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# **Lecture #16: Dynamic programming**

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School of Computer Science and Engineering  
Kyungpook National University (KNU)

Woo-Jeoung Nam



# Agenda

- **Dynamic programming**
  - Rod cutting problem
  - Fibonacci sequence problem
  - Knapsack problem
  - Longest common subsequence problem
  - **Sequence alignment**
    - Global alignment (Needleman-Wunsch algorithm)
    - Local alignment (Smith-waterman algorithm)



# The Longest Common Sequence (LCS) Problem

- In many applications, we are required to search for a word within a document
- Such problems are normally referred to the string pattern matching problem

Hello World!

| • || | | || | •

Hallo World?

- The LCS is the longest common subsequence between sequence X and Y



# The Longest Common Sequence (LCS) Problem

$$c[i, j] = \begin{cases} 0 & \text{If } i = 0, \text{ or } j = 0 \\ c[i - 1, j - 1] + 1 & \text{If } i, j > 0 \text{ and } x_i = y_i \\ \max(c[i, j - 1], c[i - 1, j]) & \text{If } i, j > 0 \text{ and } x_i \neq y_i \end{cases}$$

		<i>j</i>	0	1	2	3	4	5	6
		$y_j$		<i>B</i>	<i>D</i>	<i>C</i>	<i>A</i>	<i>B</i>	<i>A</i>
<i>i</i>	$x_i$								
0		0	0	0	0	0	0	0	0
1	<i>A</i>	0	0	0	0	1	1	1	1
2	<i>B</i>	0	1	1	1	1	2	2	2
3	<i>C</i>	0	1	1	2	2	2	2	2
4	<i>B</i>	0	1	1	1	1	3	3	3
5	<i>D</i>	0	1	2	2	2	3	3	3
6	<i>A</i>	0	1	2	2	3	3	4	4
7	<i>B</i>	0	1	2	2	3	4	4	4



# The Longest Common Sequence (LCS) Problem

		$j$	0	1	2	3	4	5	6
		$y_j$		$B$	$D$	$C$	$A$	$B$	$A$
$i$	$x_i$								
0		0	0	0	0	0	0	0	0
1	$A$	0	↑ 0	↑ 0	↑ 0	↖ 1	← 1	↖ 1	
2	$B$	0	↖ 1	↖ 1	← 1	↑ 1	↖ 2	← 2	
3	$C$	0	↑ 1	↑ 1	↖ 2	← 2	↑ 2	↑ 2	
4	$B$	0	↖ 1	↑ 1	↑ 2	↑ 2	↖ 3	← 3	
5	$D$	0	↑ 1	↖ 2	↑ 2	↑ 2	↑ 3	↑ 3	
6	$A$	0	↑ 1	↑ 2	↑ 2	↖ 3	↑ 3	↖ 4	
7	$B$	0	↖ 1	↑ 2	↑ 2	↑ 3	↖ 4	↑ 4	↑ 4

$$LCS = \langle B, C, B, A \rangle$$

What can we say about the “4” in  $c[7,6]$ ?

Does the LCS have some meaning in real life?

E.g.,  
 $X = \text{Hello World}$ ,  $Y = \text{Yexxo Wxrly}$   
 $Z = \text{eorl}$



# The Longest Common Sequence (LCS) Problem

- The LCS edit distance is not well applicable to real world problems
- Furthermore, every match or mismatch has the same edit distance
  - 0 for mismatch
  - 1 for match

$$c[i, j] = \begin{cases} 0 & \text{If } i = 0, \text{ or } j = 0 \\ c[i-1, j-1] + 1 & \text{If } i, j > 0 \text{ and } x_i = y_i \\ \max(c[i, j-1], c[i-1, j]) & \text{If } i, j > 0 \text{ and } x_i \neq y_i \end{cases}$$

- However, in many cases, even if two characters mismatch, some mismatches are more meaningful than others (and thus, the mismatch score shall differ)
- Based on the classical LCS problem, it can be extended to real world problems with slight modifications
- 즉, 불일치한다고 해서 무의미한것은 아니다

↓  
실제 real world 에서는 의미가 있다.



# The Longest Common Sequence (LCS) Problem

Hello World!  
| • | | | | | •  
Hallo World?

171,476 words in Oxford  
English Dictionary

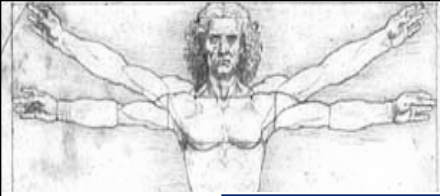


→ 원본 문자열

AATGCGACCTGA  
| • | • | | | | | • |  
AGTCCGACCTCA

3.3 billion characters in a single  
cell nucleus. Huge number of  
possible words.

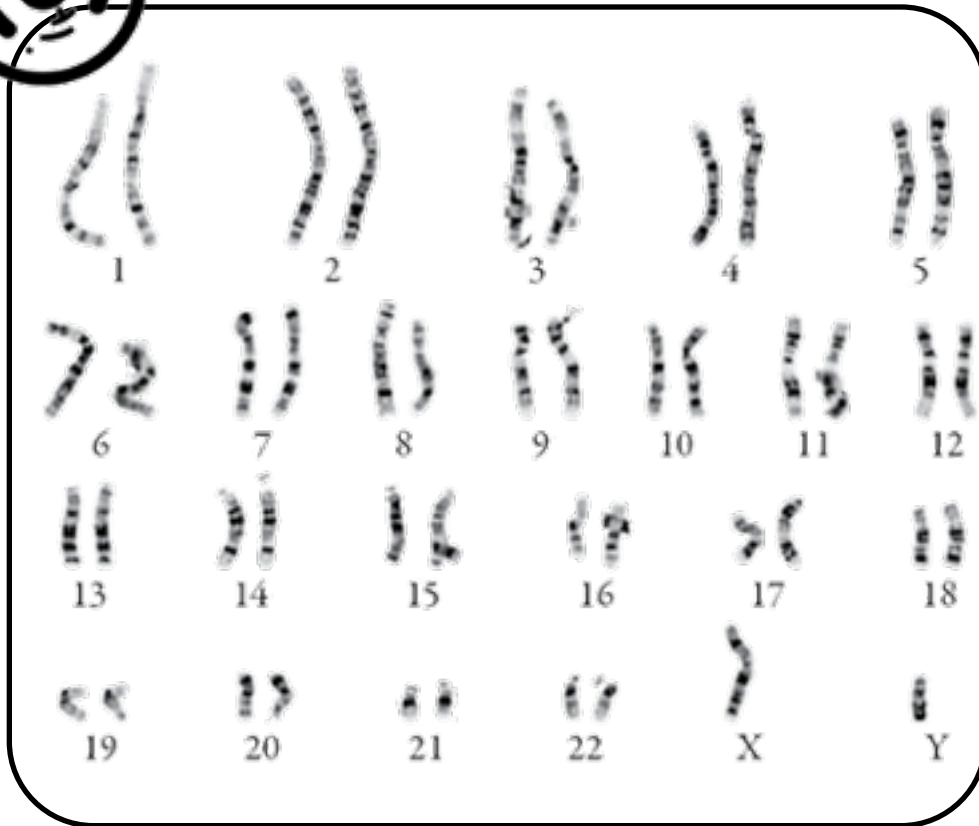




# The Longest Common Sequence (LCS) Problem



A single cell



Data



Chr 1

...

Chr Y

AATGCGACCTGA...AATGCGACCTGATTCAGACAG

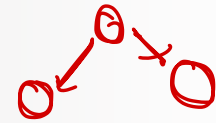
Length = 3.2 billion characters

Is "GCGAC" in the string?





# Sequence alignment



## Homologs



➤ 공통된 조상을 가지면서 비슷한 DNA sequence를 가지는 종들

- Orthologs: 다른 종임에도 불구하고 공통된 조상으로부터 종 분화(species divergence)로 인해 유사한 DNA sequence를 가지는 경우
- Paralogs: 같은 종에서 gene duplication으로 인해 유사한 DNA sequence를 가짐에도 서로 기능이 다른 경우

## Analogs

➤ 공통된 조상이 없는데도 불구하고 비슷한 DNA sequence를 가지는 종들

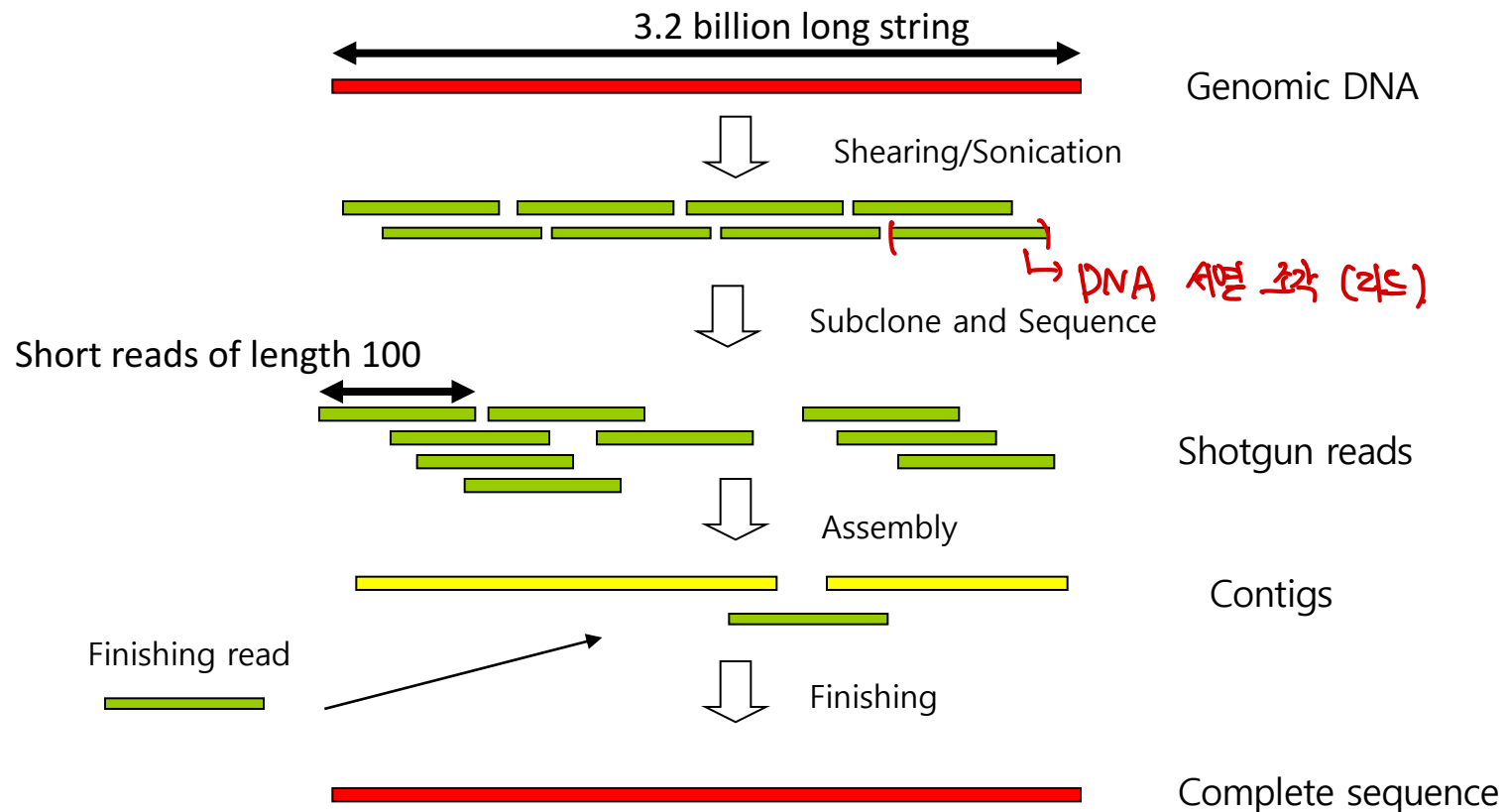
- Sequence Alignment는 기능이 알려진 DNA, protein의 sequence와 기능이 알려지지 않은 sequence의 유사도를 계산하여 기능을 유추하거나 공통의 조상을 찾는데 활용

영가서면  
발음.



# The Assembly Problem

- DNA 시퀀싱에서의 어셈블리 문제는 많은 수의 짧은 DNA 서열 조각(리드)을 사용하여 원래의 긴 DNA 서열을 정확하게 재구성하는 과정





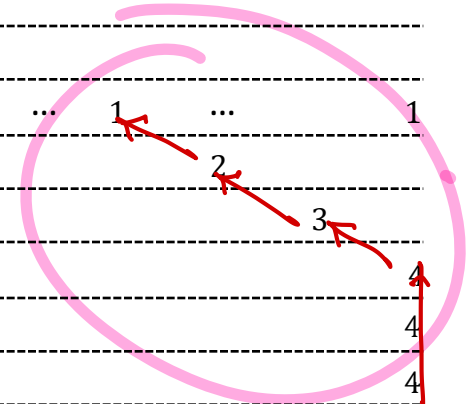
# The Longest Common Sequence (LCS) Problem Revisited

- What is the LCS if the X and Y sequences were as below?

➤  $X = \text{XXX}\underline{\text{ABB}}\text{XXX}\underline{\text{ABB}}\text{XXX}\underline{\text{ABBA}}$

➤  $Y = \text{YYYABBA}\text{YY}$

	$j$	0	1	2	3	4	5	6	...	.	.	$m-1$	$m$
$i$			X	X	A	B	B	A	...	A	B	B	A
0		0	0		...								
1	Y	0	0										
2	Y	0	0										
3	A				1	1	1	1	...	1	...		1
4	B	:				2							
5	B						3						
6	A							4					
7	Y												4
8	Y												4














- |     | $j$ | 0 | 1   | 2   | 3   | 4   | 5   | 6   | ... | .   | .   | $m-1$ | $m$ |   |
|-----|-----|---|-----|-----|-----|-----|-----|-----|-----|-----|-----|-------|-----|---|
| $i$ |     |   | $X$ | $X$ | $A$ | $B$ | $B$ | $A$ | ... | $A$ | $B$ | $B$   | $A$ |   |
| 0   |     | 0 | 0   |     | ... |     |     |     |     |     |     |       |     |   |
| 1   | $Y$ | 0 | 0   |     |     |     |     |     |     |     |     |       |     |   |
| 2   | $Y$ | 0 | 0   |     |     |     |     |     |     |     |     |       |     |   |
| 3   | $A$ |   |     |     | 1   | 1   | 1   | 1   | ... |     | 1   | ...   | 1   |   |
| 4   | $B$ | : |     |     |     | 2   |     |     |     |     |     | 2     |     |   |
| 5   | $B$ |   |     |     |     |     | 3   |     |     |     |     |       | 3   |   |
| 6   | $A$ |   |     |     |     |     |     | 4   |     |     |     |       |     | 4 |
| 7   | $Y$ |   |     |     |     |     |     |     |     |     |     |       |     | 4 |
| 8   | $Y$ |   |     |     |     |     |     |     |     |     |     |       |     | 4 |



-  **Z=ABBA**

- What if we want to find all “ABBA” occurrences?

	$j$	0	1	2	3	4	5	6	...	.	.	$m-1$	$m$	
			$X$	$X$	$A$	$B$	$B$	$A$	...	$A$	$B$	$B$	$A$	
$i$														
0		0	0	...										
1	$Y$	0	0											
2	$Y$	0	0											
3	$A$				1	1	1	1	...		1	...	1	
4	$B$	:				2						2		
5	$B$						3						3	
6	$A$							4						4
7	$Y$													4
8	$Y$													4



# The Longest Common Sequence (LCS) Problem Revisited

- This time, let's observe what the LCS if the X and Y sequences were as below?
  - X="XXXHELLOXXXHELLOXXXHELLO"
  - Y="YYHEIIOYY"
  - Z="HEO"
- But are we interested in "HEO"? In other words, do we want to accept Z as a result?



# Meaningful string pattern matching

- As we have observed in the previous LCS problem, we can actually score a common subsequence, so that it has some meaningful value
- The scoring scheme is key
- Modifications to LCS → 점수를 매기는 방식이 핵심이다.
  1. A different backtracking scheme
  2. A different scoring scheme (the previous scored each match by 1)
- Such modification leads to a slightly different problem referred to as the sequence alignment problem



# The Sequence Alignment Problem

- Objective: Given a string X and Y, find the best alignment that maximizes the alignment score (= minimize mismatches)
- An alignment is scored by three criteria

- Match: when  $x_i = y_j$
- Mismatch: when  $x_i \neq y_j$
- Gap: when  $x_i$  or  $y_j$  aligned to a gap

↓  
alignment score를 최대로 하는  
방법으로..

↓  
(gap) 을 주는 위치에 대해서  
score가 바뀐다!

- For sequences

- X="Hello\_World"
- Y="Hwllo\_qWorld"

score 1 ----  
score -1 ●  
score -2 ●

X=    Hello\_World□  
      | | | | |  
      ● | | | | ● ● ● ● ● ● ●  
Y=    Hwllo\_qWorld

$$\text{Alignment score} = 5 \cdot 1 + 6 \cdot -1 + 1 \cdot -2 = -3$$





# The Sequence Alignment Problem

- Can we find another alignment that maximizes the alignment score?
- Other alignments can be found by 1) adding gaps, 2) refining the matching score matrix

X= Hello\_World    
| | | | |  
Y= Hwllo\_qWorld

Alignment score =  $5 \cdot 1 + 6 \cdot -1 + 1 \cdot -2 = -3$

X= Hello\_   World  
| | | | |  
Y= Hwllo\_qWorld

Alignment score =  $10 \cdot 1 + 1 \cdot -1 + 1 \cdot -2 = \underline{\underline{7}}$

- During the alignment, we have choices
  - if  $x_i$  and  $y_j$  match, we are happy
  - otherwise, we have to decide whether to align to a mismatch or a gap (which maximizes the alignment score?)



# Two types of sequence alignment algorithms

X= ATCGGCTAGGAACACGACGAGCAGCT

Y= GTGCCGCTGGATGAGTGGTCAGTCTG

(Match score=1, mismatch score=0)

## Global alignment

- Needleman-Wunsch[1]
- 두 시퀀스의 전체 길이에 걸쳐 최적의 정렬을 찾는 문제
- 전체 구조가 유사한 시퀀스들 사이의 정렬에 적합

→ 서열정렬 알고리즘 → 여러 DNA 혹은 단백질 서열을 정렬하는 알고리즘

두 시퀀스 전체에서 최상의 정렬 (서열 길이가 비슷)

ATCG-GCTAGGAACACGACG-AGCA-GCT  
| | | | | | | | | | | | | | |  
GTGCCGCTGG-ATGAGTGGTCAGCTTG  
Alignment score=10

→ 완전 정렬

## Local alignment

- Smith-Waterman algorithm [2]
- 두 시퀀스에서 가장 유사한 부분 서열(subsequence)을 찾는 문제
- 길이가 다르거나 일부 구조만 유사한 시퀀스들 사이의 정렬에 적합

→ 부분 유사성!

---GCT---GCT---  
| | | | |  
---GCT---GCT---  
Alignment score=6

Which is better?

전체 서열이 유사하지 않더라도

두 시퀀스에서 가장 유사한 부분은 중요



# Global alignment

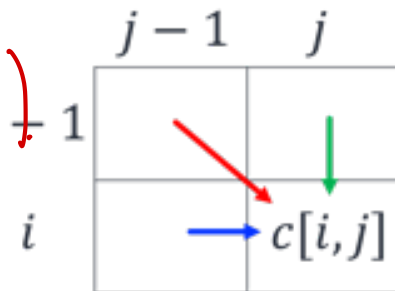
- Global alignment is based on the LCS problem
- However, instead of searching for a LCS, we now search for some subsequence that yields the highest alignment score
- This is still an optimization problem, since we aim to find an alignment that gives us the maximum alignment score (also retains the optimal substructure)
- Again, for sequences  $X = \langle A, B, C, B, D, A, B \rangle$  and  $Y = \langle B, D, C, A, B, A \rangle$ , the matrix  $C$  will be computed similar to the LCS problem
- However, the scoring scheme is different
  - (S and P will be discussed later)

$$\frac{c[i-1, j-1] + S}{c[i][j-1]}, \text{ match}$$

$$c[i, j] = \max \begin{cases} c[i-1, j-1] + S \\ c[i, j-1] + P \\ c[i-1, j] + P \end{cases}$$

gap is a penalty  
의 중재이다.

$c[i, j] =$





# Global alignment

- For sequences  $X = \langle GATTACA \rangle$  and  $Y = \langle GCATGCG \rangle$ , find the best alignment  $Z$ .
- Scoring scheme

G A T T A C A  
G C A T G C G

$$c[i, j] = \max \begin{cases} c[i-1, j-1] + S(i, j) \\ c[i, j-1] + P \\ c[i-1, j] + P \end{cases}$$

$$S(i, j) = \begin{cases} 1 & x_i = y_j \\ -1 & \text{else} \end{cases}$$

$$P = -1 \rightarrow \text{gap}$$

## 1. Initialize

$$c[i, 0] = i * P \text{ for } \forall i$$

$$c[0, j] = j * P \text{ for } \forall j$$

## 2. Compute $c[i, j]$ , for $\forall i, j$

		j							
		0	1	2	3	4	5	6	7
		$y_j$	G	C	A	T	G	C	G
0	$x_i$	0	-1	-2	-3	-4	-5	-6	-7
1	G	-1	0	-1	-1	-2	-3	-4	-5
2	A	-2	0	0	1	0	-1	-2	-3
3	T	-3	-1	-1	0	2	1	0	-1
4	T	-4	-2	-2	-1	1	1	0	-1
5	A	-5	-3	-3	-1	0	0	0	-1
6	C	-6	-4	-2	-2	-1	-1	1	0
7	A	-7	-5	-3	-1	-2	-2	0	0

初始化

GC  
G-  
-1 = 0



# Global alignment

- For each computation of  $c[i, j]$  we fill the backtrack matrix in parallel

$$c[i, j] = \max \begin{cases} c[i-1, j-1] + S(i, j) \\ c[i, j-1] + P \\ c[i-1, j] + P \end{cases}$$

Handwritten red notes:  $-A$ ,  $-A$ ,  $2$ , and arrows pointing to the three cases in the max function.

- In case of tie, priority is as (not mandatory)

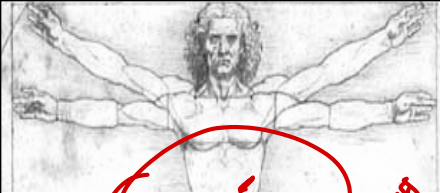
- Left  $\leftarrow$
- Up  $\uparrow$
- Diagonal  $\nearrow$



Handwritten red note:  $0$ 를 정한 생각할 수 있다.

j	0	1	2	3	4	5	6	7
$y_j$	G	C	A	T	G	C	G	
i	0	$x_i$						
0	0	-1	-2	-3	-4	-5	-6	-7
1	-1	1	0	-1	-2	-3	-4	-5
2	-2	0	0	1	0	-1	-2	-3
3	-3	-1	-1	0	2	1	0	-1
4	-4	-2	-2	-1	1	1	0	-1
5	-5	-3	-3	-1	0	0	0	-1
6	-6	-4	-2	-2	-1	-1	1	0
7	-7	-5	-3	-1	-2	-2	0	0

Handwritten red arrows and boxes highlight specific cells and paths in the matrix, including a pink box around the cell (1,2) and a pink arrow pointing from (1,2) to (2,3).



# Global alignment

Handwritten notes at the top right:   
 $G C A T G$   
 $G - A T -$

Handwritten notes in a red circle:   
 $G C A T$   
 $G - A T$

		j	0	1	2	3	4	5	6	7
		y <sub>j</sub>	G	C	A	T	G	C	G	
i	x <sub>i</sub>	0	0	-1	-2	-3	-4	-5	-6	-7
1	G	-1	1	0	-1	-2	-3	-4	-5	
2	A	-2	0	1	0	-1	-2	-3	-4	
3	T	-3	-1	0	1	0	-1	-2	-3	
4	T	-4	-2	-1	0	1	0	-1	-2	
5	A	-5	-3	-2	-1	0	-1	-2	-3	
6	C	-6	-4	-3	-2	-1	0	-1	-2	
7	A	-7	-5	-4	-3	-2	-1	0	-1	

Handwritten notes in a pink box:   
 $G C A T G C G$   
 $G - - - - -$

Handwritten note:  $A-G = -5$

Handwritten note in a circle:  $P = -1$

Handwritten note: mismatch.

Handwritten note: optimal substructure.

Backtrack starts here

Alignment diagram:  
 $G-ATTACA$   
 $GCATG-CG$

Alignment score=0

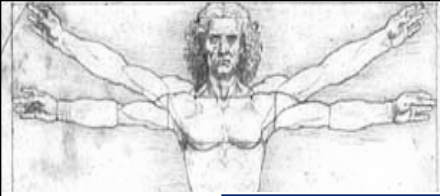
Handwritten notes:  $A C$ ,  $G C$

seq 1 → G - A T T A C A

seq 2 → G C A - T G

gap gap

Q



# Local alignment

- Similar to global alignment, but has a different objective
- Global alignment tries to align the entire sequences X and Y

X= Hello\_ World  
Y= Hwlllo\_qWorld

... sequence 는 정렬해 줘

- “Local” alignment tries to find smaller subsequence in X that well aligns to a subsequence in Y

X= Hello\_ World  
Y= [ ] Wor [ ]

gaps





# Local alignment

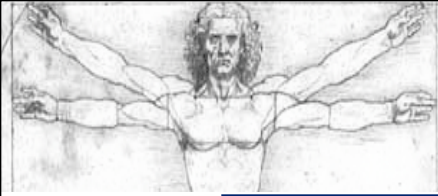
- Given two sequences X and Y, find two subsequences, x and y, whose alignment has the **highest score** amongst all subsequence pairs

X= Hello\_ World  
Y= Hwlllo\_qWorld

Global alignment (indicated by a red bracket and the word "global" written in red cursive).

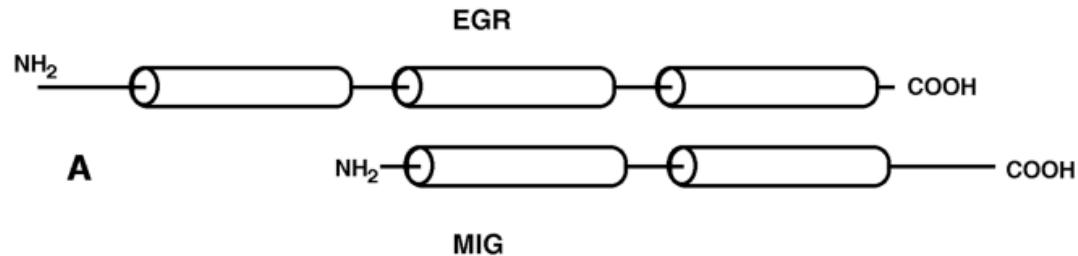
X= Hello\_ World  
Y= [ ] Wor [ ]

Local alignment (indicated by red dots in sequence X and red boxes in sequence Y, with arrows pointing to the boxes labeled "gaps").



# Local alignment - meaning

- Given two sequences X and Y, find two subsequences, x and y, whose alignment has the **highest score** amongst all subsequence pairs



EGR4_HUMAN	KA	[FACPVESCVRSFARSDELNRHLRIH]	TGHKP	[FQCRICLRNFSRSDHLTSHVRTH]	TGEKP	[FACDV--CGRRFARSDEKKRHSKVH]
EGR4_RAT	KA	[FACPVESCVRTFARSDELNRHLRIH]	TGHKP	[FQCRICLRNFSRSDHLTTHVRTH]	TGEKP	[FACDV--CGRRFARSDEKKRHSKVH]
EGR3_HUMAN	RP	[HACPAEGCDRRFSRSDELTRHLRIH]	TGHKP	[FQCRICMRFSRSDHLTTHIRTH]	TGEKP	[FACEF--CGRKFARSDEKRHAKIH]
EGR3_RAT	RP	[HACPAEGCDRRFSRSDELTRHLRIH]	TGHKP	[FQCRICMRFSRSDHLTTHIRTH]	TGEKP	[FACEF--CGRKFARSDEKRHAKIH]
EGR1_HUMAN	RP	[YACPVESCDRRFSRSDELTRHIRIH]	TGQKP	[FQCRICMRNFSRSDHLTTHIRTH]	TGEKP	[FACDI--CGRKFARSDEKRHTKIH]
EGR1_MOUSE	RP	[YACPVESCDRRFSRSDELTRHIRIH]	TGQKP	[FQCRICMRNFSRSDHLTTHIRTH]	TGEKP	[FACDI--CGRKFARSDEKRHTKIH]
EGR1_RAT	RP	[YACPVESCDRRFSRSDELTRHIRIH]	TGQKP	[FQCRICMRNFSRSDHLTTHIRTH]	TGEKP	[FACDI--CGRKFARSDEKRHTKIH]
EGR1_BRARE	RP	[YACPVETCDRRFSRSDELTRHIRIH]	TGQKP	[FQCRICMRNFSRSDHLTTHIRTH]	TGEKP	[FACEI--CGRKFARSDEKRHTKIH]
EGR2_RAT	RP	[YPCPAEGCDRRFSRSDELTRHIRIH]	TGHKP	[FQCRICMRNFSRSDHLTTHIRTH]	TGEKP	[FACDY--CGRKFARSDEKRHTKIH]
EGR2_XENLA	RP	[YPCPAEGCDRRFSRSDELTRHIRIH]	TGHKP	[FQCRICMRNFSRSDHLTTHIRTH]	TGEKP	[FACDY--CGRKFARSDEKRHTKIH]
EGR2_MOUSE	RP	[YPCPAEGCDRRFSRSDELTRHIRIH]	TGHKP	[FQCRICMRNFSRSDHLTTHIRTH]	TGEKP	[FACDY--CGRKFARSDEKRHTKIH]
EGR2_HUMAN	RP	[YPCPAEGCDRRFSRSDELTRHIRIH]	TGHKP	[FQCRICMRNFSRSDHLTTHIRTH]	TGEKP	[FACDY--CGRKFARSDEKRHTKIH]
EGR2_BRARE	RP	[YPCPAEGCDRRFSRSDELTRHIRIH]	TGHKP	[FQCRICMRNFSRSDHLTTHIRTH]	TGEKP	[FACDF--CGRKFARSDEKRHTKIH]
MIG1_KLULA	--	[-----]	---RP	[YVCPICQRGFHRLEHQTRHIRTH]	TGERP	[HACDFPGCSKRFSRSDELTRHRRIH]
MIG1_KLUMA	--	[-----]	---RP	[YMCPICHRGFHRLEHQTRHIRTH]	TGERP	[HACDFPGCAKRFSRSDELTRHRRIH]
MIG1_YEAST	--	[-----]	---RP	[HACPICHRAFHRLHQTRHMRTH]	TGEKP	[HACDFPGCVKRFSRSDELTRHRRIH]
MIG2_YEAST	--	[-----]	---RP	[FRCDTCHRGFHRLEHKKRHLRTH]	TGEKP	[HHCAFPGCGKSFSRSDELKRHMRTH]
		[	:*	[. * * * * : * * * * ]		**** *
						***** .** :

<https://www.cs.cmu.edu/~02710/Lectures/SeqAlign2015.pdf>



# Local alignment

<Insert, Delete, Match>

- For sequences  $X = \langle GATTACA \rangle$  and  $Y = \langle GCATGCG \rangle$ , find the best alignment  $Z$
- Scoring scheme

$$c[i, j] = \max \begin{cases} c[i-1, j-1] + S(i, j) \\ c[i, j-1] + P \\ c[i-1, j] + P \end{cases}$$

A small change, but huge difference  
 → 문자면이 바뀌면 +1

$$S(i, j) = \begin{cases} 1 & x_i = y_j \\ 0 & \text{else} \end{cases}$$

$$P = -1$$

Global alignment

vs

$$c[i, j] = \max \begin{cases} c[i-1, j-1] + S(i, j) \\ c[i, j-1] + P \\ c[i-1, j] + P \end{cases}$$

(Arrows: red arrow pointing up-right, blue arrow pointing left, green arrow pointing up)



# Local alignment – Procedure

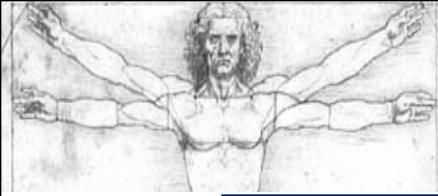
- For sequences  $X = \langle GATTACA \rangle$  and  $Y = \langle GCATGCG \rangle$ , find the best alignment  $Z$
1. Initialize
    - $c[i, 0] = 0$  for  $\forall i$
    - $c[0, j] = 0$  for  $\forall j$
  2. Compute  $c[i, j]$ , for  $\forall i, j$
  3. Terminate backtrack (two options)
    - For both options, stop backtrack at a cell where  $c[i, k] = 0$  and return alignment
    - Option 1) Find the **best alignment** and trace back OR
      - start backtrack at  $\max(c[i, j])$
    - Option 2) Find optimal or suboptimal alignments  $>$  threshold score and trace back to return all **good alignments**
      - Start backtrack at each  $(i, j)$  where  $c[i, j] > \text{threshold}$



# Local alignment – Procedure

- For each computation of  $c[i, j]$  we fill the backtrack matrix in parallel
- Here, STOP means to 1) stop the current backtrack and terminate OR 2) continue to find other alignments  $>$  threshold score
- For option 1), the backtrack start at  $\max(c[i, j])$  (from right to left)
- For option 2), the backtrack starts at every  $(i, j)$  where  $c[i, j] > \text{threshold}$

$$c[i, j] = \max \begin{cases} 0 & \text{STOP} \\ c[i-1, j-1] + S(i, j) & \nearrow \\ c[i, j-1] + P & \leftarrow \\ c[i-1, j] + P & \uparrow \end{cases}$$



# Local alignment – Procedure

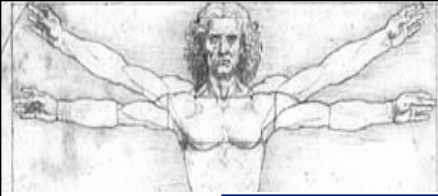
$$c[i, j] = \max \begin{cases} 0 \\ c[i-1, j-1] + S(i, j) \\ c[i, j-1] + P \\ c[i-1, j] + P \end{cases}$$

Match score = 10  
Gap penalty = -15

## ■ The score matrix $C$

		H	E	A	G	A	W	G	H	E	E	V	V	H	E	A
P	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	10	0	10	0	0	0	0	0	0	0	0	0	10
W	0	0	0	0	0	0	20	5	0	0	0	0	0	0	0	0
H	0	10	0	0	0	0	5	10	15	0	0	0	0	10	0	0
E	0	0	20	5	0	0	0	0	0	25	10	0	0	0	20	5
A	0	0	5	30	15	10	0	0	0	10	15	0	0	0	5	30
E	0	0	10	15	20	5	0	0	0	10	20	5	0	0	10	15

A W G  
 A W H



# Local alignment (option 1)

Score matrix

		H	E	A	G	A	W	G	H	E	E	V	V	H	E	A
		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
P		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
A		0	0	0	10	0	10	0	0	0	0	0	0	0	0	10
W		0	0	0	0	0	20	5	0	0	0	0	0	0	0	0
H		0	10	0	0	0	0	5	10	15	0	0	0	10	0	0
E		0	0	20	5	0	0	0	0	25	10	0	0	0	20	5
A		0	0	5	30	15	10	0	0	10	15	0	0	0	5	30
E		0	0	10	15	20	5	0	0	10	20	5	0	0	10	15

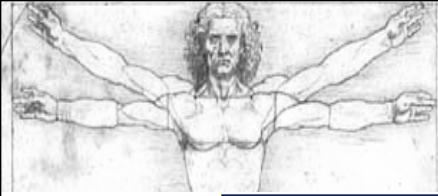
$$\max(c[i,j]) = \{c[6,3], c[6,15]\} = 30$$

Backtrack matrix

		H	E	A	G	A	W	G	H	E	E	V	V	H	E	A
		-	-	-	-	-	-	-	-	-	-	-	-	-	-	-
P		-	-	-	-	-	-	-	-	-	-	-	-	-	-	-
A		-	-	-	↖	-	↖	-	-	-	-	-	-	-	-	↖
W		-	-	-	-	↖	↖	←	-	-	-	-	-	-	-	-
H		-	↖	-	-	-	↑	↖	↖	←	-	-	-	↖	-	-
E		-	-	↖	←	-	-	-	↑	↖	←	-	-	-	↖	←
A		-	-	↑	↖	←	↖	-	-	↑	↖	←	-	-	↑	↖
E		-	-	↖	↑	↖	←	↖	-	↖	↖	←	-	-	↖	↑

A W G  
A W -

A W G  
A W H



# Local alignment (option 1)

Score matrix

		H	E	A	G	A	W	G	H	E	E	V	V	H	E	A
	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
P	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
A	0	0	0	10	0	10	0	0	0	0	0	0	0	0	0	10
W	0	0	0	0	0	0	20	5	0	0	0	0	0	0	0	0
H	0	10	0	0	0	0	5	10	15	0	0	0	0	10	0	0
E	0	0	20	5	0	0	0	0	0	25	10	0	0	0	20	5
A	0	0	5	30	15	10	0	0	0	10	15	0	0	0	5	30
E	0	0	10	15	20	5	0	0	0	10	20	5	0	0	10	15

$\max(c[i,j])=c[6,15]$

Backtrack matrix

		H	E	A	G	A	W	G	H	E	E	V	V	H	E	A
	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-
P	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-
A	-	-	-	↖	-	↖	-	-	-	-	-	-	-	-	-	↖
W	-	-	-	-	↖	-	↖	←	-	-	-	-	-	-	-	-
H	-	↖	-	-	-	-	↑	↖	↖	←	-	-	-	↖	-	-
E	-	-	↖	←	-	-	-	-	↑	↖	←	-	-	-	↖	←
A	-	-	↑	↖	←	↖	-	-	-	↑	↖	←	-	-	↑	↖
E	-	-	↖	↑	↖	←	↖	-	-	↖	↖	←	-	-	↖	↑

Alignment=

HEA
HEA

Alignment score=30





# Local alignment (option 2)

Score matrix

		H	E	A	G	A	W	G	H	E	E	V	V	H	E	A
		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
P		0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
A		0	0	0	10	0	10	0	0	0	0	0	0	0	0	10
W		0	0	0	0	0	20	5	0	0	0	0	0	0	0	0
H		0	10	0	0	0	5	10	15	0	0	0	0	10	0	0
E		0	0	20	5	0	0	0	0	25	10	0	0	0	20	5
A		0	0	5	30	15	10	0	0	10	15	0	0	0	5	30
E		0	0	10	15	20	5	0	0	10	20	5	0	0	10	15

Threshold=20

Candidates={c[6,15], c[5,9], c[6,3]}

3개

Backtrack matrix

		H	E	A	G	A	W	G	H	E	E	V	V	H	E	A
		-	-	-	-	-	-	-	-	-	-	-	-	-	-	-
P		-	-	-	-	-	-	-	-	-	-	-	-	-	-	-
A		-	-	-	↖	-	↖	-	-	-	-	-	-	-	-	↖
W		-	-	-	↖	-	↖	-	-	-	-	-	-	-	-	-
H		-	↖	-	-	-	↑	↖	-	-	-	-	-	↖	-	-
E		-	↖	↖	↖	-	-	-	↖	-	-	-	-	↖	↖	↖
A		-	-	↑	↖	↖	-	-	-	↑	↖	↖	-	-	↑	↖
E		-	-	↑	↖	↖	↖	-	-	↖	↖	↖	-	-	↖	↑

Alignment 1= HEA ||| (30)  
HEA

Alignment 2= EH-WA || || (25)  
EHGWA

Alignment 3= HEA ||| (30)  
HEA

HEA  
|||  
HEA (30)

AW-HE  
AWGHE

HEA  
|||  
HEA (30)



# Global vs Local alignment

Global alignment  
backtrack matrix

		H	E	A	G	A	W	G	H	E	E	V	V	H	E	A
	-	←	←	←	←	←	←	←	←	←	←	←	←	←	←	←
P	↑	↖	←	←	←	←	←	←	←	←	←	←	←	←	←	←
A	↑	↑	↖	↖	←	←	←	←	←	←	←	←	←	←	←	←
W	↑	↑	↑	↑	↖	←	↖	←	←	←	←	←	←	←	←	←
H	↑	↖	←	↑	↑	↖	↑	↖	↖	←	←	←	←	←	←	←
E	↑	↑	↖	←	←	←	↑	↑	↑	↖	←	←	←	←	←	←
A	↑	↑	↑	↖	←	←	←	←	↑	↑	↖	←	←	←	←	↖
E	↑	↑	↑	↑	↖	←	←	←	←	↑	↖	←	←	←	↖	←

$$c[i,j] = \max \begin{cases} c[i-1,j-1] + S(i,j) \\ c[i,j-1] + P \\ c[i-1,j] + P \end{cases}$$

P-A--W-HEA---E-  
.-|--| -|| .---| -  
HEAGAWGHEEVVHEA

Local alignment  
backtrack matrix

		H	E	A	G	A	W	G	H	E	E	V	V	H	E	A
	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-
P	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-	-
A	-	-	-	↖	-	↖	-	-	-	-	-	-	-	-	-	↖
W	-	-	-	-	↖	-	↖	←	-	-	-	-	-	-	-	-
H	-	↖	-	-	-	-	↑	↖	↖	←	-	-	-	↖	-	-
E	-	-	↖	←	-	-	-	-	↑	↖	←	-	-	-	↖	←
A	-	-	↑	↖	←	↖	-	-	-	↑	↖	←	-	-	↑	↖
E	-	-	↖	↑	↖	←	↖	-	-	↖	↖	←	-	-	↖	↑

$$c[i,j] = \max \begin{cases} 0 \\ c[i-1,j-1] + S(i,j) \\ c[i,j-1] + P \\ c[i-1,j] + P \end{cases}$$

HEA  
|||  
HEA

13      2      2  
~~35~~   ~~79~~      → 가려 줄 수 없으면 ①

0	<u>1</u>	<u>2</u>	<u>3</u>	<u>4</u>	<u>5</u>
0	0	1	<u>2</u>	2	