



Chap. 10

Other Models of Turing Machines

10.11 이론은- STM과 power 같 아?
?

M 이 STM에서 simulate. 하는 과정

Agenda of Chapter 10

- Minor variations on the Turing machine theme
 - Turing machines with a Stay-option
 - Turing machines with multiple tracks
 - Turing machines with semi-infinite tape
 - The off-line Turing machine

- Turing machines with more complex storage
 - Multitape Turing machines
 - Multidimensional Turing machines

- Nondeterministic Turing machines
- A Universal Turing machine

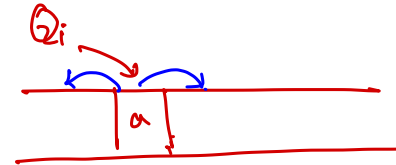
power가 동등함.
 \Leftrightarrow STM 이. SSTM은 변환함

Turing Machine with minor variations (1/3)

□ Turing Machine with a stay-option

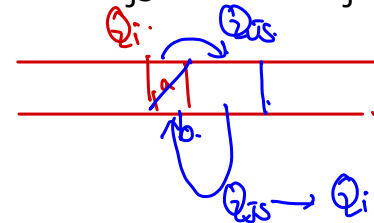
- $\delta: Q \times \Gamma \rightarrow Q \times \Gamma \times \{L, R, S\}$
- simulation with standard TM

$$\delta(q_i, a) = (q_j, b, S) \Leftrightarrow \delta(q_i, a) = (q_{js}, b, R), \delta(q_{js}, c) = (q_j, c, L) \text{ for all } c$$

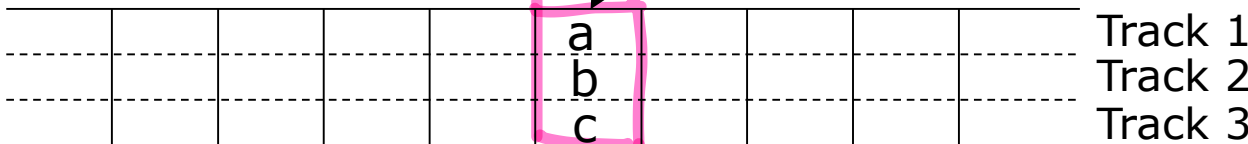


"q장에서 대충으로 봤을 것"

$$\delta: Q \times \Gamma \rightarrow Q \times \Gamma \times \{L, R\}$$



□ Turing Machine with multiple tracks



- String on three tracks can be considered as a symbol

$$\delta: Q \times \Gamma \times \Gamma \times \Gamma \rightarrow Q \times \Gamma \times \Gamma \times \Gamma \times \{L, R\}$$

$$\delta(Q_i, a, b, c) \rightarrow (Q_j, def, L)$$

(STM은 대충 simulation.)

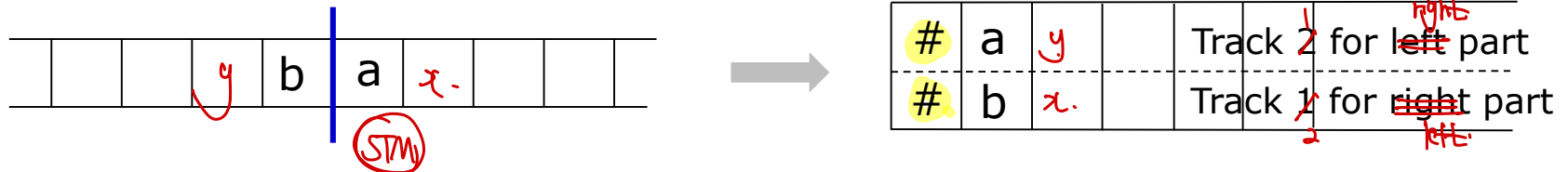
abc
 4bit symbol.

4bit tape alphabet: Γ_{13} 정도

Turing Machine with minor variations (2/3)

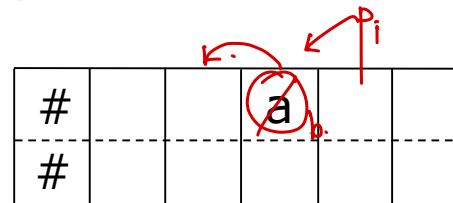
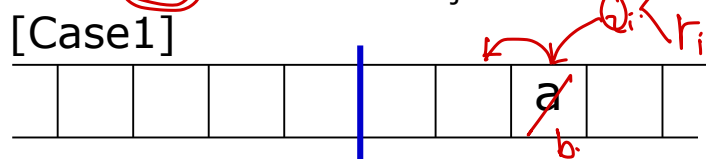
□ Turing machine with semi-infinite tape

- Tape is unbounded only in one direction
- Semi-infinite tape with 2-tracks can simulate Standard TM.

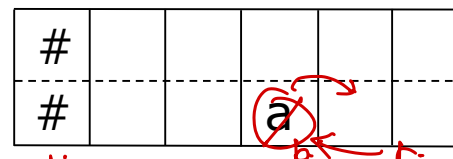
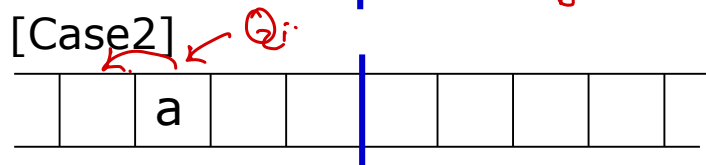


– Simulation of STM using semi-STM

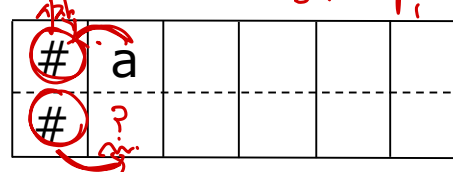
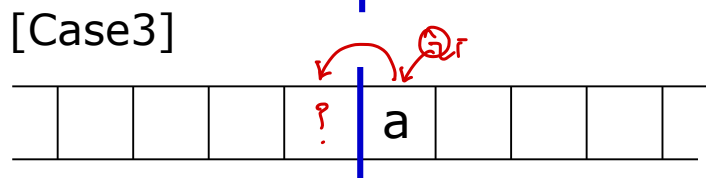
STM: $\delta(q_i, a) = (q_j, b, L)$



$$\delta(p_i, (a, *)) = (p_j, (b, *), L)$$



$$\delta(r_i, (*, a)) = (r_j, (*, b), R)$$



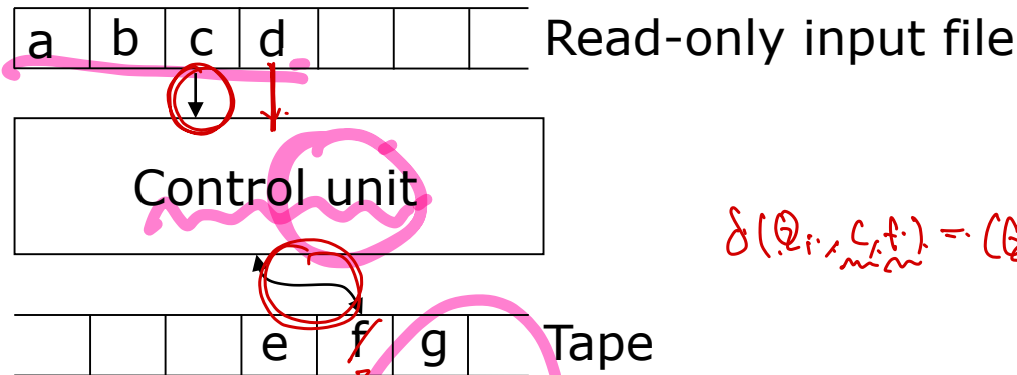
$$\delta(p_j, (\#, \#)) = (r_j, (\#, \#), R)$$

$$\delta(r_j, (\#, \#)) = (p_j, (\#, \#), R)$$

Turing Machine with minor variations (3/3)

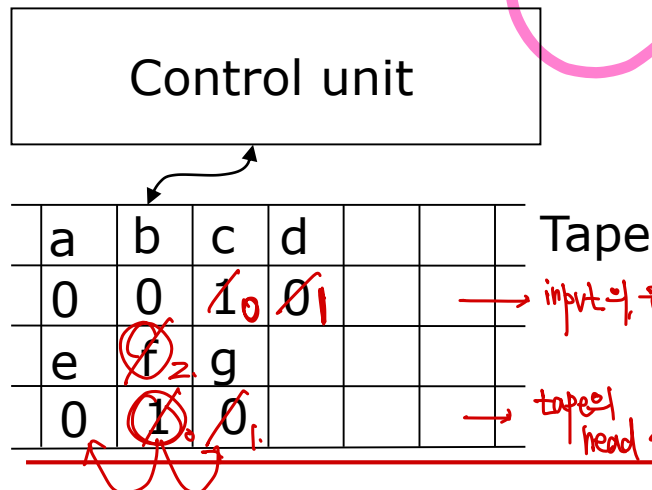
Off-line Turing machine

- Turing machine with input file



$$\delta(Q_i, c, f) = (Q_j, z, L)$$

- Standard TM with 4 tracks can simulate off-line TM.



Operations:

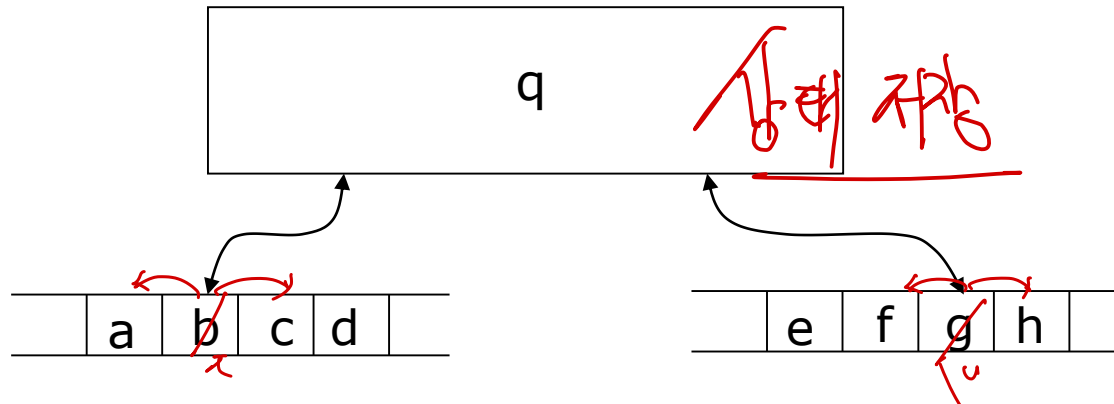
- Find 1 on track 2.
- Read input symbol, save state, change symbols on track 2 → input을 읽고 input을 쓴다
- Find 1 on track 4 → tape의 head 찾기
- Read tape symbol, change state, and change symbols on track 3 & 4.

Multitape Turing Machine

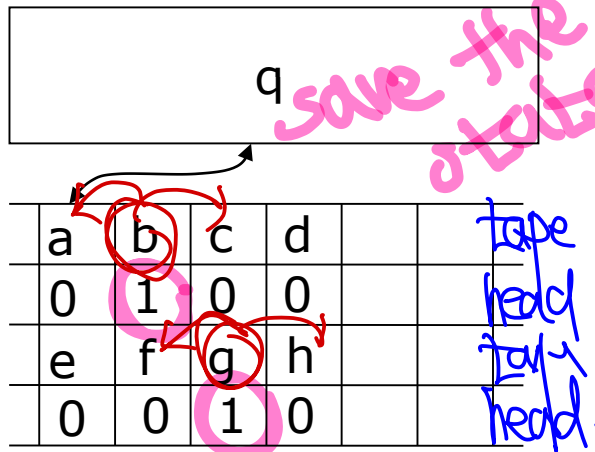
□ n-tape Turing machine

$$\delta : Q \times \Gamma^n \rightarrow Q \times \Gamma^n \times \{L, R\}^n$$

$$\delta : Q \times \Gamma_1 \times \Gamma_2 \rightarrow Q \times \Gamma_1 \times \Gamma_2 \times \{L, R\} \times \{L, R\}$$



— Standard TM with $2n$ tracks can simulate n -tape TM.

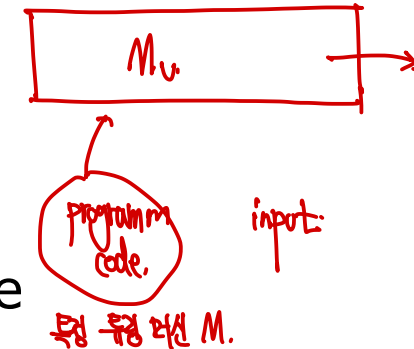


Operations:

- Find 1 on track 2
- Read tape 1 symbol, save the state, and change symbol on track 1 & 2
- Find 1 on track 4
- Read tape 2 symbol, change state, and change symbols on track 3 & 4.

Universal Turing Machine (1/2)

- Turing machine vs. Digital computer
 - Special purpose vs. general purpose.
- To overcome the objection
 - Design a reprogrammable Turing machine
- A universal Turing machine M_u
 - input : a description of any Turing Machine M and a string w
 - Simulate the computation of M on w 1 0 1 0 1 1 0 1 1 0 1
- Description of Turing machines.
 - Assumption
 - $Q = \{q_1, q_2, \dots, q_n\}$, q_1 = initial state, $F = \{q_2\}$, $\Gamma = \{a_1, a_2, \dots, a_m\}$, $a_1 = \square$
 - Encoding
 - $q_i = 11\dots 1$ ($|q_i| = i$), $a_i = 11\dots 1$ ($|a_i| = i$), 0 separator between 1's.
 - Ex) $\delta(q_1, a_2) = (q_2, a_3, L) \rightarrow \dots \underbrace{10}_{q_1} \underbrace{110}_{a_2} \underbrace{110}_{q_2} \underbrace{1110}_{a_3} \underbrace{10}_{L}$ T.M. 10110



$$\delta(q_1, a_2) = (q_2, a_3, L)$$

1 11 11 111

L = 1
R = 1

Universal Turing Machine (2/2)

- A universal Turing Machine M_u simulating an M with input w
- Consider M_u with three tapes.
 - Operations of M_u
 1. Look at tapes 2 & 3.
 2. Consults tape 1 to see what M would do.
 3. Modify tapes 2 & 3.

