\ \overrightarrow{args} \longrightarrow k'\ \overrightarrow{args}}$$

APPLYCONT-β
$(\lambda\ x.\ e)\ v \longrightarrow e\ [x \setminus v]$

APPLYR
$$\frac{\overrightarrow{args} \longrightarrow \overrightarrow{args}'}{e\ res_k\ exn_k\ \overrightarrow{args} \longrightarrow e\ res_k\ exn_k\ \overrightarrow{args}'}$$

APPLYL
$$\frac{e \longrightarrow e'}{e\ res_k\ exn_k\ \overrightarrow{args} \longrightarrow e'\ res_k\ exn_k\ \overrightarrow{args}}$$

APPLY-β
$(\lambda\ \overrightarrow{x}.\ e)\ res_k\ exn_k\ \overrightarrow{v} \longrightarrow e\ [x \setminus \overrightarrow{v}] \hookrightarrow res_k/exn_k$

SWITCH
$$\frac{e \longrightarrow e'}{\mathsf{switch}\ (e)\ arms \longrightarrow \mathsf{switch}\ (e')\ arms}$$

SWITCH-β
$\mathsf{switch}\ (v)\ [x \mapsto e] \longrightarrow e\ [x \setminus v]$

$$\overrightarrow{args} \longrightarrow \overrightarrow{args}' := \forall a, a \in \overrightarrow{args}. \exists a', a \longrightarrow^* a' \wedge a \in \overrightarrow{args}'$$

2. Rewrite Rules

FLATTENMATCH
$\mathsf{switch}\ (\mathsf{switch}\ (e_1)\ [A \mapsto e_2 : B|..])\ [B \mapsto e_2'|..] \longrightarrow \mathsf{switch}\ (e_1)\ [A \mapsto e_2'\ [B \setminus e_2]|..]$

3. Features

A wishlist of desirable inlining/semantic features to support for the validator.

3.1. **Inlining.**

- function calls
- recursive functions
- inlining (direct calls, within same function)
- cross-module inlining
- low-priority: locals

3.2. **Semantics.**

- mutable state
- exceptions
- effects (printing, etc.)
- external calls

3.3. **Primitives evaluation.**

- arithmetic evaluation: commutative and associative laws for arithmetic? It is likely that the commutative/associative laws are not necessary for the simplifier
- block-based primitives (makeblock, loading from block) The blocks have a tag, corresponding to the constructors (i.e. tag0 is the first constructor) values either are immediate tags or blocks Mutability corresponds to reference cells

  Being able to treat the block-related primitives will resolve supporting the structures below (except for arrays, which have a tricky case involving storing floating-point values. See floating-point valued array optimization)

TODO: Refactor [simplify_primitive].

3.4. **Supported structures.**

- structs
- tuples
- lists
- arrays